# Syntax and Semantics: Exercise Session 7

**Recall that.** A CFG is in Chomsky normal form if every production is of the following form

$$A \to BC$$
 or  $A \to a$ 

where A, B, C are non-terminals,  $a \in \Sigma$  and B, C are not the initial non-terminal. In addition  $S \to \varepsilon$  is permitted only if S is the initial non-terminal.

**Recall that.** A DFA  $(Q, \Sigma, \gamma, q_0, F)$  can be encoded to a PDA  $(Q, \Sigma, \Gamma, \delta, q_0, F)$  with empty stack alphabet (i.e.,  $\Gamma = \emptyset$ ) and transition function  $\delta \colon Q \times \Sigma_{\varepsilon} \times \Gamma_{\varepsilon} \to \wp(Q \times \Gamma_{\varepsilon})$  defined, for arbitrary  $s \in Q$  and  $\sigma \in \Sigma_{\varepsilon}$  as

$$\delta(q, \sigma, \varepsilon) = \{ (q', \varepsilon) \mid q' \in \gamma(q, \sigma) \}.$$

This is an alternative way to prove that context-free languages are a superset of the regular languages (i.e., any regular language is a context-free language). It is worth to recall that the converse inclusion does not hold.

#### Exercise 1.

For each of the CFGs below find an equivalent CFG in Chomsky normal form. The grammar  $G_1$  produces mathematical expressions with the alphabet  $\Sigma = \{a, +, \times, (,)\}.$ 

$$G_1 \colon E \to E + T \mid T$$
  $G_2 \colon R \to XRX \mid S$   $T \to T \times F \mid F$   $S \to aTb \mid bTa$   $T \to XTX \mid X \mid \varepsilon$   $X \to a \mid b \mid \varepsilon$ 

## Exercise 2.

Provide an equivalent PDA for the languages generated by the grammars  $G_1$  and  $G_2$  from Exercise 1

#### Exercise 3.

Construct a PDA for each of the following languages.

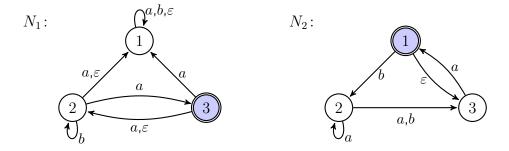
 $L_1 = \{ w \in \{0,1\}^* \mid w \text{ contains at least three 1s} \}$ 

 $L_2 = \{w \in \{0,1\}^* \mid w \text{ starts and ends with the same symbol}\}$ 

 $L_3 = \{w \in \{0,1\}^* \mid |w| \text{ is odd and 0 is its middle symbol}\}$ 

## Exercise 4.

Construct an equivalent PDA for each of the following NFAs.



# Exercise 5.

Give context-free grammars in Chomsky normal form for the following languages

$$L_4 = \{w \in \{a, b\}^* \mid w \text{ has more } a\text{'s than } b\text{'s}\}$$
  
 $L_5 = \{w \# x \in \{0, 1, \#\}^* \mid w, x \in \{0, 1\}^* \text{ and } w^R \text{ is a prefix of } x\}$