

TUTORIAL - 7 & 8

Q.1. > Write a Program Specification to compute greatest common divisor.

A.1. >

Pre-condition $\{ i_1 > 0 \text{ and } i_2 > 0 \}$

Procedure P

output

Post-condition

$0 = Z$ and such that

$$\left\{ \begin{array}{l} \text{exist } z, k_1, k_2 \quad (i_1 = z \times k_1 \text{ and } i_2 = z \times k_2) \\ \text{and} \\ \text{not exist } y \quad (i_1 = y \times k_3 \text{ and } i_2 = y \times k_4 \text{ and } y > z) \\ \text{(exist } k_3, k_4) \end{array} \right\}$$

Q.2. > Write a Program Specification to produce reverse of input sequence.

A.2. >

Precondition $\{ n > 0 \}$

Procedure

P

Input

output

Post-condition

for all $i \ (1 \leq i \leq n)$,

$O_i = I_{n-i+1}$

Q.3. > Give a logic specification for a program that reads a sequence of $n+1$ values and checks whether the first value also appears in next input n values.

A.3. >

$\{ n > 0 \}$

check Element $(n, i_0, \{ i_1, \dots, i_n \})$

{

0 , implies $(\text{exist } j \ (1 \leq j \leq n) \text{ and } i_0 = i_j)$

}

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(Extra space)

Q4.4 Give a logic specification for a program that first read two words (i.e. two sequences of alphabetic character by a blank and terminated by '#'). The second word may be null, the first must not.

Then, the program reads a sequence of other words, separated by blanks and terminated by '#', and rewrites the sequence, substituting ~~the~~ all occurrences of first word by second.

P.T.O

$\{ \begin{array}{l} n_1 \geq 2 \\ n_2 \geq 0 \end{array}$
 $\text{words}_{n_1} = \#$
 $\text{exists } j \text{ and } \text{words}_j = ' ' \text{ and } 1 \leq j \leq n_1$
 $\text{for all } j \text{ in } 1 \leq j \leq n_1 \text{ (} \text{words}_j \text{ in } '[a-zA-Z\backslash]' \text{)}$
 $\text{sentence}_{n_2} = \#$
 $\text{for all } j \text{ in } 1 \leq j \leq n_2 \text{ sentence}_j \text{ in } '[a-zA-Z\backslash]'$
 $\}$
 $\text{replaceWord}(n_1, \text{words}, n_2, \text{sentence})$
 $\{ \begin{array}{l} 0\text{-sentence} = \# \\ n_3 \end{array}$
 $n_3 \geq 0$
 $\text{first word} = \text{words} [1 \dots \text{position of } (' ')]$
 $\text{second word} = \text{words} [\text{position of } (' ') \dots n_1]$
 $\text{first word in sentence}$
 $\text{implies second word in } 0\text{-sentence}$
 $\}$

systems.

Q5> Mention advantages and disadvantages of using FSM specifications for
A5>

Advantages of FSM

- ① FSM simplicity make it easy for inexperienced developers to implement with little to no extra knowledge (low entry level)
- ② Predictability (in deterministic FSM)
- ③ Quick to design, quick to implement & quick in execution
- ④ Relatively Flexible (Number of ways to implement FSM)
- ⑤ Easy to transfer meaningful abstract representation to a coded implementation.
- ⑥ Low processor overhead & easy determination of reachability of ^{state}

Disadvantages of FSM

- ① The predictable nature of deterministic FSM can be ^{unwanted in} ~~difficult to manage~~ some domains like computer games.
- ② Finite memory
 > expressive power is limited
- ③ State explosion
 > given a number of FSMs with k_1, k_2, \dots, k_n states, their composition is a FSM with $k_1 \times k_2 \times k_3 \dots k_n$. (Exponential)
- ④ FSMs are essentially a synchronous model (at any time, a global state of the system, must be defined and a single transition must occur.

Q6. > Write Algebraic Specification for stack operations.

algebra StackOfItem

imports Boolean;

introduces

sort Stack, Item;

operations

Create: \rightarrow Stack; create() \rightarrow Stack

IsEmpty: Stack \rightarrow Boolean; push (Stack, Integer) \rightarrow Stack

Push: Stack \times Item \rightarrow Stack; pop (Stack) \rightarrow Stack

Pop: Stack \rightarrow Stack; peek (Stack) \rightarrow Integer

Top: Stack \rightarrow Item; IsEmpty (Stack) \rightarrow Boolean

constraints Create, IsEmpty, Push, Pop, Top so that
Stack generated by [Create, Push]

for all [S: stack, i: item]

Pop (create()) = Exception (empty stack)

IsEmpty (Create) = true; pop (Push (S, I)) = S

IsEmpty (Push (S, i)) = false; peek (create()) = exception (empty
Stack)

Pop (Create) = error;

Top (Create) = error; peek (Push (S, I)) = I

Pop (Push (S, i)) = S; IsEmpty (create()) = true

Top (Push (S, i)) = i; IsEmpty (Push (S, I)) = false

end StackOfItem;

Q7.7 Compare or Differentiate PetriNets with FSMs.

A7.7 We know that both PetriNets and FSM are models that represent the discrete interactions in system. Both are not possible practically but define the behavior of system.

Here the difference b/w PetriNets (PN) and Finite State machines (FSM)

- ① FSM is a model that represent how single activity change its behavior over time whereas PN represents multiactivity co-ordinates.
- ② In FSM, designer knows the starting point and how the process is going, but PN is asynchronous process in which designer does not know about starting point and how process is working. (concurrent system)
- ③ FSM occurs only on one system based on time, but PN contain co-ordinate system that includes two or three processes, one depend on the previous one. So deadlocks can be placed in PN.
- ④ FSM can be based on inputs and outputs and of two types - Mealey FSM and Moore FSM but PN is based on places (like communication roles), transition (like transformation of objects) and Tokens (like physical objects).
- ⑤ PetriNets are suited for timed processes whereas FSMs are not suited for it.