

# Syntax and Semantics:

## Exercise Session 13

### Exercise 1.

Consider the following grammar for arithmetic expressions  $\mathbb{A}$ :

$$n ::= 0 \mid s(n) \qquad a ::= n \mid a + a \mid a * a$$

Intuitively,  $s(n)$  is the successor of  $n$ , e.g.,  $s(0)$  and  $s(s(0))$  encodes 1 and 2, respectively.

- 1) Let the set of states be  $\mathbb{A}$ , the set of final states be  $\{0, s(0), s(s(0)), s(s(s(0))), \dots\}$  and consider the following *small step* transition rules for addition:

$$\begin{array}{ll} [r_1] \quad s(n_1) + n_2 \Rightarrow n_1 + s(n_2) & [r_2] \quad 0 + a \Rightarrow a \\ [r_3] \quad a_1 + a_2 \Rightarrow a_2 + a_1 & [r_4] \quad \frac{a_1 \Rightarrow a'_1}{a_1 + a_2 \Rightarrow a'_1 + a_2} \end{array}$$

Prove  $s(s(s(0))) + s(s(0)) \Rightarrow^* s(s(s(s(s(0)))))$  using  $[r_1], \dots, [r_4]$ . All intermediate computations have to be proven. (Intuitively, this corresponds to deducing  $3+2 = 5$ .)

- 2) Complete the *small step* semantics from 1) by providing the rules for multiplication.

Note: The complete semantics of  $\mathbb{A}$  is intended to be the same as for arithmetic expressions in the language **Bims** from Hans Hüttel's book but with the modified representation of numbers as explained above.

- 3) Using the transition rules from 1) and 2), prove  $s(s(0)) * s(s(0)) \Rightarrow^* s(s(s(s(0))))$ . All intermediate computations have to be proven.

**Exercise 2.**

Consider the arithmetic and boolean expressions from Chapter 4 in Hans Hüttel's book that have been extended by the existential ( $\exists$ ) and the universal ( $\forall$ ) quantifier:

$$a ::= n \mid x \mid a_1 + a_2 \mid a_1 * a_2 \mid a_1 - a_2 \mid (a_1)$$

$$b ::= a_1 = a_2 \mid a_1 < a_2 \mid \neg b_1 \mid b_1 \wedge b_2 \mid b_1 \rightarrow b_2 \mid (b_1) \mid \exists x.(b_1) \mid \forall x.(b_1)$$

1. Consider the following *big step* transition rules for the existential quantifier:

$$[\exists_1] \frac{s[x \mapsto v] \vdash b_1 \rightarrow_b \mathbf{tt}}{s \vdash \exists x.(b_1) \rightarrow_b \mathbf{tt}} \text{ FOR SOME } v \in \mathbb{Z}$$

$$[\exists_2] \frac{s[x \mapsto v] \vdash b_1 \rightarrow_b \mathbf{ff}}{s \vdash \exists x.(b_1) \rightarrow_b \mathbf{ff}} \text{ FOR ALL } v \in \mathbb{Z}$$

By combining the above rules with those from the book of Hans Hüttel, evaluate the expression  $\exists x.(x * x = \underline{4})$  in the state  $s = [x \mapsto 0]$  to its truth value.

2. Provide the *big step* transition rules for the universal quantifier.

**Exercise 3.**

Consider the language  $L'' = \{w \in \Sigma^* \mid w = a^n \text{ with } n \geq 1\}$  over the singleton alphabet  $\Sigma = \{a\}$ . With this, consider the big step transitions  $\rightarrow \subseteq L'' \times \mathbb{N}$  given by

$$[r_1] \frac{}{a \rightarrow 1} \qquad [r_2] \frac{w \rightarrow k}{aw \rightarrow k'} \quad k' = k + |aw|,$$

where  $|aw|$  denotes the length of the word  $aw$ .

1. Using the big step semantics, prove  $aa \rightarrow 2$  and  $aaa \rightarrow 6$ .
2. Using the big step semantics and induction, prove that  $a^n \rightarrow v$  with  $v = 1 \cdot 2 \cdot \dots \cdot n$  for all  $n \geq 1$ . Note that  $a^n$  is the word consisting of  $n$  symbols  $a$ .

**Exercise 4.**

Consider the language  $L'' = \{\underline{0}, \underline{1}\}^+$  of binary strings. With this, we consider the big step semantics  $\rightarrow \subseteq L'' \times \mathbb{N}$  given by

$$[r_0] \frac{}{\underline{0} \rightarrow 0}, \quad [r_1] \frac{}{\underline{1} \rightarrow 1}, \quad [r_2] \frac{w \rightarrow k'}{\underline{0}w \rightarrow k} \quad k = k', \quad [r_3] \frac{w \rightarrow k'}{\underline{1}w \rightarrow k} \quad k = 2^{|w|} + k'$$

As usual,  $|w|$  denotes the length of  $w$ , while  $\sigma w$  stands for a concatenation of  $\sigma \in \{\underline{0}, \underline{1}\}$  and  $w \in L''$ . Essentially,  $\rightarrow$  assigns each binary string its decimal value (i.e., the value represented by the string with 2 as the base of the numeral system).

- 1) Using the big step semantics, prove  $\underline{10} \rightarrow 2$  and  $\underline{110} \rightarrow 6$ .
- 2) Consider the language of ternary strings  $L''' = \{\underline{0}, \underline{1}, \underline{2}\}^+$ . Provide the big step semantics  $\rightarrow_3 \subseteq L''' \times \mathbb{N}$  that maps each ternary string to its decimal value (i.e., the value represented by the string with 3 as the base of the numeral system).