

Syntax and Semantics:

Exercise Session 6

Exercise 1.

Give context-free grammars that generate the following languages with the alphabet $\Sigma = \{0, 1\}$.

$$L_1 = \{w \in \Sigma^* \mid w \text{ contains at least three occurrences of } 1\}$$

$$L_2 = \{w \in \Sigma^* \mid w \text{ starts and ends with the same symbol}\}$$

$$L_3 = \{w \in \Sigma^* \mid \text{the length of } w \text{ is odd and its middle symbol is } 0\}$$

$$L_4 = \emptyset$$

$$L_5 = \{\varepsilon\}$$

Solution 1.

1. L_1 is recognized by a CFG with the following rules

$$S \rightarrow A1A1A1A \quad A \rightarrow \varepsilon \mid \Sigma A \quad \Sigma \rightarrow 0 \mid 1$$

2. L_2 is recognized by a CFG with the following rules

$$S \rightarrow 0 \mid 1 \mid 1A1 \mid 0A0 \quad A \rightarrow \varepsilon \mid \Sigma A \quad \Sigma \rightarrow 0 \mid 1$$

3. L_3 is recognized by a CFG with the following rules

$$S \rightarrow 0 \mid \Sigma S \Sigma \quad \Sigma \rightarrow 0 \mid 1$$

4. Remember that words recognized by a CFG have to be formed only by terminals. Therefore, L_4 can be recognized by a CFG with the following rule $S \rightarrow S$.

5. L_5 is recognized by a CFG with the following rule $S \rightarrow \varepsilon$

Exercise 2.

Consider the following CFG G with start variable A .

$$A \rightarrow X A X \mid S \quad S \rightarrow a T b \mid b T a \quad T \rightarrow X T X \mid X \mid \varepsilon \quad X \rightarrow a \mid b$$

1. Describe G formally by giving all its components.
2. Give five strings in $\mathcal{L}(G)$.
3. Give five strings not in $\mathcal{L}(G)$.
4. Which of the following derivations is allowed in G ?

$$\begin{array}{lll} T \Rightarrow aba & T \Rightarrow^* aba & XXX \Rightarrow^* aba \\ T \Rightarrow^* XX & T \Rightarrow^* XXX & S \Rightarrow^* \varepsilon \end{array}$$

Solution 2.

1. $V = \{A, X, S, T\}$; $\Sigma = \{a, b\}$; start variable A ; rules as given.
2. $\{ab, aaabb, abaa, ba, bbab\} \subseteq \mathcal{L}(G)$.
3. Let $\Sigma = \{a, b\}$. One can note that

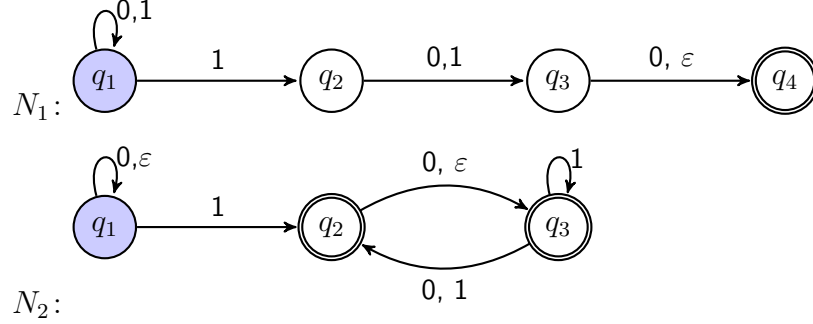
$$\mathcal{L}(G) = \{\Sigma^n x \Sigma^* y \Sigma^n \mid x \neq y, n \in \mathbb{N}\} \subseteq \Sigma^*$$

therefore, $\{\varepsilon, a, b, aa, bb\} \cap \mathcal{L}(G) = \emptyset$.

4. (a) $T \not\Rightarrow aba$
- (b) $T \Rightarrow X T X \Rightarrow a T X \Rightarrow a X X \Rightarrow a b X \Rightarrow a b a$.
- (c) $X X X \Rightarrow a X X \Rightarrow a b X \Rightarrow a b a$
- (d) $T \Rightarrow X T X \Rightarrow X X$
- (e) $T \Rightarrow X T X \Rightarrow X X X$
- (f) $S \not\Rightarrow^* \varepsilon$

Exercise 3.

Provide a CFG equivalent to the following NFAs



Solution 3.

- N_1 recognizes the language $(0 \cup 1)^* 1 (0 \cup 1) (0 \cup \varepsilon)$; a CFG for it can be

$$S \rightarrow 0S \mid 1S \mid 1A \quad A \rightarrow 0B \mid 1B \quad B \rightarrow 0 \mid \varepsilon$$

- N_2 recognizes the language $0^* 1 (0 \cup 1)^*$; a CFG for it can be

$$S \rightarrow Z1A \quad Z \rightarrow \varepsilon \mid 0Z \quad A \rightarrow \varepsilon \mid \Sigma A \quad \Sigma \rightarrow 0 \mid 1$$

Exercise 4.

Provide CFGs equivalent to each of the following regular expressions:

$$0^* 1 0^* \quad 1 \cup 0^* \emptyset^* \quad (01^+)^+$$

Solution 4.

- A CFG for $0^* 1 0^*$ can be

$$S \rightarrow Z1Z \quad Z \rightarrow \varepsilon \mid 0Z$$

- A CFG for $1 \cup 0^* \emptyset^* = 1 \cup 0^*$ can be

$$S \rightarrow 1 \mid Z \quad Z \rightarrow \varepsilon \mid 0Z$$

- A CFG for $(01^+)^+$ can be

$$S \rightarrow A \mid AS \quad A \rightarrow 0B \quad B \rightarrow 1 \mid 1B$$