

a) Since pthread\_cond\_wait is used inside critical regions (mutex locked regions) it needs to unlock mutex first to not cause a deadlock by preventing other threads to continue and waiting for a signal itself. So inside pthread\_cond\_wait it unlocks mutex, waits for a signal then locks it again.

b)

```
li $t0, 0xA000
li $t1, 0xA0A0
li $t2, 0
mv $t3, $t0
li $t5, 0
loop:
    sw $t4, 0($t3)
    bgt $t4, $t2, swap
loopc:
    addi $t3, $t3, 1
    addi $t5, $t5, 1
    bgt $t5, $t6, loop
    halt
swap:
    lw $t2, $t4
    j loopc
```

li \$t6, 10

c) for P1 only there is 12GB free ram  $n = 12/2 = 6$

$$\text{CPU ut.} = 1 - p^n \text{ where } n = 6, 1 - \left(\frac{6}{10}\right)^6 =$$

for P2 only it would be

$$\text{CPU ut.} = 1 - p^n \text{ where } n = 2$$

$$n = 12/6 = 2$$

$$1 - \left(\frac{3}{10}\right)^2 =$$

2P2s