

### FA with ε-Transitions

- We can allow <u>explicit</u> ε-transitions in finite automata
  - i.e., a transition from one state to another state without consuming any additional input symbol
  - Makes it easier sometimes to construct NFAs

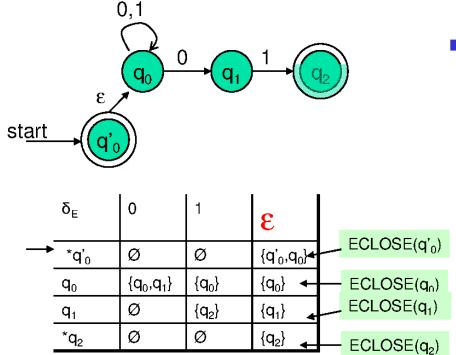
# <u>Definition:</u> $\varepsilon$ -NFAs are those NFAs with at least one explicit $\varepsilon$ -transition defined.

 ε -NFAs have one more column in their transition table



## Example of an $\varepsilon$ -NFA

 $L = \{w \mid w \text{ is empty, } \underline{or} \text{ if non-empty will end in } 01\}$ 



ε-closure of a state q, ECLOSE(q), is the set of all states (including itself) that can be reached from q by repeatedly making an arbitrary number of εtransitions.

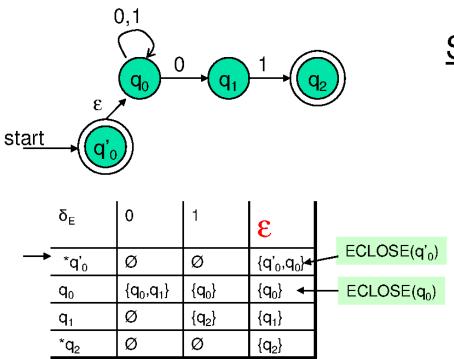
#### To simulate any transition:

Step 1) Go to all immediate destination states.

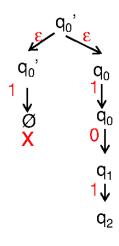
Step 2) From there go to all their  $\varepsilon$ -closure states as well.

# Example of an $\varepsilon$ -NFA

 $L = \{w \mid w \text{ is empty, or if non-empty will end in 01}\}$ 



### Simulate for w=101:

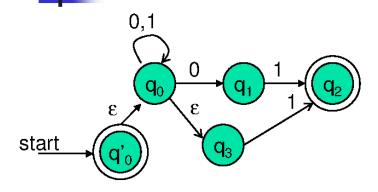


#### To simulate any transition:

Step 1) Go to all immediate destination states.

Step 2) From there go to all their  $\epsilon$ -closure states as well.

## Example of another $\varepsilon$ -NFA



	$\delta_{E}$	0	1	3
<b>→</b>	*q' <sub>0</sub>	Ø	Ø	{q' <sub>0</sub> ,q <sub>0</sub> ,q <sub>3</sub> }
	$q_0$	$\{q_0,q_1\}$	$\{q_0\}$	{q <sub>0,</sub> q <sub>3</sub> }
	$q_1$	Ø	{q <sub>2</sub> }	{q <sub>1</sub> }
	*q <sub>2</sub>	Ø	Ø	{q <sub>2</sub> }
,	$q_3$	Ø	{q <sub>2</sub> }	{q <sub>3</sub> }

#### Simulate for w=101:

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