

## Solutions Networking 2

R3

3. The process which initiates the communication is the client; the process that waits to be contacted is the server.

R12

12. As the HTTP protocol is stateless, the website needs to explicitly do something to identify users. One way is to use cookies as explained in Section 2.2.4. Cookies allow websites to track users across different sessions (i.e., when the browser is closed and re-opened). If it is enough to keep track of a user within a single session, the website can generate a unique ID at the server side and craft webpages so that such an ID is appended to all subsequent requests (for instance, at the end of the URLs).

Moreover, from [DF] you studied already `localStorage` / `sessionStorage`. Other options can be found for example here:

<https://stackoverflow.com/questions/95655/comparison-of-ways-to-maintain-state>

R20

20. A recursive DNS query allows the server to issue additional queries to other DNS servers before replying to the client. In contrast, an iterative DNS query allows the server to reply with a referral to another DNS server. The client has to explicitly query the DNS server in the referral to proceed with name resolution.

P5

### Problem 5

- a) The status code of 200 and the phrase OK indicate that the server was able to locate the document successfully. The reply was provided on Tuesday, 07 Mar 2008 12:39:45 Greenwich Mean Time.
- b) The document `index.html` was last modified on Saturday 10 Dec 2005 18:27:46 GMT.
- c) There are 3874 bytes in the document being returned.
- d) The first five bytes of the returned document are : `<!doc`. The server agreed to a persistent connection, as indicated by the `Connection: Keep-Alive` field

P10 (Bog løsninger forkerte, se korrigerede svar nedenfor)

### Problem 10

The total download time is:

- a.  $2 \cdot 100 \text{ ms} + 8 \cdot 10^3 \text{ bits} / 10^6 \text{ bits/s} + 5 \cdot (2 \cdot 100 \text{ ms} + 4 \cdot 10^5 \text{ bits} / 10^6 \text{ bits/s}) = 3.208 \text{ s}$
- b.  $2 \cdot 100 \text{ ms} + 3 \cdot (4 \cdot 10^5 \text{ bits} / 10^6 \text{ bits/s}) = 1.4 \text{ s}$
- c.  $2 \cdot 100 \text{ ms} + 4 \cdot 10^5 \text{ bits} / 10^6 \text{ bits/s} = 0.6 \text{ s}$
- d.  $2 \cdot 100 \text{ ms} + 8 \cdot 10^3 \text{ bits} / 10^6 \text{ bits/s} + 5 \cdot (4 \cdot 10^5 \text{ bits} / 10^6 \text{ bits/s}) = 2.208 \text{ s}$

Given:  $RTT=0.1\text{s}$ ,  $\text{transferrate}=10^6 \text{ bits/s}$ ,  $\text{main document}=8 \cdot 10^3 \text{ bits}$ ,  $\text{image} = 8 \cdot 50 \cdot 10^3 \text{ bits} = 4 \cdot 10^5 \text{ bits}$

- $\text{Transfertime}(1\text{k byte}) = 8 \cdot 10^3 \text{ bits} / 10^6 \text{ b/s} = 0.008 \text{ sek}$
- $\text{Transfertime}(50\text{k byte}) = 0.4 \text{ sek}$

a) We need to download the main-page first:  $(2 \cdot RTT + \text{transfertime}(1\text{kbytes})) = 0.208$   
 $+ 5 \cdot (2 \cdot RTT + \text{transfertime}(50\text{k bytes})) = 5 \cdot (0.2 + 0.4) = 3 \text{ sek}$

**alt 3.208**

b) 2 parallel connections

We need to download the main-page first:  $(2 \cdot RTT + \text{transfertime}(1\text{kbytes})) = 0.208$   
+  
5 images on 2 parallel connections require 3 series of https (img1+2, img3+4, img5):

$2 \cdot (2 \cdot RTT + \text{transfertime}(2 \cdot 50\text{kbytes})) = 2 \cdot (0.2 + 0.8) = 2 \text{ sek}$   
+  
 $1 \cdot (2 \cdot RTT + \text{transfertime}(1 \cdot 50\text{kbytes})) = 0.2 + 0.4 = 0.6 \text{ sek}$

**alt 2.808**

C) 6 parallel connections

We need to download the main-page first:  $(2 \cdot RTT + \text{transfertime}(1\text{kbytes})) = 0.208$   
+  
 $2 \cdot RTT + 5 \cdot \text{transfertime}(50 \text{ kbyte}) = 0.2 + 2 \text{ sek}$

**alt: 2.408**

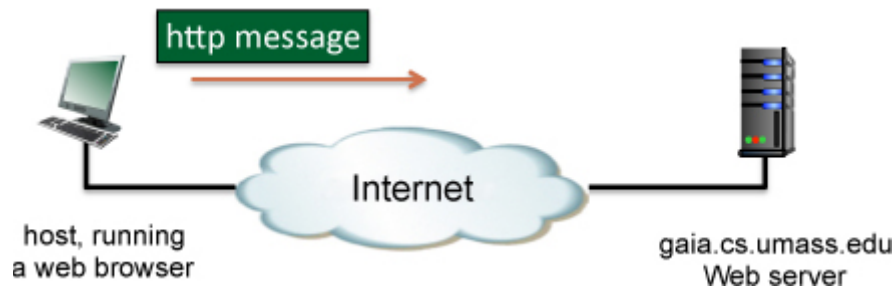
D) One persistent connection with pipelining

We need to download the main-page first:  $(2 \cdot RTT + \text{transfertime}(1\text{kbytes})) = 0.208$   
+  $RTT$  (5 requests+responses handled as 1) +  $5 \cdot \text{transfertime}(50 \text{ kbyte}) = 0.1 + 2 \text{ sek}$

**alt 2.308**

## The HTTP GET message

Consider the figure below, where a client is sending an HTTP GET message to a web server, gaia.cs.umass.edu.



Suppose the client-to-server HTTP GET message is the following:

```
GET /kurose_ross/interactive/quotation6.htm HTTP/1.0
Host: gaia.cs.umass.edu
If-Modified-Since: Mon, 23 Mar 2020 00:23:34 -0700
```

Answer the following questions:

- What is the name of the file that is being retrieved in this GET message?
- What version of HTTP is the client running?
- Does the client already have a (possibly out-of-date) copy of the requested file? Explain. If so, approximately how long ago did the client receive the file, assuming the GET request has just been issued?

### Solution:

The file being fetched is /kurose\_ross/interactive/quotation6.htm. The specific file name is quotation7.htm.

- The client is running HTTP version 1.0.
- The time indicated in the browser's If-Modified-Since header field is approximately 30 minutes ago, indicating that it has a cached copy. Therefore the server will only send a copy of the requested URL in response to this HTTP GET message if the server-side copy has been changed in the last 30 minutes.

[http://gaia.cs.umass.edu/kurose\\_ross/interactive/DNS\\_HTTP\\_delay.php](http://gaia.cs.umass.edu/kurose_ross/interactive/DNS_HTTP_delay.php)

Løsning online

## Problem 22

For calculating the minimum distribution time for client-server distribution, we use the following formula:

$$D_{cs} = \max \{NF/u_s, F/d_{min}\}$$

Similarly, for calculating the minimum distribution time for P2P distribution, we use the following formula:

$$D_{P2P} = \max \{F/u_s, F/d_{min}, NF/(u_s + \sum_{i=1}^N u_i)\}$$

Where,  $F = 15 \text{ Gbits} = 15 * 1024 \text{ Mbits}$

$$u_s = 30 \text{ Mbps}$$

$$d_{min} = d_i = 2 \text{ Mbps}$$

**Note, 300Kbps = 300/1024 Mbps.**

### Client Server

		N		
		10	100	1000
u	300 Kbps	7680	51200	512000
	700 Kbps	7680	51200	512000
	2 Mbps	7680	51200	512000

### Peer to Peer

		N		
		10	100	1000
u	300 Kbps	7680	25904	47559
	700 Kbps	7680	15616	21525
	2 Mbps	7680	7680	7680

## Practice

Eksempel (se også slides/video)

The image shows a Wireshark packet capture analysis of a DNS query. The top pane displays a list of packets, with packet 569 selected. The middle pane shows the details of packet 569, which is a DNS Standard query for stanford.edu. The bottom pane shows the raw packet data in hexadecimal and ASCII.

**Packet List:**

No.	Time	Source	Destination	Protocol	Length	Info
270	8.569482	192.168.0.1	192.168.0.181	DNS	284	Standard query response 0x44f1 A aauk-my.sharepoint.com CNAME aa
569	12.686475	192.168.0.181	192.168.0.1	DNS	72	Standard query 0x9437 A stanford.edu
570	12.686487	192.168.0.181	192.168.0.1	DNS	72	Standard query 0x9437 A stanford.edu
579	12.710650	192.168.0.1	192.168.0.181	DNS	88	Standard query response 0x9437 A stanford.edu A 171.67.215.200
596	13.037182	192.168.0.181	192.168.0.1	DNS	76	Standard query 0x598e A www.stanford.edu

**Packet Details (Packet 569):**

- Frame 569: 72 bytes on wire (576 bits), 72 bytes captured (576 bits) on interface \Device\NPF\_{D5EB09A5-5960-4374-81C3-5C4316773839}, id 0
- Ethernet II, Src: ASUSTekC\_19:22:8a (54:04:a6:19:22:8a), Dst: Icotera\_88:3a:53 (00:1e:80:88:3a:53)
- Internet Protocol Version 4, Src: 192.168.0.181, Dst: 192.168.0.1
- User Datagram Protocol, Src Port: 51815, Dst Port: 53
- Domain Name System (query)
  - Transaction ID: 0x9437
  - Flags: 0x0100 Standard query
    - 0... .. = Response: Message is a query
    - .000 0... .. = Opcode: Standard query (0)
    - ... ..0. .... = Truncated: Message is not truncated
    - ... ..1. .... = Recursion desired: Do query recursively
    - ... ..0.. .... = Z: reserved (0)
    - ... ..0 .... = Non-authenticated data: Unacceptable
  - Questions: 1
  - Answer RRs: 0
  - Authority RRs: 0
  - Additional RRs: 0
  - Queries
    - > stanford.edu: type A, class IN

**Raw Packet Data (Hex/ASCII):**

```
0000 00 1e 80 88 3a 53 54 04 a6 19 22 8a 08 00 45 00  ....:ST.  ."...E.
0010 00 3a 01 54 00 00 80 11 00 00 c0 a8 00 b5 c0 a8  .:.T.....
0020 00 01 ca 67 00 35 00 26 82 3e 94 37 01 00 00 01  ..g.5-& .->7....
0030 00 00 00 00 00 00 08 73 74 61 6e 66 6f 72 64 03  ....s tanford.
0040 65 64 75 00 00 01 00 01                edu.....
```

\*Ethernet

File Edit View Go Capture Analyze Statistics Telephony Wireless Tools Help

Apply a display filter ... <Ctrl-/>

No.	Time	Source	Destination	Protocol	Length	Info
270	8.569482	192.168.0.1	192.168.0.181	DNS	284	Standard query response 0x44f1 A aauk-my.sharepoint.com CNAME aai
569	12.686475	192.168.0.181	192.168.0.1	DNS	72	Standard query 0x9437 A stanford.edu
570	12.686487	192.168.0.181	192.168.0.1	DNS	72	Standard query 0x9437 A stanford.edu
579	12.710650	192.168.0.1	192.168.0.181	DNS	88	Standard query response 0x9437 A stanford.edu A 171.67.215.200
596	13.037182	192.168.0.181	192.168.0.1	DNS	76	Standard query 0x598e A www.stanford.edu

> Frame 579: 88 bytes on wire (704 bits), 88 bytes captured (704 bits) on interface \Device\NPF\_{D5EB09A5-5960-4374-81C3-5C4316773839}, id 0

> Ethernet II, Src: Icotera\_88:3a:53 (00:1e:80:88:3a:53), Dst: ASUSTekC\_19:22:8a (54:04:a6:19:22:8a)

> Internet Protocol Version 4, Src: 192.168.0.1, Dst: 192.168.0.181

> User Datagram Protocol, Src Port: 53, Dst Port: 51815

▼ Domain Name System (response)

Transaction ID: 0x9437

▼ Flags: 0x8180 Standard query response, No error

1... .. = Response: Message is a response

.000 0... .. = Opcode: Standard query (0)

.... .0... .. = Authoritative: Server is not an authority for domain

.... .0... .. = Truncated: Message is not truncated

.... .1... .. = Recursion desired: Do query recursively

.... .1... .. = Recursion available: Server can do recursive queries

.... .0... .. = Z: reserved (0)

.... .0... .. = Answer authenticated: Answer/authority portion was not authenticated by the server

.... .0... .. = Non-authenticated data: Unacceptable

.... .0000 = Reply code: No error (0)

Questions: 1

Answer RRs: 1

Authority RRs: 0

Additional RRs: 0

▼ Queries

> stanford.edu: type A, class IN

▼ Answers

> stanford.edu: type A, class IN, addr 171.67.215.200

[Request In: 569]

[Time: 0.024175000 seconds]

```

0000  54 04 a6 19 22 8a 00 1e 80 88 3a 53 08 00 45 00  T...". .:S..E.
0010  00 4a 00 00 40 00 40 11 b8 9c c0 a8 00 01 c0 a8  .J..@. ....
0020  00 b5 00 35 ca 67 00 36 b9 e5 94 37 81 80 00 01  ...S.g.6 ...7...
0030  00 01 00 00 00 00 08 73 74 61 6e 66 6f 72 64 03  ....s tanford.
0040  65 64 75 00 00 01 00 01 c0 0c 00 01 00 01 00 00  edu.....
0050  07 08 00 04 ab 43 d7 c8                ....C..

```