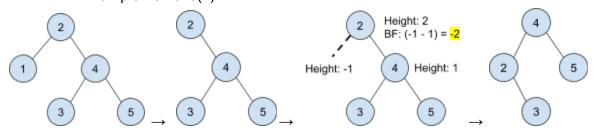
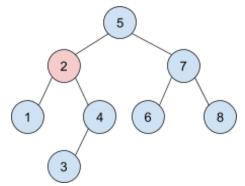
Fall Break - 10/8/18 10/10/18

#### **Topics**

- I. Exam 2 Topic List: see announcement on Canvas for full list
  - A. Coding up to (but not including) HashMaps will be fair game
- II. AVL Remove
  - A. The algorithm:
    - 1. Remove the node like you are removing from a BST (considering all the same 0, 1, and 2 child cases).
    - 2. When tracing back up from the node you removed, update heights and balance factors, and check for and perform rotations.
      - a) If this was a two child remove case, make sure to update and rotate when coming back from removing the predecessor/successor.
  - B. Example: remove(1)

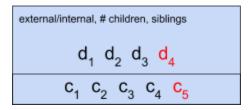


- 1. Remove like you're removing from a BST
- 2. Update the height and balance factor of the node containing 2 (the only node we traversed to find the node containing 1)
  - a) Height = max(-1, 1) + 1
  - b) Balance factor = -1 1 = -2
- 3. Since the balance factor is -2 and its right child's balance factor is not > 0, we just need to do a single left rotation on the node containing 2
- C. Example: remove(2) using predecessor
  - 1. Hint: we need to do a double rotation

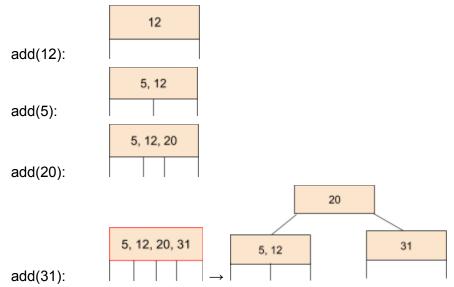


#### III. 2-4 Trees / B-Trees

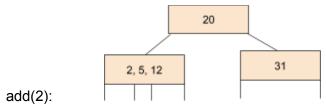
A. 2-4 nodes:



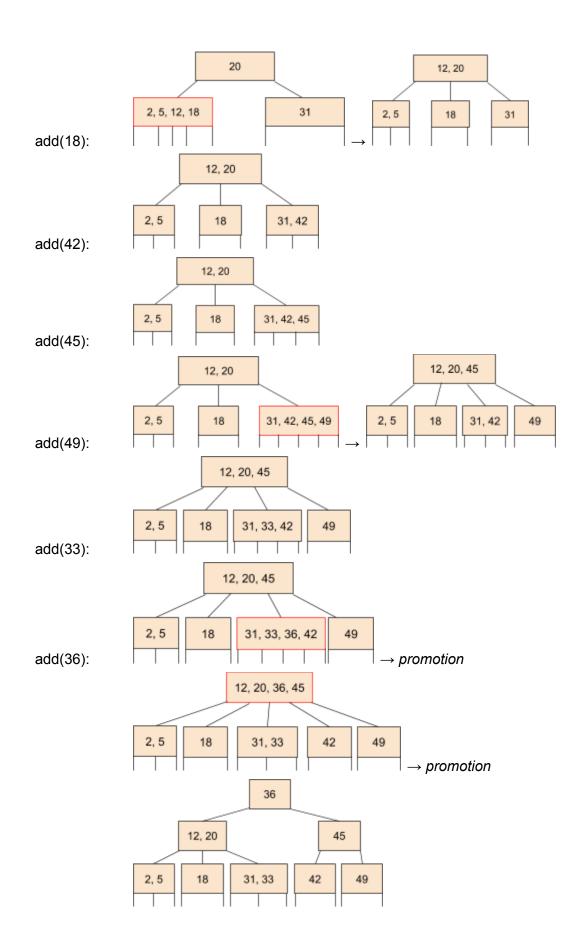
- 1. Can have 2 to 4 children (c<sub>i</sub>) and can store 1 to 3 data values (d<sub>i</sub>)
  - a) A node with m data values must have m + 1 children
  - b) Data 4 and Child 5 are in red because they are allowed to exist momentarily, but will be immediately fixed
- 2. Data is stored within a node in ascending order (d1 < d2 < d3)
- 3. Nodes can also store whether they're internal or external, the number of children they have, and references to their siblings or parent
- B. Order Property:
  - 1.  $c_1 data < d_1 < c_2 data < d_2 < c_3 data < d_3 < c_4 data < d_4 < c_5 data$
- C. Shape Property:
  - 1. Every leaf MUST reside on the same level
  - 2. Run times are guaranteed to be O(log n)
- D. Example: add 12, 5, 20, 31, 2, 18, 42, 45, 49, 33, 36



*Promotion*: occurs when we have overflow in a node; we will always promote a middle node (12 or 20), the choice is arbitrary



We always add new data to a leaf node.



#### Activities

I. Big O reality check

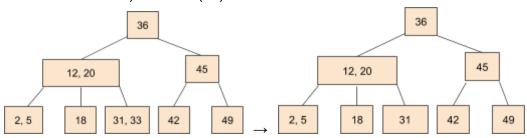
10/12/18

## <u>Announcements</u>

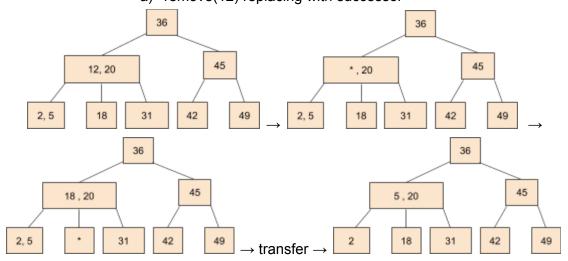
- I. Exam 2
  - A. 2 coding questions
  - B. 1 recursion tracing show your work so we can award partial credit!
  - C. Big-O multiple choice/matching
- II. Practice exams
  - A. 2 for each exam
  - B. We provide hints not answers to help you arrive at the answers yourselves
  - C. If you want your work checked, post privately to Piazza or come to office hours

## **Topics**

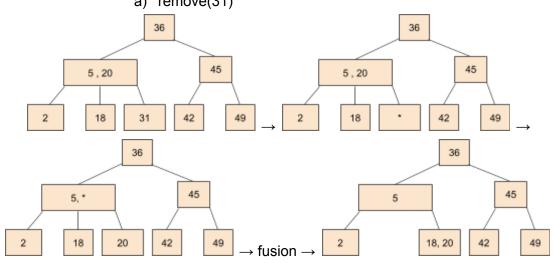
- I. 2-4 Trees
  - A. Add can cause *overflow*, we fix overflow with *promotion* (move a middle value to its parent or create a new node if the overflow is in the root)
  - B. Remove can cause underflow, we fix underflow with transfers and fusion
    - 1. Removing from a leaf with > 1 values (easiest case)
      - a) remove (33)



- 2. Removing from an internal node with a child with > 1 value
  - a) remove(12) replacing with successor



3. Removing from a leaf with a parent with multiple data a) remove(31)



\*for a flowchart of the remove cases, see Canvas Files > Recitation > 2\_4\_remove.jpg

# **Activities**

I. Tree review reality check