## MIPS REFERENCE SHEET

#### TA: Kevin Liston

There are a few special notations outlined here for reference.

Notation	Meaning	Example
{X, Y}	Concatenate the bits of X and Y together.	{10, 11, 011} = 1011011
X × Y	Repeat bit X exactly Y times.	{1, 0 × 3} = 1000
(X)[B:A]	Slice bits A through B (inclusive) out of X.	(1100110101)[4:0] = 10101
SignExt <sub>Nb</sub> (X)	Sign-extend X from N bits to 32 bits.	SignExt <sub>4b</sub> (1001) = $\{1 \times 28, 1001\}$
Mem <sub>NB</sub> (X)	Refers to the N-byte quantity in memory at byte address X.	
R[N]	Refers to the general-purpose register number N.	

### **INSTRUCTION FORMATS**

There are 3 main instruction formats in MIPS. The fields in each type are laid out in such a way that the same fields are always in the same place for each type.

Туре	31	26	25	21	20	16	15	11	10	06	05	00
R-Type	opcode		\$rs		\$rt		\$rd		shamt		funct	
I-Type	opcode		\$rs		\$rt		imm					
J-Type	opcode		addre	ss								

### **R-Type Instructions**

These instructions are identified by an opcode of 0, and are differentiated by their funct values. Except for the first 3 shift instructions, these operations only use registers. Note that in addition to arithmetic operations, these instructions also include jumps and the system call instruction.

	Instruction	RTL	Notes
00	sll \$rd, \$rt, shamt	R[\$rd] ← R[\$rt] << shamt	
02	srl \$rd, \$rt, shamt	R[\$rd] ← R[\$rt] >> shamt	Unsigned right shift
03	sra \$rd, \$rt, shamt	R[\$rd] ← R[\$rt] >> shamt	Signed right shift
04	sllv \$rd, \$rt, \$rs	R[\$rd] ← R[\$rt] << R[\$rs]	

	Instruction	RTL	Notes
06	srlv \$rd, \$rt, \$rs	R[\$rd] ← R[\$rt] >> R[\$rs]	Unsigned right shift
07	srav \$rd, \$rt, \$rs	R[\$rd] ← R[\$rt] >> R[\$rs]	Signed right shift
08	jr \$rs	PC ← R[\$rs]	R[\$rs] must be a multiple of 4
09	jalr \$rd, \$rs	tmp ← R[\$rs] R[\$rd] ← PC + 8 PC ← tmp	R[\$rs] must be a multiple of 4; Undefined if \$rs = \$rd
09	jalr \$rs	(special form of "jalr \$rd, \$	Grs" where \$rd = 31, implicitly)
12	syscall	System Call	
16	mfhi \$rd	R[\$rd] ← HI	
17	mthi \$rs	HI ← R[\$rs]	
18	mflo \$rd	R[\$rd] ← LO	
19	mtlo \$rs	LO ← R[\$rs]	
24	mult \$rs, \$rt	{HI, LO} ← R[\$rs] * R[\$rt]	Signed multiplication
25	multu \$rs, \$rt	{HI, LO} ← R[\$rs] * R[\$rt]	Unsigned multiplication
26	div \$rs, \$rt	LO ← R[\$rs] / R[\$rt] HI ← R[\$rs] % R[\$rt]	Signed division
27	divu \$rs, \$rt	LO ← R[\$rs] / R[\$rt] HI ← R[\$rs] % R[\$rt]	Unsigned division
32	add \$rd, \$rs, \$rt	R[\$rd] ← R[\$rs] + R[\$rt]	Exception on signed overflow
33	addu \$rd, \$rs, \$rt	R[\$rd] ← R[\$rs] + R[\$rt]	
34	sub \$rd, \$rs, \$rt	R[\$rd] ← R[\$rs] - R[\$rt]	Exception on signed overflow
35	subu \$rd, \$rs, \$rt	R[\$rd] ← R[\$rs] - R[\$rt]	
36	and \$rd, \$rs, \$rt	R[\$rd] ← R[\$rs] & R[\$rt]	
37	or \$rd, \$rs, \$rt	R[\$rd] ← R[\$rs]   R[\$rt]	
38	xor \$rd, \$rs, \$rt	R[\$rd] ← R[\$rs] ^ R[\$rt]	
39	nor \$rd, \$rs, \$rt	R[\$rd] ← !(R[\$rs]   R[\$rt])	
42	slt \$rd, \$rs, \$rt	R[\$rd] ← R[\$rs] < R[\$rt]	Signed comparison
43	sltu \$rd, \$rs, \$rt	R[\$rd] ← R[\$rs] < R[\$rt]	Unsigned comparison

# J-Type Instructions

These instructions are identified and differentiated by their opcode numbers (2 and 3). Jump instructions use pseudo-absolute addressing, in which the upper 4 bits of the computed address are taken relatively from the program counter.

	Instruction	RTL
02	j address	PC ← {(PC + 4)[31:28], address, 00}
03	jal address	R[31] ← PC + 8 PC ← {(PC + 4)[31:28], address, 00}

### I-Type Instructions

These instructions are identified and differentiated by their opcode numbers (any number greater than 3). All of these instructions feature a 16-bit immediate, which is sign-extended to a 32-bit value in every instruction (except for the and, or, and xor instructions which zero-extend and the lui instruction in which it does not matter). Branch instructions also effectively multiply the immediate by 4, to get a byte offset.

	Instruction	RTL	Notes
04	beq \$rs, \$rt, imm	if(R[\$rs] = R[\$rt]) PC ← PC + 4 + SignExt <sub>18b</sub> ({imm, 00})	
05	bne \$rs, \$rt, imm	if(R[\$rs] != R[\$rt]) PC ← PC + 4 + SignExt <sub>18b</sub> ({imm, 00})	
06	blez \$rs, imm	if(R[\$rs] <= 0) PC ← PC + 4 + SignExt <sub>18b</sub> ({imm, 00})	Signed comparison
07	bgtz \$rs, imm	if(R[\$rs] > 0) PC ← PC + 4 + SignExt <sub>18b</sub> ({imm, 00})	Signed comparison
08	addi \$rt, \$rs, imm	R[\$rt] ← R[\$rs] + SignExt <sub>16b</sub> (imm)	Exception on signed overflow
09	addiu \$rt, \$rs, imm	R[\$rt] ← R[\$rs] + SignExt <sub>16b</sub> (imm)	
10	slti \$rt, \$rs, imm	R[\$rt] ← R[\$rs] < SignExt <sub>16b</sub> (imm)	Signed comparison
11	sltiu \$rt, \$rs, imm	R[\$rt] ← R[\$rs] < SignExt <sub>16b</sub> (imm)	Unsigned comparison
12	andi \$rt, \$rs, imm	R[\$rt] ← R[\$rs] & {0 × 16, imm}	
13	ori \$rt, \$rs, imm	R[\$rt] ← R[\$rs]   {0 × 16, imm}	
14	xori \$rt, \$rs, imm	R[\$rt] ← R[\$rs] ^ {0 × 16, imm}	
15	lui \$rt, imm	R[\$rt] ← {(imm)[15:0], 0 × 16}	
32	lb \$rt, imm(\$rs)	$R[\$rt] \leftarrow SignExt_{8b}(Mem_{1B}(R[\$rs] + SignExt_{16b}(imm)))$	
33	lh \$rt, imm(\$rs)	R[\$rt] ← SignExt <sub>16b</sub> (Mem <sub>2B</sub> (R[\$rs] + SignExt <sub>16b</sub> (imm)))	Computed address must be a multiple of 2

	Instruction	RTL	Notes
34	lw \$rt, imm(\$rs)	R[\$rt] ← Mem <sub>4B</sub> (R[\$rs] + SignExt <sub>16b</sub> (imm))	Computed address must be a multiple of 4
36	lbu \$rt, imm(\$rs)	$R[\$rt] \leftarrow \{0 \times 24, Mem_{1B}(R[\$rs] + SignExt_{16b}(imm))\}$	
37	lhu \$rt, imm(\$rs)	$R[\$rt] \leftarrow \{0 \times 16, Mem_{2B}(R[\$rs] + SignExt_{16b}(imm))\}$	Computed address must be a multiple of 2
40	sb \$rt, imm(\$rs)	$Mem_{1B}(R[\$rs] + SignExt_{16b}(imm)) \leftarrow (R[\$rt])[7:0]$	
41	sh \$rt, imm(\$rs)	$Mem_{2B}(R[$rs] + SignExt_{16b}(imm)) \leftarrow (R[$rt])[15:0]$	Computed address must be a multiple of 2
43	sw \$rt, imm(\$rs)	Mem <sub>4B</sub> (R[\$rs] + SignExt <sub>16b</sub> (imm)) ← R[\$rt]	Computed address must be a multiple of 4