



COMPUTER ENGINEERING DEPARTMENT,
SVNIT, SURAT

Automata and Formal Languages - CS208

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* Required

AFL-Unit Test 1

Which one of the following regular expressions represents the language: the set of all binary strings having two consecutive 0s and two consecutive 1s? *

- (A) $(0+1)^*0011(0+1)^* + (0+1)^*1100(0+1)^*$
(B) $(0+1)^*(00(0+1)^*11 + 11(0+1)^*00)(0+1)^*$
(C) $(0+1)^*00(0+1)^* + (0+1)^*11(0+1)^*$
(D) $00(0+1)^*11 + 11(0+1)^*00$

- ☐ D
☒ A
☐ C
☐ B

Eliminate left recursion from the grammar: *

$E \rightarrow E+T \mid T$

$E \rightarrow TE'$
 $E' \rightarrow +TE' \mid \epsilon$

☒ Option 1

$E \rightarrow E'$
 $E' \rightarrow +TE'$

☐ Option 2

$E \rightarrow +TE'$
 $E' \rightarrow TE'$

☐ Option 3

None of the above

☐ Option 4

The given grammar is of type *

$XYZ \rightarrow ba$

☒ Type 0 Grammar

☐ Type 2 Grammar

☐ Type 1 Grammar

☐ Type 3 Grammar

Total number of useless symbols in given grammar G is *

$S \rightarrow ABC \mid BaB$
 $A \rightarrow aA \mid BaC \mid aaa$
 $B \rightarrow bBb \mid a$
 $C \rightarrow CA \mid AC$

☐ 0

☒ 1

☐ 2

☐ 3

Consider the Following regular expressions *

$r1 = 1(0 + 1)^*$
 $r2 = 1(1 + 0)^+$
 $r3 = 11^*0$

What is the relation between the languages generated by the regular expressions above ?

☐ $L(r1) \supseteq L(r2)$ and $L(r2) \subseteq L(r3)$

☒ $L(r1) \supseteq L(r2)$ and $L(r2) \supseteq L(r3)$

☐ $L(r1) \supseteq L(r3)$ and $L(r2) \subseteq L(r1)$

☐ $L(r1) \subseteq L(r2)$ and $L(r1) \subseteq L(r3)$

Context Free Languages is closed under *

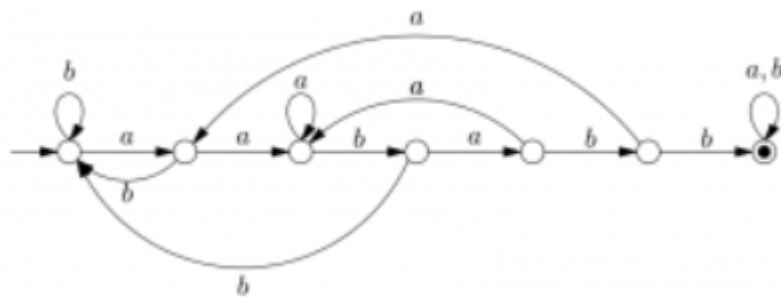
☐ Union

☐ Concatenation



- ☒ Both (a) and (b)
- ☐ None of the above

Consider the following Deterministic Finite Automata and state which statement is true. *



- ☐ It only accepts strings with prefix as "aababb"
- ☐ It only accepts strings with substring as "aababb"
- ☐ It only accepts strings with suffix as "aababb"
- ☒ All of the above

How many unit productions are there in following grammar: *

$S \rightarrow S+T / T$
 $T \rightarrow T * F / F$
 $F \rightarrow (S) / a$

- ☐ 1
- ☒ 2
- ☐ 3
- ☐ 4

Ambiguity can be removed using *

- ☐ Associativity of Operators



- ☐ Precedence of Operators
- ☐ Separate the Productions
- ☒ All of the above

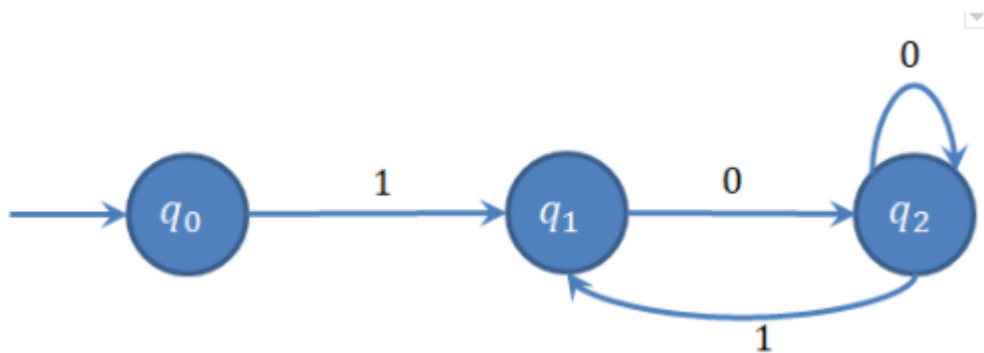
According to the 5-tuple representation i.e. $FA = \{Q, \Sigma, \delta, q, F\}^*$

Statement 1: $q \in Q^*$;

Statement 2: $F \in Q$

- ☐ Statement 1 is false, Statement 2 is true
- ☐ Statement 1 may be true, Statement 2 is false
- ☒ Statement 1 is true, Statement 2 is false
- ☐ Statement 1 is false, Statement 2 may be true

In given automata, q_0 =initial state and q_2 =final state .What is $\delta^*(101)$? *



- ☐ q_2
- ☐ q_0
- ☒ q_1
- ☐ q_1, q_2

In some programming language, L denotes the set of letters and D denotes the set of digits. An identifier is permitted to be a letter followed by any number of letters or digits. The regular expression that defines an identifier is *

☐ $L D^*$

- ☐ $(L.D)^*$
- ☐ $(L + D)^*$
- ☐ $L (L.D)$
- ☒ $L (L + D)^*$

Is the given grammar ambiguous? *

$S \rightarrow SaSbS \mid SbSaS \mid \epsilon$

- ☐ Given grammar is not CFG
- ☐ Insufficient Data
- ☐ No
- ☒ Yes

Are the given productions in chomsky normal form? *

$A \rightarrow a \mid BCD \mid cd$

- ☐ Yes
- ☐ Can't say
- ☐ Given grammar is not CFG
- ☒ No

There is moore machine wit m states and n outputs.If we convert this moore machine to mealy machine then maximum number of states possible in mealy machine is: *

- ☒ $m-1$
- ☐ $m+1$
- ☐ $m*n$
- ☐ m



Is this production in CFG? *

$SA \rightarrow ab$

- ☐ Insufficient Data
- ☐ Given production is not satisfying the grammar rules.
- ☐ Yes
- ☒ No

Consider the context-free grammar and which of the following terminal strings has more than one parse tree when parsed according to the grammar? *

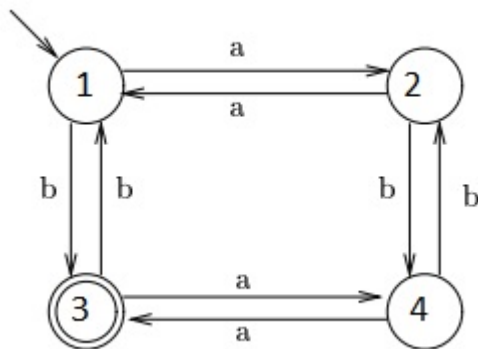
$E \rightarrow E + E$
 $E \rightarrow (E * E)$
 $E \rightarrow id$

- ☐ $((id * id + id) * id)$
- ☒ $id + id + id + id$
- ☐ $(id * (id * id)) + id$
- ☐ $id + (id * (id * id))$

Which two of the following four regular expressions are equivalent? (i). $(00)^*$ $(\epsilon + 0)$ (ii). $(00)^*$ (iii). 0^* (iv). $0(00)^*$ *

- ☐ (iii) and (iv)
- ☐ (i) and (ii)
- ☒ (i) and (iii)
- ☐ (ii) and (iii)

This DFA represents *



- ☐ Set of all strings over $\{a,b\}$ which contains odd a and odd b
- ☐ Set of all strings over $\{a,b\}$ which contains even a and even b
- ☐ Set of all strings over $\{a,b\}$ which contains odd a and even b
- ☒ Set of all strings over $\{a,b\}$ which contains even a and odd b

The language corresponding to the regular expression $(0 + 1)^*(10)$ is *

- ☐ Always ends with 0
- ☒ Always ends with 10
- ☐ Start with 0
- ☐ All of the above

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