

corresponding to

1. a tour is ^{corresponding to} a cycle and if we delete an edge it's a spanning tree. $MST \leq cycle \leq Tour \leq OPT$

2. Total sum is even. Sum of evens is even.

→ Sum of odds is even. → Size of S is even.

3. the TSP with only S is $\leq OPT$. Now

Since the size is even, alternating

will give us $W(M_1) + \overset{W}{W}(M_2) \leq OPT$ So

$$W(M_1) + W(M_2) \leq \frac{OPT}{2}$$

4. S is nodes with odd degree and M adds

one edge to each node, so \forall nodes even degree

→ \exists Euler circuit, $Cost(T) + Cost(M)$

$$\leq OPT + \frac{OPT}{2}$$

5. $TSP \leq Euler \leq \frac{3}{2} OPT$. ~~Because in some cases~~

~~xxxxxxxxxxxx~~

6. \square