

TUTORIAL - 4

NFA : NON-DETERMINISTIC FINITE AUTOMATA

UI9CS012

[BHAGYA VINOD RANA]

Q1. Construct a NFA to accept strings over alphabet $\Sigma = \{0, 1\}$

a) The third symbol from right is '0'.

 $\Sigma = \{0, 1\}$ [Input symbol] $L = \text{Language} = \{011, 010, 000, 001, 0011, \dots\}$ $Q = \text{No. of finite states} = n+1 = 3+1 = 4$ ($n = \text{size of min 'L'}$)
String

(A) Transition Table

(State)		0	1
Initial \rightarrow	q_0	$\{q_0, q_1\}$	$\{q_0\}$
	q_1	$\{q_2\}$	$\{q_2\}$
	q_2	$\{q_3\}$	$\{q_3\}$
Final \times	q_3	ϕ	ϕ

(B) Transition Diagram

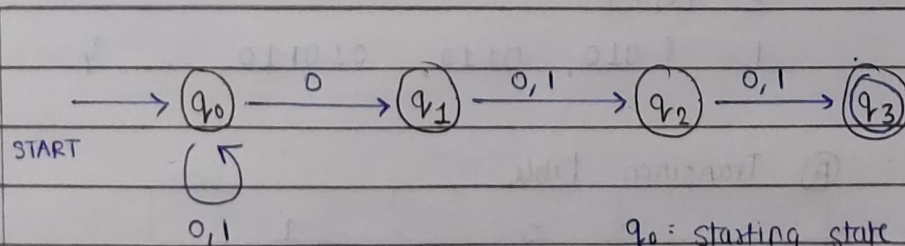


Fig (1). (a)

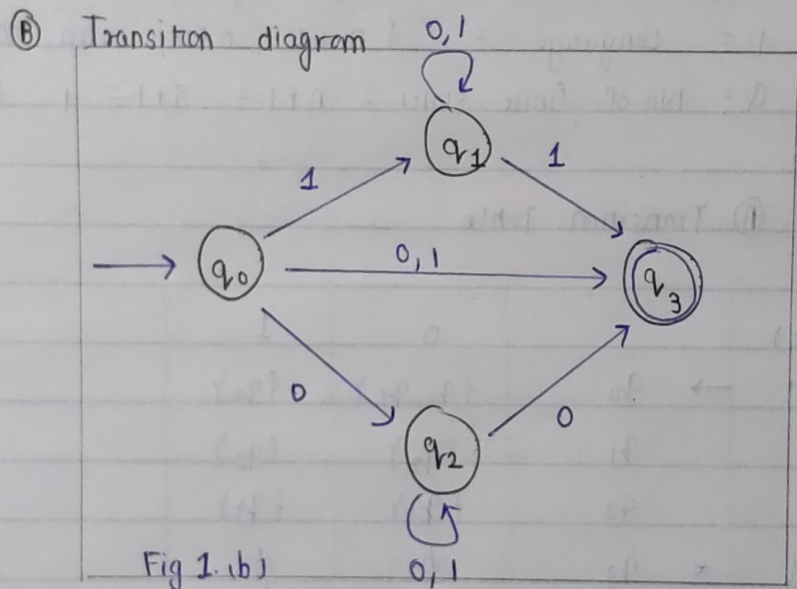
 $q_0 = \text{starting state}$ $q_3 = \text{final state}$

b) The first and last digits are same

 $\Sigma = \{0, 1\}$ $L = \{0, 1, 010, 101, 000, 111, \dots\}$

1.6) (A) Transition table

	0	1
→ q_0	$\{q_2, q_3\}$	$\{q_1, q_3\}$
q_1	$\{q_1\}$	$\{q_1, q_3\}$
q_2	$\{q_2, q_3\}$	$\{q_2\}$
* q_3	ϕ	ϕ



1.c) Start with 01 and end with 10

$$\Sigma = \{0, 1\}$$

$$L = \{010, 0110, 010110, \dots\}$$

(A) Transition table

	0	1
Start → q_0	$\{q_1\}$	ϕ
q_1	ϕ	$\{q_2, q_3\}$
q_2	$\{q_2\}$	$\{q_2, q_3\}$
q_3	$\{q_4\}$	ϕ
final * q_4	ϕ	ϕ

1(c) (B) Transition Diagram

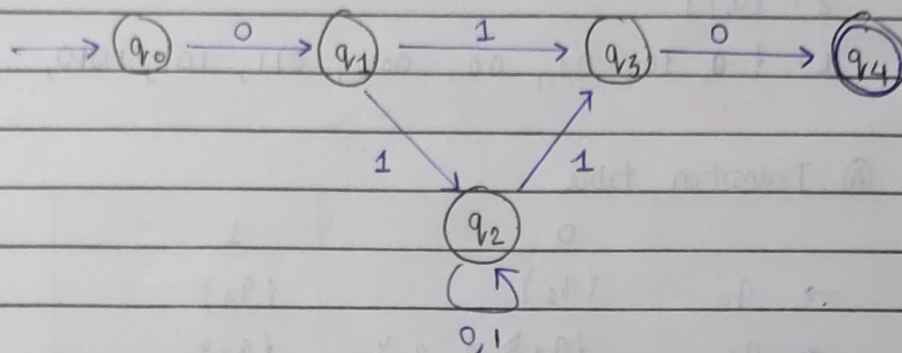


Fig 1.(c)

1) d) Ending with 110

$$\Sigma = \{0, 1\}$$

$$L = \{110, 0110, 1110, 01110, \dots\}$$

$$\text{No. of States} = n+1 = 3+1 = 4$$

$$\{n = \text{min}^m \text{ string in } L\}$$

(A) Transition Table

	0	1
q_0	$\{q_0\}$	$\{q_0, q_2\}$
q_1	\emptyset	$\{q_2\}$
q_2	$\{q_3\}$	\emptyset
q_3	\emptyset	\emptyset

(B) Transition Diagram

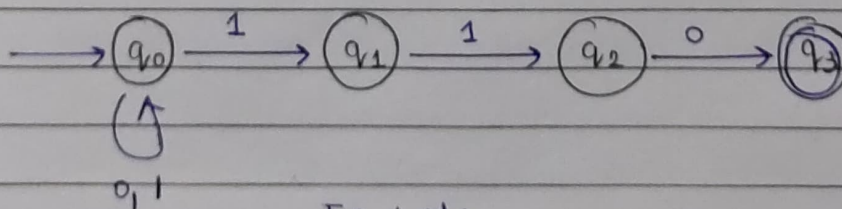


Fig 1.(d)

(1)(e) 0 as one of last three characters in string

$$\Sigma = \{0, 1\}$$

$$L = \{0, 10, 01, 00, 001, 011, 101, 1010, \dots\}$$

(A) Transition table

	0	1
$\rightarrow q_0$	$\{q_1\}$	$\{q_4\}$
* q_1	$\{q_1, q_2, q_3\}$	$\{q_2\}$
* q_2	$\{q_2, q_1, q_3\}$	$\{q_3\}$
* q_3	$\{q_1, q_3, q_2\}$	$\{q_4\}$
q_4	$\{q_1, q_2, q_3\}$	$\{q_4, q_0\}$

(B) Transition Diagram

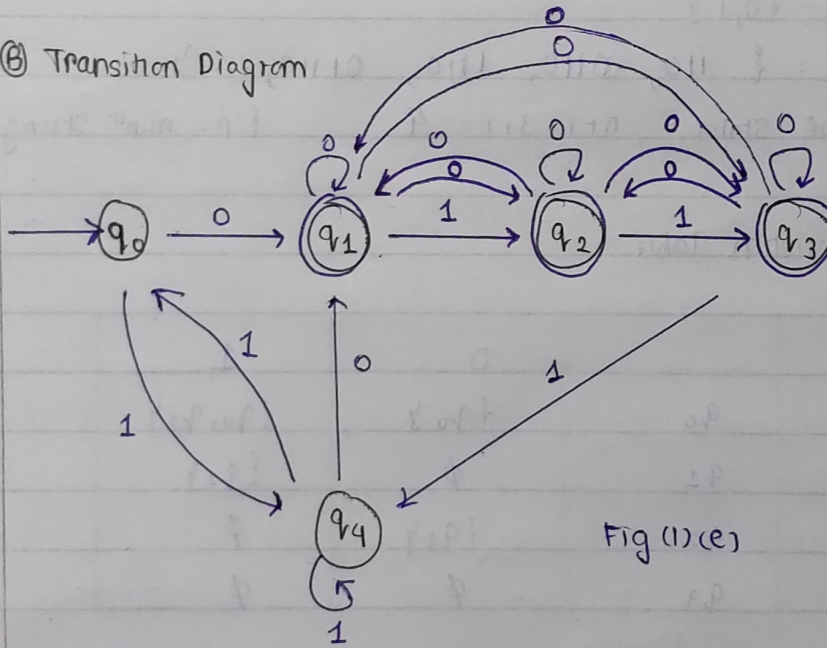


Fig (1)(e)

2.7 Design a NFA to accept strings over alphabets $\Sigma = \{a, b\}$ ending with aba. Now use this NFA to construct DFA to accept the same set of strings.

2.8 NFA : $Q = \{q_0, q_1, q_2, q_3\}$ start = q_0
 $\Sigma = \{a, b\}$ final = q_3 } state

(A) Transition table

	a	b
q_0	$\{q_0, q_1\}$	$\{q_0\}$
q_1	\emptyset	$\{q_2\}$
q_2	$\{q_3\}$	\emptyset
q_3	\emptyset	\emptyset

(B) Transition Diagram

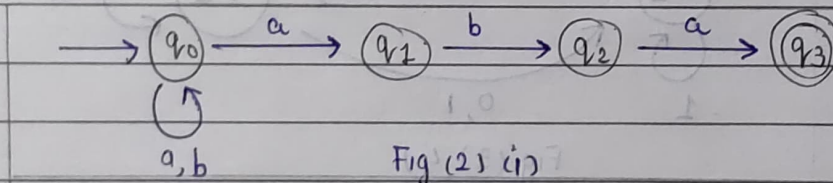


Fig (2) (i)

DFA : $Q = \{q_0, q_1, q_2, q_3\}$ $q_0 = \text{start}$
 $\Sigma = \{a, b\}$ $q_3 = \text{end}$

(A) Transition Table

	a	b
q_0	q_1	q_0
q_1	q_1	q_2
q_2	q_3	q_0
q_3	q_1	q_2

(B) Transition Diagram

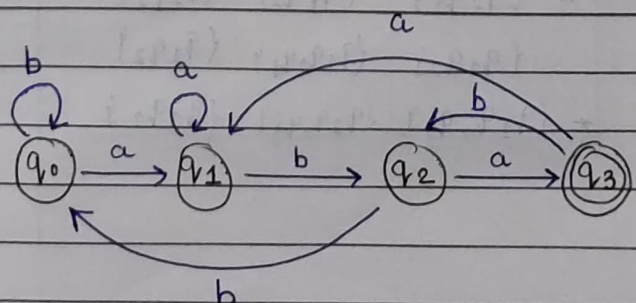


Fig (2) (ii)

3. > let $M = \{q_1, q_2, q_3\}, \{0, 1\}, \{q_1\}, \{q_3\}$ NFA

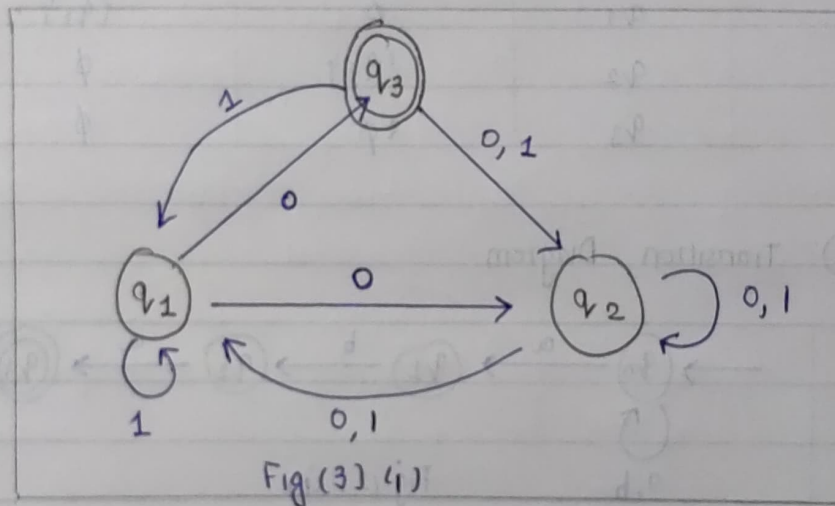
Transition Table Given :

	0	1
q_1	$\{q_2, q_3\}$	$\{q_1\}$
q_2	$\{q_1, q_2\}$	$\{q_1, q_2\}$
q_3	$\{q_2\}$	$\{q_1, q_2\}$

Construct Transition Diagram corresponding to above NFA.

Also convert NFA \rightarrow equivalent DFA.

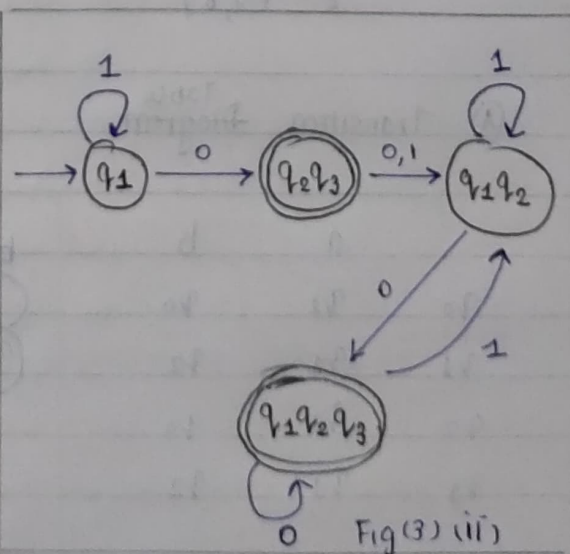
⑧ Transition Diagram



⑨ Transition Table for DFA

State	0	1
$\rightarrow \{q_1\}$	$\{q_2, q_3\}$	$\{q_1\}$
$* \{q_2, q_3\}$	$\{q_1, q_2\}$	$\{q_1, q_2\}$
$\{q_3, q_2\}$	$\{q_2, q_3\}$	$\{q_1, q_2\}$
$* \{q_1, q_2, q_3\}$	$\{q_1, q_2, q_3\}$	$\{q_1, q_2\}$

⑩ Transition Diagram



4. Construct the transition diagram from given transition table for NFA.
Convert NFA to equivalent DFA.

(A) Transition Table

States	0	1
$\rightarrow q_0$	$\{q_2\}$	\emptyset
q_1	\emptyset	$\{q_0, q_2\}$
$* q_2$	$\{q_0, q_1\}$	$\{q_0\}$

(B) Transition Diagram

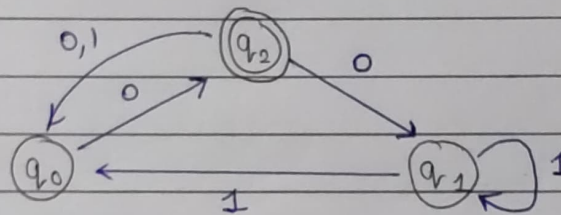


Fig 4 (i)

(A2) Transition table for DFA

States	0	1
$\rightarrow q_0$	$\{q_2\}$	$\{q_3\}$
$* \{q_2\}$	$\{q_0, q_1\}$	$\{q_0\}$
$\{q_0, q_1\}$	$\{q_2\}$	$\{q_0, q_1\}$
$\{q_3\}$	$\{q_3\}$	$\{q_3\}$

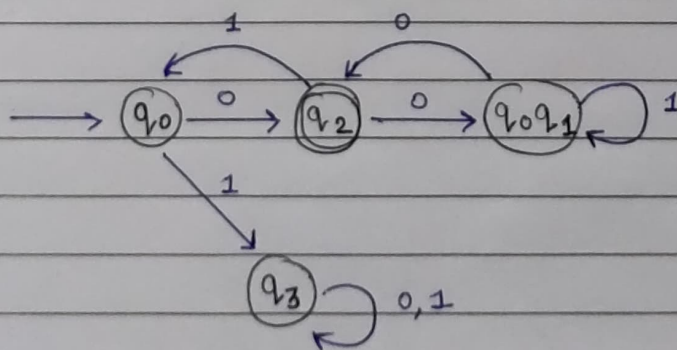


Fig (4) (ii)

SUBMITTED BY:

BHAQVA VINOD RANA

UI9CS012

B-TECH C IInd year