Topics in Database Theory – Homework 3

September 17, 2023

1 The AGM Bound

- 1. (0 points)
 - (a) Consider the following query:

$$Q(x, y, z, u, v, w) = R(x, y) \land S(y, z) \land T(y, u) \land K(u, v) \land M(x, w)$$

Assume that $|R| = |S| = |T| = |K| = |M| \le N$.

- i. Find the maximum size of the output to the query Q
- ii. Find a worst-case database instance where the query Q has the bound you found above.
- (b) Consider the query:

$$Q(x,y,z,u) = \!\! R(x,y) \wedge S(y,z) \wedge T(z,u) \wedge K(u,x)$$

Suppose the four relations have cardinalities N_1, N_2, N_3, N_4 .

Give a formula that represents a tight upper bound on |Q|. Your formula should use the cardinalities N_1, N_2, N_3, N_4 and operations like $+, \times, /, \hat{}$, max, for example $\max(N_1/N_2, N_3^{3/2} + N_4)$ (not a real answer).

(c) Consider the same query as above, and repeat your answer for the case when y is a key in S:

$$Q(x,y,z,u) = \!\! R(x,y) \wedge S(\underline{y},z) \wedge T(z,u) \wedge K(u,x)$$

2 Information Inequalities

- 2. (0 points)
 - (a) Consider the following query:

$$Q(x, y, z, u) = R(x, y, z) \wedge S(y, z, u) \wedge T(z, u, x) \wedge K(u, x, y)$$

Prove that the following inequalities hold:

$$|Q| \le (|R| \cdot |S| \cdot |T| \cdot |K|)^{1/3}$$

$$|Q| \le |R| \cdot \max(\deg_S(u|yz))$$

$$|Q| \le |T| \cdot \max(\deg_K(y|ux))$$

(b) Consider the following query:

$$Q(x, y, z, u, v, w) = R(x, y, z) \land S(z, u, v) \land T(v, w, x)$$
$$\land A(y, z, u) \land B(u, v, w) \land C(w, x, y)$$

Prove the following inequality:

$$|Q| \leq \sqrt{|R| \cdot |S| \cdot |T| \cdot \max(\mathsf{deg}_A(y|zu)) \cdot \max(\mathsf{deg}_B(u|vw)) \cdot \max(\mathsf{deg}_C(w|xy))}$$

(c) Prove the following inequality:

$$h(xyz) + h(zuv) + h(vwx) + h(yuw) + h(y|x) + h(z|y) + h(u|z) + h(v|u) + h(w|v) + h(x|w) \ge 3h(xyzuvw)$$

More details about information inequalities can be found in [1].

References

[1] D. Suciu. Applications of information inequalities to database theory problems. In *LICS*, pages 1–30, 2023.