



Indices:

- $i = (1, n)$: *packing orders*
- $j = (1, m)$: *line numbers*
- $t = (1, T)$: *time slots, given in terms of a base time interval ΔT e. g. 15 minutes*

Decision Variables:

- $x(i, j, t,) :=$ *packing order i starts on line j at time t* (binary)
- $y(i, k, j) :=$ *setup between orders i and k on line j* (binary)
- $b(i, k) :=$ *batch indicator, order i and k batched together* (binary)
- $prod(i, t) :=$ *number of units produced at time t* (positive integer)
- $u(j) :=$ *line in use* (binary)

Decision Variables (OTIF tracking)

- $late(i) :=$ (binary)
- $lateness(i) :=$ (positive integer)

- $workers_{used}(t) :=$ total workers actually working at time t (integer $[0, W]$)
- $workers_{max} :=$ maximal workers used in any time slot (integer $[0, W]$)
- $workers_{min} :=$ minimum workers used in any time slot (integer $[0, W]$)

Decision Variables (WIP tracking)

- $inv(i) :=$ number of packing order i in stock (positive integer)
- $ship(i, t) :=$ packing order i is shipped at time t (binary)
- $time_{ship}(i) :=$ time slot when order i is shipped (integer $[0, T]$)
- $has_{inv}(i, t) :=$ order i has available inventory at time t (binary)
- $time_{start}(i) :=$ actual start time for packing order i (integer $[1, T]$)
- $time_{completion}(i) :=$ actual completion time for packing order i (integer $[1, T]$)
- $time_{flow}(i) :=$ time from start to completion for packing order i (non-negative integer)



Input parameters:

- $p(i, j) :=$ processing time for packing order i on line j (positive real, dividable by base time interval)
- $s(i, k, j) :=$ setup time for packing order j and k on line j (positive real, dividable by base time interval)
- $a(w, t) :=$ worker w is available at time t (binary)
- $inv_0(i) :=$ initial inventory stock for packing order i (positive integer)
- $\alpha :=$ reserved capacity (positive real less than 1)

Input parameters (OTIF):

- $due(i) :=$ due date for packing order i (integer $[1, T]$)
- $priority(i) :=$ priority weight for order i (positive integer, e.g. $[1, 100]$)

Constraints

- One assignment: Each order assigned to only one line, one time

$$\sum_j \sum_t x(i, j, t) = 1 \quad \forall i$$

- Line capacity: No overlap of order on the same line

$$\sum_i \sum_{t \leq \tau < t + p(i, j) + \sum_k y(i, k, j) * s(i, k, j) * (1 - b(i, k))} x(i, j, t) \leq u(j) \quad \forall j, \forall \tau$$

- Reserved line capacity

$$\sum_i \sum_j \sum_{\tau} \sum_{t \leq \tau < t + p(i, j) + \sum_k y(i, k, j) * s(i, k, j) * (1 - b(i, k))} x(i, j, t) < (1 - \alpha) * m * T$$



Constraints

- Line in use

$$\sum_i \sum_{t \leq \tau < t+p(i,j)} x(i,j,t) \leq u(j) \quad \forall \tau$$

- Each order ships exactly once

$$\sum_t ship(i,t) = 1 \quad \forall i$$

- Shipping time

$$time_{ship}(i) = \sum_t t * ship(i,t) \quad \forall i$$



Constraints (OTIF)

- Start time

$$\text{time}_{\text{start}}(i) = \sum_j \sum_t t * x(i, j, t) \quad \forall i$$

- Completion time

$$\text{time}_{\text{complete}}(i) = \sum_j \sum_t (t + p(i, j)) * x(i, j, t) \quad \forall i$$

- Lateness

$$\text{lateness}(i) \geq \text{time}_{\text{complete}}(i) - \text{due}(i) \quad \forall i$$

$$\text{lateness}(i) \geq 0 \quad \forall i$$

- Identify late orders (forces $\text{late}(i) = 1$ if order is late)

$$\text{lateness}(i) \leq T * \text{late}(i) \quad \forall i$$

$$\text{time}_{\text{complete}}(i) > \text{due}(i) - T(1 - \text{late}(i)) \quad \forall i$$



Constraints (Shipping)

- Cannot ship before order complete

$$time_{ship}(i) \geq time_{complete}(i) \quad \forall i$$

- Has inventory

$$inv(i, t) \geq due(i, t) - M * (1 - has_{inv}(i, t)) \quad \forall i \forall t$$

$$inv(i, t) \leq M * has_{inv}(i, t) \quad \forall i \forall t$$

- Can only ship when inventory is sufficient

$$ship(i, t) \leq has_{inv}(i, t) \quad \forall i \forall t$$



Constraints (WIP)

- Flow time

$$time_{flow}(i) = time_{ship}(i) - time_{start}(i) \quad \forall i$$


- Number of packing orders produced

$$prod(i, t) = \sum_j x(i, j, t - p(i, j))$$

- Inventory balance equation

$$inv(i, t) = inv(i, t - 1) + prod(i, t) - ship(i, t) \quad \forall i, \forall t > 0$$

$$inv(i, 0) = inv_0(i) \quad \forall i$$



Constraints (Workforce)

- Active workers per time slot

$$\text{workers}_{\text{used}}(\tau) = \sum_i \sum_j \sum_{t \leq \tau < t+p(i,j)} x(i,j,t) \quad \forall \tau$$

- Maximum workforce

$$\text{workers}_{\text{max}} \geq \text{workers}_{\text{used}}(t) \quad \forall t$$

- Minimum workforce

$$\text{workers}_{\text{used}}(t) \geq \text{workers}_{\text{min}}$$

- Workforce range

$$\text{workforce}_{\text{range}} = \text{workers}_{\text{max}} - \text{workers}_{\text{min}}$$

Objective function

Minimization objective function:

$$f = \alpha * otif + \beta * wip_{obj} + \gamma * workforce + \delta * total_{not_utilized}$$

where

- OTIF term (heavy penalty for late order but also penalizes according to lateness)

$$otif = \sum_i priority(i) * (7 * late(i) + 3 * lateness(i))$$

- WIP term

$$wip_{obj} = \sum_t inv(t)$$

- Workforce term (reduces worker variance, deviations from target workforce and changes from one time instant to another)

$$workforce = workers_{range}$$

- Number of lines fully utilized

$$total_{not_utilized} = \sum_j u(j)$$