

CSE 321

Homework 4

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1)

000000000000000000... -> n

For an n sized 0 sequence, in worst scenario it will take $n - 3$ comparisons to verify if the sequence contains a sequence like "0010" since the length of the searched text is 4.

...000000

The last possible sequence is the last four characters in terms of character length, so we do not check the last three. Since the character that breaks the pattern is the third one (0010), there will be three characters comparisons for each n sized sequence character which is equal to $3 \cdot (n - 3) = 3n - 9$.

The worst-case input pattern of length 3 is 001 since the character that breaks the pattern is the last one.

2)

Since the starting point and traveling in forward or reverse order of the same point sequence does not affect total distance travelled, we will divide permutation of 5 by 5 and 2 respectively as,

$$\frac{P(5,5)}{5.2} = \frac{5!}{5.2} = 4.3 = 12$$

$$\mathbf{A \rightarrow B \rightarrow E \rightarrow D \rightarrow C \rightarrow A = 16}$$

$$\mathbf{A \rightarrow D \rightarrow C \rightarrow B \rightarrow E \rightarrow A = 16}$$

$$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow A = 22$$

$$A \rightarrow B \rightarrow C \rightarrow E \rightarrow D \rightarrow A = 25$$

$$A \rightarrow B \rightarrow D \rightarrow C \rightarrow E \rightarrow A = 21$$

$$A \rightarrow E \rightarrow D \rightarrow E \rightarrow C \rightarrow A = 27$$

$$A \rightarrow B \rightarrow E \rightarrow C \rightarrow D \rightarrow A = 19$$

$$A \rightarrow C \rightarrow B \rightarrow D \rightarrow E \rightarrow A = 27$$

$$A \rightarrow C \rightarrow B \rightarrow E \rightarrow D \rightarrow A = 22$$

$$A \rightarrow C \rightarrow D \rightarrow B \rightarrow E \rightarrow A = 18$$

$$A \rightarrow C \rightarrow E \rightarrow B \rightarrow D \rightarrow A = 21$$

$$A \rightarrow D \rightarrow B \rightarrow C \rightarrow E \rightarrow A = 24$$

3)

```
PROCEDURE log2(n)
  IF n <= 1 THEN
    RETURN 0
  ELSE THEN
    RETURN 1 + log2(n / 2)
  ENDIF
END
```

$$T(n) = T(n / 2) + 1, n > 1, A(1) = 0$$

$$T(n) = \log_2 n \in Q(\log n)$$

4)

```
PROCEDURE findBottle(arr, avrWeight)
    IF arr1.size == 1 THEN
        RETURN arr[0]
    ENDIF

    totalWeightL = totalWeight(arr, 0, arr.size / 2)
    totalWeightR = totalWeight(arr, arr.size / 2 + 1, arr.size)

    IF totalWeightL != avrWeight * (arr.size / 2) THEN
        RETURN getXth(arr[0 : arr.size / 2])
    ELSE THEN
        RETURN getXth(arr[arr.size / 2 + 1 : arr.size])
    ENDIF
END

PROCEDURE totalWeight(arr, start, end)
    totalWeight = 0
    FOR i IN range(start, end)
        totalWeight = totalWeight + arr[i]
    ENDFOR
    RETURN totalWeight
END
```

In this part I implemented an algorithm that divides the bottles into two halves, calculates their total weight and checks which half does not meet “bottle count * average bottle weight” requirement and then calls the recursive function again for that incorrect half. The best, worst and average case complexities of the algorithm is $O(\log n)$ if we assume, we can get the total weight of a half in constant time.

5)

```
PROCEDURE getXth(arr1, arr2, x)
    IF arr1.size == 0 THEN
        RETURN arr2[x]
    ENDIF
    IF arr2.size == 0 THEN
        RETURN arr1[x]
    ENDIF

    mid1 = (arr1.size - 1) / 2
    mid2 = (arr2.size - 1) / 2

    IF mid1 + mid2 < x THEN
        IF arr1[mid1] > arr2[mid2] THEN
            RETURN getXth(arr1, arr2[mid2 + 1 : arr2.size], x - mid2 - 1)
        ELSE THEN
            RETURN getXth(arr1[mid1 + 1 : arr1.size], arr2, x - mid1 - 1)
        ENDIF
    ELSE THEN
        IF arr1[mid1] > arr2[mid2] THEN
            RETURN getXth(arr1[0 : mid1], arr2, x)
        ELSE THEN
            RETURN getXth(arr1, arr2[0 : mid2], x)
        ENDIF
    ENDIF
END
```

Worst-case -> $O(\log n + \log m)$

n -> arr1.size,

m -> arr2.size