

1. Proposition: 命题, 必须为 truth or false, 不能不确定. 通常用 p, q, r, s 表示, 大都是尽量用多, 细的命题
2. 命题操作: 逆: \neg , 异或: \oplus , 交: \wedge , 并: \vee , implication: \rightarrow , biconditional: \leftrightarrow . $p \leftrightarrow q$ iff p, q 同为 true 或同为 false
3. tautology: 恒为 true 的组合命题; contradiction: 恒为 false 的组合命题
4. p and q are logically equivalent if $p \leftrightarrow q$ 为 tautology, 记作 $p \equiv q$
5. Associative laws: $(p \vee q) \vee r \equiv p \vee (q \vee r)$, $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$. De Morgan's laws: $\neg(p \wedge q) \equiv \neg p \vee \neg q$, $\neg(p \vee q) \equiv \neg p \wedge \neg q$
- Distributive laws: $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$, $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$. Absorption laws: $p \vee (p \wedge q) \equiv p$, $p \wedge (p \vee q) \equiv p$
- Negation laws: $p \vee \neg p \equiv \top$, $p \wedge \neg p \equiv \bot$
- Useful law: $p \rightarrow q \equiv \neg p \vee q$

6. Predicate Logic: make statements with variables: $P(x)$. 如

7. 全称量词: \forall , $\forall x$, 对所有 x . 存在量词: \exists , $\exists x$, 存在一个 x .

8. The argument form with premises P_1, P_2, \dots, P_n and conclusion q is valid if $(P_1 \wedge P_2 \wedge \dots \wedge P_n) \rightarrow q$ is a tautology.

9. Rules of Inference for propositional logic: Modus ponens: $\frac{p \rightarrow q, p}{q}$; Modus tollens: $\frac{p \rightarrow q, \neg q}{\neg p}$; Hypothetical syllogism: $\frac{p \rightarrow q, q \rightarrow r}{p \rightarrow r}$; Addition: $\frac{p}{p \vee q}$

Disjunctive syllogism: $\frac{p \vee q, \neg p}{q}$; Simplification: $\frac{p \wedge q}{p}$; Conjunction: $\frac{p, q}{p \wedge q}$; Resolution: $\frac{p \vee q, \neg p \vee r}{q \vee r}$

The intended meaning of the predicate symbol: B ('is a baby'), I ('is illogical'), D ('is despised'), C ('is a crocodile') and M ('can manage'). Domain of x : All human. Domain of y : All animals.

- (a) $\forall x(B(x) \rightarrow I(x))$
 (b) $\forall x(\exists y(C(y) \wedge M(x, y)) \rightarrow \neg D(x))$
 (c) $\forall x(I(x) \rightarrow D(x))$
 (d) $\forall x(B(x) \rightarrow \exists y(C(y) \wedge \neg M(x, y)))$ or $\exists y(C(y) \wedge \forall x(B(x) \rightarrow \neg M(x, y)))$

逻辑与证明

10. 证明方法: direct proof, proof by contrapositive 证逆否命题, proof by cases: 分类讨论, proof of equivalence: 证等价命题.

11. Set 集合: 不考虑元素顺序, 元素可重复, 形如 $\{x | x \text{ has property } P\}$

12. Cardinality: 集的势, 表示集合中不同的元素的数量, 记作 $|S|$, S 为一集合

13. Power set: 幂集, the power set of S is the set of all subsets of the set S , 记作 $P(S)$.

14. Tuple: 元组, 考虑元素顺序, 如 (a_1, a_2, \dots, a_n)

15. Cartesian Product: 笛卡尔积 of 集合 A, B , 记作 $A \times B$, $A \times B = \{(a, b) | a \in A \wedge b \in B\}$

16. 集合操作: $A \cup B = \{x | x \in A \vee x \in B\}$, $A \cap B = \{x | x \in A \wedge x \in B\}$, $\bar{A} = \{x \in U | x \notin A\}$, $A - B = \{x | x \in A \wedge x \notin B\}$.

17. Principle of inclusion-exclusion: $|A \cup B| = |A| + |B| - |A \cap B|$.

18. 函数: Let A and B be two sets. A function from A to B , denoted by $f: A \rightarrow B$, is an assignment of exactly one element of B to each element of A .

相关证明

- One-to-one (injective) function:
 - A function f is called one-to-one or injective if and only if $f(x) = f(y)$ implies $x = y$ for all x, y in the domain of f .
- Onto (surjective) function:
 - A function f is called onto or surjective if and only if for every $b \in B$ there is an element $a \in A$ such that $f(a) = b$.
- One-to-one (bijective) correspondence
 - One-to-one and onto

Inverse function: Let f be a one-to-one correspondence (bijection) from the set A to the set B . The inverse function of f is the function that assigns to an element b belonging to B the unique element a in A such that $f(a) = b$.

Let f be a function from B to C and let g be a function from A to B . The composition of the functions f and g , denoted by $f \circ g$, is defined by $(f \circ g)(x) = f(g(x))$.

The floor function assigns a real number x the largest integer that is $\leq x$, denoted by $\lfloor x \rfloor$. E.g., $\lfloor 3.5 \rfloor = 3$.

The ceiling function assigns a real number x the smallest integer that is $\geq x$, denoted by $\lceil x \rceil$. E.g., $\lceil 3.5 \rceil = 4$.

集合与函数

Suppose that $f: A \rightarrow B$.

To show that f is injective	Show that if $f(x) = f(y)$ for all $x, y \in A$, then $x = y$
To show that f is not injective	Find specific elements $x, y \in A$ such that $x \neq y$ and $f(x) = f(y)$
To show that f is surjective	Consider an arbitrary element $y \in B$ and find an element $x \in A$ such that $f(x) = y$
To show that f is not surjective	Find a specific element $y \in B$ such that $f(x) \neq y$ for all $x \in A$

A set that is either finite or has the same cardinality as the set of positive integers \mathbb{Z}^+ is called countable.

If there is a one-to-one function from A to B , the cardinality of A is less than or equal to the cardinality of B , denoted by $|A| \leq |B|$.

Theorem: If there is a one-to-one correspondence between elements in A and B , then the sets A and B have the same cardinality.

Theorem: If A and B are sets with $|A| \leq |B|$ and $|B| \leq |A|$, then $|A| = |B|$.

A sequence is a function from a subset of the set of integers (typically the set $\{0, 1, 2, \dots\}$ or $\{1, 2, 3, \dots\}$) to a set S .

We use the notation a_n to denote the image of the integer n . $\{a_n\}$ represents the ordered list $\{a_1, a_2, a_3, \dots\}$

19. $f(x)$ is $O(g(x))$ if $|f(x)| \leq C|g(x)|$ when $x > k$; $f(x)$ is $\Omega(g(x))$ if $|f(x)| \geq C|g(x)|$ when $x > k$
 $f(x)$ is $\Theta(g(x))$ if $f(x)$ is $O(g(x))$ and $f(x)$ is $\Omega(g(x))$

算法复杂度

20. 可除性: Divisibility: We say that a divides b if there is an integer c such that $b = ac$, or equivalently b/a is an integer.

余数计算

• If a, b, c are integers, where $a \neq 0$, such that $a|b$ and $a|c$, then $a|(mb + nc)$ whenever m and n are integers.

Corollary: Let m be a positive integer and let a and b be integers. Then,

$$(a + b) \bmod m = ((a \bmod m) + (b \bmod m)) \bmod m$$

$$ab \bmod m = ((a \bmod m)(b \bmod m)) \bmod m$$

同余关系

Congruence Relation: If a and b are integers and m is a positive integer, then a is congruent to b modulo m if m divides $a - b$, denoted by $a \equiv b \pmod{m}$.

Theorem: Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then

$$a + c \equiv b + d \pmod{m}$$

$$ac \equiv bd \pmod{m}$$

The integers a and b are congruent modulo m if and only if there is an integer k such that

$$a = b + km.$$

21. 质数: prime. A integer $p > 1$ is a prime if the only positive factors of p are 1 and p .
若 n 为 composite 数 (非质数), 则 n has a prime divisor less than or equal to \sqrt{n} .

22. 最大公因数 greatest common divisor: 最大整数 d such that $d|a, d|b$, $\mathbb{P} d = \gcd(a, b)$

若 $u = p_1^{a_1} p_2^{a_2} \dots p_n^{a_n}$, $b = p_1^{b_1} p_2^{b_2} \dots p_n^{b_n}$, 则 $\gcd(a, b) = p_1^{\min(a_1, b_1)} p_2^{\min(a_2, b_2)} \dots p_n^{\min(a_n, b_n)}$

23. 最小公倍数 least common multiple: 最小整数 d such that $a|d, b|d$, $\mathbb{P} d = \text{lcm}(a, b)$

若 $u = p_1^{a_1} p_2^{a_2} \dots p_n^{a_n}$, $b = p_1^{b_1} p_2^{b_2} \dots p_n^{b_n}$, 则 $\text{lcm}(a, b) = p_1^{\max(a_1, b_1)} p_2^{\max(a_2, b_2)} \dots p_n^{\max(a_n, b_n)}$

24. Euclidean Algorithm 欧几里德算法 计算 \gcd :

For two integers 287 and 91, we want to find $\gcd(287, 91)$.

$$\text{Step 1: } 287 = 91 \cdot 3 + 14$$

$$\text{Step 2: } 91 = 14 \cdot 6 + 7$$

$$\text{Step 3: } 14 = 7 \cdot 2 + 0$$

25. 若 $\gcd(a, b) = 1$, 则称 a, b are relatively prime

$$\gcd(287, 91) = \gcd(91, 14) = \gcd(14, 7) = 7$$

26. Bezout's Theorem: If a and b are positive integers, then there exist integers s and t such that

$$\gcd(a, b) = sa + tb.$$

Lemma: If a, b, c are positive integers such that $\gcd(a, b) = 1$ and $a|bc$, then $a|c$.

Lemma: If p is prime and $p|a_1 a_2 \dots a_n$, then $p|a_i$ for some i .

27. 线性同余:

A congruence of the form $ax \equiv b \pmod{m}$, where m is a positive integer, and a and b are integers, and x is a variable, is called a linear congruence.

The solutions to a linear congruence $ax \equiv b \pmod{m}$ are all integers x that satisfy the congruence.

$$x \equiv ab \pmod{m}$$

28. 模逆元:

Modular Inverse: An integer \bar{a} such that $\bar{a}a \equiv 1 \pmod{m}$ is said to be an inverse of a modulo m .

存在条件

When does inverse exist?

Theorem: If a and m are relatively prime integers and $m > 1$, then an inverse of a modulo m exists. The inverse is unique modulo m . That is,

Solve the congruence $ax \equiv b \pmod{m}$ by multiplying both sides by \bar{a} .

$$x \equiv \bar{a}b \pmod{m}.$$

例题

How to find inverses?

Using extended Euclidean algorithm:

Example: Find an inverse of 101 modulo 4620. That is, find \bar{a} such that $\bar{a} \cdot 101 \equiv 1 \pmod{4620}$.

With extended Euclidean algorithm, we obtain $\gcd(a, b) = sa + tb$, i.e., $1 = -35 \cdot 4620 + 1601 \cdot 101$. It tells us that -35 and 1601 are Bezout coefficients of 4620 and 101. We have

$$1 \bmod 4620 = 1601 \cdot 101 \bmod 4620.$$

Thus, 1601 is an inverse of 101 modulo 4620.

generating functions:

TABLE 1 Useful Generating Functions.

$G(x)$	a_k
$(1+x)^n = \sum_{k=0}^n C(n, k)x^k$ $= 1 + C(n, 1)x + C(n, 2)x^2 + \dots + x^n$	$C(n, k)$
$(1+ax)^n = \sum_{k=0}^n C(n, k)a^k x^k$ $= 1 + C(n, 1)ax + C(n, 2)a^2 x^2 + \dots + a^n x^n$	$C(n, k)a^k$
$(1+x^r)^n = \sum_{k=0}^n C(n, k)x^{rk}$ $= 1 + C(n, 1)x^r + C(n, 2)x^{2r} + \dots + x^{rn}$	$C(n, k/r)$ if $r k$; 0 otherwise
$\frac{1-x^{n+1}}{1-x} = \sum_{k=0}^n x^k = 1 + x + x^2 + \dots + x^n$	1 if $k \leq n$; 0 otherwise
$\frac{1}{1-x} = \sum_{k=0}^{\infty} x^k = 1 + x + x^2 + \dots$	1
$\frac{1}{1-ax} = \sum_{k=0}^{\infty} a^k x^k = 1 + ax + a^2 x^2 + \dots$	a^k
$\frac{1}{1-x^r} = \sum_{k=0}^{\infty} x^{rk} = 1 + x^r + x^{2r} + \dots$	1 if $r k$; 0 otherwise
$\frac{1}{(1-x)^2} = \sum_{k=0}^{\infty} (k+1)x^k = 1 + 2x + 3x^2 + \dots$	$k+1$
$\frac{1}{(1-x)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)x^k$ $= 1 + C(n, 1)x + C(n+1, 2)x^2 + \dots$	$C(n+k-1, k) = C(n+k-1, n-1)$
$\frac{1}{(1+x)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)(-1)^k x^k$ $= 1 - C(n, 1)x + C(n+1, 2)x^2 - \dots$	$(-1)^k C(n+k-1, k) = (-1)^k C(n+k-1, n-1)$
$\frac{1}{(1-ax)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)a^k x^k$ $= 1 + C(n, 1)ax + C(n+1, 2)a^2 x^2 + \dots$	$C(n+k-1, k)a^k = C(n+k-1, n-1)a^k$
$e^x = \sum_{k=0}^{\infty} \frac{x^k}{k!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots$	$1/k!$
$\ln(1+x) = \sum_{k=1}^{\infty} \frac{(-1)^{k+1}}{k} x^k = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots$	$(-1)^{k+1}/k$