

Dr. Hakim Mellah Department of Computer Science & Software Engineering Concordia University, Montreal, Canada

Lecture 7: Synchronization Examples

COMP 346: Operating Systems winter 2020

These slides has been extracted, modified and updated from original slides of:

Operating System Concepts, 10th Edition, by: Silberschatz/Galvin/Gagne, published by John Wiley & Sons

Classical Problems of Synchronization

Classical problems used to test newly-proposed synchronization schemes

- Bounded-Buffer Problem
- Readers and Writers Problem

Bounded-Buffer Problem

- **>**n buffers, each can hold one item
- ➤ Semaphore **mutex** initialized to the value 1
- ➤ Semaphore **full** initialized to the value 0
- Semaphore **empty** initialized to the value n

Bounded Buffer Problem (Cont.)

➤ The structure of the producer process

```
do {
     /* produce an item in next produced */
   wait(empty);
   wait(mutex);
     /* add next produced to the buffer */
   signal(mutex);
   signal(full);
} while (true);
```

Bounded Buffer Problem (Cont.)

☐ The structure of the consumer process

```
Do {
   wait(full);
   wait(mutex);
    /* remove an item from buffer to next consumed */
       . . .
    signal(mutex);
    signal(empty);
    /* consume the item in next consumed */
 } while (true);
```

Readers-Writers Problem

- > A data set is shared among a number of concurrent processes
 - Readers only read the data set; they do **not** perform any updates
 - Writers can both read and write
- ➤ Problem allow multiple readers to read at the same time
 - Only one single writer can access the shared data at the same time
- ➤ Several variations of how readers and writers are considered all involve some form of priorities
- ➤ Shared Data
 - **❖** Data set
 - Semaphore rw_mutex initialized to 1
 - Semaphore mutex initialized to 1
 - Integer read_count initialized to 0

Readers-Writers Problem (Cont.)

> The structure of a writer process

Readers-Writers Problem (Cont.)

➤ The structure of a reader process

```
do {
     wait(mutex);
      read count++;
      if (read count == 1)
       wait(rw mutex);
    signal(mutex);
      /* reading is performed */
        . . .
    wait(mutex);
      read count--;
      if (read count == 0)
    signal(rw mutex);
    signal(mutex);
} while (true);
```

Readers-Writers Problem Variations

- ▶First variation no reader kept waiting unless writer has permission to use shared object
- ➤ **Second** variation once writer is ready, it performs the write

 ASAP
- ➤ Both may have starvation leading to even more variations
- ➤ Problem is solved on some systems by kernel providing reader-writer locks