

Design and Analysis of Algorithms (CS206)

Assignment - 3

U19CS012

1. Given the following algorithms, answer the questions.

- Merge sort: Sorting Problem

Input: A Sequence of Unsorted 'n' numbers, a_1, a_2, \dots, a_n

Output: A Permutation (Reordering) (a_1', a_2', \dots, a_n') of Input Sequence such that $a_1' \leq a_2' \leq \dots \leq a_n'$

1. Write a program to sort an array, *arr*, consisting *n* numbers using the divide and conquer approach - Use only merge sort.

(1) The divide step should split the array into two (nearly) equal sub-arrays.

(2) The divide step should split the array into three (nearly) equal sub-arrays

1.1. (T) Write pseudocodes to design the algorithms for above mentioned computational problem. Both algorithms should sort the data by dividing them into two and three (nearly) equal sub-arrays respectively.

(1) The divide step should split the array into two (nearly) equal sub-arrays.

A.) MergeSort2 Function

```
• Merge_Sort2(arr, low, high)

1. if low < high
2.   int mid = low + (high-low)/2
   // Call this Function to Recursively Divide into Smaller Sub-array [l,m]
3.   Merge_Sort2(arr, low, mid);
   // Call this Function to Recursively Divide into Smaller Sub-array [m+1,h]
4.   Merge_Sort2(arr, mid + 1, high);
   // Merge the Both Sorted Array
5.   Merge2(arr, low, mid, high);
6. return
```

B.) Merge2 Function

```
// To Merge Two Sorted Array
• Merge2(arr, low, mid, high)
// Create a Temp Array of size high-low+1
1. tmp(high - low + 1, 0);
// Crawlers for Temp
2. i = low, j = mid + 1, k = 0;

3. while i <= mid AND j <= high
4.     if (arr[i] <= arr[j])
5.         tmp[k] = arr[i];
6.         k++;
7.         i++;
8.     else
9.         tmp[k] = arr[j];
10.        k++;
11.        j++;

// Remaining Elements in Second Interval
12. while i <= mid
13.     tmp[k] = arr[i];
14.     i++;
15.     k++;

// Remaining Elements in Second Interval
16. while j <= high
17.     tmp[k] = arr[j];
18.     j++;
19.     k++;

// Copy Temp Array to Original Array
20. for i = low to high
21.     arr[i] = tmp[i - low];
```

(2) The divide step should split the array into three (nearly) equal sub-arrays

A.) MergeSort3 Function

```
• Merge_Sort3(vll &arr, ll low, ll high)

// BASE CASE : 1 Element
1.  if (high - low < 2)
2.      return;

3.  mid1 = low + ((high - low) / 3);
4.  mid2 = low + 2 * ((high - low) / 3) + 1;

// Call this Function to Recursively Divide into Smaller Sub-array [l,m1)
5.  Merge_Sort3(arr, low, mid1);
// Call this Function to Recursively Divide into Smaller Sub-array [m1,m2)
6.  Merge_Sort3(arr, mid1, mid2);
// Call this Function to Recursively Divide into Smaller Sub-array [m1,high)
7.  Merge_Sort3(arr, mid2, high);
// Merge the Both Sorted Array
8.  Merge3(arr, low, mid1, mid2, high);
9.  return;
```

B.) Merge3 Function

```
// To Merge Two Sorted Array
• Merge3(arr, low, mid1, mid2, high)
// Create a Temp Array of size high-low+1
  tmp(high - low + 1, 0);

// Crawlers for Temp
  i = low, j = mid1, k = mid2, id = 0;

  while (i < mid1 && j < mid2 && k < high)
      if (arr[i] < arr[j])
          if (arr[i] < arr[k])
              tmp[id++] = arr[i++];
          else
              tmp[id++] = arr[k++];
      // arr[j] < arr[i]
      else
          if (arr[j] < arr[k])
              tmp[id++] = arr[j++];
          else
              tmp[id++] = arr[k++];
```

```

// i & j
while ((i < mid1) && (j < mid2))
    if (arr[i] < arr[j])
        tmp[id++] = arr[i++];
    else
        tmp[id++] = arr[j++];

// j & k
while ((j < mid2) && (k < high))
    if (arr[j] < arr[k])
        tmp[id++] = arr[j++];
    else
        tmp[id++] = arr[k++];

// i & k
while ((i < mid1) && (k < high))
    if (arr[i] < arr[k])
        tmp[id++] = arr[i++];
    else
        tmp[id++] = arr[k++];

// Copy Remaining Elements
// [0,mid1)
while (i < mid1)
    tmp[id++] = arr[i++];

// [mid1,mid2)
while (j < mid2)
    tmp[id++] = arr[j++];

// [mid2,high)
while (k < high)
    tmp[id++] = arr[k++];

// Copy Temp Array to Original Array
for (i = low; i < high; i++)
    arr[i] = tmp[i - low];

```

1.2. (T) Analyze the time complexity of both algorithms (split the array into two and three sub-arrays) using the recursion tree method.

A.) Merge Sort Analysis by Dividing into Two Parts

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Merge-sort (A, p, r)

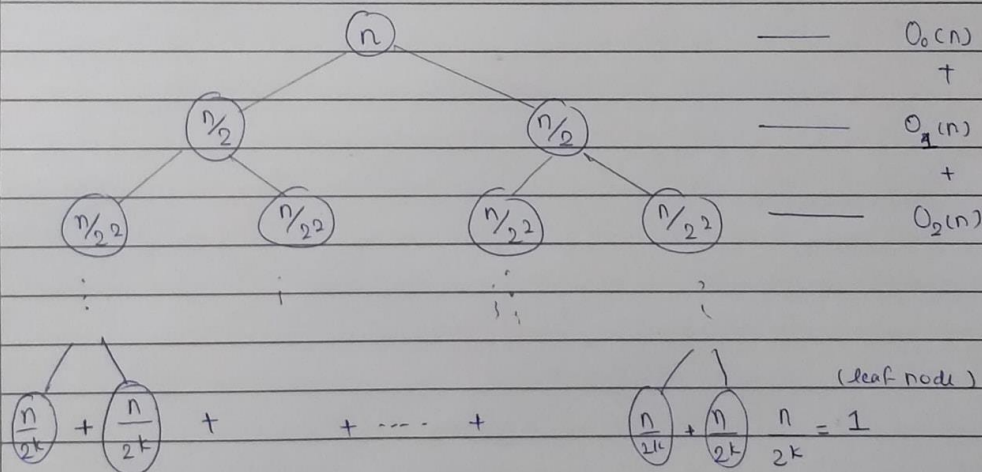
// T(n) (Assume)

(*) if $p < r$ + c
 $q = \lfloor (p+r)/2 \rfloor$ + c
 Merge-sort (A, p, q) + T(n/2)
 Merge-sort (A, q+1, r) + T(n/2)
 Merge (A, p, q, r) + $\Theta(n)$

$$T(n) = 2 \cdot T\left(\frac{n}{2}\right) + \Theta(n)$$

(\because in function merge
it take linear
time)

Recursion Tree



$$O\left(\sum_{i=0}^k 2^i \cdot \frac{n}{2^i}\right) = O\left(\sum_{i=0}^k n\right) = O(k \cdot n)$$

$$= O(\log_2(n) \times n)$$

[$k = \log_2(n)$]

\therefore Time complexity for Merge sort = $\Theta(n \log_2(n))$

B.) Merge Sort Analysis by Dividing into Three Parts

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Following the similar logic
(for 3 parts)

(*) Merge-sort3 (arr, low, high) // T(n)

if high - low < 2 + c

return

mid1 = low + ((high - low) / 3); + c

mid2 = low + 2 * ((high - low) / 3); + c

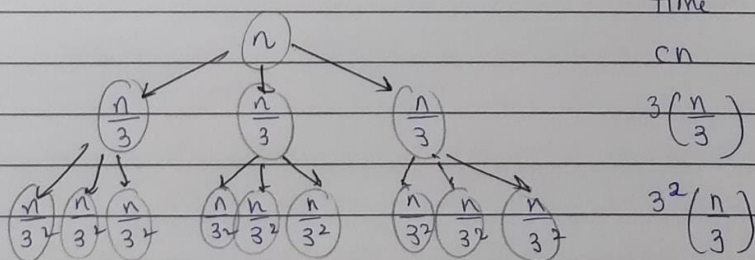
Merge-sort3 (arr, low, mid1); $T(n/3)$

Merge-sort3 (arr, mid1, mid2); $T(n/3)$

Merge-sort3 (arr, mid2, high); $T(n/3)$

Merge (arr, low, mid1, mid2, high); $O(n)$

$$T(n) = 3 T\left(\frac{n}{3}\right) + O(n)$$



leaf Node

$$\frac{n}{3^k} \dots \dots \dots \frac{n}{3^k}$$

$$\frac{n}{3^k} = 1 \quad [k = \log_3 n]$$

$$O\left(\sum_{i=0}^k 3^i \frac{n}{3^i}\right) = O\left(\sum_{i=0}^k n\right) = O(k \cdot n)$$

$$= O(\log_3 cn) \times n$$

∴ Time complexity for Merge sort = $O(n \log_3 cn)$

1.3. (L) Provide the details of Hardware/Software you used to implement algorithms and to measure the time.

Hardware Details of My Laptop:

PARAMETER	LAPTOP CONFIGURATION
Operating System	Microsoft Windows 10.0.19042
Processor	Intel(R) Core(TM) i5-10210U [Core i5 10th Gen]
CPU	1.60GHz, 2112 Mhz, 4 Core(s), 8 Logical Processor(s)
System Type	x64-based PC [64 Bit]
RAM	8.00 GB
Hard Drive/SSD	512 GB SSD

Software Used:

PARAMETER	LAPTOP CONFIGURATION
Code Editor	Visual Studio Code [Version 1.52]
Compiler	gcc (MinGW.org GCC-8.2.0-5) 8.2.0
Time	Measured using chrono Library in C++
Programming Language Used	C++

1.4. (L) Submit the code (complete programs).

A.) *Merge Sort Program* by Dividing into Two Parts

```
// HEADERS AND NAMESPACE
#include <bits/stdc++.h>
// INSTEAD OF ALL THESE
#include <iostream>
// For Creating File
#include <fstream>
#include <vector>
// For set - precision
#include <iomanip>
// For Time Calculation
#include <chrono>
// For File Name and Output File Name
```

```

#include <string>

using namespace std;
using namespace std::chrono;

// COMMONLY USED TYPES
typedef long long ll;
typedef vector<ll> vll;

// Basic Algorithm Implementation of Merge Sort
// To Merge Two Sorted Array
void merge(vll &arr, ll low, ll mid, ll high)
{
    // Create a Temp Array
    vll tmp(high - low + 1, 0);

    // Crawlers for Temp
    ll i = low, j = mid + 1, k = 0;

    while (i <= mid && j <= high)
    {
        if (arr[i] <= arr[j])
        {
            tmp[k] = arr[i];
            k++;
            i++;
        }
        else
        {
            tmp[k] = arr[j];
            k++;
            j++;
        }
    }

    // Remaining Elements in Second Interval
    while (i <= mid)
    {
        tmp[k] = arr[i];
        i++;
        k++;
    }

    // Remaining Elements in Second Interval
    while (j <= high)
    {
        tmp[k] = arr[j];
        j++;
        k++;
    }
}

```



```

    // Copy Temp Array to Original Array
    for (i = low; i <= high; i++)
    {
        arr[i] = tmp[i - low];
    }
}

// Real Merge Sort Function
void merge_sort(vll &arr, ll low, ll high)
{
    if (low < high)
    {
        ll mid = low + (high - low) / 2; // To Avoid Overflow
        // Call this Function to Recursively Divide into Smaller Sub-array [l,m]
        merge_sort(arr, low, mid);
        // Call this Function to Recursively Divide into Smaller Sub-array [m+1,h]
        merge_sort(arr, mid + 1, high);
        // Merge the Both Sorted Array
        merge(arr, low, mid, high);
    }

    return;
}

int main()
{
    // For Read & Write from "Input File" and Return Output to "Output" File
    freopen("output.txt", "w", stdout);

    // EDIT THIS FILE NUMBER , LIMIT and Number of Times File Runs
    int file_no = 1;
    int limit = 10;
    int each_file_runs = 2;

    for (; file_no <= limit; file_no++)
    {
        string inp_file = "File";
        string num = to_string(file_no);
        string ext = ".txt";
        inp_file += num;
        inp_file += ext;

        ifstream File;
        File.open(inp_file);

        vector<ll> arr;

        ll number, idx = 0;

```

```

while (!File.eof())
{
    File >> number;
    arr.push_back(number);
}

ll Best_Duration = 0, Worst_Duration = 0, Average_Duration = 0;
auto start = high_resolution_clock::now();
auto end = high_resolution_clock::now();
auto time_taken = duration_cast<nanoseconds>(end - start);
ll n1 = arr.size();
for (int f = 0; f < each_file_runs; f++)
{
    // -----AVERAGE CASE [ $O(n^2)$ ]-----

    start = high_resolution_clock::now();
    // Function Here
    merge_sort(arr, 0, arr.size() - 1);
    // Function Ends here
    end = high_resolution_clock::now();
    time_taken = duration_cast<nanoseconds>(end - start);
    Average_Duration += time_taken.count();

    // -----BEST CASE [ $O(n^2)$ ]-----
    // The Array is Already Sorted from Average Case, So it Becomes out Best Case
    // sort(arr.begin(), arr.end());
    start = high_resolution_clock::now();
    // Function Here
    merge_sort(arr, 0, arr.size() - 1);
    // Function Ends here
    end = high_resolution_clock::now();
    time_taken = duration_cast<nanoseconds>(end - start);
    Best_Duration += time_taken.count();

    // -----WORST CASE [ $O(n^2)$ ]-----
    // This will Reverse the Sorted Array, Therefore we will Get the Worst Case

    reverse(arr.begin(), arr.end());
    // sort(arr.begin(), arr.end(), greater<ll>());
    start = high_resolution_clock::now();
    // Function Here
    merge_sort(arr, 0, arr.size() - 1);
    // Function Ends here
    end = high_resolution_clock::now();
    time_taken = duration_cast<nanoseconds>(end - start);
    Worst_Duration += time_taken.count();
}

cout << "-----" << endl;
cout << inp_file << endl;

```

```

        cout << "AVERAGE CASE : ";
        double avg = (double)Average_Duration / (double)each_file_runs;
        avg *= 1e-9;
        cout << fixed << avg << setprecision(9);
        cout << " seconds" << endl;
        cout << "BEST CASE : ";
        double best = (double)Best_Duration / (double)each_file_runs;
        best *= 1e-9;
        cout << fixed << best << setprecision(9);
        cout << " seconds" << endl;
        cout << "WORST CASE : ";
        double worst = (double)Worst_Duration / (double)each_file_runs;
        worst *= 1e-9;
        cout << fixed << worst << setprecision(9);
        cout << " seconds" << endl;
    }

    return 0;
}

```

B.) Merge Sort Program by Dividing into Three Parts

```

// HEADERS AND NAMESPACE
#include <bits/stdc++.h>
// INSTEAD OF ALL THESE
#include <iostream>
// For Creating File
#include <fstream>
#include <vector>
// For set - precision
#include <iomanip>
// For Time Calculation
#include <chrono>
// For File Name and Output File Name
#include <string>

using namespace std;
using namespace std::chrono;

// COMMONLY USED TYPES
typedef long long ll;
typedef vector<ll> vll;

// Basic Algorithm Implementation of Merge Sort
// To Merge Two Sorted Array
void merge(vll &arr, ll low, ll mid1, ll mid2, ll high)
{
    // Create a Temp Array

```

```

vll tmp(high - low + 1, 0);

// Crawlers for Temp
ll i = low, j = mid1, k = mid2, id = 0;

while (i < mid1 && j < mid2 && k < high)
{
    if (arr[i] < arr[j])
    {
        if (arr[i] < arr[k])
        {
            tmp[id++] = arr[i++];
        }
        else
        {
            tmp[id++] = arr[k++];
        }
    }
    // arr[j] < arr[i]
    else
    {
        if (arr[j] < arr[k])
        {
            tmp[id++] = arr[j++];
        }
        else
        {
            tmp[id++] = arr[k++];
        }
    }
}

// i & j
while ((i < mid1) && (j < mid2))
{
    if (arr[i] < arr[j])
    {
        tmp[id++] = arr[i++];
    }
    else
    {
        tmp[id++] = arr[j++];
    }
}

// j & k
while ((j < mid2) && (k < high))
{
    if (arr[j] < arr[k])
    {
        tmp[id++] = arr[j++];
    }
}

```

```

    }
    else
    {
        tmp[id++] = arr[k++];
    }
}
// i & k
while ((i < mid1) && (k < high))
{
    if (arr[i] < arr[k])
    {
        tmp[id++] = arr[i++];
    }
    else
    {
        tmp[id++] = arr[k++];
    }
}

// Copy Remaining Elements
// [0,mid1)
while (i < mid1)
{
    tmp[id++] = arr[i++];
}
// [mid1,mid2)
while (j < mid2)
{
    tmp[id++] = arr[j++];
}
// [mid2,high)
while (k < high)
{
    tmp[id++] = arr[k++];
}

// Copy Temp Array to Original Array
for (i = low; i < high; i++)
{
    arr[i] = tmp[i - low];
}
}

// Real Merge Sort Function
void merge_sort(vll &arr, ll low, ll high)
{
    // BASE CASE : 1 Element
    if (high - low < 2)
    {
        return;
    }
}

```

```

}

ll mid1 = low + ((high - low) / 3);
ll mid2 = low + 2 * ((high - low) / 3) + 1;

// Call this Function to Recursively Divide into Smaller Sub-array [l,m1)
merge_sort(arr, low, mid1);
// Call this Function to Recursively Divide into Smaller Sub-array [m1,m2)
merge_sort(arr, mid1, mid2);
// Call this Function to Recursively Divide into Smaller Sub-array [m1,high)
merge_sort(arr, mid2, high);
// Merge the Both Sorted Array
merge(arr, low, mid1, mid2, high);

return;
}

int main()
{
    // For Read & Write from "Input File" and Return Output to "Output" File
    freopen("output.txt", "w", stdout);

    // EDIT THIS FILE NUMBER , LIMIT and Number of Times File Runs
    int file_no = 1;
    int limit = 10;
    int each_file_runs = 3;

    for (; file_no <= limit; file_no++)
    {
        string inp_file = "File";
        string num = to_string(file_no);
        string ext = ".txt";
        inp_file += num;
        inp_file += ext;

        ifstream File;
        File.open(inp_file);

        vector<ll> arr;

        ll number, idx = 0;
        while (!File.eof())
        {
            File >> number;
            arr.push_back(number);
        }

        ll Best_Duration = 0, Worst_Duration = 0, Average_Duration = 0;
        auto start = high_resolution_clock::now();
        auto end = high_resolution_clock::now();
    }
}

```



```

auto time_taken = duration_cast<nanoseconds>(end - start);
ll n1 = arr.size();
for (int f = 0; f < each_file_runs; f++)
{
    // -----AVERAGE CASE [ $O(n^2)$ ]-----

    start = high_resolution_clock::now();
    // Function Here
    merge_sort(arr, 0, arr.size());
    // Function Ends here
    end = high_resolution_clock::now();
    time_taken = duration_cast<nanoseconds>(end - start);
    Average_Duration += time_taken.count();

    // -----BEST CASE [ $O(n^2)$ ]-----
    // The Array is Already Sorted from Average Case, So it Becomes out Best Case
    // sort(arr.begin(), arr.end());
    start = high_resolution_clock::now();
    // Function Here
    merge_sort(arr, 0, arr.size() - 1);
    // Function Ends here
    end = high_resolution_clock::now();
    time_taken = duration_cast<nanoseconds>(end - start);
    Best_Duration += time_taken.count();

    // -----WORST CASE [ $O(n^2)$ ]-----
    // This will Reverse the Sorted Array, Therefore we will Get the Worst Case

    reverse(arr.begin(), arr.end());
    // sort(arr.begin(), arr.end(), greater<LL>());
    start = high_resolution_clock::now();
    // Function Here
    merge_sort(arr, 0, arr.size() - 1);
    // Function Ends here
    end = high_resolution_clock::now();
    time_taken = duration_cast<nanoseconds>(end - start);
    Worst_Duration += time_taken.count();
}

cout << "-----" << endl;
cout << inp_file << endl;
cout << "AVERAGE CASE : ";
double avg = (double)Average_Duration / (double)each_file_runs;
avg *= 1e-9;
cout << fixed << avg << setprecision(9);
cout << " seconds" << endl;
cout << "BEST CASE : ";
double best = (double)Best_Duration / (double)each_file_runs;
best *= 1e-9;
cout << fixed << best << setprecision(9);

```

```
    cout << " seconds" << endl;
    cout << "WORST CASE    : ";
    double worst = (double)Worst_Duration / (double)each_file_runs;
    worst *= 1e-9;
    cout << fixed << worst << setprecision(9);
    cout << " seconds" << endl;
}

return 0;
}
```

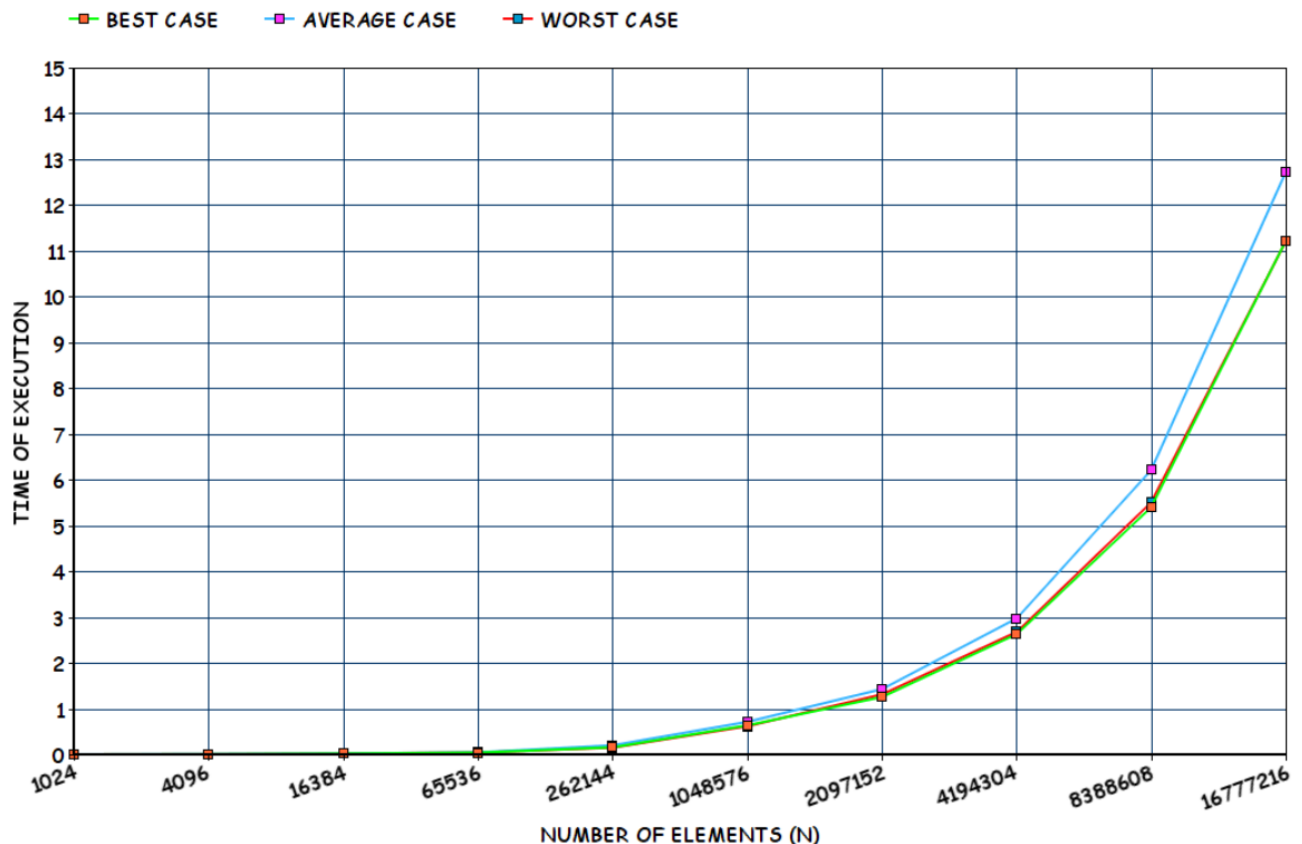
1.5. (L) Measure the best-case time, average-case time and worst-case time of the above two algorithms for all ten files (Assignment 1). Plot a graph.

A.) *Merge Sort Program* by Dividing into Two Parts

FILE	No. of Elements	BEST CASE [in sec]	AVERAGE CASE [in sec]	WORST CASE [in sec]
1	$1024 = 2^{10}$	0.000000000	0.000519000	0.000000000
2	$4096 = 2^{12}$	0.003997500	0.008046500	0.003997500
3	$16384 = 2^{14}$	0.017086500	0.017533500	0.017251500
4	$65536 = 2^{16}$	0.037438500	0.043737500	0.042738500
5	$262144 = 2^{18}$	0.154819000	0.198323000	0.147636500
6	$1048576 = 2^{20}$	0.634084000	0.711899500	0.616180000
7	$2097152 = 2^{21}$	1.255310000	1.427379000	1.311771000
8	$4194304 = 2^{22}$	2.629683500	2.973717500	2.677120500
9	$8388608 = 2^{23}$	5.407570500	6.219963500	5.513057500
10	$16777216 = 2^{24}$	11.222652500	12.724297000	11.212041500

*Worst Case = Reverse Sorted Array

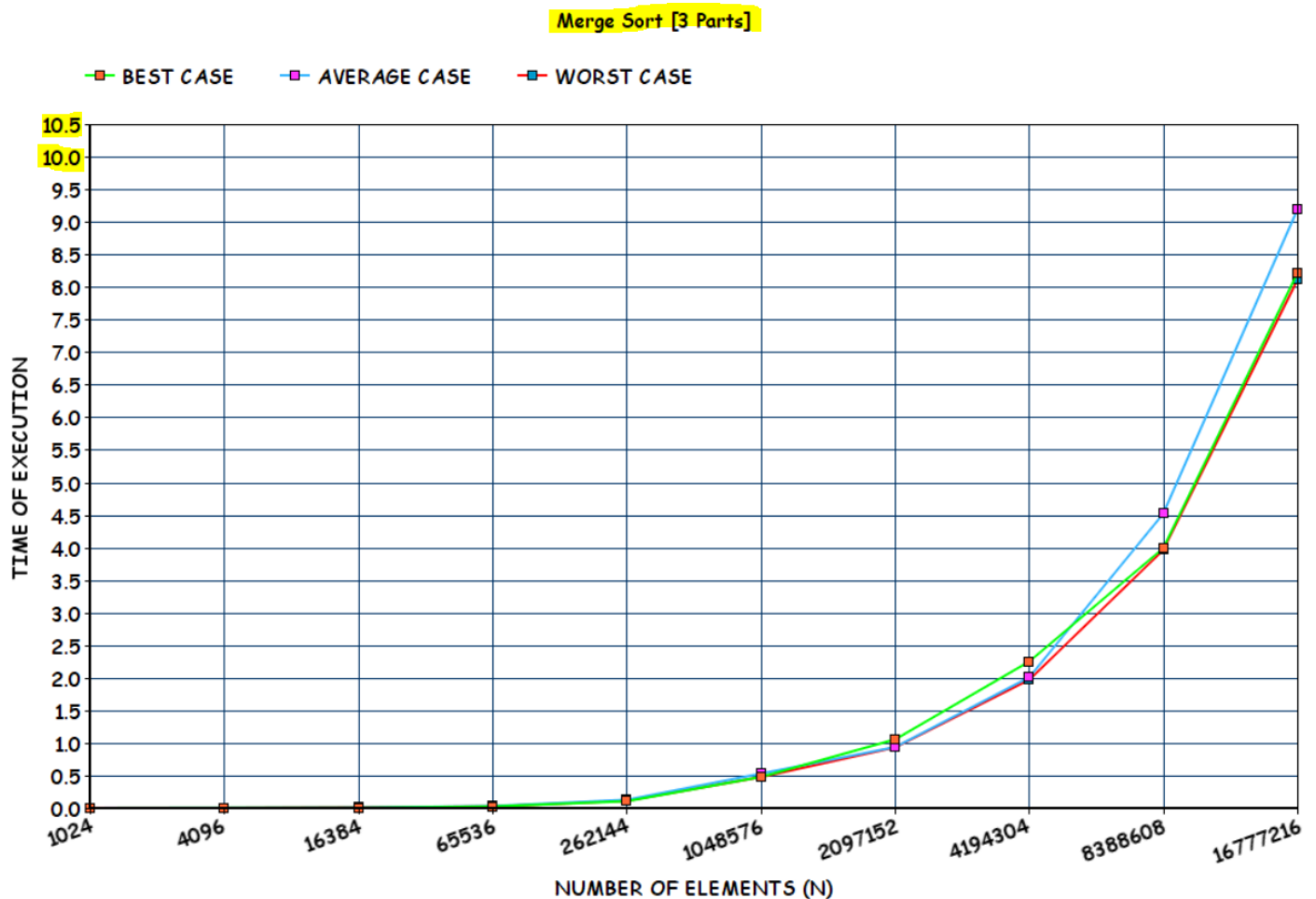
Merge Sort [2 Parts]



B.) Merge Sort Program by Dividing into Three Parts

FILE	No. of Elements	BEST CASE [in sec]	AVERAGE CASE [in sec]	WORST CASE [in sec]
1	$1024 = 2^{10}$	0.000000000	0.000332	0.001713000
2	$4096 = 2^{12}$	0.003036000	0.002839667	0.001595667
3	$16384 = 2^{14}$	0.003331667	0.007095667	0.010413333
4	$65536 = 2^{16}$	0.024336000	0.030064667	0.025325667
5	$262144 = 2^{18}$	0.105656333	0.126667667	0.111899000
6	$1048576 = 2^{20}$	0.477917333	0.528992333	0.476208667
7	$2097152 = 2^{21}$	1.054969667	0.936536667	0.931854000
8	$4194304 = 2^{22}$	2.252637667	2.009962000	1.972199000
9	$8388608 = 2^{23}$	3.994126333	4.524697667	3.966805333
10	$16777216 = 2^{24}$	8.216317333	9.190697000	8.110807000

*Worst Case = Reverse Sorted Array

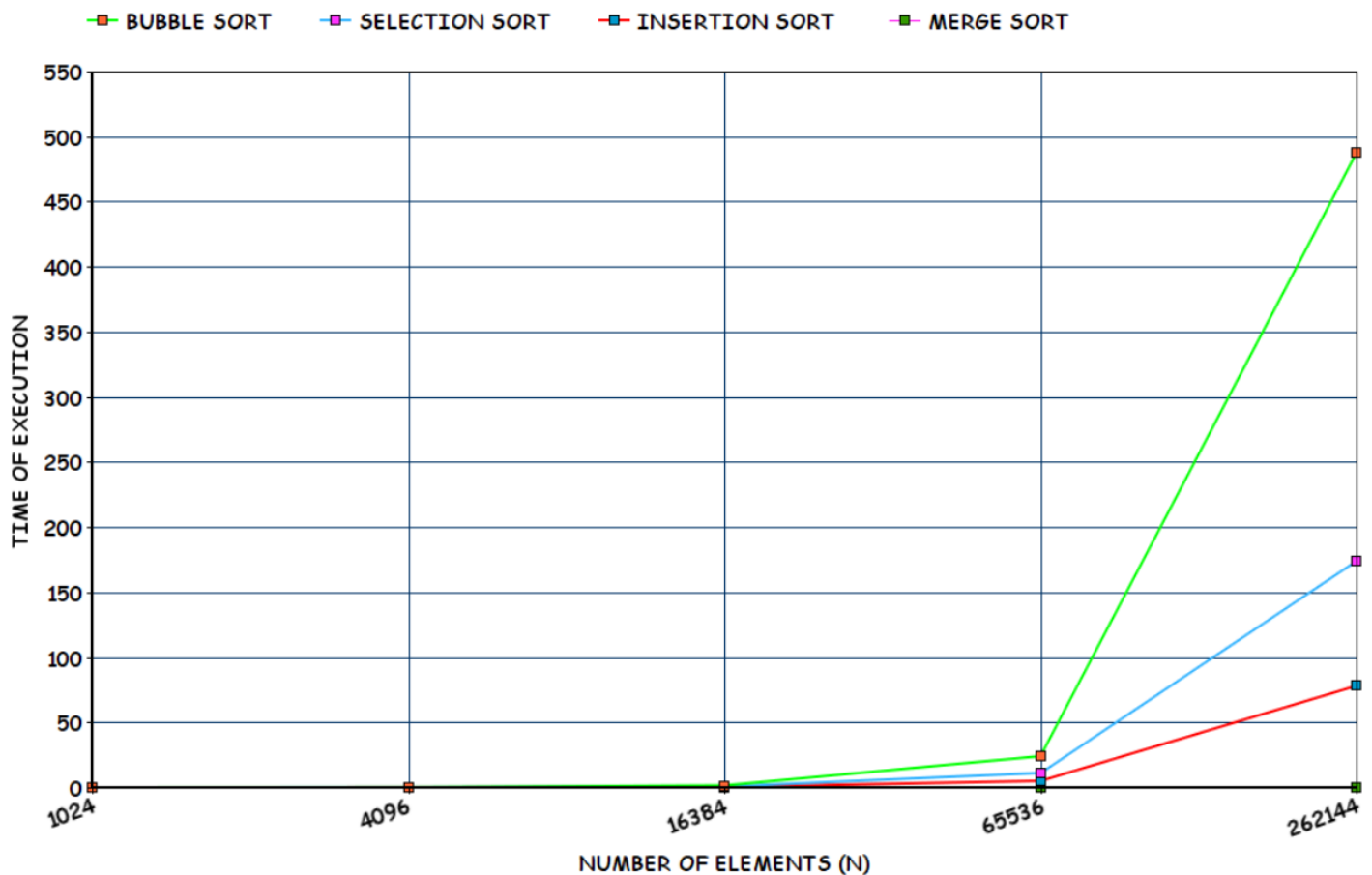


1.7. (L) Compare the average-case performance of bubble sort, selection sort, insertion sort, and merge sort for all ten files. Plot a graph.

AVERAGE CASE

File	No. Of Elements	Bubble Sort	Selection Sort	Insertion Sort	Merge Sort
1	2^{10}	0.01099400	0.002992000	0.000000000	0.000519000
2	2^{12}	0.111313000	0.050864000	0.038989000	0.008046500
3	2^{14}	1.501978000	0.669212000	0.345732500	0.017533500
4	2^{16}	24.017980000	11.059123000	4.971833000	0.043737500
5	2^{18}	487.501293000	174.081319000	78.119213000	0.198323000

AVERAGE CASE

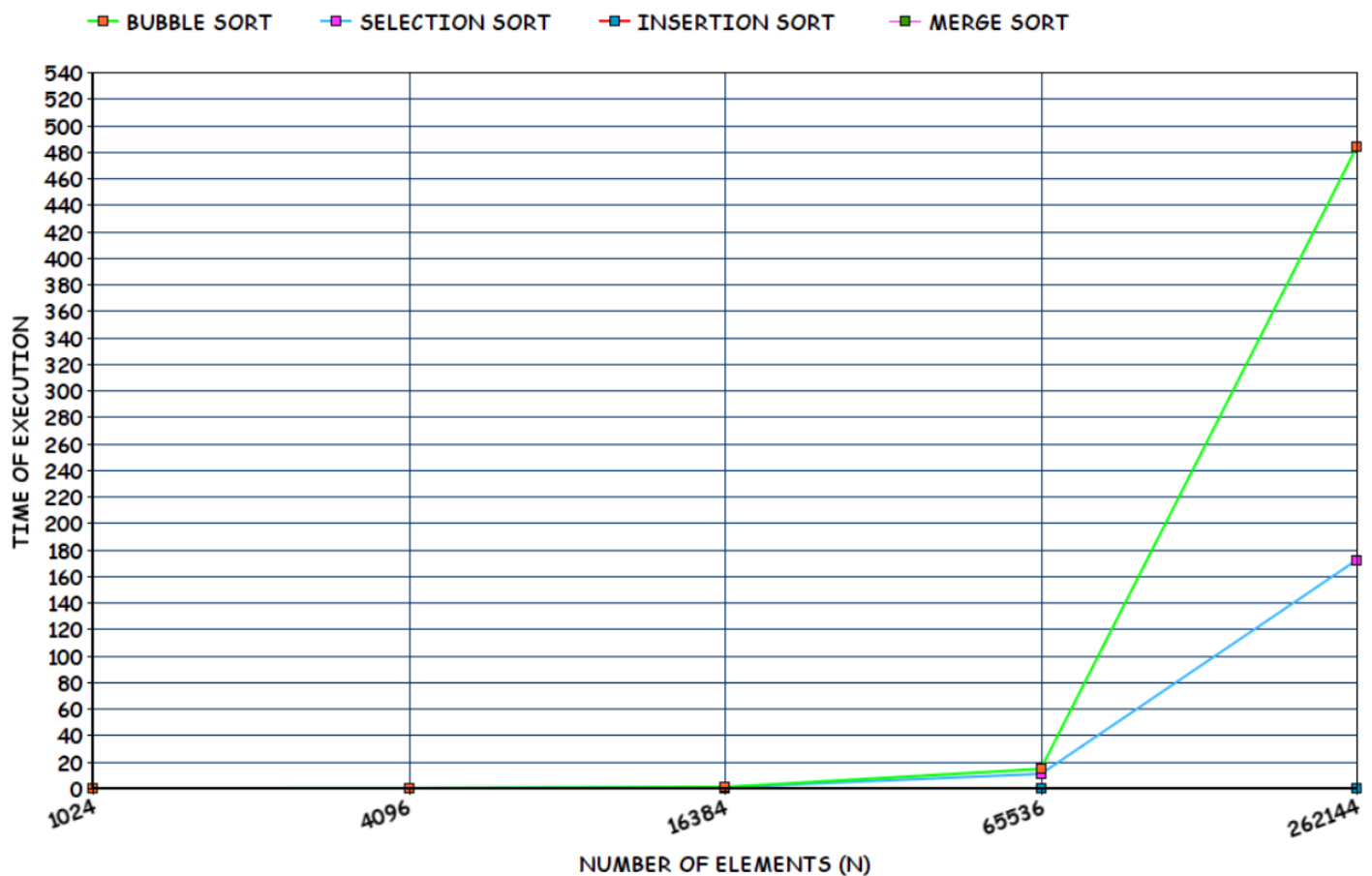


1.6. (L) Compare the best-case performance of bubble sort, selection sort, insertion sort, and merge sort for all ten files. Plot a graph.

BEST CASE

File	No. Of Elements	Bubble Sort	Selection Sort	Insertion Sort	Merge Sort
1	2^{10}	0.008177500	0.001993000	0.000000000	0.000000000
2	2^{12}	0.057657500	0.048870000	0.000000000	0.003997500
3	2^{14}	0.913967500	0.667217000	0.005000000	0.017086500
4	2^{16}	14.563756500	10.864665000	0.000000000	0.037438500
5	2^{18}	483.525254500	172.253910000	0.001998500	0.154819000

BEST CASE



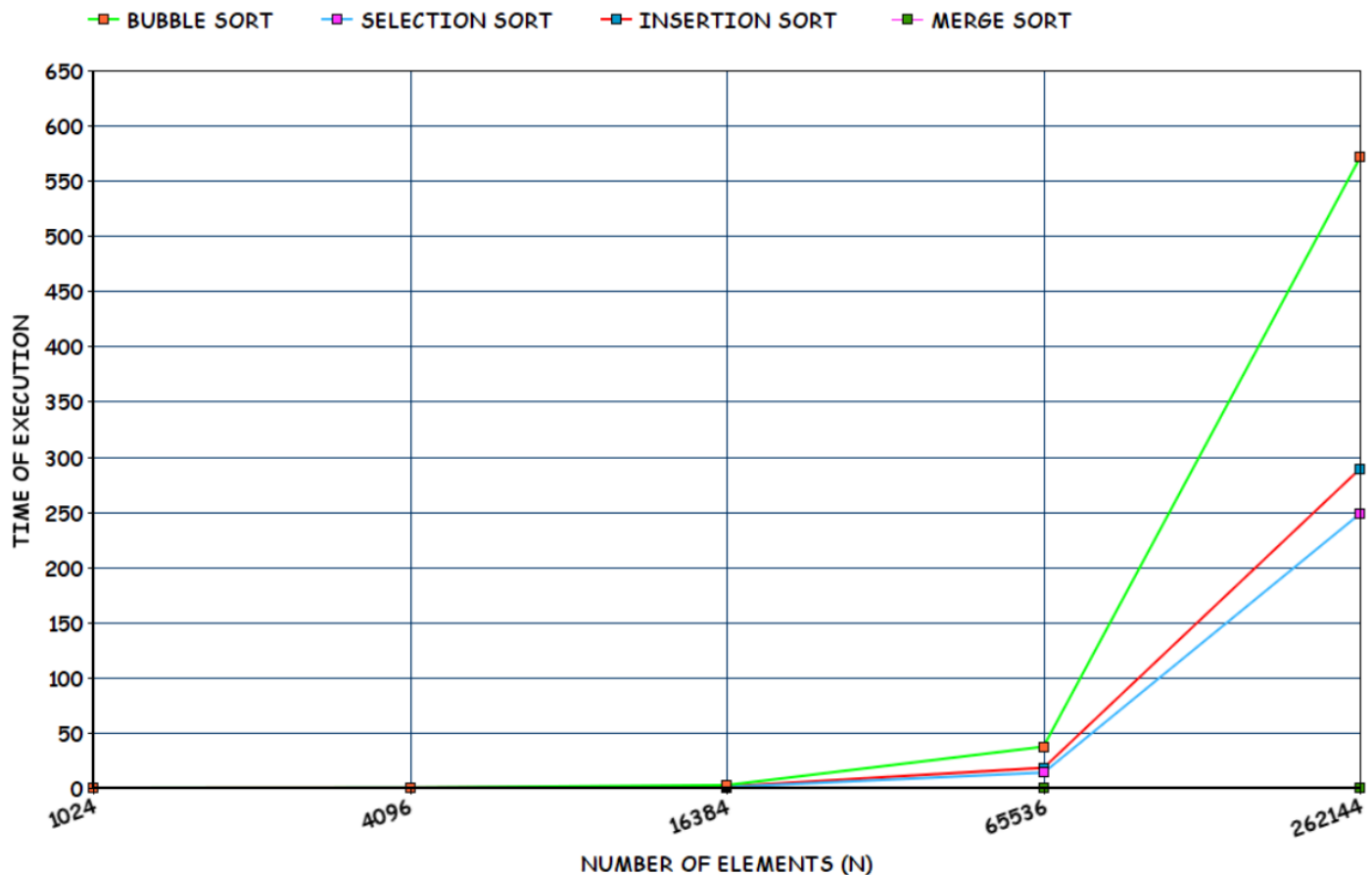
1.8. (L) Compare the worst-case performance of bubble sort, selection sort, insertion sort, and merge sort for all ten files. Plot a graph.

WORST CASE

File	No. Of Elements	Bubble Sort	Selection Sort	Insertion Sort	Merge Sort [Rev Sort]
1	2 ¹⁰	0.014323500	0.003026000	0.010075500	0.000000000
2	2 ¹²	0.142320000	0.045876000	0.086874000	0.003997500
3	2 ¹⁴	2.449665000	0.743014000	1.208551500	0.017251500
4	2 ¹⁶	37.231916500	13.791273000	18.156971000	0.042738500
5	2 ¹⁸	571.021661500	248.630703000	289.277434500	0.147636500

After File 5 Onwards, It would take a Minimum of 2 hrs for Each File Execution. So Avoided Executing for Rest of the Files.

WORST CASE



Note: Real Worst Case is Calculated in Following Manner:

The worst case of merge sort will be the one where merge sort will have to do **maximum number of comparisons**.

So I will try building the worst case in bottom up manner:

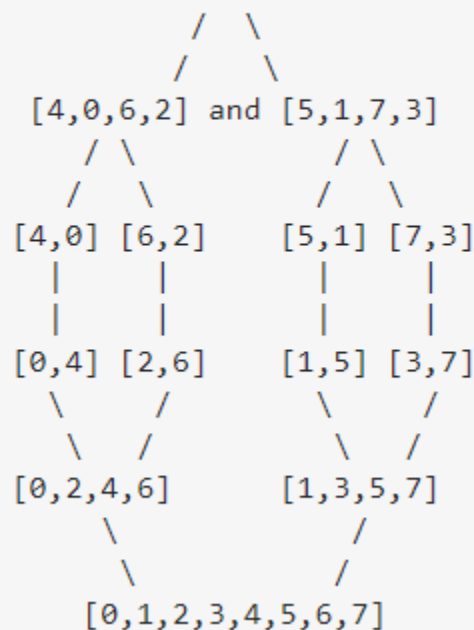
1. Suppose the array in final step after sorting is $\{0,1,2,3,4,5,6,7\}$
2. For worst case the array before this step must be $\{0,2,4,6,1,3,5,7\}$ because here left subarray = $\{0,2,4,6\}$ and right subarray = $\{1,3,5,7\}$ will result in maximum comparisons.
(Storing alternate elements in left and right subarray)

Reason: Every element of array will be compared atleast once.

3. Applying the same above logic for left and right subarray for previous steps : For array $\{0,2,4,6\}$ the worst case will be if the previous array is $\{0,4\}$ and $\{2,6\}$ and for array $\{1,3,5,7\}$ the worst case will be for $\{1,5\}$ and $\{3,7\}$.
4. Now applying the same for previous step arrays: *For worst cases:* $\{0,4\}$ must be $\{4,0\}$, $\{2,6\}$ must be $\{6,2\}$, $\{1,5\}$ must be $\{5,1\}$, $\{3,7\}$ must be $\{7,3\}$. Well if you look clearly this step is **not necessary** because if the size of set/array is 2 then every element will be compared atleast once even if array of size 2 is sorted.

Applying Merge Sort using Divide and Conquer

Input array arr[] = [4,0,6,2,5,1,7,3]



MERGE SORT ALL CASE [THEORATICAL CALCULATION]

FILE	NUMBER OF ELEMENTS	NO OF OPERATIONS $O(N \cdot \log_2(N))$	APPROX TIME TAKEN [OP/ 10^8]
FILE 1	$1024 = 2^{10}$	$1024 \cdot 10$	0.0001024
FILE 2	$4096 = 2^{12}$	$4096 \cdot 12$	0.00049152
FILE 3	$16384 = 2^{14}$	$16384 \cdot 14$	0.00229376
FILE 4	$65536 = 2^{16}$	$65536 \cdot 16$	0.01048576
FILE 5	$262144 = 2^{18}$	$262144 \cdot 18$	0.04718592
FILE 6	$1048576 = 2^{20}$	$1048576 \cdot 20$	0.20971520
FILE 7	$2097152 = 2^{21}$	$2097152 \cdot 21$	0.44040192
FILE 8	$4194304 = 2^{22}$	$4194304 \cdot 22$	0.92274688
FILE 9	$8388608 = 2^{23}$	$8388608 \cdot 23$	1.92937984
FILE 10	$16777216 = 2^{24}$	$16777216 \cdot 24$	4.02653184

For Bubble Sort, Selection Sort and Insertion Sort.

BEST CASE [THEORATICAL CALCULATION]

FILE	NUMBER OF ELEMENTS	NO OF OPERATIONS [CASE] = $O(N)$	APPROX TIME TAKEN [OP/ 10^8]
FILE 1	$1024 = 2^{10}$	1024	0.00001024
FILE 2	$4096 = 2^{12}$	4096	0.00004096
FILE 3	$16384 = 2^{14}$	16384	0.00016384
FILE 4	$65536 = 2^{16}$	65536	0.00065536
FILE 5	$262144 = 2^{18}$	262144	0.00262144
FILE 6	$1048576 = 2^{20}$	1048576	0.01048576
FILE 7	$2097152 = 2^{21}$	2097152	0.02097152
FILE 8	$4194304 = 2^{22}$	4194304	0.04194304
FILE 9	$8388608 = 2^{23}$	8388608	0.08388608
FILE 10	$16777216 = 2^{24}$	16777216	0.16777216

WORST/AVERAGE CASE [THEORATICAL CALCULATION]

FILE	NUMBER OF ELEMENTS	NO OF OPERATIONS [CASE] = $O(N^2)$	APPROX TIME TAKEN [OP/ 10^8]
FILE 1	$1024 = 2^{10}$	2^{20}	0.0104 seconds = 0.01 sec
FILE 2	$4096 = 2^{12}$	2^{24}	0.167 seconds = 0.16 sec
FILE 3	$16384 = 2^{14}$	2^{28}	2.684 seconds = 2.6 sec
FILE 4	$65536 = 2^{16}$	2^{32}	43 seconds = 43 sec
FILE 5	$262144 = 2^{18}$	2^{36}	687 seconds = 11 mins
FILE 6	$1048576 = 2^{20}$	2^{40}	10995 seconds = 3 hrs 3 mins
FILE 7	$2097152 = 2^{21}$	2^{42}	43980 seconds = 12 hrs 13 mins
FILE 8	$4194304 = 2^{22}$	2^{44}	175922 seconds = 2 days 52 hrs 2 mins
FILE 9	$8388608 = 2^{23}$	2^{46}	703687 seconds = 8 days 3 hrs 28 mins
FILE 10	$16777216 = 2^{24}$	2^{48}	2814750 seconds = 32 days 13 hrs 52 mins

CONCLUSION:

Bubble sort: repeatedly compare neighbor pairs and swap if necessary.

Selection sort: repeatedly pick the smallest element to append to the result.

Insertion sort: repeatedly add new element to the sorted result.

Sorting Algorithm	Time Complexity			Space Complexity
	Best Case	Average Case	Worst Case	Worst Case
Bubble Sort	$O(N)$	$O(N^2)$	$O(N^2)$	$O(1)$
Selection Sort	$O(N^2)$	$O(N^2)$	$O(N^2)$	$O(1)$
Insertion Sort	$O(N)$	$O(N^2)$	$O(N^2)$	$O(1)$

Merge Sort:

- 1.) If it is only one element in the list it is already sorted, return.
- 2.) **Divide** the list recursively into two halves until it can no more be divided.
- 3.) **Merge** the smaller lists into new list in sorted order.

- ✓ Merge Sort is useful for sorting linked lists.
- ✓ Merge Sort is a stable sort which means that the same element in an array maintain their original positions with respect to each other.
- ✓ Overall time complexity of Merge sort is $O(n \log n)$.
It is more efficient as it is in worst case also the runtime is $O(n \log n)$
- ✓ The space complexity of Merge sort is $O(n)$. [Not In-Place]
- ✓ This means that this algorithm takes a lot of space and may slower down operations for the last data sets.

SUBMITTED BY:

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