Delays

- Each instruction passes through different combinations of Fetch, Memory Read, and Memory Write cycles.
- Knowing the combinations of cycles, one can calculate how long such an instruction would require to complete.
- B for Number of Bytes
- M for Number of Machine Cycles
- T for Number of T-State.

Delays

- Knowing how many T-States an instruction requires, and keeping in mind that a T-State is one clock cycle long, we can calculate the time using the following formula:
- Delay = No. of T-States / Frequency
- For example a "MVI" instruction uses 7 T-States. Therefore, if the Microprocessor is running at 2 MHz, the instruction would require 3.5 μSeconds to complete.

Delay loops

- We can use a loop to produce a certain amount of time delay in a program.
- The following is an example of a delay loop:

	MVI C, FFH	7 T-States
LOOP	DCR C	4 T-States
	JNZ LOOP	10 T-States

- The first instruction initializes the loop counter and is executed only once requiring only 7 T-States.
- The following two instructions form a loop that requires 14 T-States to execute and is repeated 255 times until C becomes 0

Delay Loops

- We need to keep in mind though that in the last iteration of the loop, the JNZ instruction will fail and require only 7 T-States rather than the 10.
- Therefore, we must deduct 3 T-States from the total delay to get an accurate delay calculation.
- To calculate the delay, we use the following formula:
 Tdelay= TO+ TL

Tdelay= total delay, TO= delay outside the loop, TL= delay of the loop

TO is the sum of all delays outside the loop.

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Clock frequency of the system f = 2 MHz Clock period $T = 1/f = 1/2 \times 10^{-6} = 0.5 \mu s$ Time to execute MVI = 7 T-states × 0.5 = 3.5 μs

The time delay in the loop T_L with 2 MHz clock frequency is calculated as

$$TL = (T \times Loop T-states \times N10)$$

where T_L = Time delay in the loop

T = System clock period

 N_{10} = Equivalent decimal number of the hexadecimal count loaded in the delay register

$$T_L = (0.5 \times 10^{-6} \times 14 \times 255)$$

= 1785 \mus
\approx 1.8 ms

$$T_{LA} = T_L - (3 \text{ T-states} \times \text{Clock period})$$

= 1785.0 µs - 1.5 µs = 1783.5 µs

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$$Total Delay = \frac{Time to execute instructions}{outside loop} + \frac{Time to execute}{loop instructions}$$

$$T_D = T_O + T_{LA}$$

= $(7 \times 0.5 \ \mu s) + 1783.5 \ \mu s = 1787 \ \mu s$
 $\approx 1.8 \ ms$

Using a Register Pair as a Loop Counter

- Using a single register, one can repeat a loop for a maximum count of 255 times.
- It is possible to increase this count by using a register pair for the loop counter instead of the single register.
 - A minor problem arises in how to test for the final count since DCX and INX do not modify the flags.
 - However, if the loop is looking for when the count becomes zero, we can use a small trick by ORing the two registers in the pair and then checking the zero flag.

Using a Register Pair as a Loop Counter

 The following is an example of a delay loop set up with a register pair as the loop counter.

LXI B, 1000H 10 T-States

LOOP DCX B 6 T-States

MOV A, C 4 T-States

ORA B 4 T-States

JNZ LOOP 10 T-States

Label	Opcode	Operand	Comments	T-states
	LXI	В,2384Н	;Load BC with 16-bit count	10
LOOP:	DCX	В	;Decrement (BC) by one	6
	MOV	A,C	;Place contents of C in A	4
	ORA	В	;OR (B) with (C) to set Zero flag	4
	JNZ	LOOP	;If result ≠ 0, jump back to LOOP	10/7

The time delay in the loop is calculated as in the previous example. The loop includes four instructions: DCX, MOV, ORA, and JNZ, and takes 24 clock periods for execution. The loop is repeated 2384H times, which is converted to decimals as

$$2384H = 2 \times (16)^3 + 3 \times (16)^2 + 8 \times (16)^1 + 4(16^0)$$

= 9092₁₀

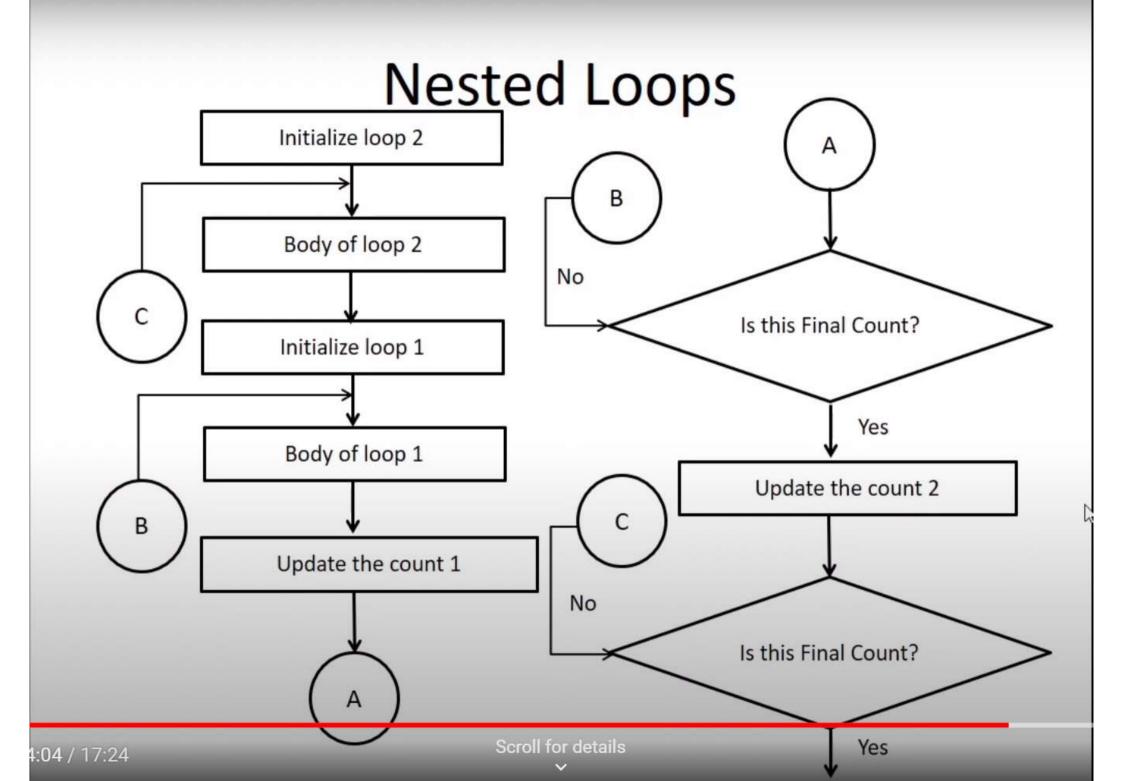
If the clock period of the system = $0.5 \mu s$, the delay in the loop T_L is

$$T_L = (0.5 \times 24 \times 9092_{10})$$

 $\approx 109 \text{ ms (without adjusting for the last cycle)}$
Total Delay $T_D = 109 \text{ ms} + T_O$
 $\approx 109 \text{ ms (The instruction LXI adds only 5 } \mu \text{s.)}$

Nested Loops

- Nested loops can be easily setup in Assembly language by using two registers for the two loop counters and updating the right register in the right loop.
- In the figure, the body of loop2 can be before or after loop1.



	MVI B,38H	7T
LOOP2:	MVI C,FFH	7T
LOOP1:	DCR C	4T
	JNZ LOOP1	10/7T
	DCR B	4T
	JNZ LOOP2	10/7T

The delay in LOOP1 is $T_{L1} = 1783.5 \,\mu s$. These calculations are shown in Section 8.1.1. We can replace LOOP1 by T_{L1} , as shown in Figure 8.3(b). Now we can calculate the delay in LOOP2 as if it is one loop; this loop is executed 56 times because of the count (38H) in register B:

$$T_{L2} = 56(T_{L1} + 21 \text{ T-states} \times 0.5 \text{ }\mu\text{s})$$

= 56(1783.5 \text{ }\mu s + 10.5 \text{ }\mu s)
= 100.46 \text{ ms}

Nested Loops for Delay

 Instead (or in conjunction with) Register Pairs, a nested loop structure can be used to increase the total delay produced.

	MVI B, 10H	7 T-States
LOOP2	MVI C, FFH	7 T-States
LOOP1	DCR C	4 T-States
	JNZ LOOP1	10 T-States
	DCR B	4 T-States
	JNZ LOOP2	10 T-States

Delay Calculation of Nested Loops

- The calculation remains the same except that it the formula must be applied recursively to each loop.
 - Start with the inner loop, then plug that delay in the calculation of the outer loop.
- Delay of inner loop
 - TO1= 7 T-States
- MVI C, FFH instruction
 - TL1= (255 X 14) -3 = 3567 T-States
- 14 T-States for the DCR C and JNZ instructions repeated
 255 times (FF16= 25510) minus 3 for the final JNZ

Increasing the delay

- The delay can be further increased by using register pairs for each of the loop counters in the nested loops setup.
- It can also be increased by adding dummy instructions (like NOP) in the body of the loop.