

Example 2.18 List the ordered pairs in the relation represented by the digraph given in Fig. 2.19. Also use the graph to prove that the relation is a partial ordering. Also draw the directed graphs representing R^{-1} and \bar{R} .

The ordered pairs in the relation are $\{(a, a), (a, c), (b, a), (b, b), (b, c), (c, c)\}$.

Since there is a loop at every vertex, the relation is reflexive.

Though there are edges $b - a$, $a - c$ and $b - c$, the edges $a - b$, $c - a$ and $c - b$ are not present in the digraph. Hence the relation is antisymmetric.

When edges $b - a$ and $a - c$ are present in the digraph, the edge $b - c$ is also present (for example). Hence the relation is transitive.

Hence the relation is a partially ordering. The digraph of R^{-1} is got by reversing the directions of the edges (Fig. 2.20). The digraph of \bar{R} contains the edges (a, b) , (c, a) , and (c, b) as shown in Fig. 2.21.

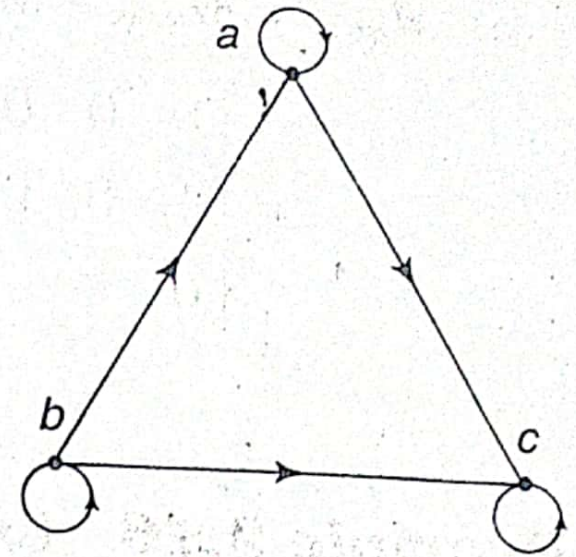


Fig. 2.19

HASSE DIAGRAMS FOR PARTIAL ORDERINGS

The simplified form of the digraph of a partial ordering on a finite set that contains sufficient information about the partial ordering is called a **Hasse diagram**, named after the twentieth-century mathematician Helmut Haasse.

The simplification of the digraph as a Hasse diagram is achieved in three ways:

- (i) Since the partial ordering is a **reflexive relation**, its digraph has loops at all vertices. We need not show these loops since they must be present.
- (ii) Since the partial ordering is **transitive**, we need not show those edges that must be present due to transitivity. For example, if $(1, 2)$ and $(2, 3)$ are edges in the digraph of a partial ordering, $(1, 3)$ will also be an edge due to transitivity. This edge $(1, 3)$ need not be shown in the corresponding Hasse diagram.
- (iii) If we assume that all edges are **directed upward**, we need not show the directions of the edges.

Thus the Hasse diagram representing a partial ordering can be obtained from its digraph, by removing **all the loops**, by removing **all edges** that are present due to **transitivity** and by **drawing each edge without arrow** so that its initial vertex is below its terminal vertex.

For example, let us construct the Hasse diagram for the partial ordering $\{(a, b) \mid a \leq b\}$ on the set $\{1, 2, 3, 4\}$ starting from its digraph. (Fig. 2.15)

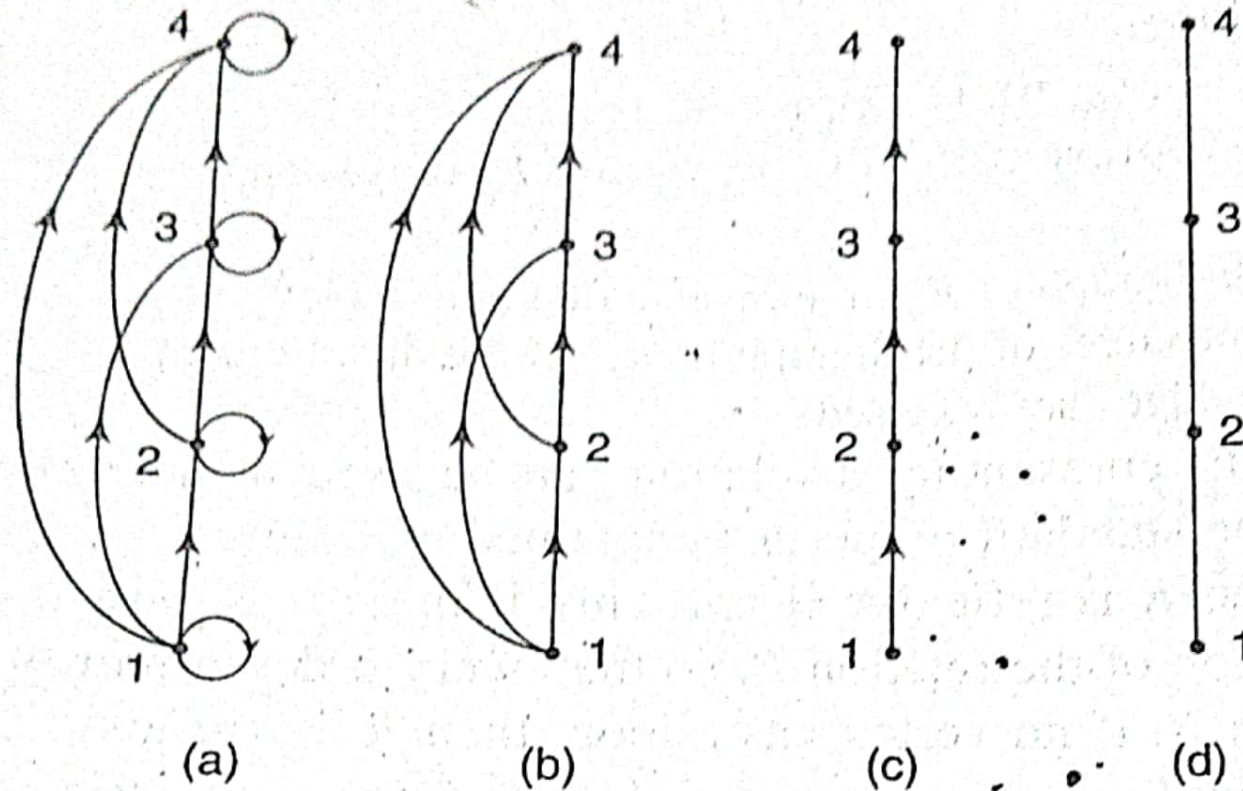
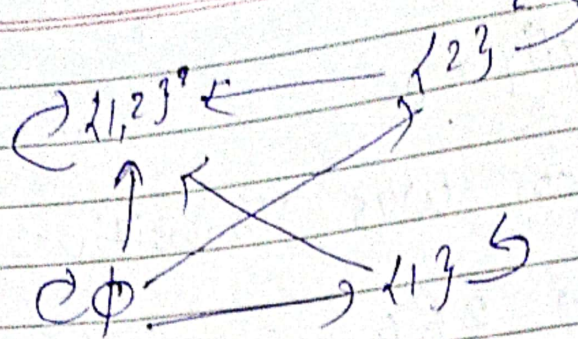


Fig. 2.15

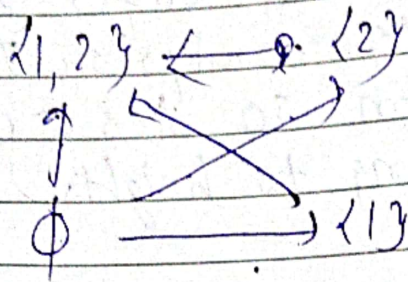
ex $A = \{1, 2\}$

$$P(A) = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\}$$

$$R = \{\{\emptyset, \{1\}\}, \{\emptyset, \{2\}\}, \{\emptyset, \{1, 2\}\}, \{\{1\}, \{1\}\}, \{\{1\}, \{1, 2\}\}, \\ \{\{2\}, \{2\}\}, \{\{2\}, \{1, 2\}\}, \{\{1, 2\}, \{1, 2\}\}\}$$

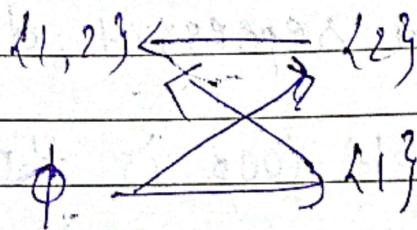


1) Remove Self loop (Reflexive)

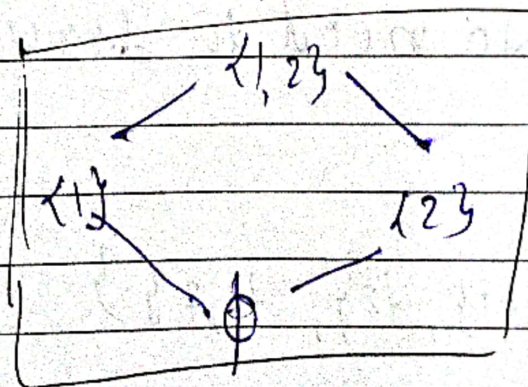


2) Remove Transitive edge.

$\phi R 1 \ \& \ 1 R \{1, 2\} \rightarrow \phi R \{1, 2\}$
 $\phi R 2 \ \& \ 2 R \{1, 2\} \rightarrow \phi R \{1, 2\}$



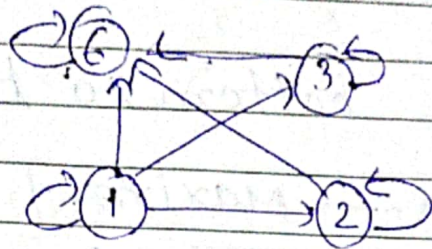
3) Remove direction



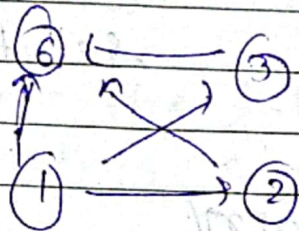
HASS Diagram

Ex. $A = \{1, 2, 3, 6\}$ and $R = a \text{ divides } b$.

$\rightarrow R = \{(1,1), (1,2), (1,3), (1,6), (2,2), (2,6), (3,3), (3,6), (6,6)\}$



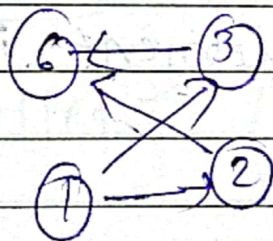
1) Remove self loop. (Reflexive)



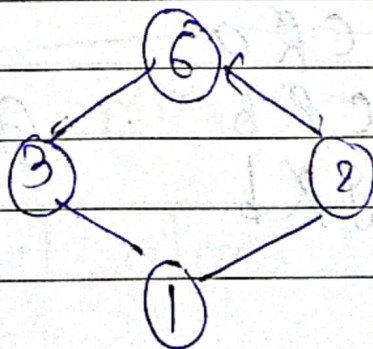
2) Remove Transitive edge.

$1R2, 2R6 \Rightarrow 1R6$

$1R3, 3R6 \Rightarrow 1R6$



3) Remove dissections.



\Rightarrow Hasse diagram