

Directed hypergraph connectivity augmentation by hyperarc re-orientations

Combinatorial Optimization and Graph Theory

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- Connectivity problems, characterisations
- Hypergraphs

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 - ▶ Exhibiting a sequence of orientations such that :
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 - The next orientation in the sequence can be obtained from the previous one by flipping exactly one arc.
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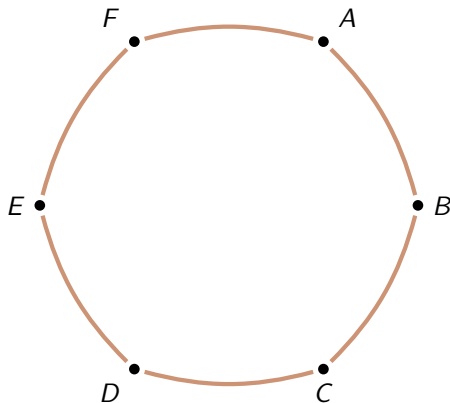
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Side note : This article generalise the results of **Ito and al.**, as directed graphs are special case of hypergraphs.

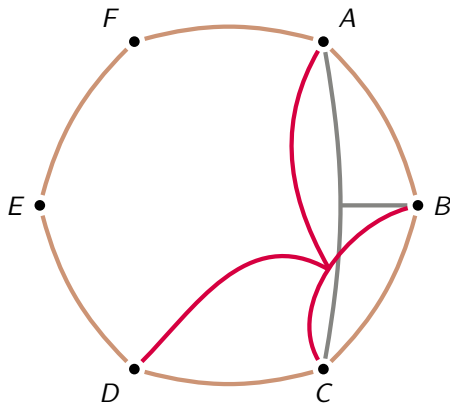
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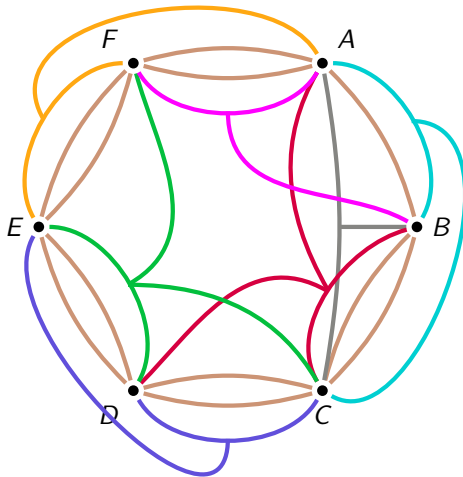


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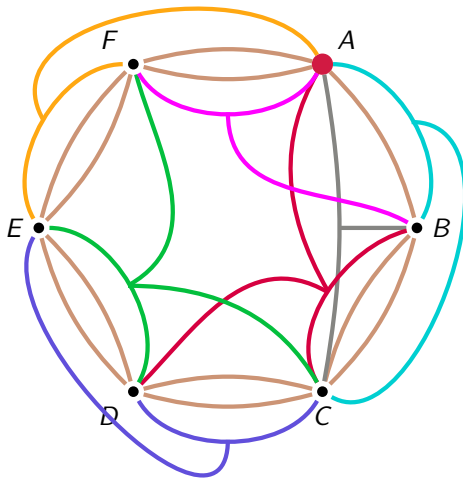
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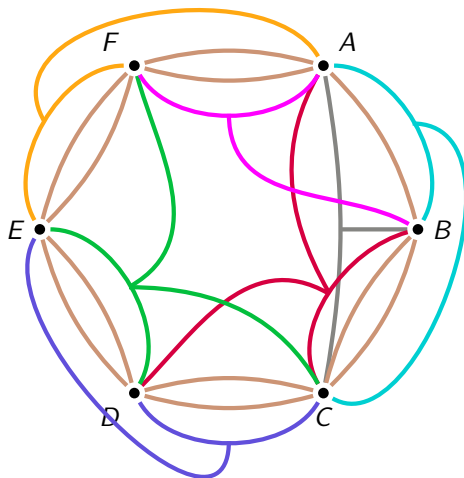


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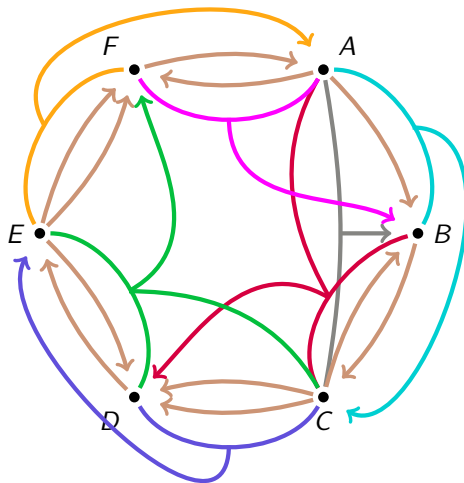
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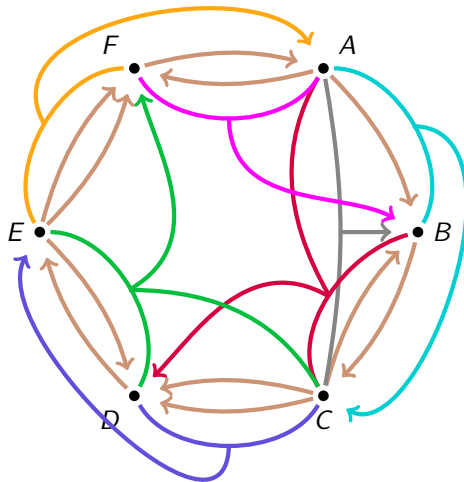


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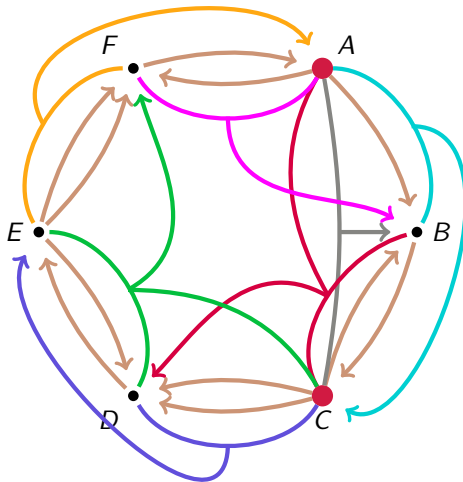
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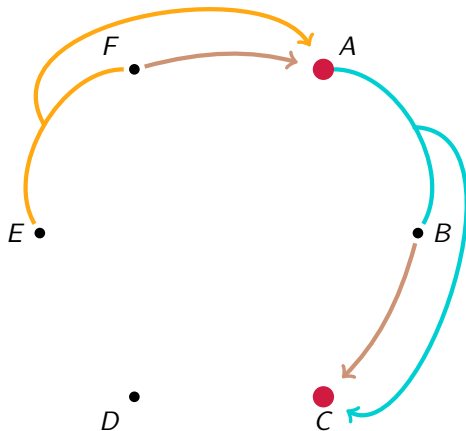
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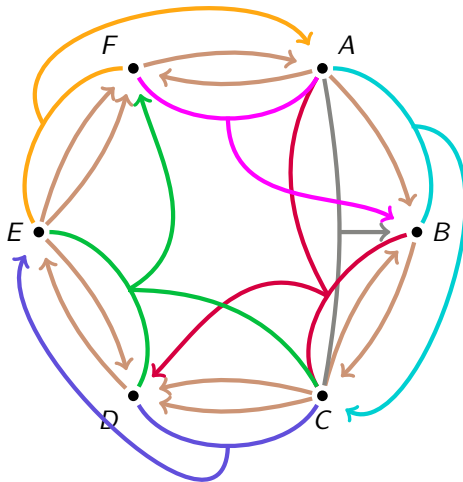
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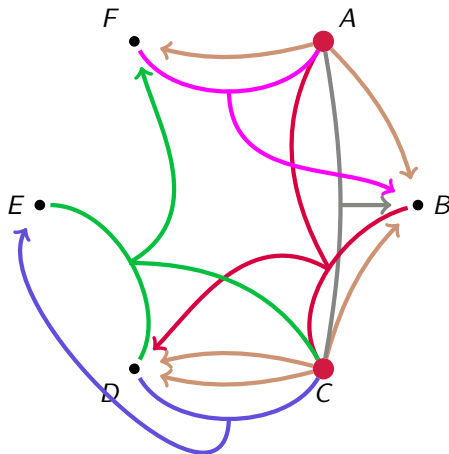
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Hyperarc-connectivity and (k, k) -partition connected hypergraphs

- $\vec{\mathcal{H}}$ is k -hyperarc-connected, if, $\forall \emptyset \neq X \subsetneq V$, $d_{\vec{\mathcal{H}}}^+(X) \geq k$.
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We use a result of Frank : \mathcal{H} is (k, k) -partition-connected if and only if it admits a k -hyperarc-connected orientation.

Main result

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Let $\mathcal{H} = (V, E)$ be a $(k + 1, k + 1)$ -partition-connected hypergraph and $\vec{\mathcal{H}}$ is a k -hyperarc connected orientation of \mathcal{H} . Then there exists a sequence of hypergraphs $(\vec{\mathcal{H}}_i)_{i \in 0 \dots \ell}$ such that $\vec{\mathcal{H}}_{i+1}$ is obtained from $\vec{\mathcal{H}}_i$ by reorienting exactly one hyperarc and $\lambda(\vec{\mathcal{H}}_{i+1}) \geq \lambda(\vec{\mathcal{H}}_i)$ and $\lambda(\vec{\mathcal{H}}_\ell) = k + 1$. Such a sequence of orientations can be obtained with $\ell \leq |V|^3$ and found in polynomial time (in the size of \mathcal{H}).

Main result

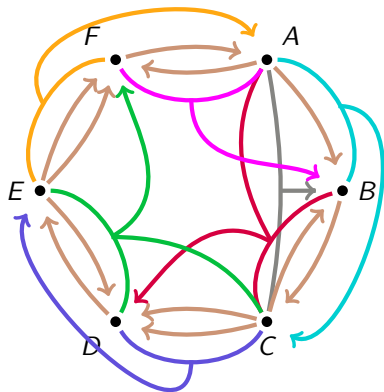
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Generalization of **Ito and al.**, as digraphs are special cases of hypergraphs.

"High-Level"-running of the algorithm

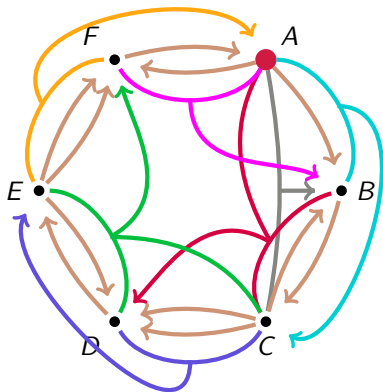
Our algorithm will provide a 3-hyperarc-connected orientation of \mathcal{H} , starting from a 2-hyperarc-connected.



- 1 Take r in $V(\mathcal{H})$.
- 2 Compute sets of vertices.
- 3 Stopping Criteria
- 4 Select a set R (cf. 2.)
- 5 Find an admissible (s, t) -hyperpath in R to reorient
- 6 Reorient the corresponding hyperpath.
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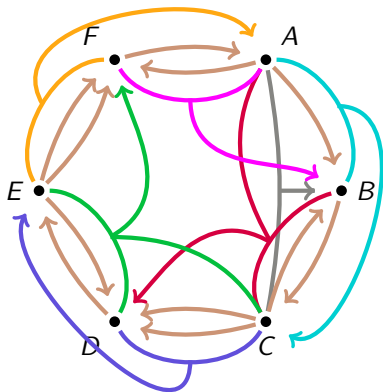
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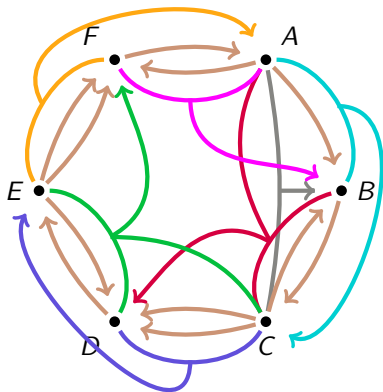
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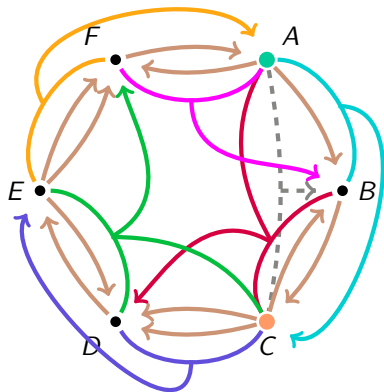
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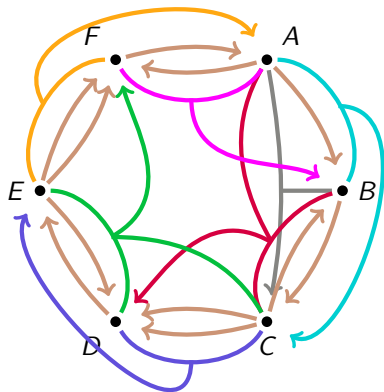
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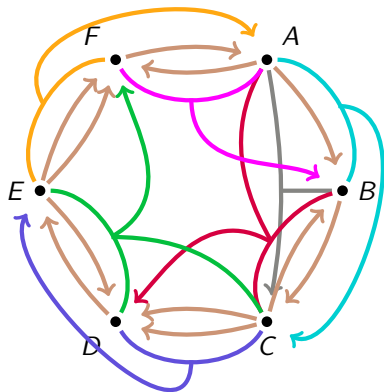
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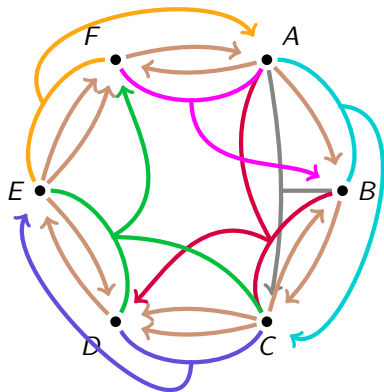
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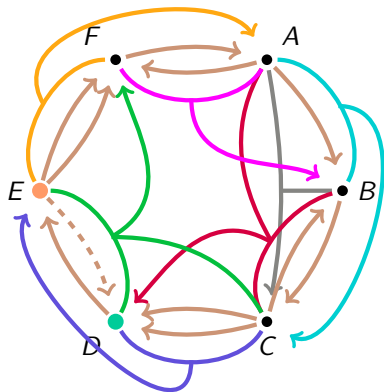
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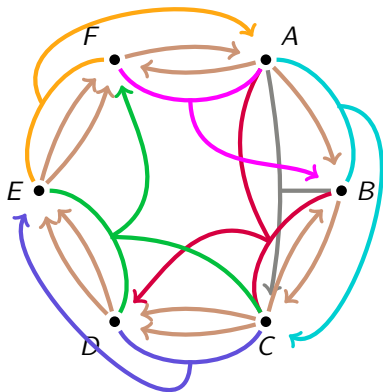
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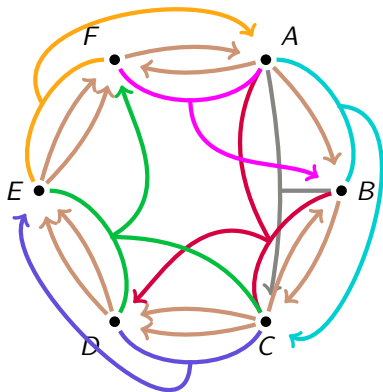
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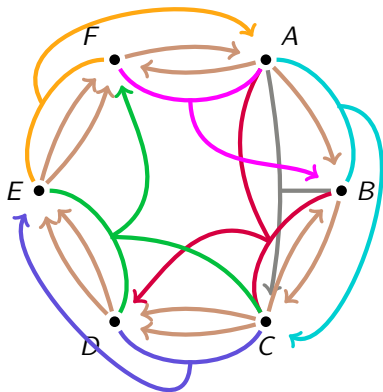
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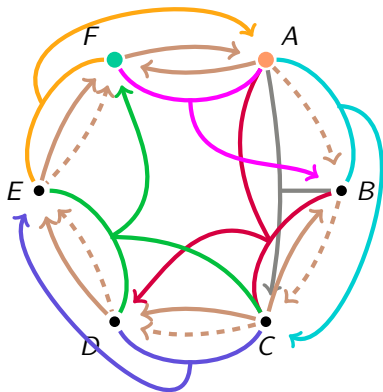
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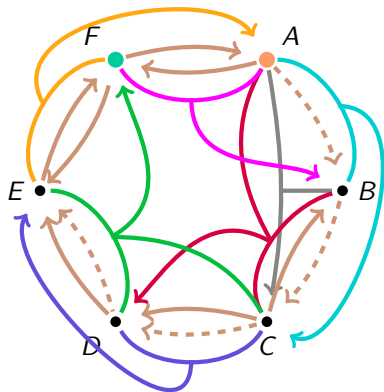
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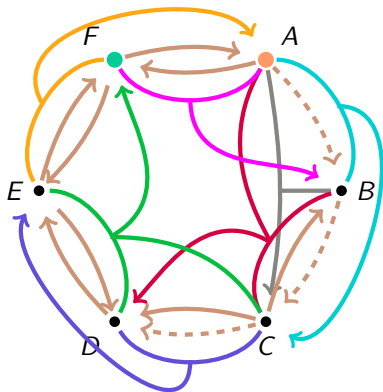
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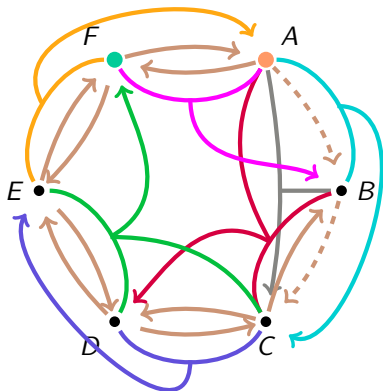
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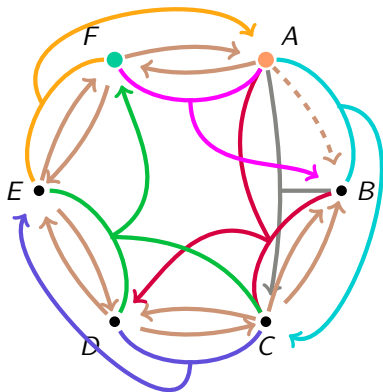
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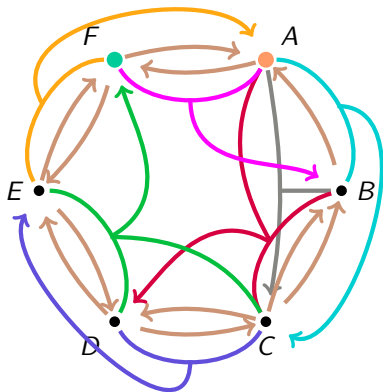
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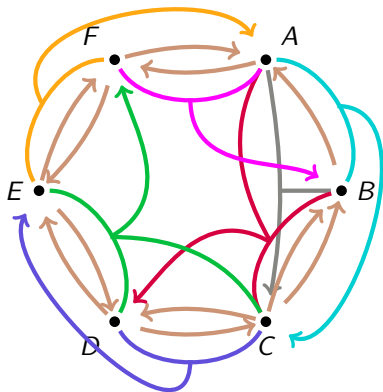
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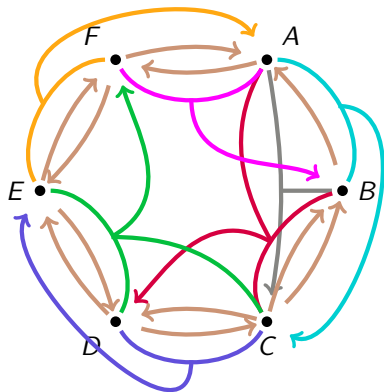
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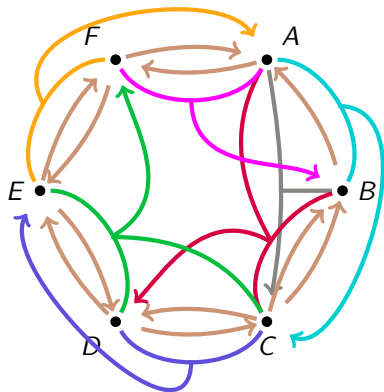
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Finding *admissible* (s, t) -hyperpaths in R

- Crucial segment of the algorithm.
- Three criterion for P to be an admissible (s, t) -hyperpath in R :
 - ① s is a safe source in $S \subseteq R$, t is a safe sink in $T \subseteq R$.
 - ② Reorient each hyperarc, **one by one**, does not decrease the hyperarc-connectivity.
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What are **safe sources** and **safe sinks** ?

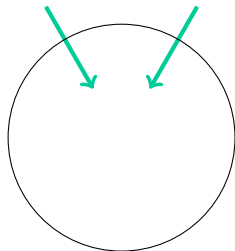
A brief detour...

Tight and Dangerous sets

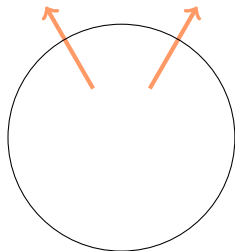
Remainder of the algorithm :

- Input : A k -hyperarc-connected orientation of a $(k + 1, k + 1)$ -partition-connected hypergraph.
- Output : A $k + 1$ -hyperarc-connected hypergraph.

Tight and Dangerous sets



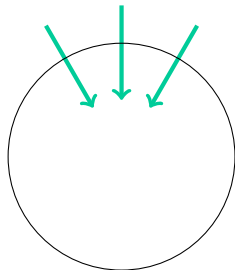
In-Tight sets



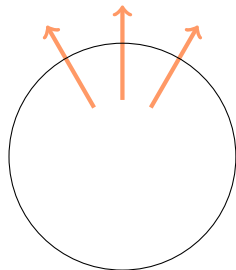
Out-Tight sets

- $\mathcal{T}_- = \{X \subseteq V - r, d^-(X) = k\} \cup \{V\}$
- $\mathcal{T}_+ = \{X \subseteq V - r, d^+(X) = k\} \cup \{V\}$
- $\mathcal{D}_- = \{X \subseteq V - r, d^-(X) = k + 1\}$
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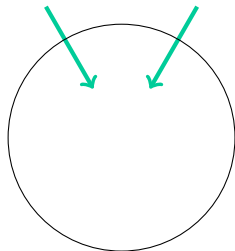
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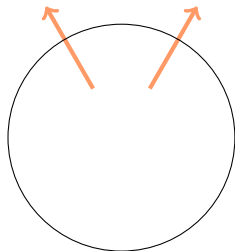
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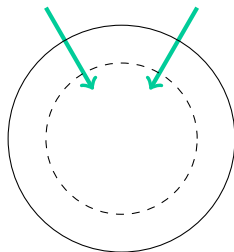
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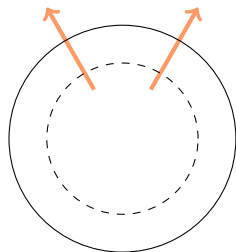
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Tight and Dangerous sets



Minimal In-Tight sets



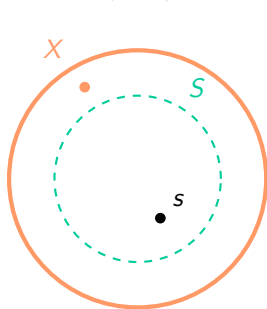
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Safe Sources and Safe Sinks

Definitions are symmetric (but proofs are not).

- For $S \in \mathcal{M}_-$, s is a safe source in S if :
 - For every $s \in X \in \mathcal{T}_+$, we have $S \subsetneq X$.
 - For every $s \in X \in \mathcal{D}_+$ such that $S \setminus X \neq \emptyset$, there exists $Y \in \mathcal{T}_+$ such that $s \notin Y \subsetneq X$.



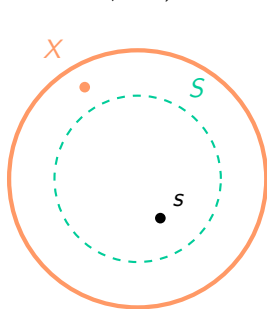
Condition (a)

Finding a safe sink t in $T \in \mathcal{M}_+$ can be done by checking each vertex if they correspond to the definition.

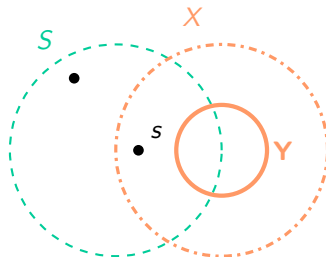
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Condition (a)



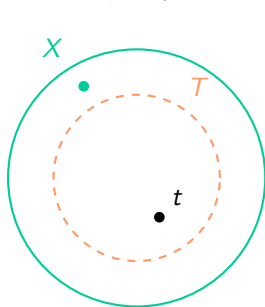
Condition (b)

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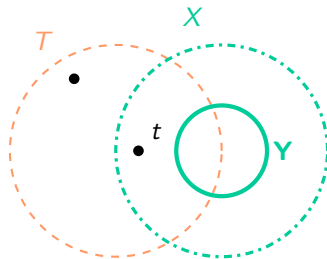
Safe Sources and Safe Sinks

Definitions are symmetric (but proofs are not).

- For $T \in \mathcal{M}_+$, t is a safe sink in T if :
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 - d For every $t \in X \in \mathcal{D}_-$ such that $T \setminus X \neq \emptyset$, there exists $Y \in \mathcal{T}_-$ such that $t \notin Y \subsetneq X$.



Condition (c)



Condition (d)

Finding a safe sink t in $T \in \mathcal{M}_+$ can be done by checking each vertex if they correspond to the definition.

Existence of a safe source (*a safe sink*)

Lemma 10

$\forall S \in \mathcal{M}_-$, there is a safe source $s \in S$.

Likewise,

Lemma 11

$\forall T \in \mathcal{M}_+$, there is a safe sink $t \in T$.

Quick sketch of a proof for Lemma 10 :

- Let $S \in \mathcal{M}_-$.
- Considering a family of vertex sets (χ) that cover as many vertices of S as possible, but using as little as vertex sets possible.
- We can prove that, under given assumptions, χ cannot cover every vertex of S .
- Vertices that are not covered by χ are "potential" safe sources, the last part of the proof is verifying that they are effectively safe sources.

Finding *admissible* (s, t) -hyperpaths in R

- Three criterion for P to be an admissible (s, t) -hyperpath in R :
 1. s is a safe source in $S \subseteq R$, t is a safe sink in $T \subseteq R$.
 - The hyperpath obtained does not leave given sets.
 2. Reorient each hyperarc, **one by one**, does not decrease the hyperarc-connectivity.
 - At the end, we will prove part of this statement.
 3. Let $\vec{\mathcal{H}}$ the hypergraph obtained after reorientation of P .
 - \mathcal{M}' : Inclusion-wise minimal members of $\mathcal{M}'_- \cup \mathcal{M}'_+$
 - Either $|\mathcal{M}'| < |\mathcal{M}|$, either $|\mathcal{M}'| = |\mathcal{M}|$ and \mathcal{M}' covers more vertices than \mathcal{M} .

Point 3. is the stopping criteria for the main algorithm :

- $\mathcal{M} = \{V\}$ implies both $\mathcal{M}_- = \{V\}$ and $\mathcal{M}_+ = \{V\}$.
- $\mathcal{T}_- = \{X \subseteq V - r, d^-(X) = k\} \cup \{V\}$
- $\mathcal{T}_+ = \{X \subseteq V - r, d^+(X) = k\} \cup \{V\}$
- \mathcal{M}_- : Inclusion-wise minimal members of \mathcal{T}_-
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- Finally, if $\lambda(\vec{\mathcal{H}}) \geq k$ and $\mathcal{T}_- = \mathcal{T}_+ = \{V\}$, $\vec{\mathcal{H}}$ is $(k+1)$ -hyperarc-connected.

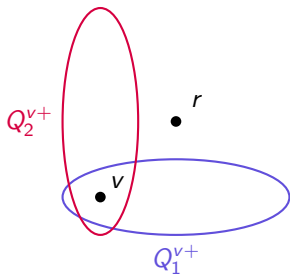
Introduction of Q_+^v

Definition of Q_+^v

Consider the sets of \mathcal{T}_+ containing v . Q_+^v is **the** minimal (inclusion-wise) one.

Unicity of Q_+^v :

If it exists, Q_+^v is unique.



Let Q_1^{v+} , Q_2^{v+} verifying the above definition.

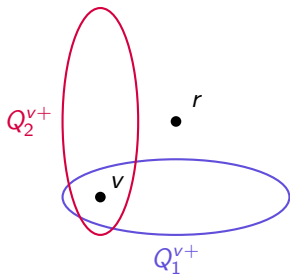
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By definition, $Q_1^{v+} \not\subseteq Q_2^{v+}$ and $Q_2^{v+} \not\subseteq Q_1^{v+}$.

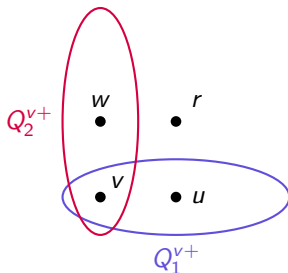
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Denote $u \in Q_1^{v+} \setminus Q_2^{v+}$, $w \in Q_2^{v+} \setminus Q_1^{v+}$.

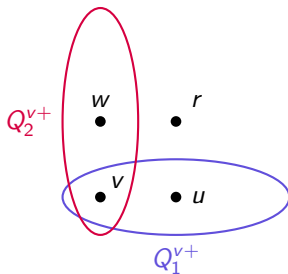
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As $r \notin Q_1^{v+}, Q_2^{v+}$, both are crossing sets.

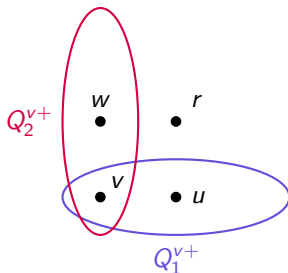
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By submodularity, if $X, Y \in \mathcal{T}_+$, both $X \cup Y \in \mathcal{T}_+$ and $X \cap Y \in \mathcal{T}_+$.

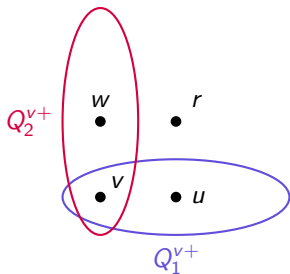
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$Q_1^{v+} \cap Q_2^{v+}$ is smaller (inclusion-wise) than Q_1^{v+} and Q_2^{v+} .