Operating System (CS301)

Assignment - 7 **U19CS012**

1. Consider this code example for allocating and releasing processes:

```
1. #define MAX_PROCESSES 255
2.
3. int numberOfProcesses = 0;
int allocateProcess()
7. {
    int newPid;
    if (numberOfProcesses == MAX_PROCESSES)
            return -1;
11.
        else
12.
        {
13.
14.
            ++numberOfProcesses;
15.
            return newPid;
16.
        }
17. }
18.
19. /*the implementation of exit() calls this function */
20. void releaseProcess()
21. {
22.
        --numberOfProcesses;
23.
24. }
```

A) Identify the race condition(s). Be specific — refer to the code.

Race condition occurs when <u>multiple threads</u> read and write the <u>same variable</u> i.e. they have access to some shared data and they try to change it at the same time. In such a scenario threads are "racing" each other to access/change the data.

Data Involved in Race Condition: The Variable "numberOfProcesses"

The Location in the Code where the Race Condition occurs:

```
Line Number 14 - 14. ++number 0fProcesses;

Line Number 23 - 23. --number 0fProcesses;
```

The code that **decrements** number_of_processes and the code that **increments** number_of_processes are the statements that could be involved in Race Conditions.

B) Assume that you have a mutex lock named mutex with the operations acquire () and release(). Indicate where in the code above that the locking/unlocking needs to be placed to prevent the race condition(s).

Method 1

```
int numberOfProcesses = 0;
/*the implementation of fork() calls this function */
int allocateProcess()
    try
       mutex.acquire();
        int newPid;
        if (numberOfProcesses == MAX_PROCESSES)
            return -1;
        else
        { /*allocate necessary process resources */
            ++numberOfProcesses;
            return newPid;
        }
    }
    finally
        mutex.release();
    }
void releaseProcess()
{
    try
    {
        mutex.acquire();
        /*release process resources */
        --numberOfProcesses;
    finally
       mutex.release();
```

Method 2

The programmer can place the two operations at the very first and end of a function call as well

The acquire() function should be placed in the beginning of function call.

Whereas, release() operation call should be placed just before the end of function call.

2. Consider how to implement a mutex lock using an <u>Atomic Hardware</u>

<u>Instruction</u>. Assume that the following structure defining the mutex lock is available:

```
typedef struct {
int unavailable;
} lock;
```

```
(unavailable == 0) indicates that the lock is <u>available</u> & (unavailable == 1) indicates that the lock is unavailable.
```

Using this struct, illustrate how the following functions can be implemented using the test_and_set() instruction and compare_and_swap() instructions:

- void acquire(lock *mutex)
- void release(lock *mutex)

Be sure to include any initialization that may be necessary.

```
typedef struct
{
    int unavailable;
} lock;

// (1) Initialization
void init(lock *mutex)
{
    // unavailable == 0 -> lock is available,
    // unavailable == 1 -> lock unavailable
    mutex->unavailable = 0;
}
```

```
int test_and_set(int *target)
    int rv = *target;
    *target = true;
    return rv;
void acquire(lock *mutex)
    while (test_and_set(&mutex->unavailable) != 0)
void release(lock *mutex)
    mutex->unavailable = 0;
int compare_and_swap(int *value, int expected, int new_value)
    int temp = *value;
    if (*value == expected)
        *value = new_value;
    return temp;
void acquire(lock *mutex)
    while (compare_and_swap(&mutex->unavailable, 0, 1) != 0)
    return;
void release(lock *mutex)
    mutex->unavailable = 0;
```

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