Modeling caches in gem5



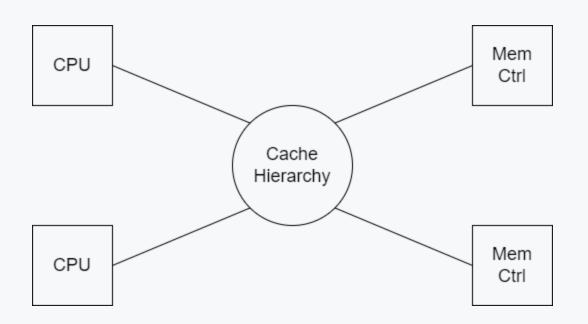
Cache Hierarchy in gem5

One of the main types of components in gem5 is the **cache hierarchy**.

In the standard library, the cache hierarchy has the Processor (with multiple cores) on one side and the Memory on the other side.

Between the cores and the caches (and the memory controllers and caches) are Ports.

Ports allow models in gem5 to send *Packets* to each other (more on this in <u>Modeling memory</u> <u>objects in gem5: Ports</u>).





Caches

Caches are used to store copies of data that are frequently accessed.

One of gem5's main goals is to model multicore systems, so caches are a key part of the system.

Much of the cache model's details and complexity come from implementing cache coherence.



Types of caches in gem5

There are two types of cache models in gem5:

- 1. Classic Cache: Simplified, faster, and less flexible
- 2. **Ruby**: Models cache coherence in detail

This is a historical quirk of the combination of *GEMS* which had Ruby and *m5* whose cache model we now call "classic" caches.

Ruby is a highly-detailed model with many different coherence protocols (specified in a language called "SLICC")

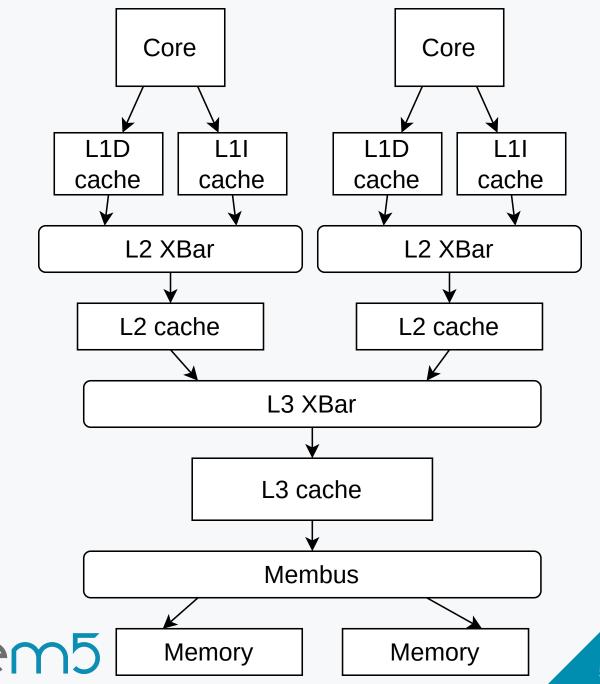
More on Ruby in Modeling Cache Coherence in gem5 in the Summer 2024 bootcamp.

Classic caches are simpler and faster, but less flexible and detailed. The coherence protocol is not parameterized and the hierarchies and topologies are fixed.



Classic Cache: Coherent Crossbar

Each crossbar can connect *n* cpu-side ports and *m* memory-side ports.



Exercise 1: Measure the cache bandwidth and latency

Use the Private L1, Private L2 cache hierarchy to measure the bandwidth and latency of the cache-based system.

Use a Dual channel DDR4 memory.
Set the cache sizes to 32KiB and 256KiB for L1 and L2 respectively.

Use the traffic generator from the previous section to generate traffic.

Questions for exercise 1

A. What is the bandwidth and latency for

- 1. The L1 cache?
- 2. The L2 cache?
- 3. The memory?

B. What is the hit/miss ratios for the cache in your tests?



Step 1: Use a classic cache hierarchy

Use the PrivateL1PrivateL2CacheHierarchy.

See PrivateL1PrivateL2CacheHierarchy for details.

Think about what the value for $[max_addr]$ should be for the generator to test L1, L2, and memory.

Think about whether you want to use the random or linear generator.



Step 1: Answer

```
board = TestBoard(
    generator=RandomGenerator(
        num_cores=1, max_addr=max_addr,
        rd_perc=75, duration="1ms"
    cache_hierarchy=PrivateL1PrivateL2CacheHierarchy(
        11d_size="32KiB",
        11i_size="32KiB",
        12_size="256KiB",
    memory=DualChannelDDR4_2400(),
    clk_freq="3GHz",
```

You want to get 100% hits in the cache when testing each level and as close to 0% when you are testing the next level or memory.

Use a max address of 16384 to test the L1 cache, 131072 to test the L2 cache, and 1048576 to test the memory.

Use the random generator to maximize the hits in the cache you are testing.

Looking at "missRate::total" will show that there's 100% hits in the level we are evaluating.



Creating custom cache hierarchies

The standard library will not have all components that you will need.

Our next goal is to take a model from gem5 and wrap it in the standard library API.

Specifically, we will create a three-level cache hierarchy with private L1s, private L2s, and a shared L3 connected to multiple memory channels.

Before we dive into that exercise, we need to cover a few gem5 topics.



Classic cache models

The classic cache consists of two main models, the Cache and the XBar.

Cache

This is an object that represents a cache in gem5. It has a number of parameters that can be set to configure the cache.

The cache has a cpu_side and a mem_side that can be connected to other components.

XBar

This is a crossbar that can connect multiple ports together. It has a cpu_side_ports and a mem_side_ports that can be connected to other components.

The CPU-side and memory-side can connect to multiple objects.



gem5 Ports

Ports

Ports are the way that gem5 models communicate with each other. They are used to send packets between models.

We will see more about ports in the <u>Modeling memory objects in gem5: Ports</u> section.

Using ports

In gem5, to connect two ports together, you just need to use the = operator.

```
cache = Cache()
xbar = XBar()

cache.mem_side = xbar.cpu_side_ports
```



Specializing gem5 models

In gem5, you can specialize a model by specifying the parameters that you want to change.

For instance, if you want to create a cache with a size of 32 KiB and an associativity of 4, you can do the following:

```
class MyCache(Cache):
    def __init__(self):
        super().__init__()
        self.size = "32KiB"
        self.assoc = 4
```

You can override any of the *parameters* in the Cache class this way.



Exercise 2: Create a three level cache hierarchy

Create a three level cache hierarchy with private L1s, private L2s, and a shared L3 connected to multiple memory channels as shown on the previous slide. Run your tests with all cores.

Note: For some of these steps, you'll just have to take our word for it. There are many things that have "legacy" reasons.

Questions for exercise 2:

What is the bandwidth and latency for

- 1. The L1 cache?
- 2. The L2 cache?
- 3. The L3 cache?
- 4. The memory?



Step 2: declare the hierarchy and add a membus

(Note: Step 1 is exercise 1.)

Open exercises/02-Using-gem5/04-cache-hierarchies/three_level.py

We are creating a new class which will implement an abstract class AbstractCacheHierarchy. By implementing this class, we can use it in the standard library as a cache hierarchy in our test board (or even in a CPU-based board).

The constructor is already provided.

Add a membus to the constructor

This will be your first time using a "raw" gem5 model (i.e., not something in the stdlib).

For this, use a SystemXBar with a width of 64.

See SystemXBar in gem5/src/mem/xbar.py.



Step 2: Answer

```
self.membus = SystemXBar(width=64)
```

This will be what connects the cache to memory.

We will make it 64 Bytes wide (as wide as the cache line) so that it's maximum bandwidth.



Step 3: Implement the hierarchy interface

The board needs to be able to get the port to connect to memory.

You will be implementing some of the abstract functions from AbstractCacheHierarchy.

See AbstractCacheHierarchy for details.

You need to implement get_mem_side_port and get_cpu_side_port in this step.



Step 3: Answer

```
def get_mem_side_port(self):
    return self.membus.mem_side_ports
```

The "cpu_side_port" is used for coherent IO access from the board.

```
def get_cpu_side_port(self):
    return self.membus.cpu_side_ports
```



Step 4: Implement incorporate_cache

The main function is <u>incorporate_cache</u> which is called by the board after the <u>Processor</u> and <u>Memory</u> are ready to be connected together.

In the [incorporate_cache] function, we will create the caches and connect them together.

First, (given below) we will connect the system port.

Then, we will connect the memory ports to the memory bus.

You can get the list of ports for memory from the MemorySystem object.

Then, we can connect all of these ports to the memory side of memory bus.

Then, we will create the L3 crossbar based on the <u>L2 crossbar</u> parameters.

```
def incorporate_cache(self, board):
    board.connect_system_port(self.membus.cpu_side_ports)
```



Step 4: Answer

```
# Connect the memory system to the memory port on the board.
for _, port in board.get_memory().get_mem_ports():
    self.membus.mem_side_ports = port

# Create an L3 crossbar
self.13_bus = L2XBar()
```

Here, we connect all of the memory ports to the membus. Note that for gem5 ports, to connect things together, you just have to use the = operator.

We are using a "L2XBar" for the L3 crossbar because "L2XBar" is just a crossbar that has been configured to not be the system xbar.



Step 5: Create the core clusters

Let's move away from the [incorporate_cache] function for a moment and create the core clusters.

Our goal is that each core cluster should have a core connected to its private caches (L1I, L1D, L2) and the L2.

We'll then connect the L2 to the L3 via the L3 crossbar.

You'll be modifying the [_create_core_cluster] function.



Step 5: More details

A SubSystem is a special object that can hold multiple objects (with a shared clock/voltage domain).

- 1. Create a SubSystem object in the _create_core_cluster function.
- 2. Create the L1I and L1D caches as sub-objects of the cluster subsystem.
- 3. Create an L2 cache as a sub-object of the cluster subsystem.
- 4. Create an L2 crossbar (L2XBar) as a sub-object of the cluster subsystem.
- 5. Use the core.connect_icache and core.connect_dcache functions to connect the L1I and L1D caches to the core.
- 6. Connect the L1I and L1D caches to the L2 crossbar.
- 7. Connect the L2 cache to the L2 crossbar.
- 8. Connect the L2 cache to the L3 bus.

Use the L1DCache, L1ICache and L2Cache objects from the standard library.

Note: there is extra code in this function. Ignore this for now.



Step 5: Answer

Since each core is going to have many private caches, let's create a cluster. In this cluster, we will create L1I/D and L2 caches, the L2 crossbar and connect things.

```
def _create_core_cluster(self, core, 13_bus, isa):
   cluster = SubSystem()
    cluster.l1dcache = L1DCache(size=self._l1d_size, assoc=self._l1d_assoc)
    cluster.l1icache = L1ICache(
        size=self._l1i_size, assoc=self._l1i_assoc, writeback_clean=False
    core.connect_icache(cluster.l1icache.cpu_side)
    core.connect_dcache(cluster.l1dcache.cpu_side)
    cluster.12cache = L2Cache(size=self._12_size, assoc=self._12_assoc)
    cluster.12_bus = L2XBar()
    cluster.l1dcache.mem_side = cluster.l2_bus.cpu_side_ports
    cluster.l1icache.mem_side = cluster.l2_bus.cpu_side_ports
    cluster.12cache.cpu_side = cluster.12_bus.mem_side_ports
    cluster.12cache.mem_side = 13_bus.cpu_side_ports
```

Full-system specific things

You have been given some code to set up other caches, interrupts, etc. that are needed for full system simulation in x86 and Arm.

You can ignore this for now.

```
cluster.iptw_cache = MMUCache(size="8KiB", writeback_clean=False)
cluster.dptw_cache = MMUCache(size="8KiB", writeback_clean=False)
core.connect_walker_ports(
    cluster.iptw_cache.cpu_side, cluster.dptw_cache.cpu_side
)

# Connect the caches to the L2 bus
cluster.iptw_cache.mem_side = cluster.12_bus.cpu_side_ports
cluster.dptw_cache.mem_side = cluster.12_bus.cpu_side_ports
if isa == ISA.X86:
    int_req_port = self.membus.mem_side_ports
    int_resp_port = self.membus.cpu_side_ports
    core.connect_interrupt(int_req_port, int_resp_port)
else:
    core.connect_interrupt()

return cluster
```



Step 6: Create a cluster for each core

Back to [incorporate_cache]...

Now that we have the cluster, we can create the clusters.

Create a list of clusters, one for each of the cores in the processor.



Step 6: Answer

```
self.clusters = [
    self._create_core_cluster(
        core, self.13_bus, board.get_processor().get_isa()
    )
    for core in board.get_processor().get_cores()
]
```



Step 7: Create an L3 cache

The standard library has L1 and L2 caches, but not an L3 cache. So, we need to create an L3 cache.

Make the L3 cache a subclass of the [Cache] object. Allow it to have a configurable size and associativity.

For the other parameters, use 20 cycles for the tag and data latency, 1 cycle for the response latency, 20 MSHRs, 12 targets per MSHR, and set writeback_clean to False. Set the "clusivity" to "mostly_incl". In other words, this will be a mostly inclusive cache.



Step 7: Answer

```
class L3Cache(Cache):
    def __init__(self, size, assoc):
        super().__init__()
        self.size = size
        self.assoc = assoc
        self.tag_latency = 20
        self.data_latency = 20
        self.response_latency = 1
        self.mshrs = 20
        self.tgts_per_mshr = 12
        self.writeback clean = False
        self.clusivity = "mostly_incl"
```



Step 8: Connect the L3 cache

Again, back to [incorporate_cache]...

Create an instance of your L3 cache object, pass through the size and associativity, then connect the L3 cache to the L3 crossbar and the memory bus.

Also, add this code at the end of the incorporate_cache function. (Trust me.)

```
if board.has_coherent_io():
    self._setup_io_cache(board)
```



Step 8: Answer

```
self.13_cache = L3Cache(size=self._13_size, assoc=self._13_assoc)

# Connect the L3 cache to the system crossbar and L3 crossbar
self.13_cache.mem_side = self.membus.cpu_side_ports
self.13_cache.cpu_side = self.13_bus.mem_side_ports
```



Step 9: Re run the test from Exercise 1 with 3 levels

Go back to the test-cache.py file that you created and import your new cache hierarchy.

Extend this to also be able to test the performance of the L3 cache.



Answers to questions

What is the bandwidth and latency for

- 1. The L1 cache?
- 2. The L2 cache?
- 3. The L3 cache?
- 4. The memory?



Classic Cache: Parameters

- src/mem/cache/Cache.py
 - src/mem/cache/cache.cc
 - src/mem/cache/noncoherent_cache.cc

Parameters:

- size
- associativity
- number of miss status handler register (MSHR) entries
- prefetcher
- replacement policy

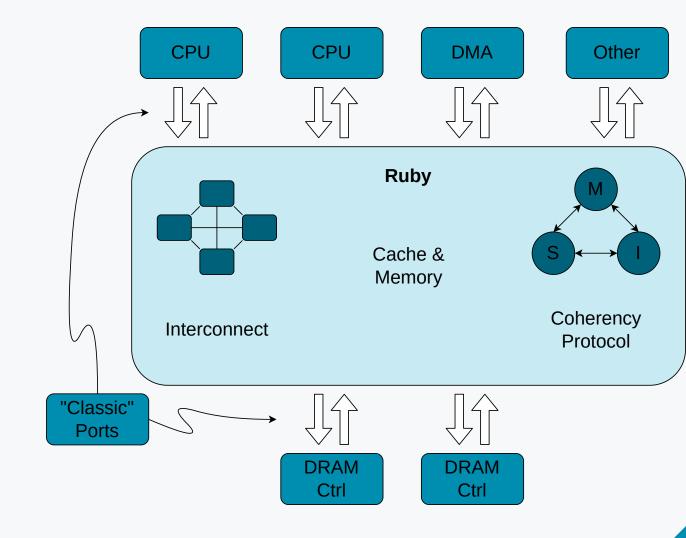


Ruby



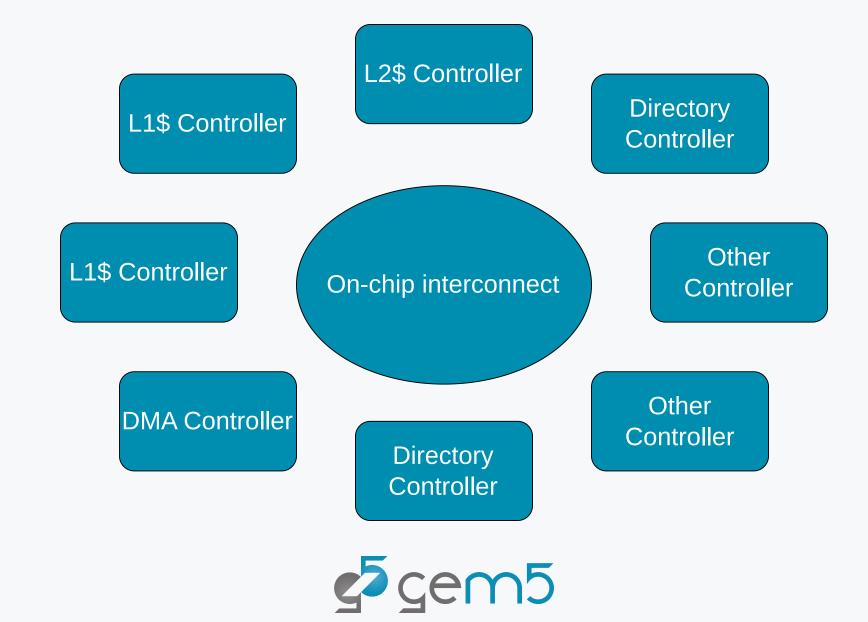
Ruby Cache

- 1. Coherence Controller
- 2. Caches + Interface
- 3. Interconnect





Ruby



Ruby Components

- Controller models (cache controller, directory controller)
- Controller topology (Mesh, all-to-all, etc.)
- Network models
- Interface (classic ports)

Ruby Cache: Controller Models

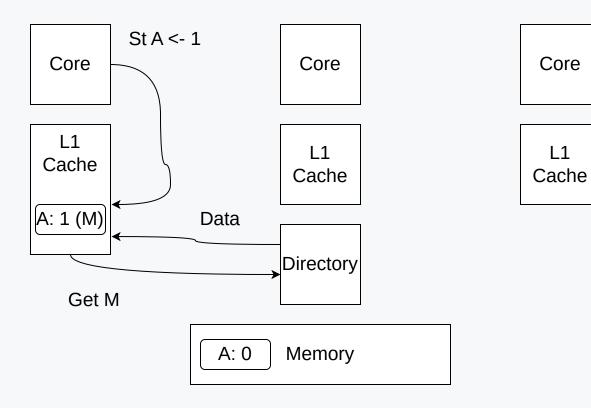
Code for controllers is "generated" via SLICC compilers

You can see much more detail in Modeling Cache Coherence in gem5 (Not covered this week).



Cache Coherence Reminder

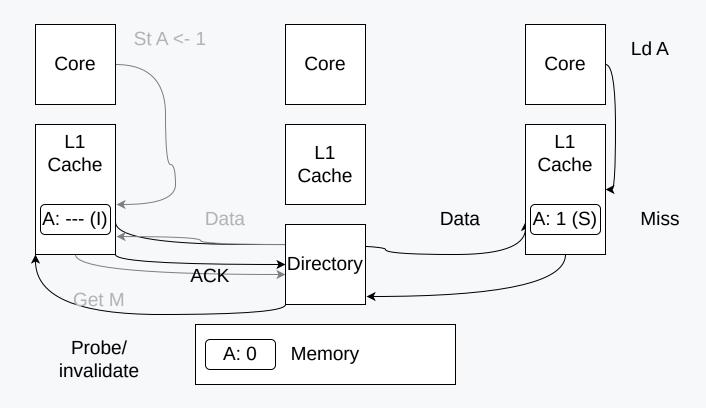
Single-Writer Multiple-Reader (SWMR) invariant





Cache Coherence Reminder

Single-Writer Multiple-Reader (SWMR) invariant





Ruby enables detailed cache coherence models

- MSI: Modified, Shared, Invalid
- From Nagarajan, Sorin, Hill, and Wood. <u>A Primer on Memory Consistency and Cache Coherence</u>.
- Excerpt of 8.2 download

TABLE 8.1: MSI Directory Protocol—Cache Controller												
	load	store	replacement	Fwd-GetS	Fwd-GetM	Inv	Put-Ack	Data from Dir (ack=0)	Data from Dir (ack>0)	Data from Owner	Inv-Ack	Last-Inv-Ack
I	send GetS to Dir/IS ^D	send GetM to Dir/IM ^{AD}										
IS^{D}	stall	stall	stall Text			stall		-/S		-/S		
IM^{AD}	stall	stall	stall	stall	stall			-/M	-/IM ^A	-/M	ack	
IM^A	stall	stall	stall	stall	stall						ack	-/M
S	hit	send GetM to Dir/SM ^{AD}	send PutS to Dir/SI ^A			send Inv-Ack to Req/I						
SM ^{AD}	hit	stall	stall	stall	stall	send Inv-Ack to Req/IM ^{AD}		-/M	-/SM ^A	-/M	ack	
SM^A	hit	stall	stall	stall	stall						ack	-/M
M	hit	hit	send PutM+data to Dir/MI ^A	send data to Req and Dir/S	send data to Req/I							
MI ^A	stall	stall	stall	send data to Req and Dir/SI ^A	send data to Req/II ^A		-/I					
SIA	stall	stall	stall			send Inv-Ack to Req/II ^A	-/I					
$\Pi^{\mathbf{A}}$	stall	stall	stall				-/I					



Events

states

SLICC Original Purpose

• Create these tables

Actual output!

	<u>Load</u>	<u>Store</u>	Replacement	<u>FwdGetS</u>	<u>FwdGetM</u>	Inv	PutAck	DataDirNoAcks	<u>DataDirAcks</u>	<u>DataOwner</u>	InvAck	<u>LastInvAck</u>	
Ī	a aT gS pQ / IS D	a aT gM pQ / IM AD											Ī
IS D	<u>z</u>	<u>z</u>	<u>z</u>			<u>z</u>		wd dT xLh pR/S		wd dT xLh pR / S			IS D
IM AD	Z	<u>z</u>	Z	<u>z</u>	<u>z</u>			wd dT xSh pR/M	wd sa pR / IM A	wd dT xSh pR/M	<u>da pR</u>		IM AD
IM A	Z	<u>z</u>	Z	<u>z</u>	<u>z</u>						<u>da pR</u>	dT xSh pR / M	IM A
<u>S</u>	<u>Lh pQ</u>	aT gM pQ/SM AD	pS/SIA			iaR d pF / I							<u>s</u>
SM AD	<u>Lh pQ</u>	<u>z</u>	Z	<u>z</u>	<u>z</u>	iaR pF / IM AD		wd dT xSh pR/M	wd sa pR / SM A	wd dT xSh pR/M	<u>da pR</u>		SM AD
SM A	<u>Lh pQ</u>	<u>z</u>	Z	<u>z</u>	<u>z</u>						<u>da pR</u>	dT xSh pR / M	SM A
<u>M</u>	<u>Lh pQ</u>	Sh pQ	pM/MIA	cdR cdD pF/S	cdR d pF/I								<u>M</u>
MI A	<u>z</u>	<u>z</u>	Z	cdR cdD pF/SIA	cdR pF/IIA		<u>d pF / I</u>						MI A
SI A	<u>z</u>	<u>z</u>	Z			<u>iaR pF</u> / <u>II A</u>	<u>d pF / I</u>						SI A
II A	<u>z</u>	<u>z</u>	<u>z</u>				<u>d pF / I</u>						II A
	<u>Load</u>	<u>Store</u>	Replacement	<u>FwdGetS</u>	<u>FwdGetM</u>	Inv	PutAck	DataDirNoAcks	DataDirAcks	<u>DataOwner</u>	InvAck	LastInvAck	



Ruby vs Classic

Ruby Caches

- Detailed cache coherence
- Any topology
- Network models
- Hierarchy fixed to protocol
 - Except CHI
- Very complex
- Slow

Classic Caches

- Easy to configure
- Fast
- Limited to a hierarchy of crossbars
- Cannot configure the protocol
- Cheats for coherence



Homework: Use a detailed coherence protocol

You will be using the Ruby cache hierarchy to model a cache coherence protocol.

You will investigate false sharing by comparing statistics with different applications.

See the homework for more details.



Summary

- Cache hierarchies are a key part of gem5
- Classic caches are simpler and faster
- Classic caches are straightforward to configure and use
- Ruby caches are more detailed and can model cache coherence
- We can use Ruby to compare different cache behaviors

