

SEMINAR 1

P1. Bank account

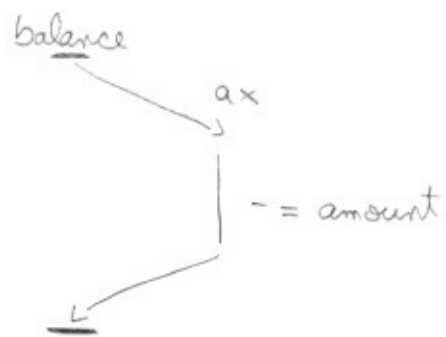
```
transfer(Account *src, Account *dest, int amount)
{
    int;
    int;
```

```
(*src) -= amount;
(*dest) += amount;
```

```
mutex global; lock();
src -> mtx.lock();
dest -> mtx.lock();
mtxglobal.unlock();
src -> balance -= amount;
dest -> balance += amount;
src -> mtx.unlock();
dest -> mtx.unlock();
```

struct Account {

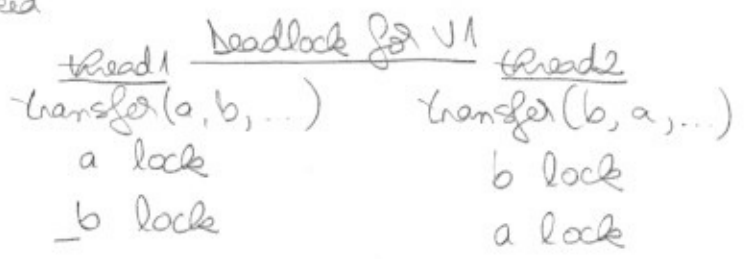
```
int balance; // initial value
mutex mtx; // balance is constant
while mtx is unlocked
```



C++
• mutex lock()
unlock()

C#/Java
lock(o)

→ every object is a mutex

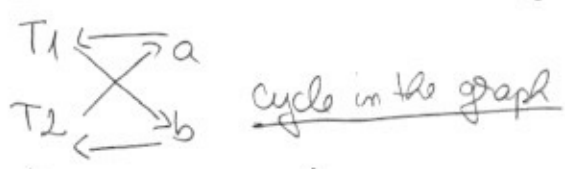


V2 Deadlock free → cât mai puține operații înainte de lock/unlock

```
src -> mtx.lock()
src -> balance -= amount
src -> mtx.unlock()

dest -> mtx.lock()
dest -> balance += amount
dest -> mtx.unlock()
```

Deadlock → circular dependency, two threads are waiting one for the other, consuming same resources



Dependency graph



multiplexes;



```
if (src->id < dest->id) {
    src->mtx.lock();
    dest->mtx.lock();
}
else {
    dest->mtx.lock();
    src->mtx.lock();
}
```

```
struct Account {
    int balance;
    mutex mtx;
    int id;
}
```



```
int sum() {
    int s = 0;
    for (int i = 0; i < accounts.size(); i++) {
        s += accounts[i].balance;
    }
    return s;
}
```

vector

accounts[i].mtx.lock();
accounts[i].mtx.unlock();

acc 1 → 10
acc 2 → 10

T1 → doing sum

sum()

acc[0].lock()

S += acc[0].balance

acc[0].unlock()

T2 → transfer

transfer(0, 1, 5)

acc[0].lock()

acc[1].lock()

acc[0] -= 5

acc[1] += 15

acc[0].unlock()

acc[1].unlock()

acc[1].lock

S += acc[1].balance

5 units are counted twice → wrong value for the sum

CORRECT

```
int sum() {  
    int s = 0;  
    for (int i = 0; i < accounts.size(); i++) {  
        accounts[i].mtx.lock();  
        s += accounts[i].balance;  
    }  
    for (int i = 0; i < accounts.size(); i++) {  
        accounts[i].mtx.unlock();  
    }  
    return s;  
}
```

the balance is constant for every version

Change transfer method so that a check for lock is done on resources:

try-lock → in C++

```
src → mutex.lock()  
while (!dest → mutex.try-lock()) {  
    src → mtx.unlock();  
    sleep(1);  
    src → mtx.lock();  
}
```

T1

a.lock()

T2

b.lock()

false ← b.try-lock()

a.unlock()

a.try-lock() → false

sleep()

b.unlock()

sleep()

a.lock()

b.lock
then it repeats forever

livelock

it's unlikely to happen