

# Information Systems Security (SOEN321)

## A Brief introduction to number Theory

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# Why do we need to learn number theory

- ◆ Many public key systems are based on number theoretic ideas
- ◆ Examples include
  - RSA (based on the difficulty of factoring)
  - Diffie-Hellman (based on the difficulty of solving the discrete log problem)
  - El-Gamal (based on discrete log)
  - etc.
- ◆ We will not dive deep into number theory. We will only learn what is necessary to understand these systems

# Modular Arithmetic

## ◆ Definition:

$$a \bmod n = r \Leftrightarrow \exists q \text{ s.t. } a = q \times n + r, \text{ where } 0 \leq r \leq n-1$$

## ◆ Example:

- $7 \bmod 3 = 1$
- $-7 \bmod 3 = 2$

## ◆ Definition (Congruence)

$$a \equiv b \bmod n \Leftrightarrow a \bmod n = b \bmod n$$

# Groups

## ◆ Definition:

A group  $(G, *)$  is a set  $G$  on which a binary operation  $*$  is defined which satisfies the following axioms:

- Closure: For all  $a, b \in G$ ,  $a * b \in G$
- Associative: For all  $a, b, c \in G$ ,  $(a * b) * c = a * (b * c)$
- Identity:  $\exists e \in G$  s.t. for all  $a \in G$ ,  $a * e = a = e * a$
- Inverse: For all  $a \in G$ ,  $\exists a^{-1} \in G$  s. t.  $a * a^{-1} = a^{-1} * a = e$

# Example

◆  $(Z_n, +)$ , where

- $Z_n = \{0, 1, \dots, n-1\}$ ,
- $a + b = a + b \bmod n$

◆  $(Z_p^*, \times)$ , where

- $Z_p^* = \{1, \dots, p-1\}$
- $a \times b = a \times b \bmod p$

◆ Abelian Group

- A group  $(G, *)$  is called an abelian group if  $*$  is a commutative operation:
- Commutative: For all  $a, b \in G$ ,  $a * b = b * a$ .

# More explicit

## ◆ Example: $(\mathbb{Z}_7, +)$

- elements:  $\{0, 1, 2, 3, 4, 5, 6\}$
- $5 + 4 = 2$
- 0 is identity element
- every element  $x$  has an inverse  $y$  such that  $x + y = 0$ , what is the inverse of 3?
  - ◆  $3 + 4 = 0$ , thus the inverse of 3 is 4.

## ◆ Example: $(\mathbb{Z}_7, \times)$

- elements:  $\{1, 2, 3, 4, 5, 6\}$
- what is the identity element?
- what is the inverse of 3?
  - ◆  $3 \times 1 = 3, 3 \times 2 = 6, 3 \times 3 = 2, 3 \times 4 = 5, 3 \times 5 = 1, 3 \times 6 = 4$

# Fields

- ◆ a field  $F$  is a commutative ring in which all non-zero elements have multiplicative inverses
- ◆ Finite fields: if  $F$  contains a finite number of elements
  - the order must be of type  $p^m$  with  $m \geq 1$
  - for every order  $p^m$  there is a unique (up to isomorphism)  $GF(p^m)$
- ◆ Example: The field:  $F_p = (\mathbb{Z}_p, +, \times)$
- ◆ Example:  $F_7$ 
  - solve  $3x + 1 = 6$
  - $3x = 5$
  - $x = 3^{-1} \times 5 = 5 \times 5 = 4$

# Useful tools

- ◆ Extended Euclidian Algorithm
- ◆ Euclid's Theorem
- ◆ Chinese reminder theorem



# Euclidian Algorithm

```
gcd (b, a)
  d=a; t=b
  while t  $\neq$  0 {
    w  $\leftarrow$  d mod t
    d  $\leftarrow$  t
    t  $\leftarrow$  w
  }
  return d
```

# Towards Extended Euclidian Algorithm

- ◆ **Theorem:** Given integers  $a, b > 0$  and  $a > b$ , then  $d = \gcd(a, b)$  is the least positive integer that can be represented as  $ax + by$ ,  $x, y$  integer numbers.
- ◆ Corollary: if  $a$  and  $b$  are relative prime, then there exist  $x$  and  $y$  such that  $ax + by = 1$ .
- ◆ In other words,  $ax \bmod b = 1$ .
- ◆ How to find such  $x$  and  $y$ ?

# Example

Find  $\gcd(143, 111)$

$$143 = 1 \times 111 + 32$$

$$111 = 3 \times 32 + 15$$

$$32 = 2 \times 15 + 2$$

$$15 = 7 \times 2 + 1$$

$$\gcd(143, 111) = 1$$

$$32 = 143 - 1 \times 111$$

$$15 = 111 - 3 \times 32$$

$$= 4 \times 111 - 3 \times 143$$

$$2 = 32 - 2 \times 15$$

$$= 7 \times 143 - 9 \times 111$$

$$1 = 15 - 7 \times 2$$

$$= 67 \times 111 - 52 \times 143$$

# Another example

◆ Example: determine  $\gcd(803; 121)$

◆ Solution

$$803 = 121 \cdot 6 + 77$$

$$121 = 77 \cdot 1 + 44$$

$$77 = 44 \cdot 1 + 33$$

$$44 = 33 \cdot 1 + 11$$

$$33 = 11 \cdot 3 + 0$$

$$\rightarrow 11 = \gcd(803, 121)$$

# Extended Euclidian Algorithm

```
x=1; y=0; d=a; r=0; s=1; t=b;
while (t>0) {
    q =  $\lfloor d/t \rfloor$ 
    u=x-qr; v=y-qs; w=d-qt
    x=r;    y=s;    d=t
    r=u;    s=v;    t=w
}
return (d, x, y)
```

# Example

$\gcd(803; 121)$

$$803 = 121 \times 6 + 77$$

$$121 = 77 \times 1 + 44$$

$$77 = 44 \times 1 + 33$$

$$44 = 33 \times 1 + 11$$

$$33 = 11 \times 3 + 0$$

$$\rightarrow 11 = \gcd(803, 121)$$

$$11 = 44 - 1 \times 33$$

$$= 44 - 1 \times (77 - 1 \times 44) = 2 \times 44 - 1 \times 77$$

$$= 2 \times (121 - 1 \times 77) - 1 \times 77 = 2 \times 121 - 3 \times 77$$

$$= 2 \times 121 - 3 \times (803 - 6 \times 121) = 20 \times 121 - 3 \times 803$$

# Example

$$m = 841$$

$$160^{-1} \text{ in } \mathbb{Z}_{841} = ?$$

$\gcd(841, 160) = 1$ ; indeed:

$$841 = 5 \cdot 160 + 41$$

$$160 = 3 \cdot 41 + 37$$

$$41 = 1 \cdot 37 + 4$$

$$37 = 9 \cdot 4 + 1$$

$$1 = 37 - 9 \cdot 4 = 37 - 9 \cdot (41 - 1 \cdot 37)$$

$$= 10 \cdot 37 - 9 \cdot 41 = 10 \cdot (160 - 3 \cdot 41) - 9 \cdot 41$$

$$= 10 \cdot 160 - 39 \cdot 41 = 10 \cdot 160 - 39 \cdot (841 - 5 \cdot 160)$$

$$= 205 \cdot 160 - 39 \cdot 841$$

$$\leadsto 160 \cdot 205 = 1 \pmod{841} \iff 160^{-1} = 205 \pmod{841}$$

# The Euler Phi Function

## ◆ Definition

Given an integer  $n$ ,  $\Phi(n) = |Z_n^*|$  is the number of all numbers  $a$  such that  $0 < a < n$  and  $a$  is relatively prime to  $n$  (i.e.,  $\gcd(a, n) = 1$ ).

◆ **Theorem:** If  $\gcd(m, n) = 1$ ,  $\Phi(mn) = \Phi(m) \Phi(n)$

Let  $p$  be prime,  $e, m, n$  be positive integers

1)  $\Phi(p) = p - 1$

2)  $\Phi(p^e) = p^e - p^{e-1}$

3) If  $n = p_1^{e_1} p_2^{e_2} \dots p_k^{e_k}$  then

$$\Phi(n) = n \left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \dots \left(1 - \frac{1}{p_k}\right)$$



# Fermat's (Little) Theorem

Let  $p$  be a prime

For each integer  $a$  one has:

$$a^p = a \pmod{p}$$

and if  $\gcd(p, a) = 1$ :

$$a^{p-1} = 1 \pmod{p}$$

**Example:**

$p = 11$ ,  $a = 2$  compute  $2^{10} \pmod{11}$ :

$$2^4 = 16 = 5 \pmod{11}$$

$$2^8 = 5^2 = 25 = 3 \pmod{11}$$

$$2^{10} = 3 \cdot 2^2 = 12 = 1 \pmod{11}$$

# Euler's Theorem

## ◆ Euler's Theorem

◆ Given integer  $n > 1$ , such that  $\gcd(a, n) = 1$  then  $a^{\Phi(n)} \equiv 1 \pmod{n}$

## ◆ Corollary

◆ Given integer  $n > 1$ , such that  $\gcd(a, n) = 1$  then  $a^{\Phi(n)-1} \pmod{n}$  is a multiplicative inverse of  $a \pmod{n}$ .

## ◆ Corollary

- Given integer  $n > 1$ ,  $x$ ,  $y$ , and a positive integers with  $\gcd(a, n) = 1$ . If  $x \equiv y \pmod{\Phi(n)}$ , then  $a^x \equiv a^y \pmod{n}$ .

# Linear Equation Modulo $n$

$$ax \equiv b \pmod{n}$$

- ◆ If  $\gcd(a, n) = 1$ , the equation has a unique solution  $x = a^{-1} b \pmod{n}$

# Chinese Remainder Theorem (CRT)

◆ Theorem:

Let  $n_1, n_2, \dots, n_k$  be integers s.t.  $\gcd(n_i, n_j) = 1$ ,  
 $i \neq j$ .

$$x \equiv a_1 \pmod{n_1}$$

$$x \equiv a_2 \pmod{n_2}$$

...

$$x \equiv a_k \pmod{n_k}$$

There exists a unique solution modulo  
 $n = n_1 n_2 \dots n_k$

Let  $y_i = (n/n_i)^{-1} \pmod{n_i}$

Then the solution is given by  $x = \sum_i x_i y_i (n/n_i) \pmod{n}$

## Example:

$$x \equiv 5 \pmod{7}$$

$$x \equiv 3 \pmod{11}$$

$$x \equiv 10 \pmod{13}$$

- $n_1=7, n_2=11, n_3=13, n=1001$
- $m_1=143, m_2=91, m_3=77$
- $y_1=143^{-1} \pmod{7} = 3^{-1} \pmod{7} = 5$
- $y_2=91^{-1} \pmod{11} = 3^{-1} \pmod{11} = 4$
- $y_3=77^{-1} \pmod{13} = 12^{-1} \pmod{13} = 12$
- $x=(5 \times 143 \times 5 + 3 \times 91 \times 4 + 10 \times 77 \times 12) \pmod{1001}$   
 $= 13907 \pmod{1001} = 894$

# Famous Number Theory Problems

FACTORING	Given $n$ , find a factor of $n$
RSAP	find $m$ such that $m^e = c \bmod n$
QRP	if $a$ is a quadratic residue mod $n$ , decide whether $a$ is a QR or not.
SQROOT	find $x$ such that $x^2 = a \bmod n$
DLP	find $x$ such that $g^x = y \bmod p$
GDLP	DLP on a finite cyclic group $G$
DHP	given $g^a \bmod p$ , $g^b \bmod p$ , find $g^{ab} \bmod p$
GDHP	DHP on a finite cyclic group $G$
SUBSETSUM	given $\{a_1, \dots, a_n\}$ and $s$ , find subset of $a_j$ that sums to $s$