# **Tutorial 7 (extra)**

#### Exercise 1 (highly recommended)

Consider the following claim that some of you used during the second test.

Claim: If a language is co-recognizable then it cannot be decidable.

Does this claim hold or not? Give arguments for your answer.

## **Exercise 2 (for further practice on the simplest reductions)**

- 1. Prove that  $HALT_{TM} \leq_m A_{TM}$ . (Note that your task is to find a mapping reduction in the opposite direction than the one provided in the book in Example 5.24 on page 236).
- 2. Prove that the language

$$EPSILON_{TM} \stackrel{\text{def}}{=} \{ \langle M \rangle \mid M \text{ is a TM and } \epsilon \in L(M) \}$$

is undecidable. Is  $EPSILON_{TM}$  recognizable? Is  $EPSILON_{TM}$  co-recognizable?

3. Prove that  $E_{TM}$  is undecidable. First, define the problem and then provide either the standard reduction or mapping reduction from a suitable undecidable problem.

### **Exercise 3 (for even further practice)**

Prove that the problem

$$INFINITE_{TM} \stackrel{\text{def}}{=} \{ \langle M \rangle \mid M \text{ is a TM and } L(M) \text{ contains infinitely many strings } \}$$

is undecidable.

### Exercise 4 (only for "feinschmeckers")

The credit for this exercise goes to Morten Dahl and Morten Kühnrich.

Assume that  $L \subseteq \{0,1\}^*$  is an undecidable language. Prove that  $L' \stackrel{\text{def}}{=} L \cup F$  remains undecidable for any finite language  $F \subseteq \{0,1\}^*$ . Is this the case also if we allow F to be an infinite language?

**Hint:** Prove the undecidability claim by reduction from L to L'. The details of the proof are rather delicate and in some sense a part of the proof is non-constructive.