

## UNIT - 2

# Concept of Boolean Algebra

### Introduction:

→ In 1854 George Boolean introduced a systematic treatment of logic and developed for this purpose of an algebraic system Now called Boolean Algebra.

→ In 1938 C.E. Shannon introduced a two-valued Boolean Algebra called a switching Algebra

### Fundamental Postulates of Boolean Algebra

S.NO	Postulates	Comments
1.	→ Result of each operator is either 0 or 1	$1, 0 \in B$
2.	a) $0+0=0, 0+1=1, 1+0=1$ b) $1 \cdot 1 = 1; 0 \cdot 1 = 0, 1 \cdot 0 = 0$	Identity element "0" for "+" and "1" for $\cdot$
3.	a) $(A+B) = (B+A)$ b) $(A \cdot B) = (B \cdot A)$	Commutative law
4.	a) $A \cdot (B+C) = (A \cdot B) + (A \cdot C)$ b) $A + (B \cdot C) = (A + B) \cdot (A + C)$	Distributive law
5.	a) $A + \bar{A} = 1, 0 + \bar{0} = 1, 1 + \bar{1} = 1$ b) $A \cdot \bar{A} = 0, \text{ if } \bar{0} \cdot \bar{0} = 0 \cdot 1 = 0$ $1 \cdot \bar{1} = 1 \cdot 0 = 0$	Complement

$$[\bar{0} = 1]$$

$$\bar{1} = 0]$$

## Basic theorems and properties

### Duality:

→ The principle of Duality theorem says that starting with a boolean relation, we can derive another boolean relation by

1. changing each "OR" sign to an "AND" sign

2. changing each "AND" sign to an "OR" sign

3. Any 0 or 1 operating in the expression of boolean function.

Ex: Dual of relation  $A + \bar{A} = 1$  is  $A \cdot \bar{A} = 0$

$$A + \bar{A} = 1$$

$$A \cdot \bar{A} = 0$$

### Theorems:

#### Laws of Boolean algebra:

##### Boolean Expression

$$\Rightarrow A + 1 = 1$$

$$A \cdot 0 = 0$$

$$\Rightarrow A + 0 = A$$

$$A \cdot 1 = A$$

$$\bar{\bar{A}} = A$$

$$\Rightarrow A + A = A$$

$$A \cdot A = A$$

$$\bar{A} + A$$

$$\Rightarrow A = A$$

$$\bar{A} + \bar{A} = A$$

##### B.A law (or) Rule

Annulation

Identity

Idempotent

Double Negation

## Boolean Expression

$$\Rightarrow A + \bar{A} = 1$$

Complement

$$A \cdot \bar{A} = 0$$

Commutative law

$$\Rightarrow A + B = B + A$$

$$A \cdot B = B \cdot A$$

De-morgan

$$\Rightarrow \overline{A+B} = \bar{A} \cdot \bar{B}$$

$$\overline{AB} = \bar{A} + \bar{B}$$

Boolean Algebra functions

Function	Description	Expression
1.	NULL	0
2.	Identity	1
3.	INPUT A	A
4.	INPUT B	B
5.	NOT A	$\bar{A}$
6.	NOT B	$\bar{B}$
7.	A AND B (AND)	$A \cdot B$
8.	A AND NOT B	$A \cdot \bar{B}$
9.	NOT A AND B	$\bar{A} \cdot B$
10.	NOT AND (NAND)	$\overline{A \cdot B}$ $A \cdot A = 0$ $A \cdot A \cdot A = 0$
11.	A OR B (OR)	$A + B$ $A + A \cdot B$ $A + \bar{B}$
12.	A OR NOT B	$A + \bar{B}$

Function	Description	Expression
13.	NOT A OR B (OR)	$A + B$
14.	NOT OR (NOR)	$\overline{A + B}$
15.	Exclusive OR (X-OR)	$AB + \overline{A}B$
16.	Exclusive NOR (X-NOR)	$\overline{AB} + \overline{\overline{A}B}$

### Theorems:

① a)  $A + A = A$

b)  $A \cdot A = A$

② a)  $A + 1 = 1$

b)  $A \cdot 0 = 0$

③ a)  $\overline{\overline{A}} = A$

④ a)  $A + AB = A$

b)  $A(A+B) = A$

⑤ a)  $A + \overline{A}B = A + B$

b)  $A(\overline{A} + B) = AB$

① Proof:

L.H.S

$$\begin{aligned}
 A + A &= (A+A) \cdot (1) \\
 &= (A+A)(A+\overline{A}) \\
 &= AA + A\overline{A} + AA + A\overline{A} \\
 &= AA + A\overline{A} \\
 &= A(A+\overline{A}) \\
 &= A // 
 \end{aligned}$$

$$\left\{ \begin{array}{l} \text{if } A = 0; 0+0=0 \\ \text{if } A = 1; 1+1=1 \end{array} \right.$$

$$[1 = A + \overline{A}]$$

② Proof:

L.H.S

$$\begin{aligned}
 A \cdot A &= A \cdot A + (0) \\
 &= A \cdot A + A \cdot \overline{A} \\
 &= A(CA + \overline{A}) \\
 &= A(A + \overline{A}) \\
 &= A //
 \end{aligned}$$

$$\left\{ \begin{array}{l} A \cdot A = A \\ A = 0 \Rightarrow 0 \cdot 0 = 0 \\ A = 1 \Rightarrow 1 \cdot 1 = 1 \end{array} \right.$$

$$[A \cdot \overline{A} = 0]$$

$$② @ A+1=1$$

Proof:  $A+1 = 1 \cdot A + 1$

$$\begin{aligned} &= (A+\bar{A}) \cdot (A+1) \\ &= A \cdot A + A + \bar{A} \cdot A + \bar{A} \\ &\quad \text{(using } A \cdot A = A \text{ and } \bar{A} \cdot A = 0) \\ &= A + A + 0 + \bar{A} \\ &= A + A + \bar{A} \\ &= A + \bar{A} \\ &= 1 \end{aligned}$$

$$\begin{aligned} A \cdot \bar{A} &= 0 \\ A + \bar{A} &= 1 \end{aligned}$$

$$\left\{ \begin{array}{l} \text{if } A=0; 0+1=1 \\ A=1; 1+1=1 \end{array} \right.$$

$$(b) A \cdot 0 = 0$$

Proof:  $A \cdot 0 = 1 \cdot A \cdot (A \cdot \bar{A})$

$$\begin{aligned} &= (A + \bar{A}) \cdot A \cdot (A \cdot \bar{A}) \\ &= A \cdot A \cdot \bar{A} \\ &= A \cdot 0 \\ &= 0 \end{aligned}$$

$$\begin{aligned} A \cdot 1 &= A \\ A \cdot A \cdot \bar{A} &= A \\ A \cdot 0 &= 0 \\ 1 \cdot 1 \cdot 0 &= 0 \\ A \cdot 0 &= 0 \\ A = 0; 0 \cdot 0 = 0 &= 0 \\ A = 1; 1 \cdot 0 = 0 &= 0 \end{aligned}$$

$$③ \bar{\bar{A}} = A$$

Proof: If  $A=0; \bar{0} = \bar{1} = 0$   $\bar{0} = 1$   
 $A=1; \bar{1} = \bar{0} = 1$   $\bar{1} = 0$

$$④ @ A+AB=A$$

Proof:  $A+AB = A(1+B)$   $[ \because 1+B=1 ]$

$$\begin{aligned} &= A(1) \\ &= A \end{aligned}$$

$$(b) A(A+B)=A$$

Proof:  $A(A+B) = (A+AB)(A+B)$

$$\begin{aligned} &= AA+AB+ABA+AB \cdot B \\ &= A+AB+A \cdot AB+AB \\ &= A+AB+AB+AB \end{aligned}$$

$$= A(1+B) + AB$$

$$= A + AB$$

$$= A(1+B)$$

$$= A \quad //$$

(5)(a)  $A + \bar{A}B = A + B$

proof:  $A + \bar{A}B = (A + \bar{A})(A + B)$

$$= A + B \quad [\because A + \bar{A} = 1]$$

(b)  $A(\bar{A} + B) = AB$

proof:  $A(\bar{A} + B) = A \cdot \bar{A} + AB$

$$= 0 + AB$$

$$\therefore A(\bar{A} + B) = AB \quad //$$

\*\* De-Morgan's theorem:

(i)  $\bar{AB} = \bar{A} + \bar{B}$

(ii)  $\bar{A+B} = \bar{A} \cdot \bar{B}$

(iii)  $\bar{AB} = \bar{A} + \bar{B}$

A	B	$\bar{A}$	$\bar{B}$	$\bar{AB}$	$\bar{A} + \bar{B}$
0	0	1	1	1	1
0	1	1	0	1	1
1	0	0	1	0	1
1	1	0	0	0	0

(iv)  $\bar{A+B} = \bar{A} \cdot \bar{B}$

A	B	$\bar{A}$	$\bar{B}$	$\bar{A} \cdot \bar{B}$	$\bar{A+B}$
0	0	1	1	1	1
0	1	1	0	0	0
1	0	0	1	0	0
1	1	0	0	0	0

→ Consensus theorem :-

→ In simplification of Boolean Expression, an expression of the form  $AB + \bar{A}C + BC$ . The term  $BC$  is redundant and can be eliminated to form the equivalent expression  $AB + \bar{A}C$ .

→ The theorem is used for this simplification is known as Consensus theorem. and it is stated as

$$AB + \bar{A}C + BC = AB + \bar{A}C$$

Proof :-

$$L.H.S \Rightarrow AB + \bar{A}C + BC = AB + \bar{A}C + BC(1)$$

$$= AB + \bar{A}C + BC(A + \bar{A})$$

$$= AB + \underline{\bar{A}C} + \underline{BCA} + \underline{BC\bar{A}}$$

$$= ABC(1 + C) + \bar{A}C(1 + B)$$

$$= AB + \bar{A}C$$

\* Solve the given expression using Consensus theorem

$$(Q) \bar{A}\bar{B} + AC + B\bar{C} + \bar{B}C + AB$$

$$\Rightarrow \bar{A}\bar{B} + AC + B\bar{C} + \bar{B}C(1) + AB \quad [A + \bar{A} = 1]$$

$$\Rightarrow \bar{A}\bar{B} + AC + B\bar{C} + \bar{B}C(A + \bar{A}) + AB$$

$$\Rightarrow \underline{\bar{A}\bar{B}} + \underline{AC} + \underline{B\bar{C}} + \underline{\bar{B}CA} + \underline{\bar{B}C\bar{A}} + AB$$

$$\Rightarrow \bar{A}\bar{B}(1 + C) + AC(1 + \bar{B}) + B\bar{C} + AB \quad [HC = 1, 1 + B = 1]$$

$$\Rightarrow \bar{A}\bar{B} + AC + B\bar{C} + AB$$

$$\Rightarrow \bar{A}\bar{B} + AC + B\bar{C} + ABC(1)$$

$$\Rightarrow \bar{A}\bar{B} + AC + B\bar{C} + AB(C + \bar{C})$$

$$\Rightarrow \bar{A}\bar{B} + \underline{AC} + \underline{B\bar{C}} + AB\bar{C} + \underline{AB\bar{C}}$$

$$\Rightarrow \bar{A}\bar{B} + AC(1 + B) + B\bar{C}(1 + A) \quad [1 + B = 1, 1 + A = 1]$$

$$\Rightarrow A\bar{B} + A\bar{C} + B\bar{C}$$

Dual of consensus theorem

The Dual form of consensus theorem is stated as

$$(A+B)(\bar{A}+C)(B+C) = (A+B)(\bar{A}+C)$$

$$(A\bar{A} + A\bar{C} + \bar{A}B + BC)(B+C) = (A\bar{A} + A\bar{C} + \bar{A}B + BC)$$

$$(0 + A\bar{C} + \bar{A}B + BC)(B+C) = (0 + A\bar{C} + \bar{A}B + BC) \quad [A \cdot \bar{A} = 0]$$

$$ABC + \bar{A}B \cdot B + BC \cdot B + AC \cdot C + \bar{A}BC + BC \cdot C = AC + \bar{A}B + BC \quad [BC + BC = BC]$$

$$BC(A+B) + \bar{A}B(B+C)$$

$$ABC + \bar{A}B \cdot B + BC + AC + \bar{A}BC + BC = AC + \bar{A}B + BC$$

$$\bar{A}B(1+C) + BC(1+A) + AC = AC + \bar{A}B + BC$$

$$\boxed{\bar{A}B + BC + AC = AC + \bar{A}B + BC}$$

Boolean function (or) switching function :-

Boolean equations are constructed by connecting the boolean constants and variables with the boolean operation.

Boolean expressions are known as boolean formulas we use

This Boolean Expressions to describe boolean functions.

Boolean expression to describe

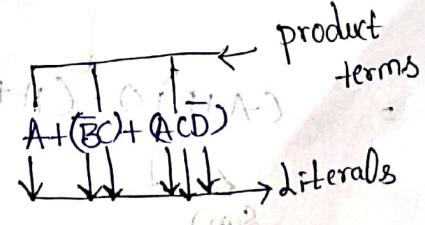
for example: If Boolean Expression  $(A+B)C$  is used to describe the function of  $f$ , then boolean function is written as,  $f(A, B, C) = (A+B)C$

$$f = (A+B)C$$

Let us consider the whole four variable boolean function.

Product-terms :-

$$f(A, B, C, D) = A + (BC) + (ACD)$$



Sum-Terms:

$$f(A, B, C, D) = (B + \bar{D}) \cdot (A + \bar{B} + C) \cdot (\bar{A} + C)$$

Literals

sum terms

→ The literals and terms are arranged in the form as in two

1. sum of product (SOP)

2. product of sum (POS)

1. sum of products:

→ The sum of product is also called Disjunctive normal form (DNF)

Disjunctive normal formula

→ The words sum & product are described from the symbolic representation "OR" and "AND" function ( $+$  and  $\cdot$ )

Ex-1

$$f(A, B, C) = AB\bar{C} + A\bar{B}\bar{C}$$

sum terms

products

Ex-2

$$f(P, Q, R) = \overline{PQ} + \overline{QR} + \overline{RS}$$

sum terms

products

2. Product of sum:

→ The product of sum is also called conjunctive normal form (CNF) conjunctive normal formula.

→ A product of sum is any group of sum terms ANDed

together.

Ex-3

$$f(A, B, C) = (A + B + C) \cdot (\bar{A} + \bar{B} + \bar{C})$$

product  
sum

$$f(P, Q, R) = (P+Q) \cdot (R+P) \cdot Q$$

↓ Sum      ↓ product

\* Canonical form  $\Leftrightarrow$  (standard SOP and POS forms)

→ Basically Canonical forms are two types

1. Standard SOP (or) minterm Canonical form

2. Standard POS (or) maxterm Canonical form

1. Standard SOP  $\Leftrightarrow$  (minterm Canonical form)

$f(A, B, C) = (A\bar{B}C) + (\bar{A}BC) + (\bar{A}\bar{B}C)$

→ The SOP is given by

Each product term is consistent all literals in either completed form or uncompleted form

2. Standard POS  $\Leftrightarrow$  (maxterm)

$f(A, B, C) = (A+\bar{B}+C) \cdot (\bar{A}+B+C) \cdot (\bar{A}+\bar{B}+C)$

Each sum term consists of all literals in the completed form

Each sum term consists of all literals in the completed form

\* Converting SOP to standard SOP (or) POS form

Steps to convert SOP to standard SOP form

1. find the missing literal in each product term if any

2. "AND" each product term having missing literals which term 1 formed by OR'ing the literals and its complement

3. Expand the terms by applying Distributive law and recorded

the literals in the product terms omitting (or) removing repeated product

4. Reduce the expression because  $(A+A) = A$

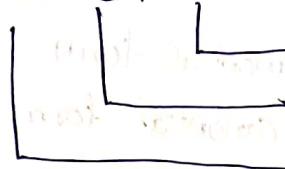
→ because  $(A+A) = A$

Ex: Convert the given expression in standard SOP form

$$f(A, B, C) = AC + AB + BC$$

Step 1:  $f(A, B, C)$  = find the missing literal in each product term

$$f(A, B, C) = AC + AB + BC$$



A literal is missing

C literal missing

B literal missing

Step 2: AND product terms with (missing literal + complement)

$$f(A, B, C) = AC(1) + AB(1) + BC(A + \bar{A})$$

$$\begin{aligned} A + \bar{A} &= 1 \\ B + \bar{B} &= 1 \\ C + \bar{C} &= 1 \end{aligned}$$

$$= AC(B + \bar{B}) + AB(C + \bar{C}) + BC(A + \bar{A})$$

Step 3: Expand the term and recorded all the terms

Expand :-

$$f(A, B, C) = ACB + A\bar{C}\bar{B} + ABC + AB\bar{C} + BCA + B\bar{C}A$$

Recorded :-

$$f(A, B, C) = ABC + A\bar{B}C + ABC + A\bar{B}\bar{C} + A\bar{B}C + \bar{A}BC$$

Step 4: omitting the repeated product term

$$f(A, B, C) = \underline{ABC} + \underline{A\bar{B}C} + \underline{ABC} + \underline{A\bar{B}\bar{C}} + \underline{A\bar{B}C} + \underline{\bar{A}BC}$$

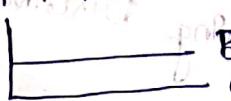
$$f(A, B, C) = ABC + A\bar{B}C + A\bar{B}\bar{C} + \bar{A}BC$$

Ex: Convert the given expression in standard SOP form

$$f(A, B, C) = A + ABC$$

Step 1: find

$$f(A, B, C) = A + ABC$$



B literal missing

C literal missing

Step 2: AND product term with (missing literal + complement)

$$f(A, B, C) = AC(1) + ABC$$

$$= A(B + \bar{B})(C + \bar{C}) + ABC$$

Step 3:

$$f(A, B, C) = AB + \overline{A}B (BC + \overline{B}C + \overline{B}\overline{C}) + ABC$$

$$= \underline{ABC} + AB\overline{C} + A\overline{B}C + A\overline{B}\overline{C} + \underline{ABC}$$

Step 4: Omitting the repeated product term

$$f(A, B, C) = ABC + AB\overline{C} + A\overline{B}C + A\overline{B}\overline{C}$$

Convert POS to standard POS form:

Step 5:

1. find the missing literals in each sum term if any.
2. OR each sum term having missing literal formed by ANDing the literal and its complement
3. Expand the terms by applying distributive law and recorded the literal in the sum terms
4. Reduce the Expression by omitting repeated sum term if any

$$A \cdot A = A$$

Ex: convert the given expression in standard POS form

$$f(A, B, C) = (A+B) \cdot (B+C) \cdot (A+C)$$

$$[A \cdot \overline{A} = 0]$$

$$[B \cdot \overline{B} = 0]$$

$$[C \cdot \overline{C} = 0]$$

Step 1: find the missing literal in each product

$$f(A, B, C) = (A+B) \cdot (B+C) \cdot (A+C)$$

$\swarrow$  B literal missing       $\searrow$  A literal missing       $\swarrow$  C literal missing

Step 2: OR product term with (missing literal + complement)

$$f(A, B, C) = [(A+B) + (C \cdot \overline{C})] \cdot [(B+C) + (A \cdot \overline{A})] \cdot [(A+C) + (B \cdot \overline{B})]$$

Step 3: expand the term and recorded the terms

Expand:

$$f(A, B, C) = [(A+B+C) \cdot (A+B+\overline{C})] \cdot [(B+C+A) \cdot (B+C+\overline{A})] \cdot [(A+C+B) \cdot (A+C+\overline{B})]$$

Recorded:

$$f(A, B, C) = [ \underline{A+B+C} \cdot \underline{A+B+\bar{C}} ] \cdot [ \underline{A+B+C} \cdot \underline{\bar{A}+B+C} ] \cdot [ \underline{A+B+C} \cdot \underline{\bar{A}+\bar{B}+C} ]$$

Step 4: omitting the repeated product term

$$f(A, B, C) = (A+B+C)(A+B+\bar{C})(\bar{A}+B+C)(\bar{A}+\bar{B}+C)$$

Ex: Convert the given expression in standard pos forms

$$f(A, B, C) = (A) \cdot (A+B+C)$$

Step 1: find the missing literal in each product

$$f(A, B, C) = A \cdot (A+B+C)$$

Step 2: B, C literals missing

Step 2: OR product term with (missing literal + complement)

$$f(A, B, C) = [A + (B \cdot \bar{B}) + (\bar{C} \cdot \bar{C})] \cdot (A+B+C)$$

Step 3: Expand the term and recorded the terms

$$\text{Expand: } [A + (B+C) + (B\bar{C}) + (\bar{B}+C) + (\bar{B}\bar{C})] \cdot (A+B+C)$$

$$= (A+B+C) \cdot (A+B+\bar{C}) \cdot (A+\bar{B}+C) \cdot (A+\bar{B}+\bar{C}) \cdot (A+\bar{B}+C)$$

Step 4: omitting the repeated product term

$$f(A, B, C) = (A+B+C) \cdot (A+B+\bar{C}) \cdot (A+\bar{B}+C) \cdot (A+\bar{B}+\bar{C})$$

$$= [(A+B+C) \cdot (A+B+\bar{C})] \cdot [(A+\bar{B}+C) \cdot (A+\bar{B}+\bar{C})]$$

$$= [(A+B+C) \cdot (A+B+\bar{C})] \cdot [(A+\bar{B}+C) \cdot (A+\bar{B}+\bar{C})]$$

$$= [(A+B+C) \cdot (A+B+\bar{C})] \cdot [(A+\bar{B}+C) \cdot (A+\bar{B}+\bar{C})]$$

# M-Notations: (minterms and maxterms)

Decimal No.	Binary Numbers (4x3)	Minterm (m) (SOP)	Maxterm (M) (POS)
0	000	$\bar{A}\bar{B}\bar{C}$ (m <sub>0</sub> )	A + B + C (M <sub>0</sub> )
1	001	$\bar{A}\bar{B}C$ (m <sub>1</sub> )	A + B + $\bar{C}$ (M <sub>1</sub> )
2	010	$\bar{A}B\bar{C}$ (m <sub>2</sub> )	A + $\bar{B}$ + C (M <sub>2</sub> )
3	011	$\bar{A}BC$ (m <sub>3</sub> )	A + $\bar{B}$ + $\bar{C}$ (M <sub>3</sub> )
4	100	A $\bar{B}\bar{C}$ (m <sub>4</sub> )	$\bar{A}$ + B + C (M <sub>4</sub> )
5	101	A $\bar{B}C$ (m <sub>5</sub> )	$\bar{A}$ + B + $\bar{C}$ (M <sub>5</sub> )
6	110	A B $\bar{C}$ (m <sub>6</sub> )	$\bar{A}$ + $\bar{B}$ + C (M <sub>6</sub> )
7	111	A B C (m <sub>7</sub> )	$\bar{A}$ + $\bar{B}$ + $\bar{C}$ (M <sub>7</sub> )

On the left

Minterm = Complement of maxterm

Example:  $f(A, B, C) = \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}C + \bar{A}BC + ABC$   
 $= m_0 + m_1 + m_3 + m_6$   
 $= \sum m(0, 1, 3, 6)$

Ex:  $f(A, B, C) = (A + B + \bar{C}) \cdot (A + \bar{B} + \bar{C}) \cdot (\bar{A} + \bar{B} + C)$   
 $= M_1 \cdot M_3 \cdot M_6$   
 $= \prod M(1, 3, 6)$

To find sum of product form to the given table

A B C	y
0 0 0	0
0 0 1	0
0 1 0	1
0 1 1	1
1 0 0	0

[ABC] 98 a  
 input variables  
 [SOP, k 1'8]

A · B · C	y
0 0 1	0
1 1 0	1
1 1 1	0

$$f(A, B, C) = \overline{A}B\overline{C} + \overline{A}B\overline{C} + A\overline{B}\overline{C}$$

$$= m_2 + m_3 + m_6$$

$$= \sum m(2, 3, 6)$$

To find the product of sum from the given table.

[Product of  
sum is zero's]

A + B + C	y
0 0 0	1
0 0 1	1
0 1 0	0
0 1 1	1
1 0 0	1
1 0 1	0
1 1 0	1
1 1 1	1

$$f(A, B, C) = (A + \overline{B} + \overline{C}) \cdot (\overline{A} + B + \overline{C})$$

[Adding 9 & 0]

$$= M_2 \cdot M_5$$

$$= \Pi M(2, 5)$$

0	0	0
0	0	0
1	0	0
0	1	0
1	1	0
0	0	1

## Algebraic Simplifications:

$$1. A \cdot \bar{A}C = 0$$

$$2. ABCD + ABD = ABD$$

$$3. A(A+B) = A$$

$$4. AB + ABC + ABC(D+E) = AB$$

$$5. XY + XYZ + XY\bar{Z} + \bar{X}YZ$$

$$6. \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}C$$

$$7. ABC + A\bar{B}C + ABC$$

$$8. A + \bar{A}B + A\bar{B} = A+B$$

$$\underline{\underline{Ex:}} \quad \overline{\bar{A}\bar{B} + \bar{A} + AB}$$

$$\Rightarrow \overline{\bar{A} + \bar{B} + \bar{A} + AB}$$

$$\Rightarrow \overline{\bar{A} + \bar{B} + AB}$$

$$\Rightarrow \overline{\bar{A} + (\bar{B}+A) \cdot (B+\bar{B})}$$

$$\Rightarrow \overline{\bar{A} + \bar{B} + A \cdot (1)}$$

$$\Rightarrow \overline{(A+\bar{A}) + \bar{B}}$$

$$\Rightarrow \overline{1 + \bar{B}}$$

$$[\bar{A}\bar{B} = \bar{A} + \bar{B}]$$

$$[\bar{B} + (AB) = \bar{B}A + (B \cdot B)]$$

$$= [(\bar{B}+A) \cdot (\bar{B}+\bar{B})]$$

$$[\bar{B} + B = 1]$$

$$[\bar{1} = 0]$$

$$\underline{\underline{Ex:}} \quad \overline{B + \bar{B} + A}$$

$$\Rightarrow \overline{B + 0}$$

$$\textcircled{1} \quad A \cdot \bar{A}C = 0$$

$$(A \cdot \bar{A})C = 0$$

$$(0)C = 0$$

$$A \cdot \bar{A}C = 0$$

$$\overline{\bar{A}\bar{B} + \bar{A} + AB}$$

$$\overline{\bar{A} + \bar{B} + AB}$$

$$\overline{A + AB + B\bar{B}}$$

$$\overline{A + AB}$$

$$\overline{B = (\bar{A}+A) \cdot (A+\bar{B})}$$

$$\overline{B = \frac{1}{1+B}}$$

$$\textcircled{2} \quad ABCD + ABD = ABD$$

$$[H + C = 1]$$

$$\Rightarrow ABD(1+c)$$

$\Rightarrow \text{ABD}$

$$\Rightarrow ABCD + \cancel{ABD} = ABD$$

$$\textcircled{3} \quad A(A+B) = A$$

$$\Rightarrow A \cdot A + A \cdot B$$

$$\Rightarrow A + AB$$

$$\Rightarrow A(1+B)$$

$\Rightarrow A$

$$④ AB + ABC + ABC(D+E) = AB$$

$$\Rightarrow AB + ABC + AB(D + \cancel{DE})$$

$$\Rightarrow AB(1+c) + ABD + ABE$$

$$\Rightarrow AB + ABCD + E$$

$$\Rightarrow AB(1+D+E)$$

$$\Rightarrow AD((1+D)+E)$$

$$\Rightarrow AB \cancel{[E^{12} + E]}$$

$\Rightarrow$  ABCD

$\Rightarrow \cdot AB$

$$5. xy + xyz + xy\bar{z} + \bar{x}yz = y(x+z)$$

$$\Rightarrow x + \bar{x}h_k + 2\bar{y}h_x + (1)y$$

$$\Rightarrow 2\bar{y}x + 2\bar{y}x + 2\bar{y}x + (\bar{z} + 2)x$$

$$\Rightarrow xy + \bar{x}\bar{y} + x\bar{y} + \bar{x}y$$

$$\Rightarrow xy(z+\bar{z}) + xy(z+\bar{z}) + \bar{x}y z$$

$$xy + xy\bar{z} + \bar{x}\bar{y}\bar{z} + \bar{x}yz$$

$$xy(x+2y) + yz(x+2y)$$

$$xy + yz \quad \text{A} \cdot A$$

$$y(x+z) \parallel A$$

$$\Rightarrow xy + \bar{x}y + \bar{x}yz$$

$$\Rightarrow xy + \bar{x}yz$$

$$\Rightarrow y(x + \bar{x}z) \Rightarrow y((x+\bar{x}) \cdot (x+z)) \Rightarrow y(1) \cdot (x+z)$$

$$\Rightarrow y(x+z) //$$

$$\textcircled{6} \quad A\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}B\bar{C}$$

$$\Rightarrow \bar{A}\bar{B}\bar{C} + \bar{A}B\bar{C} + \bar{A}B\bar{C}$$

$$\Rightarrow \bar{A}\bar{B}\bar{C} + \bar{A}B(\bar{C}+C)$$

$$\Rightarrow \bar{A}\bar{B}\bar{C} + \bar{A}B$$

$$\Rightarrow \bar{A}C(\bar{B}\bar{C} + B)$$

$$\Rightarrow \bar{A}((\bar{B}+B) \cdot (\bar{C}+B))$$

$$\Rightarrow \bar{A}(1 \cdot (\bar{C}+B))$$

$$\Rightarrow \bar{A}(\bar{C}+B)$$

$$\Rightarrow \bar{A}\bar{C} + \bar{A}B$$

$$\textcircled{7} \quad ABC + A\bar{B}C + A\bar{B}\bar{C}$$

$$ABC + A\bar{B}C + A\bar{B}\bar{C}$$

$$AC(C+B) + A\bar{B}\bar{C}$$

$$AC + AB\bar{C}$$

$$AC(C+\bar{B}C)$$

$$AC((C+B) \cdot (C+\bar{C}))$$

$$AC(C+B)$$

$$AC(A+A)(C+C)$$

$$AC+AB$$

$$8. A + \bar{A}B + A\bar{B} = A+B$$

$$\Rightarrow A + \bar{A}B + A\bar{B}$$

$$\Rightarrow A(1 + \bar{B}) + \bar{A}B$$

$$\Rightarrow A(1) + \bar{A}B$$

$$\Rightarrow A + \bar{A}B$$

$$\Rightarrow (A + \bar{A})(A + B)$$

$$\Rightarrow (1)(A + B)$$

$$\Rightarrow A + B //$$

Ex: Simplify the following three variable Expression by using Boolean algebra

$$Y = \Sigma m(1, 3, 5, 7)$$

$$\text{Given Data } Y = \Sigma m(1, 3, 5, 7) \\ = m_1 + m_3 + m_5 + m_7$$

$$= \bar{A}\bar{B}C + \bar{A}BC + A\bar{B}C + ABC$$

$$= \bar{A}C(\bar{B} + B) + AC(\bar{B} + B)$$

$$= \bar{A}C(1) + AC(1)$$

$$= C\bar{A} + A$$

$$= C //$$

Ex: Simplify the following three variable Expression by using Boolean algebra

$$Y = \prod m(1, 3, 5, 7)$$

$$\text{Given that } Y = \prod m(1, 3, 5, 7)$$

$$= M_1 \cdot M_3 \cdot M_5 \cdot M_7$$

$$= (A + B + \bar{C})(A + \bar{B} + \bar{C})(\bar{A} + B + \bar{C})(\bar{A} + \bar{B} + C)$$

$$= [A \cdot A + (A \cdot \bar{B}) + (A \cdot \bar{C}) + (B \cdot A) + (B \cdot \bar{B}) + (B \cdot \bar{C}) + (\bar{C} \cdot A) + (\bar{C} \cdot \bar{B}) + (\bar{C} \cdot \bar{C})] \cdot$$

$$[ (\bar{A} \cdot \bar{A}) + (\bar{A} \cdot B) + (\bar{A} \cdot \bar{C}) + (B \cdot A) + (B \cdot \bar{B}) + (B \cdot \bar{C}) + (\bar{C} \cdot \bar{A}) + (\bar{C} \cdot \bar{C}) + (\bar{C} \cdot \bar{C}) ] \cdot$$

$$= ABC + A\bar{B}\bar{C} + A\bar{C} + \bar{A}\bar{B}\bar{C} + ABC + A\bar{B}\bar{C} + A\bar{C} + AB\bar{C} + A\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C}$$

$$+ BC + \bar{A}\bar{B}\bar{C} + B\bar{C} + ABC + A\bar{B}\bar{C} + AC + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + A\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{B}\bar{C}$$

$$+ B\bar{C} + \bar{A}\bar{C} + \bar{A}\bar{B}\bar{C} + A\bar{C} + \bar{A}\bar{B} + B\bar{C} + A\bar{C} + \bar{B}\bar{C} + \bar{C}$$

$$= ABC + A\bar{B}\bar{C} + A\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + BC + \bar{A}\bar{B}\bar{C} + B\bar{C} + \bar{A}\bar{C} + \bar{A}\bar{B} + \bar{C}$$

$$= BC(A + \bar{A}) + \bar{B}\bar{C}(HA) + \bar{C}(A + \bar{A}) + \bar{A}\bar{B}(HC) + \bar{C}(HB) + \bar{A}\bar{B}\bar{C}$$

$$= BC + \bar{B}\bar{C} + \bar{C} + \bar{A}\bar{B} + \bar{C} + \bar{A}\bar{B}\bar{C}$$

$$= \bar{A}\bar{B} + \bar{A}\bar{B}\bar{C} + \bar{C} //$$

Ex-2 Simplify the following three variable expression. Convert the expression into minterm complementary form.

$$y = \prod_M(1, 3, 5, 7)$$

Sol: The given expression is maxterm  $y = \prod_M(1, 3, 5, 7)$ .

To convert given sequences to minterm complementary form

minterm = Complementary of Maxterm

The minterm is given by  $= \sum_m(0, 2, 4, 6)$

$$\begin{aligned} y &= \sum_m(0, 2, 4, 6) = \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} \\ &= m_0 + m_2 + m_4 + m_6 = \bar{A}\bar{C}(B + \bar{B}) + \bar{A}\bar{C}(B + \bar{B}) \\ &= \bar{A}\bar{C}(B + B) + \bar{A}\bar{C}(B + B) = \bar{A}\bar{C} \\ &= \bar{A}\bar{C} + \bar{A}\bar{C} \\ &= \bar{C}(\bar{A} + A) = \bar{C} // \end{aligned}$$

(M) \* transform each of the following Canonical expression into its other Canonical form & its decimal notations.

(q)  $f(x, y, z) = \sum_m(1, 3, 5)$

(q)  $f(x, y, z) = \prod_M(0, 2, 5, 6, 7, 8, 9, 11, 12)$

(q)  $f(x, y, z) = \sum_m(1, 3, 5)$

$$f(x, y, z) = \prod_M(0, 2, 4, 6, 7)$$

$$\begin{aligned} &= M_0 \cdot M_2 \cdot M_4 \cdot M_6 \cdot M_7 (x + y + z) \cdot (x + \bar{y} + z) \cdot (\bar{x} + y + z) \\ &\quad \cdot (\bar{x} + \bar{y} + z) \cdot (\bar{x} + y + \bar{z}) \cdot (\bar{x} + \bar{y} + \bar{z}) \\ &= (x + y + z) \cdot (A\bar{B} + \bar{C}) \cdot (A + B + C) \cdot (A + B + \bar{C}) \cdot (\bar{A} + \bar{B} + \bar{C}) \end{aligned}$$

$$f(w, x, y, z) = \overline{TM}(0, 2, 5, 6, 7, 8, 9, 11, 12)$$

$$f(w, x, y, z) = \text{PI}_N(0, 8, 15, 6, 7, 8, 9, 11, 12)$$

$$= \sum m(1, 3, 4, \cancel{10}, 14, 15)$$

$$= m_1 + m_3 + \cancel{m_4} + m_{10} + m_{13} + m_{14} + m_{15}$$

$$= \bar{w}\bar{x} + \bar{z}y\bar{w} + \bar{z}y\bar{x}\bar{w} + 2\bar{y}x\bar{w} + 2\bar{y}\bar{x}\bar{w} + 2\bar{y}\bar{x}\bar{w}$$

$$= \bar{w}xz(\bar{y}+y) + wxz(y\cdot\bar{y}) + w\bar{y}z(x+\bar{x}) + \bar{w}xz\bar{y}z$$

$$= \bar{\omega} \bar{x} z + \omega x \bar{z} + \bar{\omega} \bar{y} \bar{z} + \bar{\omega} x \bar{y} \bar{z}$$

Decimal NO	Binary number	minterm	Maxterm
8	1000	$A \bar{B} \bar{C} \bar{D}$	$(M) \bar{A} + \bar{B} + \bar{C} + \bar{D}$
9	1001	$\bar{A} \bar{B} \bar{C} D$ (m9)	$\bar{A} + \bar{B} + \bar{C} + D$
10	1010	$\bar{A} \bar{B} C \bar{D}$ (m10)	$\bar{A} + B + \bar{C} + \bar{D}$
11	1011	$\bar{A} \bar{B} \bar{C} D$	$\bar{A} + B + \bar{C} + D$
12	1100	$A \bar{B} \bar{C} \bar{D}$	$\bar{A} + \bar{B} + C + \bar{D}$
13	1101	$A \bar{B} \bar{C} D$	$\bar{A} + \bar{B} + \bar{C} + D$
14	1110	$A B \bar{C} \bar{D}$	$\bar{A} + \bar{B} + C + \bar{D}$
15	1111	$A B C \bar{D}$	$\bar{A} + \bar{B} + \bar{C} + \bar{D}$

# UNIT-III

## Gate-Level Minimization

Map :-

→ The map method gives us a systematic approach for simplifying a boolean expression. The map method first proposed by Veitch and modified by Karnaugh, hence it is known as the Veitch-Karnaugh diagram or the Karnaugh map (K-map).

One-variable, two-variable, three variable and four variable

Maps :-

$$1 \text{ Variable} = 2^1 = 2 \text{ cells}$$

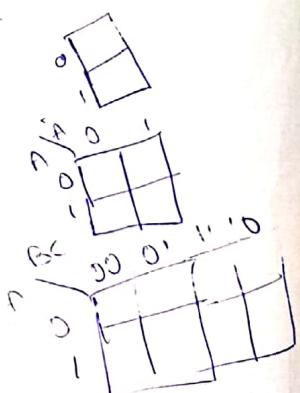
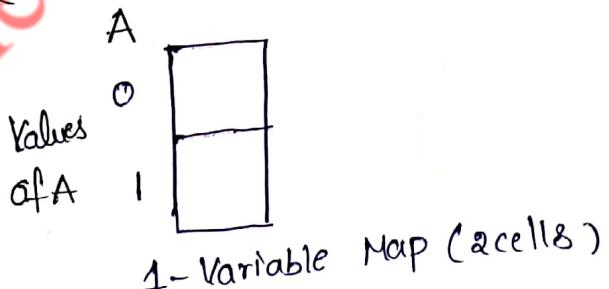
$$2 \text{ Variable} = 2^2 = 4 \text{ cells}$$

$$3 \text{ Variable} = 2^3 = 8 \text{ cells}$$

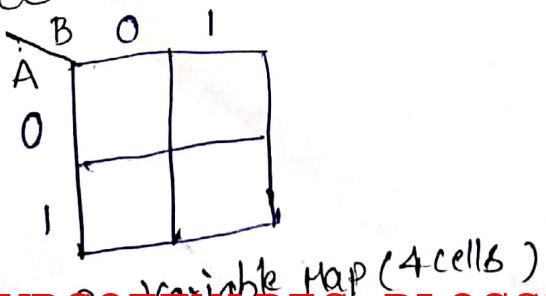
$$4 \text{ Variable} = 2^4 = 16 \text{ cells}$$

$$\begin{aligned}1 \text{ Var} &= 2^1 = 2 \text{ cells} \\2 \text{ Var} &= 2^2 = 4 \text{ cells} \\3 \text{ Var} &= 2^3 = 8 \text{ cells} \\4 \text{ Var} &= 2^4 = 16 \text{ cells}\end{aligned}$$

1 Variable K-Map :-



2-Variable K-Map :-



3-variable K-map

		BC	00	01	11	10
		A	0	1	1	0
A	B	00	0	1	1	0
		01	1	0	0	1

3-variable K-map (8 cells)

4-variable K-map

		CD		00	01	11	10
		AB	00	01	11	10	
AB	CD	00	0	1	1	0	
		01	1	0	0	1	

Values of CD

in gray code

4-variable K-map (16 cells)

1-variable

0	$\bar{A}$	$\bar{A} 0$
1	A	A 1

2-variable

		B	$\bar{B}$	B
		A	0	1
A	B	00	$\bar{A} \bar{B}$	$\bar{A} B$
		01	$A \bar{B}$	$A B$

3-variable

		BC		$\bar{B} \bar{C}$	$\bar{B} C$	$B \bar{C}$	BC
		A	00	01	11	10	
A	BC	00	$\bar{A} \bar{B} \bar{C}$	$\bar{A} \bar{B} C$	$\bar{A} B \bar{C}$	$\bar{A} B C$	
		01	$A \bar{B} \bar{C}$	$A \bar{B} C$	$A B \bar{C}$	$A B C$	

#### 4-variable

AB	CD	CD	CD	CD	CD
$\bar{A}\bar{B}$	$\bar{A}C\bar{D}$	$\bar{A}B\bar{D}$	$\bar{A}\bar{B}CD$	$A\bar{B}C\bar{D}$	$ABC\bar{D}$
	0	1	3	2	
$\bar{A}B$	$\bar{A}BC\bar{D}$	$\bar{A}B\bar{C}D$	$\bar{A}B\bar{C}\bar{D}$	$AB\bar{D}$	6
	4	5	7		
AB	$\bar{A}\bar{B}CD$	$A\bar{B}CD$	$ABC\bar{D}$	$ABC\bar{D}$	14
	12	13	15		
$A\bar{B}$	$A\bar{B}CD$	$A\bar{B}\bar{C}D$	$A\bar{B}\bar{C}\bar{D}$	$AB\bar{C}D$	10
	8	9	11		
					11 13 03 21

Representation of Truth table on k-map

two-variable k-map

→ The representation of two variable truth table on k-map is given below

A	B	y
0	0	0
0	1	1
1	0	1
1	1	0

A	B	0	1
0	0	0	1
1	1	0	0

A	B	0	1
0	0	0	1
1	1	1	0

Three-variable k-map

A	B	C	y
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	1

A	BC	00	01	11	10
0	0	0	0	0	1
1	1	1	1	1	0

A	BC	00	01	11	10
0	0	0	0	0	1
1	1	1	1	1	0

# four-variable K-maps

A	B	C	D	y
0	0	0	0	1
0	0	0	1	0
0	0	1	0	0
0	0	1	1	1
0	1	0	0	1
0	1	0	1	1
0	1	1	0	0
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	0
1	1	0	0	1
1	1	0	1	0
1	1	1	0	1
1	1	1	1	1

AB	CD	00	01	11	10
	0000	1	0	1	0
D	0001	0	1	0	1
	0011	1	1	1	0
	0010	0	0	0	1
	0100	1	1	0	1
	0101	1	0	1	0
	0111	0	1	1	1
	0110	1	1	1	0
	1100	0	0	1	1
	1101	0	1	0	1
	1111	1	1	1	1
	1110	1	1	0	0
	1000	1	0	0	0
	1001	0	1	1	1
	1011	0	1	0	0
	1000	1	0	0	0
	0100	0	0	1	0
	0101	0	1	1	1
	0111	0	1	1	0
	0110	1	1	0	1
	0010	0	0	0	1
	0011	1	1	1	0
	0001	0	1	0	1
	0000	1	0	1	0

AB	CD	CD	CD	CD
	1	0	1	0
	0	1	0	1
	1	1	1	0
	0	0	1	1
	1	0	1	0
	1	1	0	1
	0	1	1	0
	1	0	0	1
	0	0	1	1
	1	1	1	1
	0	0	0	1
	1	1	0	0
	0	1	1	1
	1	0	1	0
	0	0	0	1
	1	1	1	1
	0	0	1	0
	1	1	0	1
	0	1	1	0
	1	0	0	1
	0	0	1	1
	1	1	1	0
	0	1	0	1
	1	0	1	0
	0	0	0	1
	1	1	1	1

Plot Boolean Expression  $y = ABC\bar{t} + ABC + \bar{A}\bar{B}C$  on the k-map

A	B	C	y
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

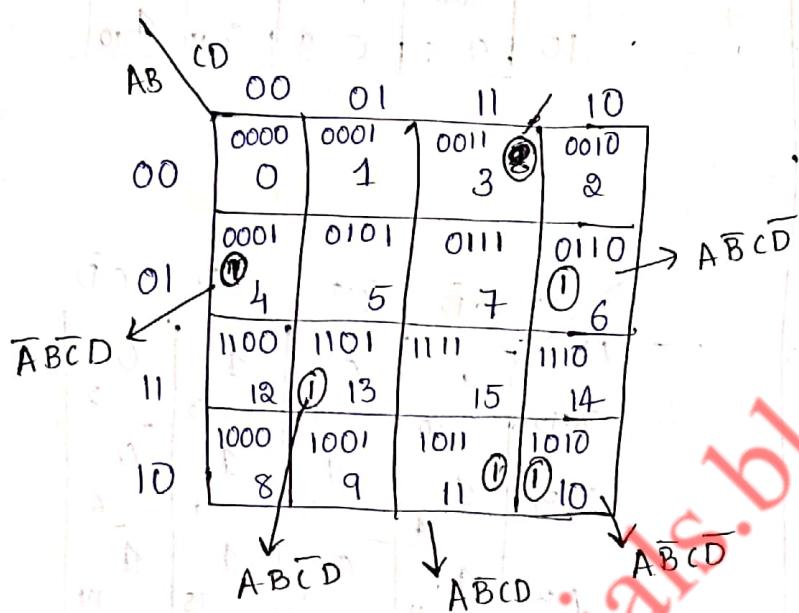
BC	00	01	11	10
0	0	1	0	0
1	0	1	1	0
0	0	0	1	1
1	0	1	0	1
0	1	1	1	0
1	0	0	0	1
0	1	0	1	0
1	1	1	0	1
0	0	1	1	0
1	1	0	0	1
0	1	1	0	1
1	0	0	1	0
0	0	0	1	1
1	1	1	1	1

Ex- plot Boolean expression  $y = \bar{A}B\bar{C}\bar{D} + A\bar{B}C\bar{D} + \bar{A}BC\bar{D} + A\bar{B}CD + ABC\bar{D}$  on the 1k-map.

$$y = \bar{A}B\bar{C}\bar{D} + A\bar{B}C\bar{D} + \bar{A}BC\bar{D} + A\bar{B}CD + ABC\bar{D}$$

$$= 0100 + 1010 + 0110 + 1011 + 1101$$

$$= m_4 + m_{10} + m_6 + m_{11} + m_{13}$$



Representation of standard pos on k-map

Three variable k-map

A	B	C	Maxterm (M)
0	0	0	$\bar{A}+\bar{B}+\bar{C}$ ( $M_0$ )
0	0	1	$\bar{A}+\bar{B}+\bar{C}$ ( $M_1$ )
0	1	0	$\bar{A}+\bar{B}+C$ ( $M_2$ )
0	1	1	$\bar{A}+\bar{B}+\bar{C}$ ( $M_3$ )
1	0	0	$\bar{A}+B+C$ ( $M_4$ )
1	0	1	$\bar{A}+B+\bar{C}$ ( $M_5$ )
1	1	0	$\bar{A}+\bar{B}+C$ ( $M_6$ )
1	1	1	$\bar{A}+\bar{B}+\bar{C}$ ( $M_7$ )

# Four Variable K-map

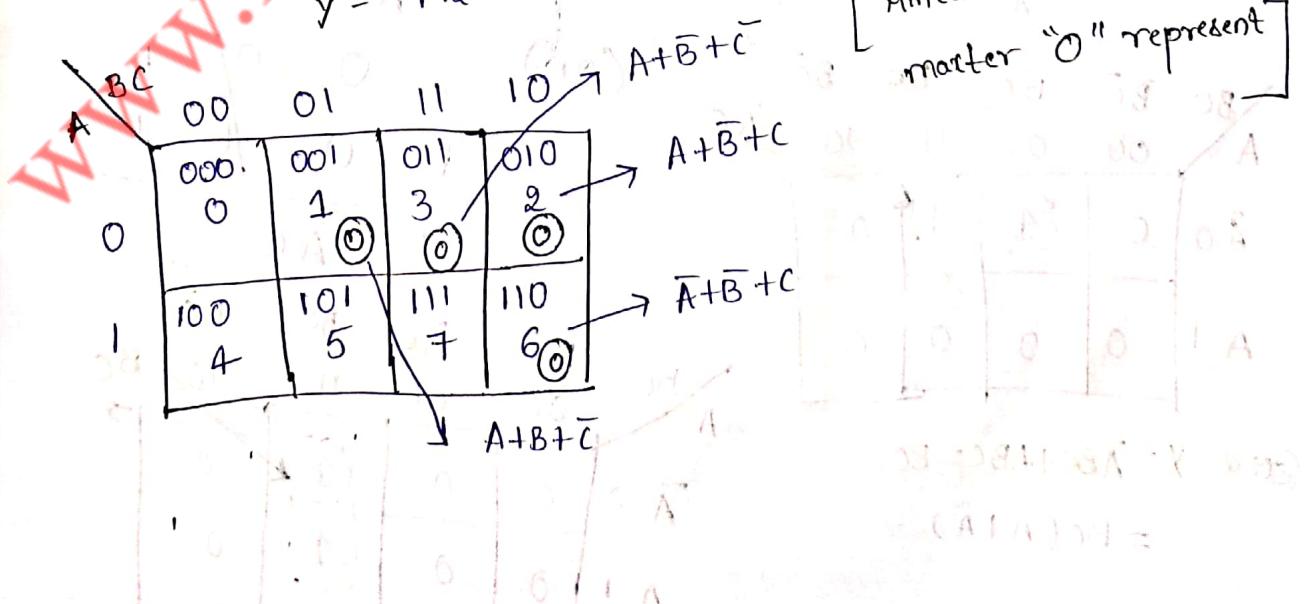
A	B	C	D	Maxterm (M)
0	0	0	0	$A+B+C+D \quad (M_0)$
0	0	0	1	$A+B+C+\bar{D} \quad (M_1)$
0	0	1	0	$A+B+\bar{C}+D \quad (M_2)$
0	0	1	1	$A+B+\bar{C}+\bar{D} \quad (M_3)$
0	1	0	0	$\bar{A}+\bar{B}+C+D \quad (M_4)$
0	1	0	1	$\bar{A}+\bar{B}+C+\bar{D} \quad (M_5)$
0	1	1	0	$\bar{A}+\bar{B}+\bar{C}+D \quad (M_6)$
0	1	1	1	$\bar{A}+\bar{B}+\bar{C}+\bar{D} \quad (M_7)$
1	0	0	0	$\bar{A}+B+C+\bar{D} \quad (M_8)$
1	0	0	1	$\bar{A}+B+C+D \quad (M_9)$
1	0	1	0	$\bar{A}+B+\bar{C}+D \quad (M_{10})$
1	0	1	1	$\bar{A}+B+\bar{C}+\bar{D} \quad (M_{11})$
1	1	0	0	$\bar{A}+\bar{B}+C+D \quad (M_{12})$
1	1	0	1	$\bar{A}+\bar{B}+C+\bar{D} \quad (M_{13})$
1	1	1	0	$\bar{A}+\bar{B}+\bar{C}+D \quad (M_{14})$
1	1	1	1	$\bar{A}+\bar{B}+\bar{C}+\bar{D} \quad (M_{15})$

Ex:- plot Boolean Expression  $y = (A+\bar{B}+C)(A+\bar{B}+\bar{C})(\bar{A}+\bar{B}+C)(\bar{A}+\bar{B}+\bar{C})$

on the K-map:

$$y = M_2 \cdot M_3 \cdot M_6 \cdot M_1$$

[Maxterm kis "1" represent  
minter "0" represent]



Ex-5 plot Boolean Expression  $y = (A+B+C+\bar{D})(A+\bar{B}+\bar{C}+\bar{D})(A+B+\bar{C}+D)$

$$(\bar{A}+\bar{B}+C+\bar{D})(\bar{A}+\bar{B}+\bar{C}+D)$$

$$Y = (A+B+C+\bar{D}) \cdot (\bar{A}+\bar{B}+\bar{C}+D) (\bar{A}+\bar{B}+\bar{C}+\bar{D}) (\bar{A}+\bar{B}+C+D)$$

$$Y = M_1 \cdot M_6 \cdot M_3 \cdot M_{13} \cdot M_{14}$$

		CD	AB	00	01	11	10	
		CD	AB	0000	0001	0011	0010	$(A+B+C+\bar{D})$
		CD	AB	00	4	0	2	$(A+\bar{B}+\bar{C}+D)$
00	00	0100	0101	0111	0110			
01	01	1100	1101	1111	1110	6	0	
11	11	1000	1001	1011	1010	14	0	$(\bar{A}+\bar{B}+\bar{C}+D)$
10	10	0100	0101	0111	0110	11	10	$\bar{A}+\bar{B}+C+\bar{D}$

Grouping cells for simplification:

Grouping two adjacent ones (pair)

$$\text{Ex-6 } y = \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}C$$

$$= \bar{A}C(\bar{B}+\bar{B})$$

$$= \bar{A}C$$

		BC	$\bar{B}\bar{C}$	$\bar{B}C$	$B\bar{C}$	$BC$
		A	00	01	11	10
$\bar{A}0$	0	0	1	1	0	
A 1	1	0	0	0	0	

(Cor.)

$$\begin{array}{r} 0 \quad 0 \\ 0 \quad 1 \\ \hline \bar{A}C \end{array}$$

		BC	$\bar{B}\bar{C}$	$\bar{B}C$	$B\bar{C}$	$BC$
		A	00	01	11	10
$\bar{A}0$	0	0	0	1	1	0
A 1	1	0	0	0	0	

$$\text{Ex-7 } Y = \bar{A}BC + ABC = BC$$

$$= BC(A + \bar{A})$$

$$= BC$$

$$\text{Ex: } ③ \quad Y = A\bar{B}\bar{C} + A\bar{B}C \\ = A\bar{C}(\bar{B} + B) \\ = A\bar{C}$$

	$\bar{B}\bar{C}$	$\bar{B}C$	$BC$	$B\bar{C}$
$\bar{A}$	0	0	0	0
A	1	1	0	1

$$\text{Ex: } ④ \quad Y = \bar{A}\bar{B}\bar{C}D + A\bar{B}\bar{C}D \\ = \bar{B}\bar{C}D(A + \bar{A}) \\ = \bar{B}\bar{C}D$$

	$\bar{C}D$	$\bar{C}D$	$\bar{C}D$	$CD$	$CD$
$\bar{A}B$	00	01	10	11	10
$\bar{A}\bar{B}$	00	11	00	00	00
$\bar{A}B$	10	00	00	00	00
$\bar{A}\bar{B}$	10	11	00	00	00

$$\text{Ex: } ⑤ \quad Y = \bar{A}\bar{B}C + \bar{A}B\bar{C} + A\bar{B}C \\ = \bar{A}\bar{B}C + \bar{A}B\bar{C} + \bar{A}BC + ABC$$

$$= \bar{A}C(\bar{B} + B) + BC(\bar{A} + A)$$

$$= \bar{A}C + BC$$

Group A  $\rightarrow \bar{A}C$

	$\bar{B}\bar{C}$	$\bar{B}C$	$BC$	$B\bar{C}$
$\bar{A}$	0	1	1	0
A	0	0	1	0

Group B  $\rightarrow BC$

[OR]

A B C

Grouping four adjacent ones (quad)

	$\bar{B}\bar{C}$	$\bar{B}C$	$BC$	$B\bar{C}$
$\bar{A}$	00	01	11	10
A	0	0	0	0
A	1	1	1	1

Ex-2

AB	CD	CD	CD	CD
$\bar{A}B$	00	01	11	10
$A\bar{B}$	00	0	1	0
$\bar{A}B$	01	0	1	0
AB	11	0	1	0
$A\bar{B}$	10	0	1	0

(b)  $y = CD$

Ex-3

AB	CD	$\bar{C}D$	$\bar{C}D$	$\bar{C}D$	$\bar{C}D$
$\bar{A}B$	00	0	0	0	0
$\bar{A}B$	01	0	1	1	0
BD	AB	0	1	1	0
AB	11	0	0	0	0
$A\bar{B}$	10	0	0	0	0

(3)  $y = BD$

Ex-4

AB	CD	$\bar{C}D$	$\bar{C}D$	CD	CD
$\bar{A}B$	00	0	0	0	0
$\bar{A}B$	01	0	0	0	0
AB	11	1	0	0	1
$A\bar{B}$	10	1	0	0	1

$y = \bar{AD}$

Ex-5

	CD	$\bar{C}D$	$\bar{B}D$	$B\bar{D}$	CD
AB	00	01	11	10	
$\bar{A}\bar{B}$	00	1	0	0	1
$\bar{A}B$	01	0	0	0	0
AB	11	0	0	0	0
$A\bar{B}$	10	1	0	0	1

$$\text{Ans: } Y = \bar{B}\bar{D}$$

$$\begin{aligned}
 & A^B C^D + B^0 C^0 \\
 & = B^0 C^0 + B^0 C^0 \\
 & = B^0 (C^0 + C^0) \\
 & = B^0 C^0
 \end{aligned}$$

Ex-6

	CD	$\bar{C}D$	$\bar{B}D$	CD	$\bar{B}\bar{D}$
AB	00	01	11	10	
$\bar{A}\bar{B}$	00	0	0	0	0
$\bar{A}B$	01	0	0	0	0
AB	11	1	1	1	1
$A\bar{B}$	10	0	1	1	1

Group 2  $\rightarrow$  AD

$$(b) Y = AB + AD + AC$$

Grouping eight adjacent ones (Octet)

	CD	$\bar{C}D$	$\bar{B}D$	CD	$\bar{B}\bar{D}$
AB	00	01	11	10	
$\bar{A}\bar{B}$	00	0	0	0	0
$\bar{A}B$	01	1	1	1	1
AB	11	1	1	1	1
$A\bar{B}$	10	0	0	0	0

$$(a) Y = B$$

	CD	$\bar{C}D$	$\bar{B}D$	CD	$\bar{B}\bar{D}$
AB	00	01	11	10	
$\bar{A}\bar{B}$	00	0	1	1	0
$\bar{A}B$	01	0	1	1	0
AB	11	0	1	1	0
$A\bar{B}$	10	0	1	1	0

$$(b) Y = D$$

$\bar{A}B$	$CD$	$\bar{C}D$	$\bar{C}D$	$CD$	$CD$
$\bar{A}B$	00	01	11	10	00
$\bar{A}B$	01	00	00	00	11
$\bar{A}B$	11	00	00	00	11
$\bar{A}B$	10	11	11	11	10

$\bar{A}B$	$CD$	$\bar{C}D$	$\bar{C}D$	$CD$	$CD$
$\bar{A}B$	00	01	11	10	00
$\bar{A}B$	01	00	00	00	11
$\bar{A}B$	11	00	00	00	11
$\bar{A}B$	10	11	11	11	10

A	B	C	D
0	0	0	0
0	0	0	1
0	0	1	1
0	0	1	0
1	0	0	0
0	0	0	1
0	0	1	1
1	0	1	0

$\overline{B}$

A	B	C	D
0	0	0	0
0	1	0	0
1	1	0	0
0	0	0	0
0	0	1	1
0	1	1	1
1	0	1	0

### Simplifications of SOP Expressions

→ from the above discussion we can outline generalized procedure to simplify Boolean Expressions as follows:

1. plot the k-map and place 1's in those cells corresponding to the 1's in the truth table or sum of product expression; place 0's in other cells.

2. check the k-map for adjacent 1's and encircle those 1's which are not adjacent to any other 1's; these are called isolated 1's.

3. check for those 1's which are adjacent to only one other 1 and encircle such pairs.

4. check for quads and octets of adjacent 1's even if it contains some 1's that have already been encircled. while doing this make sure that there are minimum no. of groups.

5. combine any pairs necessary to include any 1's that have not yet been grouped.

6. From the simplified expression by summing product terms of all groups.

Ex: minimize the expression  $Y = A\bar{B}C + \bar{A}\bar{B}C + A\bar{B}C + A\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C}$

$$\begin{aligned} Y &= 101 + 001 + 011 + 100 + 000 \\ &= m_5 + m_1 + m_3 + m_6 + m_0 \\ &= \sum m(5, 3, 4, 0) \end{aligned}$$

Step 1: The k-map for three variables and its plotted according to the given expression.

Step 2: There are no isolated 1's

Step 3: 1 in the cell 3 adjacent only to 1 in the cell 4. This pair is combined and referred to as group 1

		$\bar{B}\bar{C}$	$\bar{B}C$	$BC$	$B\bar{C}$
		00	01	11	10
$A\bar{B}$	0	000 ① 0	001 ① 1	011 ① 3	010 2
	1	100 ① 4	101 ① 5	111 7	110 6
$A$	0	000 ① 0	001 ① 1	011 ① 3	010 2
	1	100 ① 4	101 ① 5	111 7	110 6

Step 4: There are no octet, but there are a quad cells 0, 1, 4 and 5 from a quad.

This quad is combined and referred to as a group.

Step 5: All ~~the~~ have already been grouped

		$\bar{B}\bar{C}$	$\bar{B}C$	$BC$	$B\bar{C}$
		00	01	11	10
$A\bar{B}$	0	000 1	001 1	011 1	010 0
	1	100 1	101 1	111 0	110 0
$A$	0	000 1	001 1	011 1	010 0
	1	100 1	101 1	111 0	110 0

Step 6: Each group generates a term

In Expression for  $y$  - In group 1 B Variable  
is eliminated and in group 2 Variables

A and care eliminated and we get

	$\overline{BC}$	$\overline{BC}$	$BC$	$BC$
$A$	00	01	11	10
$\overline{A}$	1	1	1	0
$B$	1	0	0	0

$$y = \overline{A}c + \overline{B}$$

Ex: minimize the Expression

$$Y = \overline{A}B\overline{C}\overline{D} + \overline{A}B\overline{C}D + A\overline{B}\overline{C}\overline{D} + A\overline{B}\overline{C}D + A\overline{B}C\overline{D} + A\overline{B}CD$$

$$= 0100 + 0101 + 1100 + 1101 + 1001 + 0010$$

$$= m_4 + m_5 + m_{12} + m_{13} + m_9 + m_2$$

$$= \Sigma_m (4, 5, 12, 13, 9, 2)$$

$\overline{AB}$	$\overline{CD}$	$\overline{CD}$	$CD$	$\overline{CD}$	$G_1$	$G_2$	$G_3$
$00$	$00$	$01$	$11$	$10$	$A \quad B \quad \overline{C} \quad D$	$A \quad B \quad C \quad D$	$\overline{ABC}D$
$00$	$00$	$01$	$11$	$10$	$0 \quad 1 \quad 0 \quad 0$	$1 \quad 1 \quad 0 \quad 1$	$0 \quad 0 \quad 1 \quad 1$
$01$	$01$	$10$	$11$	$11$	$0 \quad 1 \quad 1 \quad 0$	$1 \quad 0 \quad 0 \quad 1$	$1 \quad 0 \quad 0 \quad 1$
$11$	$10$	$10$	$10$	$10$	$1 \quad 0 \quad 0 \quad 1$	$A \quad \overline{C} \quad D$	$\overline{AB}CD$
$10$	$00$	$00$	$00$	$00$	$B \bar{C}$		

## Group

Group ②

A C D

$$Y = \overline{A} \overline{B} C \overline{D} + A \overline{C} D + B \overline{C}$$

\* Simplify the logic function specified by the truth table by using the Karnaugh map method.  $y$  is the output variable and  $A, B$  and  $C$  are the input variables.

A	B	C	y
0	0	0	1
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

$$Y = \bar{A}\bar{B}C + \bar{A}B\bar{C} + A\bar{B}\bar{C} + ABC$$

$$\begin{aligned} &= m_0 + m_3 + m_4 + m_7 \\ &= \sum m(0, 3, 4, 7) \end{aligned}$$

		BC		A
		00	10	
A	0	0 4 4 1	1 1 5 7	
	1	2 3 6 7	1 1 6 7	

Group 1		Group 2	
0/0	0/0	0/0	1/1
0/0	1/1	1/1	1/1
<u><math>\bar{B}\bar{C}</math></u>		<u><math>BC</math></u>	

$$\therefore Y = \bar{B}\bar{C} + BC$$

Ex: Reduce the following function using K-map technique.

$$f(A, B, C, D) = \sum m(0, 1, 4, 8, 9, 10)$$

$$= m_0 + m_1 + m_4 + m_8 + m_9 + m_{10}$$

$$= \bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}\bar{C}D + \bar{A}\bar{B}C\bar{D} + A\bar{B}\bar{C}\bar{D} + A\bar{B}\bar{C}D + A\bar{B}C\bar{D}$$

AB	CD	00	01	11	10		Group 1	Group 2		
		2	1	3	2		A	B	C	D
00	01	1	0	1	0		1	0	0	0
01	11	4	5	7	6		1	0	0	1
11	10	12	13	15	14		1	0	0	0
10	00	8	9	11	10		1	0	0	1
		1	1	1	1					
								$\bar{B} \bar{C}$		
								Group 3		
								$\bar{A} \bar{B} \bar{C} \bar{D}$		
								0 1 0 0		
								$\bar{A} \bar{B} \bar{C} D$		

Product of sum simplification:

1. plot the k-map and place 0s in those cells corresponding to the 0s in the truth table or minterms in the product of sum expression.
  2. check the k-map for adjacent 0s and encircle those 0s which are not adjacent to any other 0s. These are called isolated 0s.
  3. check for those 0s which are adjacent to only one other 0 and encircle such pairs.
  4. check for quads and octets of adjacent 0s even if it contains some 0s that have already been encircled. while doing this make sure that there are minimum no. of groups.
  5. Combine any pairs necessary to include any 0s that have not yet been grouped.
  6. From the simplified pos expression for F by taking product of sum terms of all the groups.
- To get familiar with these steps we will solve some Examples.

Ex: minimize the Expression

$$y = (A+B+C)(A+\bar{B}+C)(\bar{A}+\bar{B}+C)(\bar{A}+B+C)(A+B+C)$$

$$= (\cancel{001})(\cancel{011})$$

$$(A+\bar{B}+C) = M_3, (\bar{A}+\bar{B}+C) = M_7$$

$$= (A+B+C) = M_1,$$

$$(\bar{A}+B+C) = M_0$$

$$(\bar{A}+B+C) = M_4,$$

Step 1: (a) shows the k-map for three variables A, B, C and minterms plotted according to given minterms.

	BC	00	01	11	10
A	0	0	0	3	2
	4	5	7	X	6
	0	0	0	0	1

Step 2: Three are no isolated 0s

Step 3: 0 in the cell 4 is adjacent only to 0 in the cell 0 and 0 in cell 7 is adjacent only to 0 in the cell 3. These two pairs are combined and referred to as group 1 and group 2 respectively.

	BC	00	01	11	10
A	0	0	1	3	2
	0	0	0	0	1
	0	0	0	0	1

Step 4: There are no quads and octets

	BC	00	01	11	10
A	0	0	1	3	2
	0	0	0	0	1
	0	0	0	0	1

Step 5: The 0 in the cell 1 can be combined with 0 in the cell 3 to be from a pair. This pair is referred to as group 3.

	BC	00	01	11	10
A	0	0	1	3	2
	0	0	0	0	1
	0	0	0	0	1

Step 6: In group 1 and in group 2, A is eliminated, where in group 3 variable B is eliminated and we get

Group 1	Group 2	Group 3	Group 3	
A	B	C	ABC	$\bar{ABC}$
0	0	0	0 11	0 01
1	0	0	1 11	0 11

$$\bar{Y} = \overline{B\bar{C}} + B\bar{C} + \bar{A}C$$

$$\bar{Y} = Y = \overline{B\bar{C}} + B\bar{C} + \bar{A}C \quad [\text{According to DeMorgan's}]$$

$$Y = \overline{\bar{B}\bar{C}} + B\bar{C} + \bar{A}C$$

$$= (\overline{\bar{B}\bar{C}}) \cdot (\overline{B\bar{C}}) \cdot (\bar{A}C)$$

$$= (\bar{B} + \bar{C}) \cdot (\bar{B} + \bar{C}) \cdot (\bar{A} + C)$$

$$Y = (B + C)(\bar{B} + \bar{C})(\bar{A} + C)$$

$$Y = \bar{B} + \bar{C} + A$$

Ex: Minimize the following expression in the pos form

$$X = (\bar{A} + \bar{B} + C + D)(\bar{A} + \bar{B} + \bar{C} + D)(\bar{A} + \bar{B} + \bar{C} + \bar{D})(\bar{A} + B + C + D)(A + \bar{B} + \bar{C} + D) \\ (\bar{A} + \bar{B} + \bar{C} + \bar{D})(A + B + C + D)(\bar{A} + \bar{B} + C + \bar{D})$$

Sol:  $(\bar{A} + \bar{B} + C + D) = M_{12}$ ,  $(\bar{A} + \bar{B} + \bar{C} + D) = M_{14}$ ,  $(\bar{A} + \bar{B} + \bar{C} + \bar{D}) = M_{15}$

$$(\bar{A} + B + C + D) = M_8, (A + \bar{B} + \bar{C} + D) = M_6, (A + \bar{B} + \bar{C} + \bar{D}) = M_7$$

$$(A + B + C + D) = M_0 \text{ and } (\bar{A} + \bar{B} + C + \bar{D}) = M_{13}$$

Step 1: Shows the k-map for four variable and it plotted according to given max terms

Step 2: There are no isolated 1's

Step 3: 0 in the cell 0 is adjacent

Only to 0 in the cell 11 this pair is combined and referred to as group 1

Step 4: There are two quads cells 12, 13, 14 and 15 form a quad

1 and cells 6, 7, 14, 15 forms a quad 2. These two quads are referred to as group 2 and group 3, respectively.

Step 5: All 0s have already been grouped

Step 6: In group 1, Variable A is eliminated

In group 2, Variable C and D are eliminated and 01

in group 3 variable A and D are eliminated.

∴ we get simplified pos Expression

		CD	C+D	C+D̄	C̄+D	C̄+D̄
		AB	00	01	11	10
A+B	00	0	1	3	2	
		0				
A+B̄	01	4	5	7	6	
		0	0	0	0	
A+B	10	12	13	15	14	
		0	0	0	0	
A+B̄	11	8	9	11	10	
		0				

		CD	Group 1		
		AB	00	01	11
	00	0	1	3	2
	01	0			
	10	4	5	7	6
	11	0	0	0	0
	11	12	13	15	14
	10	0	0	0	0
	11	8	9	11	10
	10	0			

Group 1		
A	B	CD
0	0	00
1	0	00
<u><math>\bar{B} \bar{C} \bar{D}</math></u>		

Group 2		
A	B	CD
1	1	00
1	1	01
1	1	11
1	1	10

Group 3		
A	B	CD
0	1	11
0	1	10
1	1	11
1	1	10

$$\bar{Y} = \overline{BCD} + AB + BC$$

$$y = \bar{Y} = \overline{BCD} + AB + BC$$

$$y = (B+C+D) \cdot (\bar{A}+\bar{B}) \cdot (\bar{B}+\bar{C})$$

Incompletely Specified functions (Don't Care terms or conditions)

e.g:

A	B	C	y
0	0	0	0
1	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	X
1	1	1	X

Sol:

A	BC		00	01	11	10
	00	01	11	10	11	10
0	0	1	1	1	1	0
1	4	5	X	X	X	X

Don't Care Conditions

A	B	C
0	0	1
1	1	1
1	0	1
0	1	1

Group 1

A	BC		00	01	11	10
	00	01	11	10	11	10
0	0	1	1	1	1	0
1	4	5	X	X	X	X

Group 1

$$\therefore Y = C$$

# Describing Incomplete Boolean function

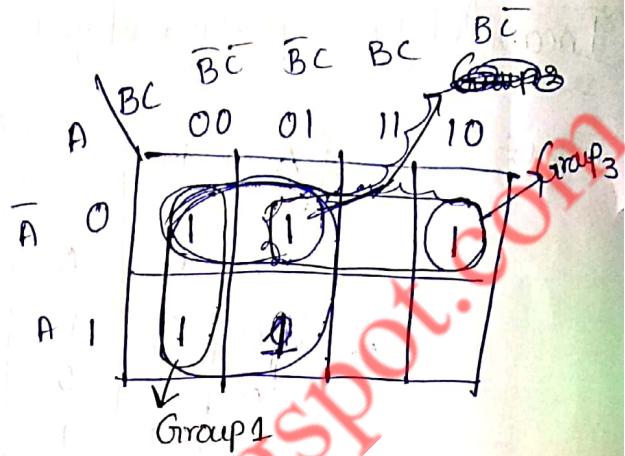
Ex:  $f(A, B, C) = \sum m(0, 2, 4) + d(1, 5)$

Q)  $f(A, B, C) = \prod M(2, 5, 7) + d(1, 3)$

Q)  $f(A, B, C) = \sum m(0, 2, 4) + d(1, 5)$

		BC	00	01	11	10
		A	0	1	3	2
0	0	1	X			1
	1	4	5	7		6

$$\begin{array}{c}
 G_1 \\
 \begin{array}{c} A \\ \diagup B \\ \diagdown C \end{array} \\
 \begin{array}{cc} 0 & 0 \\ 1 & 0 \end{array} \\
 \hline
 \overline{B} \quad \overline{C}
 \end{array}
 \quad
 \begin{array}{c}
 G_2 \\
 \begin{array}{c} A \\ \diagup B \\ \diagdown C \end{array} \\
 \begin{array}{cc} 0 & 0 \\ 0 & 1 \end{array} \\
 \hline
 \overline{A} \quad \overline{B} \quad \overline{C}
 \end{array}$$

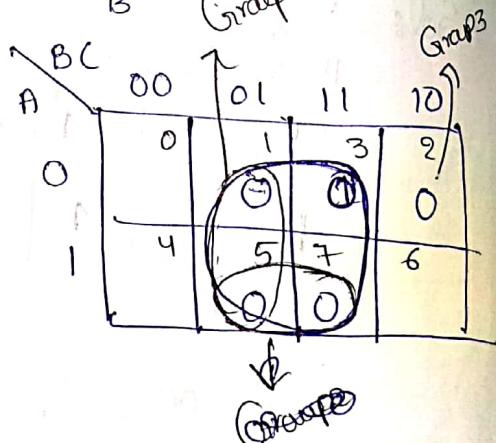


$$\begin{array}{c}
 y = \overline{B} \quad \overline{C} + \overline{A} \quad \overline{B} + \overline{B} \quad \overline{C} // \\
 \text{Group 1} \\
 \begin{array}{ccc} A & B & C \\ 0 & 0 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 1 \end{array} \\
 \hline
 \overline{B}
 \end{array}
 \quad
 \begin{array}{c}
 y = \overline{B} + \overline{B} \quad \overline{C} \\
 \text{Group 2} \\
 \begin{array}{ccc} A & B & C \\ 0 & 0 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 1 \end{array} \\
 \hline
 \overline{B}
 \end{array}$$

Q)  $f(A, B, C) = \prod M(2, 5, 7) + d(1, 3)$

		BC	00	01	11	10
		A	0	1	3	2
0	0	0	X	X		0
	1	4	5	7		6

$$\begin{array}{c}
 G_1 \\
 \begin{array}{c} A \\ \diagup B \\ \diagdown C \end{array} \\
 \begin{array}{cc} 0 & 0 \\ 1 & 0 \end{array} \\
 \hline
 \overline{B} \quad \overline{C}
 \end{array}
 \quad
 \begin{array}{c}
 G_2 \\
 \begin{array}{c} A \\ \diagup B \\ \diagdown C \end{array} \\
 \begin{array}{cc} 1 & 0 \\ 1 & 1 \end{array} \\
 \hline
 AC
 \end{array}
 \quad
 \begin{array}{c}
 G_3 \\
 \begin{array}{c} A \\ \diagup B \\ \diagdown C \end{array} \\
 \begin{array}{cc} 0 & 1 \\ 1 & 0 \end{array} \\
 \hline
 \overline{A} \quad \overline{B} \quad \overline{C}
 \end{array}$$



$$\begin{array}{c}
 G_1 \\
 \begin{array}{c} A \\ \diagup B \\ \diagdown C \end{array} \\
 \begin{array}{cc} 0 & 0 \\ 0 & 1 \\ 1 & 0 \\ 1 & 1 \end{array} \\
 \hline
 C
 \end{array}$$

$$y = \overline{A} \quad \overline{B} \quad \overline{C} + \overline{B} \quad \overline{A} \quad C //$$

Don't care conditions in logic design

	A	B	C	D	P
0	0	0	0	0	0
1	0	0	0	1	1
2	0	0	1	0	1
3	0	0	1	1	0
4	0	1	0	0	1
5	0	1	0	1	0
6	0	1	1	0	0
7	0	1	1	1	1
8	1	0	0	0	1
9	1	0	0	1	0
10	1	0	1	0	1
11	1	0	1	1	1
12	1	1	0	0	1
13	1	1	0	1	1
14	1	1	1	0	1
15	1	1	1	1	1

$$P = \Sigma m(1, 2, 4, 7, 8) + \Sigma d(10, 11, 12, 13, 14, 15)$$

Find the reduced SOP form of the following function.

$$f(A, B, C, D) = \Sigma m(1, 3, 7, 11, 15) + \Sigma d(0, 2, 4)$$

$$f(A, B, C, D) = \Sigma m(1, 3, 7, 11, 15) + \Sigma d(0, 2, 4)$$

	AB	CD	00	01	11	10
00	X	0	1	(1)	(1)	X
01	X	4	5	7	6	
11		12	13	15	14	
10		8	9	(1)	10	

	AB	CD	00	01	11	10	
00			1	1	(1)	1	G <sub>11</sub>
01			0		1		
11					1		
10						1	

↓ Group 1

↓ Group 2

G<sub>11</sub>      G<sub>2</sub>

A	B	C	D	A	B	C	D
0	0	0	0	0	0	0	0
0	0	0	1	0	0	1	1
0	0	1	1	0	1	1	1
0	1	1	1	1	1	1	1
1	1	1	1	1	1	1	1

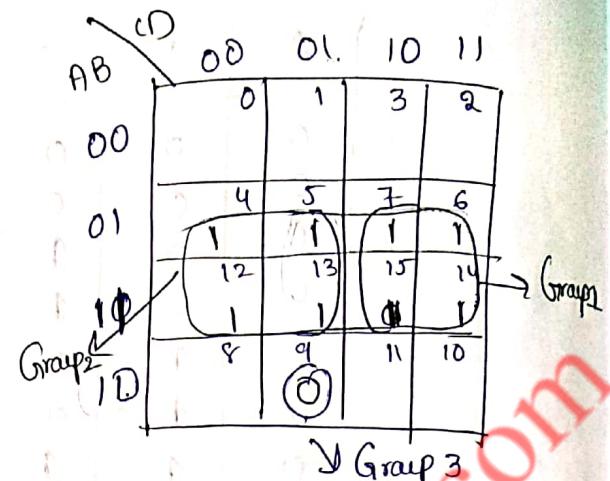
CD

$$\therefore Y = \bar{A}\bar{B} + CD$$

Ex: Reduce the following function using k-map technique

$$f(A, B, C, D) = \sum m(5, 6, 7, 12, 13) + \sum d(4, 9, 14, 15)$$

		CD	00	01	10	11
		AB	00	01	10	11
00	00	0	1	3	2	
	01	4	5	7	6	
01	00	X	1	1	1	
	01	12	13	15	14	
10	00	1	1	X	X	
	10	8	9	11	10	



$G_1$

$$\begin{array}{l} A \ B \ C \ D \\ \hline 0 \ 1 \ 0 \ 0 \\ 0 \ 1 \ 0 \ 1 \\ 1 \ 0 \ 0 \ 0 \\ 1 \ 0 \ 0 \ 1 \end{array}$$

$$\overline{B \bar{C}}$$

$G_2$

$$\begin{array}{l} A \ B \ C \ D \\ \hline 0 \ 1 \ 1 \ 0 \\ 0 \ 1 \ 1 \ 1 \\ 1 \ 1 \ 1 \ 0 \\ 1 \ 1 \ 1 \ 1 \end{array}$$

$$\overline{B \ C}$$

$$\begin{array}{l} A \ B \ C \ D \\ \hline 1 \ 0 \ 0 \ 1 \\ 1 \ 0 \ 1 \ 0 \end{array}$$

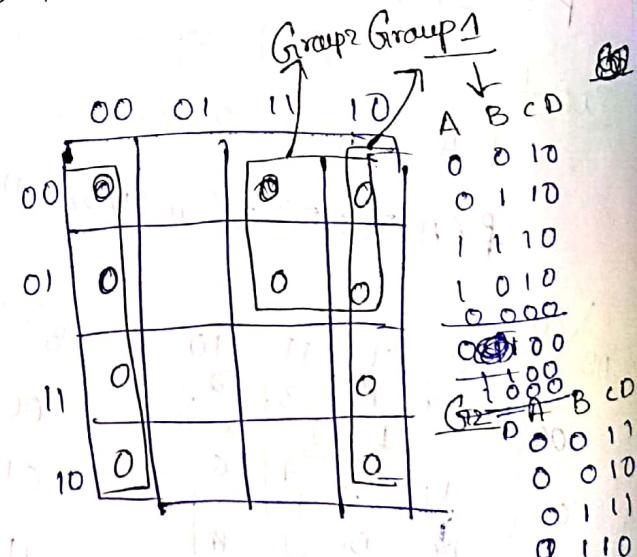
$$\overline{A \bar{B} \bar{C} \ D}$$

$$Y = A \bar{B} \bar{C} \ D + B \bar{C} + B C$$

Ex: Reduce the following function by using k-map technique

$$f(A, B, C, D) = \sum m(0, 3, 4, 7, 8, 10, 12, 14) + \sum d(2, 6)$$

		CD	00	01	11	10
		AB	00	01	11	10
00	00	0	1	0	X	2
	01	0	5	0	X	6
01	00	4	5	7	6	
	01	0	1	1	0	
11	00	12	13	15	14	
	10	0	8	9	11	10



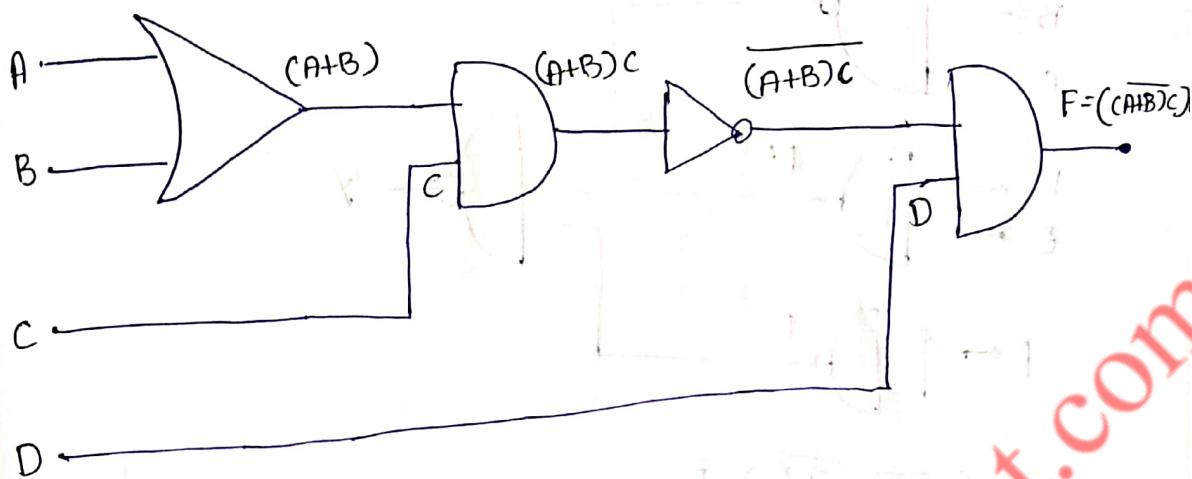
$$Y = \bar{D} + \bar{A} \ C$$

$$\begin{aligned} Y &= \bar{Y} = \overline{\bar{D} + \bar{A} \ C} \\ &= (\bar{D}) \cdot (\bar{A} \ C) \\ &= D \cdot (A + \bar{C}) \end{aligned}$$

$$D + \bar{A} \bar{C} = A + \bar{C}$$

$$* f(A, B, C, D) =$$

Ex:  $F = \overline{(A+B)}C D$  to implement logic Design?



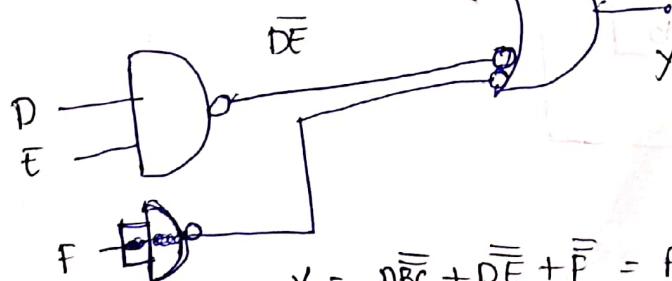
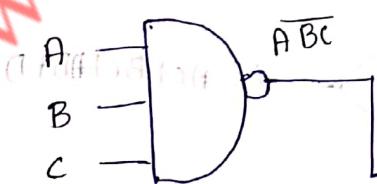
NAND to NAND Implementation:

$$Y = ABC + DEF$$

(i) AND - OR

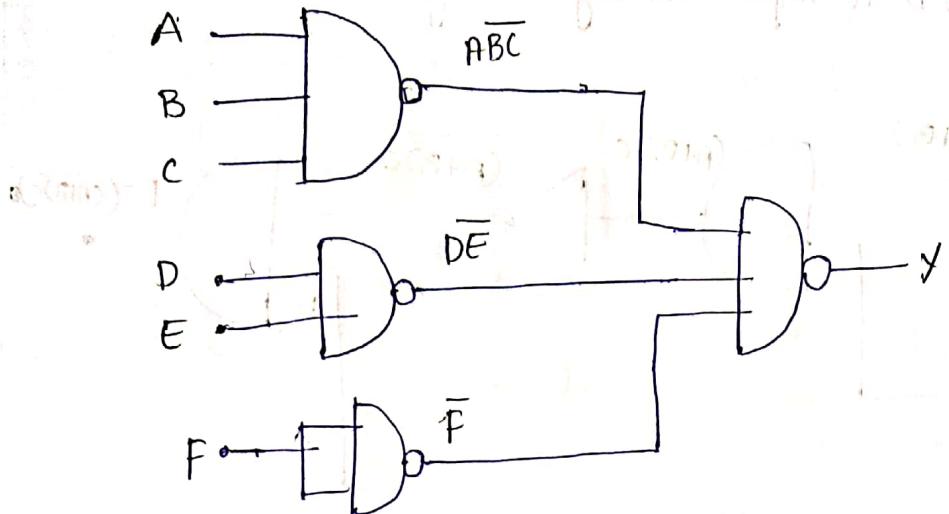


(ii) NAND - Bubble OR



$$y = \overline{ABC} + \overline{DE} + \overline{F} = ABC + DEF$$

### (iii) NAND - NAND

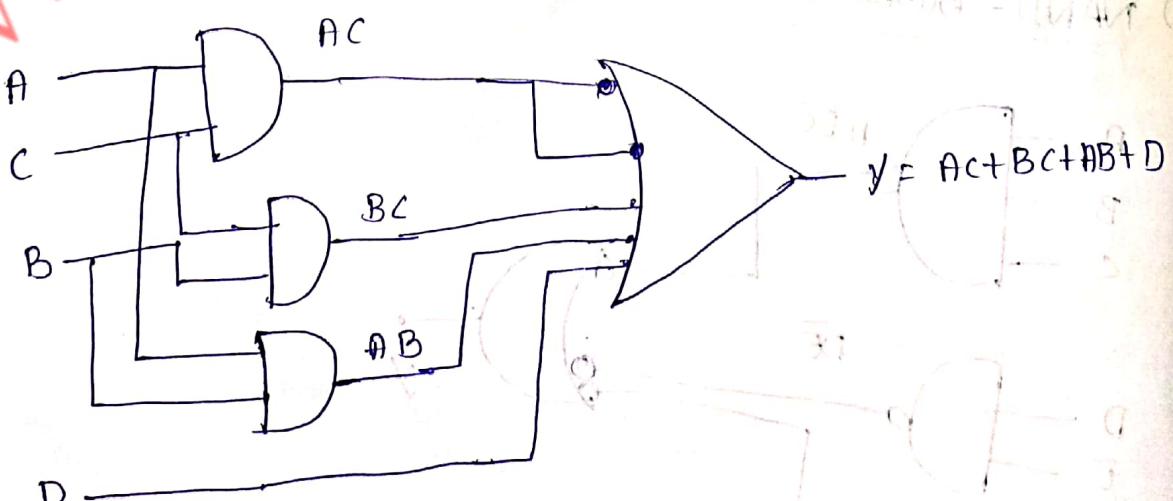


$$\begin{aligned}
 Y &= (\overline{ABC}) \cdot (\overline{DE}) \cdot \overline{F} \\
 &= (\overline{\overline{ABC}}) + (\overline{\overline{DE}}) + \overline{F} \\
 &= ABC + DE + F
 \end{aligned}$$

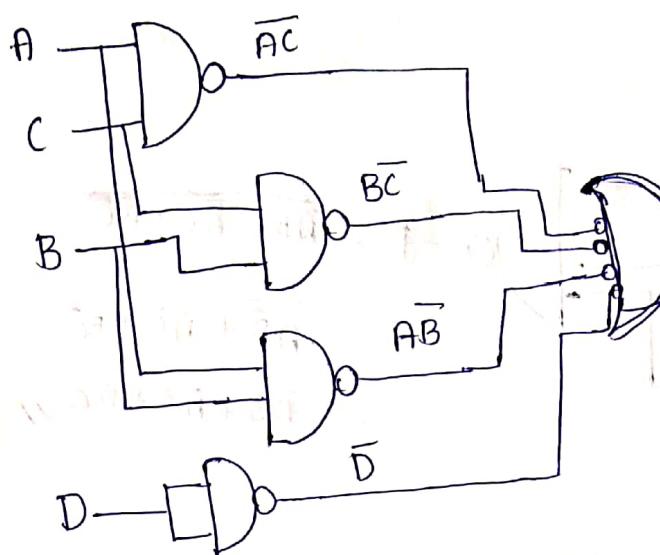
Ex:  $y = AC + ABC + \overline{ABC} + AB + D$  Implement the following Boolean function with NAND to NAND logic.

$$\begin{aligned}
 y &= AC + ABC + \overline{ABC} + AB + D \\
 &= AC + BC(A + \overline{A}) + AB + D \\
 y &= AC + BC + AB + D
 \end{aligned}$$

### (iv) AND - OR

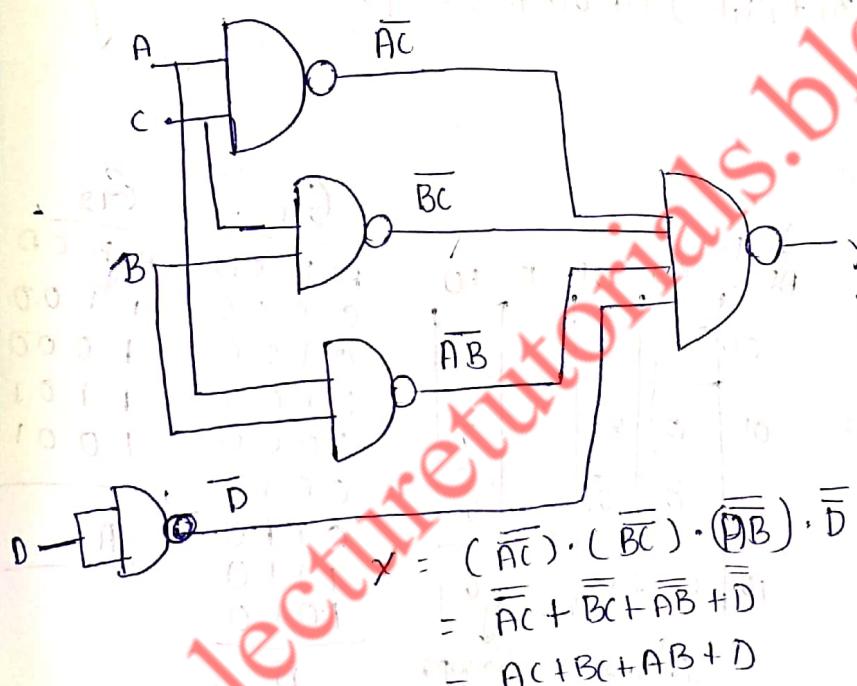


(9) NAND - Bubble OR



$$\begin{aligned}y &= \overline{\overline{AC}} + \overline{\overline{BC}} + \overline{\overline{AB}} + \overline{D} \\&= AC + BC + AB + D\end{aligned}$$

(99) NAND - NAND



$$\begin{aligned}x &= (\overline{\overline{AC}}) \cdot (\overline{\overline{BC}}) \cdot (\overline{\overline{AB}}) \cdot \overline{D} \\&= \overline{\overline{AC}} + \overline{\overline{BC}} + \overline{\overline{AB}} + \overline{D} \\&= AC + BC + AB + D\end{aligned}$$

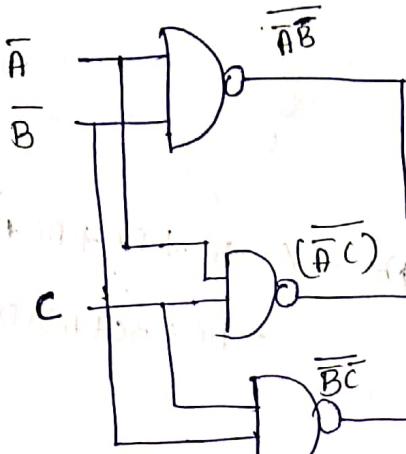
E.g. Implementation the following Boolean Expression with NAND-NAND

Logic.  $F(A, B, C) = \sum m(0, 1, 3, 5)$

		G1		G2	
		00	01	11	10
A	0	0	1	1	0
	1	4	5	7	6
		G3			

$$F(A, B, C) = \overline{A}\overline{B} + \overline{A}\overline{C} + \overline{B}\overline{C}$$

## NAND - NAND



$$F = \overline{A}\overline{B} + \overline{A}\overline{C} + \overline{B}\overline{C},$$

$$F = \overline{A}\overline{B} \cdot \overline{A}\overline{C} \cdot \overline{B}\overline{C}$$

$$F = \overline{A}\overline{B} + \overline{A}\overline{C} + \overline{B}\overline{C}$$

$$F = \overline{A}\overline{B} + \overline{A}\overline{C} + \overline{B}\overline{C} //$$

Ex: ① Find the reduced pos form of the following Equation  
 $f(A, B, C, D) = \sum m(1, 3, 7, 11, 15) + d(0, 2, 5)$  Implement and using  
 NAND logic?

CD	00	01	11	10
AB	X <sub>0</sub>	1 <sub>1</sub>	1 <sub>3</sub>	X <sub>2</sub>
00	0 <sub>4</sub>	X <sub>5</sub>	1 <sub>7</sub>	0 <sub>6</sub>
01	0 <sub>12</sub>	0 <sub>13</sub>	1 <sub>15</sub>	0 <sub>14</sub>
10	0 <sub>8</sub>	0 <sub>9</sub>	1 <sub>11</sub>	0 <sub>10</sub>

CD	00	01	11	10
AB	0 <sub>1</sub>	-	-	0 <sub>0</sub>
00	0 <sub>1</sub>	X	-	0 <sub>0</sub>
01	0 <sub>1</sub>	0 <sub>1</sub>	0 <sub>1</sub>	0 <sub>0</sub>
11	0 <sub>1</sub>	0 <sub>1</sub>	0 <sub>1</sub>	0 <sub>0</sub>
10	0 <sub>1</sub>	0 <sub>1</sub>	0 <sub>1</sub>	0 <sub>0</sub>

Group<sub>1</sub> Group<sub>2</sub>

G <sub>11</sub>	G <sub>12</sub>
A B CD	A B CD
0 0 0 0	1 1 0 0
0 1 0 0	1 0 0 0
1 1 0 0	1 1 0 1
1 0 0 0	1 0 0 1
0 0 1 0	0 0 1 0
0 1 1 0	0 1 1 0
1 1 1 0	1 1 1 0
1 0 1 0	1 0 1 0
	<u>D</u>

$$\bar{x} = \bar{D} + A\bar{C}$$

$$Y = \bar{x} = (\bar{D})(\bar{A}\bar{C})$$

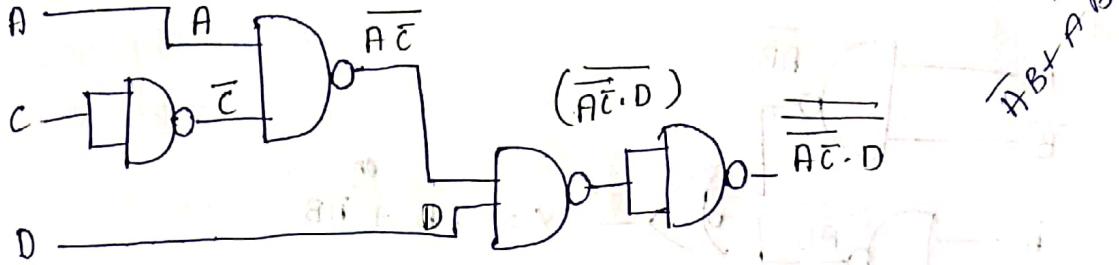
$$\therefore Y = D \cdot (\bar{A} + C)$$

Ex: ② Using k-map, determine the minimal SOP expression and realize the simplified expression using NAND logic?

$$f(w, x, y, z) = \text{ITM}(0, 2, 3, 7, 8, 9, 10)$$

$$\text{at: } f(w, x, y, z) = \sum m(1, 4, 5, 6, 11, 12, 13, 14, 15)$$

$$\therefore Y = D \cdot (\bar{A} + C)$$



$$Y = (\bar{A}\bar{C} \cdot D)$$

$$= (\bar{A} + \bar{C}) \cdot D$$

$$Y = (\bar{A} + C) \cdot D //$$

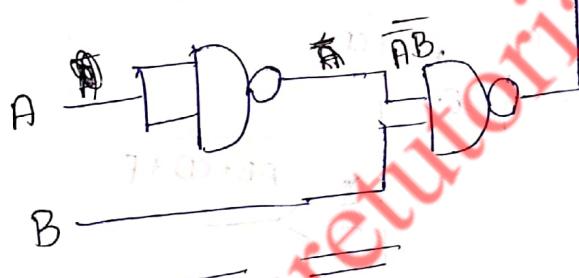
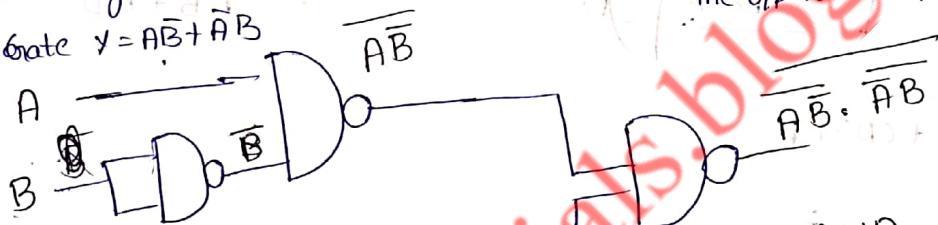
Eg: To implement X-OR gate by using NAND-NAND logic gate

$$Y = A\bar{B} + \bar{A}B$$

The X-OR gate function is two inputs and one output. The i/p variables are A and B.

The o/p variable Y

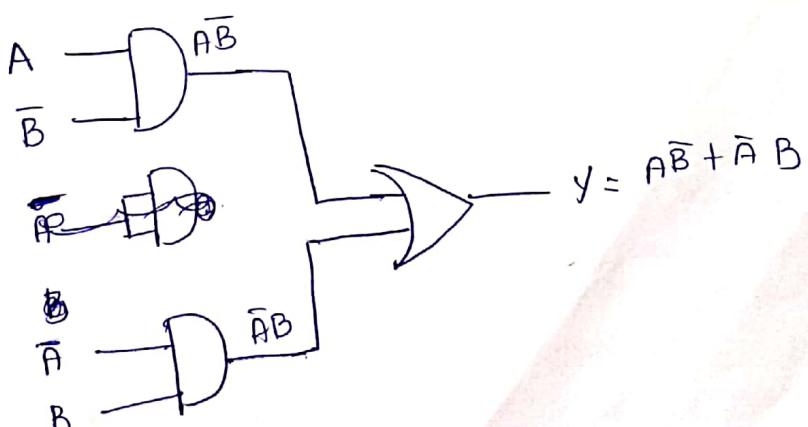
$$\text{X-OR gate } Y = A\bar{B} + \bar{A}B$$



$$Y = \bar{A}\bar{B} + \bar{A}B$$

$$Y = A\bar{B} + \bar{A}B //$$

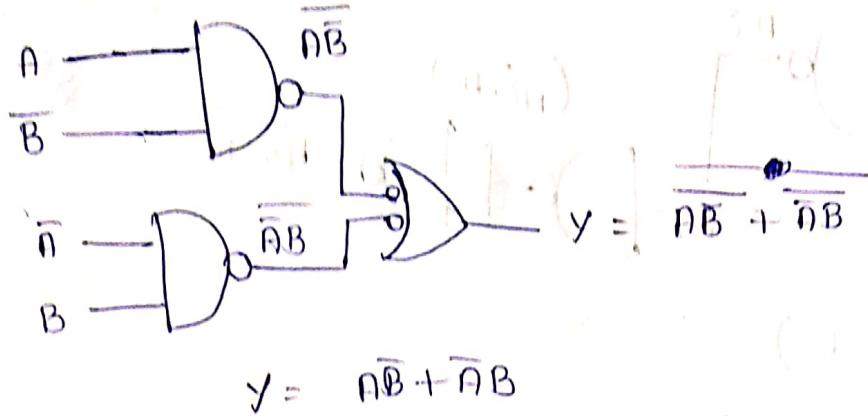
### (a) AND-OR



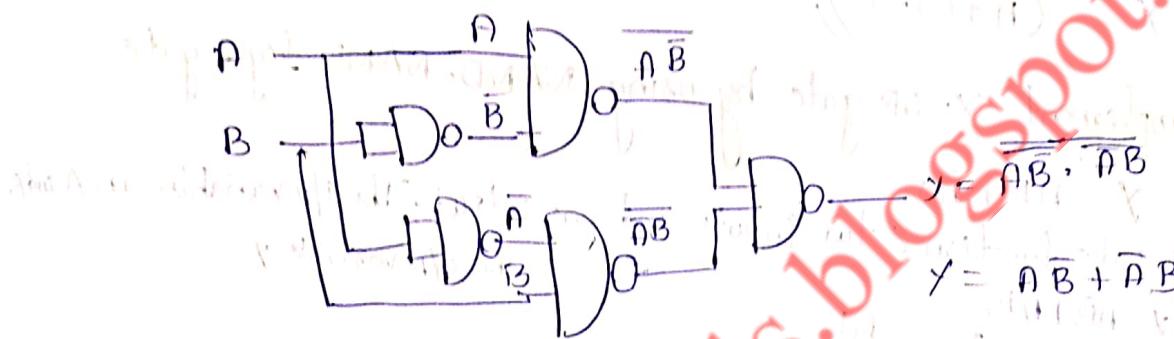
$$Y = A\bar{B} + \bar{A}B$$

## (ii) NAND - Bubbled OR

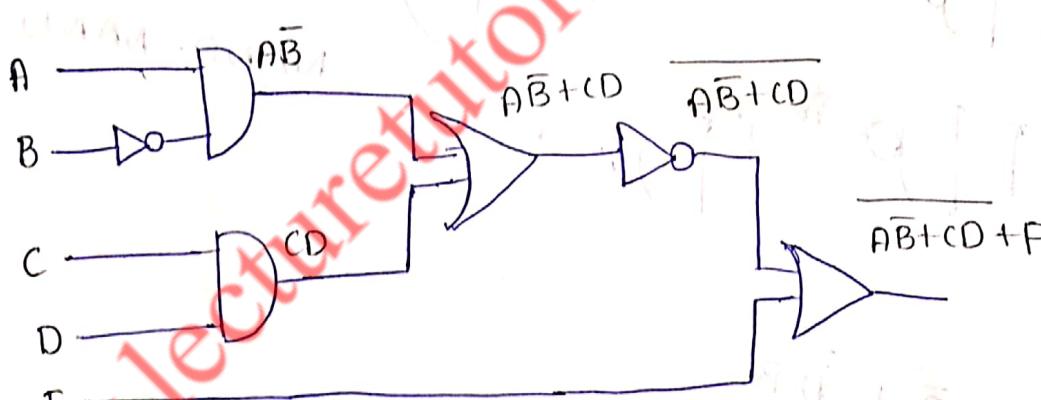
$$Y = A\bar{B} + \bar{A}B$$



## (iii) NAND-NAND



$$\text{Ex: } \overline{(AB + CD) + F}$$



Ex: Implement the following function with NAND-to-NAND logic

$$f(A, B, C, D) = \sum_m(0, 1, 3, 5)$$

$$Y = AC + BC + AB + D$$

## NOR - NOR Implementation:

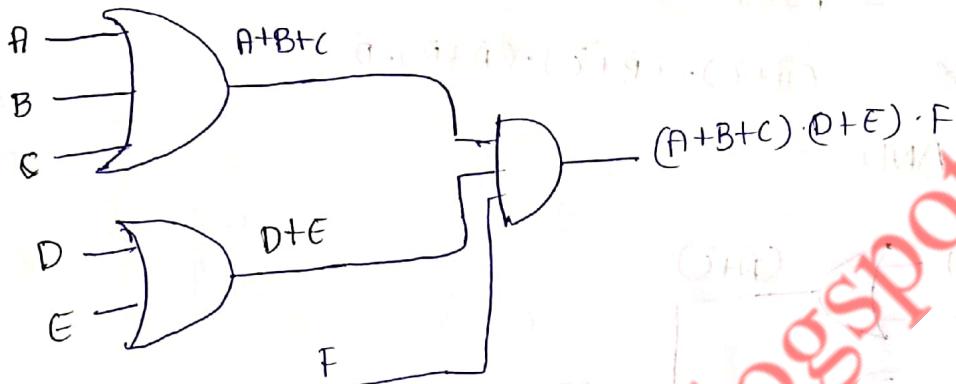
(i) OR-AND

(ii) NOR-Bubbled AND

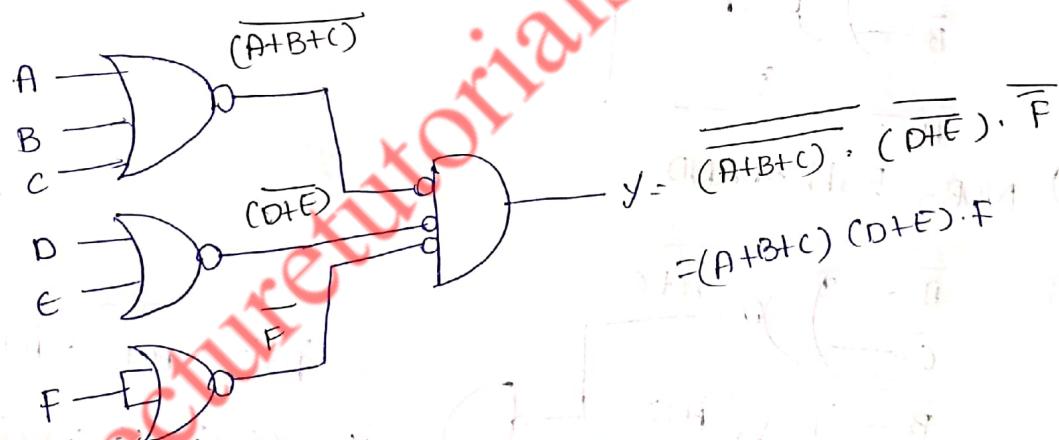
(iii) NOR-NOR

$$\text{Ex: } y = (A+B+C) \cdot (D+E) \cdot F$$

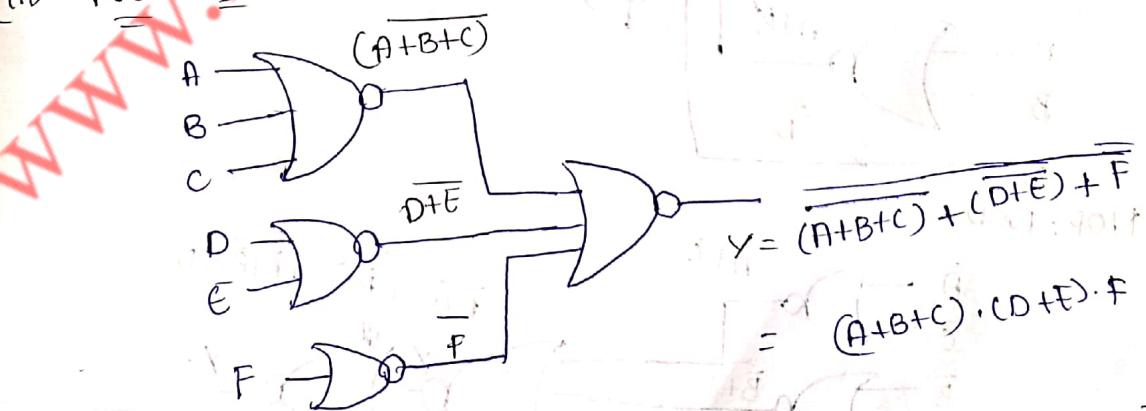
(i) OR-AND



(ii) NOR-Bubbled AND



(iii) NOR-NOR



$$\begin{aligned} \because \overline{A+B} &= \overline{A} \cdot \overline{B} \\ \overline{\overline{A}} &= A \end{aligned}$$

Ex :- Implement the following Boolean function with NOR-NOR logic

$$Y = AC + BC + AB + D$$

The given Boolean Expression

$$Y = AC + BC + AB + D$$

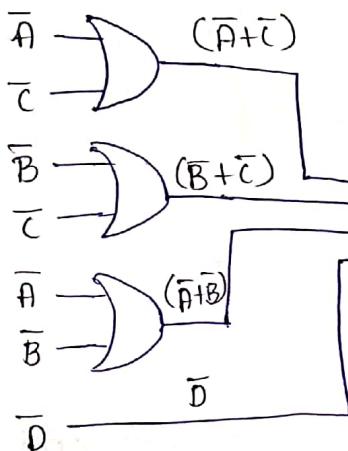
Assumption . By using duality theorem

$$\bar{Y} = \overline{AC + BC + AB + D}$$

$$= (\bar{A}\bar{C}) \cdot (\bar{B}\bar{C}) \cdot (\bar{A}\bar{B}) \cdot \bar{D}$$

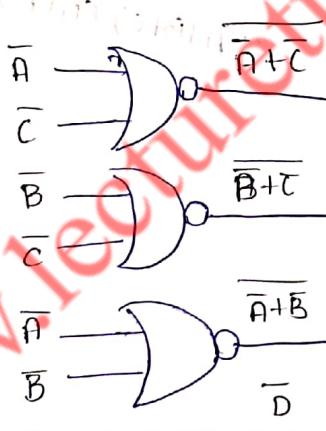
$$\bar{Y} = (\bar{A} + \bar{C}) \cdot (\bar{B} + \bar{C}) \cdot (\bar{A} + \bar{B}) \cdot \bar{D}$$

(i) OR-AND



$$\bar{Y} = (\bar{A} + \bar{C}) \cdot (\bar{B} + \bar{C}) \cdot (\bar{A} + \bar{B}) \cdot \bar{D}$$

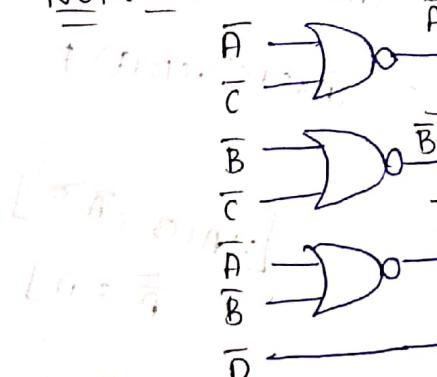
(ii) NOR-Bubbled AND



$$\bar{Y} = (\bar{A} + \bar{C}) \cdot (\bar{B} + \bar{C}) \cdot (\bar{A} + \bar{B}) \cdot \bar{D}$$

999

NOR=NOR



$$\bar{Y} = (\bar{A} + \bar{C}) + (\bar{B} + \bar{C}) + (\bar{A} + \bar{B}) + \bar{D}$$

$$Y =$$

$$\bar{A} + \bar{C} + \bar{B} + \bar{C} + \bar{A} + \bar{B} + \bar{D}$$

$$\bar{Y} = (\bar{A} + \bar{C}) \cdot (\bar{B} + \bar{C}) \cdot (\bar{A} + \bar{B}) \cdot \bar{D}$$

$$\bar{Y} = (A + C) \cdot (B + C) \cdot (A + B) \cdot D$$

$$\bar{Y} = (\bar{A} + \bar{C}) \cdot (\bar{B} + \bar{C}) \cdot (\bar{A} + \bar{B}) \cdot \bar{D} \quad [\because \bar{A} + \bar{B} = \bar{A} \cdot \bar{B} = AB]$$

$$Y = AC + BC + AB + D$$

Ex: Implement the following Boolean function with NOR-NOR logic

$$f(A, B, C) = \prod_M(0, 2, 4, 5, 6)$$

Ex: Implement the X-NOR gate by using only NOR gate ( $AB + \bar{A}\bar{B}$ )

$$\text{Ex: } f(A, B, C) = \prod_M(0, 2, 4, 5, 6)$$

A	B	C	00	01	11	10
0	0	0	0	0	1	0
1	0	0	1	0	0	1

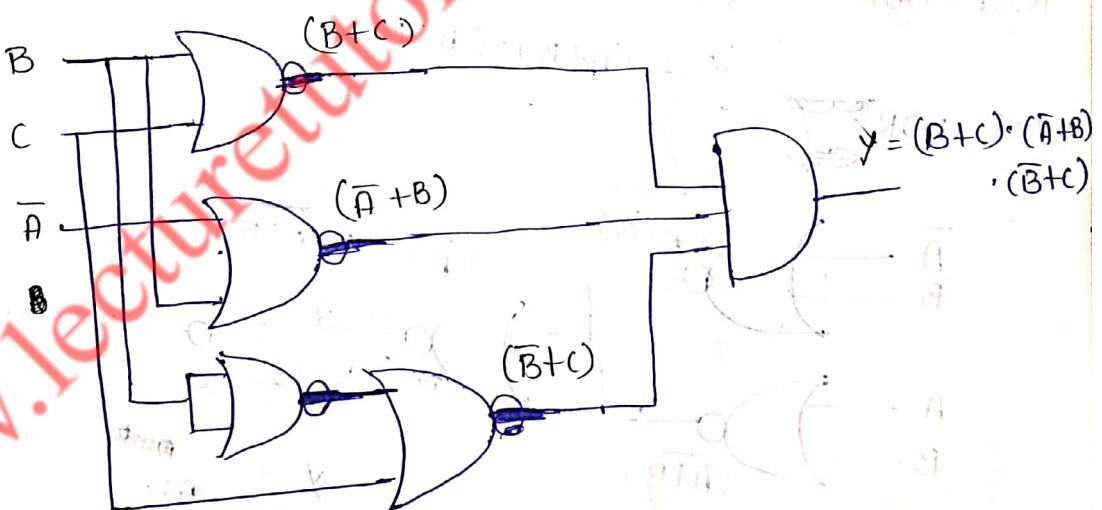
$$\begin{array}{c} G_1 \\ \hline A & B & C \\ 0 & 0 & 0 \\ 1 & 0 & 0 \\ \hline \bar{B}\bar{C} \end{array} \quad \begin{array}{c} G_2 \\ \hline A & B & C \\ 0 & 0 & 1 \\ 1 & 0 & 1 \\ \hline A\bar{B} \end{array} \quad \begin{array}{c} G_3 \\ \hline A & B & C \\ 0 & 1 & 0 \\ 1 & 1 & 0 \\ \hline B\bar{C} \end{array}$$

$$Y = \bar{B}\bar{C} + A\bar{B} + B\bar{C}$$

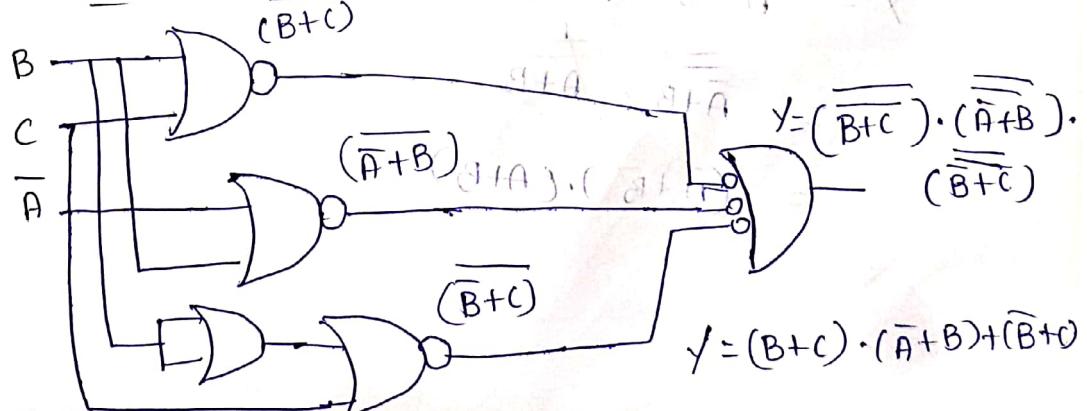
$$\bar{Y} = \bar{B}\bar{C} + A\bar{B} + B\bar{C}$$

$$Y = (\bar{B} + \bar{C}) \cdot (\bar{A} + \bar{B}) \cdot (\bar{B} + \bar{C})$$

### (i) OR-AND

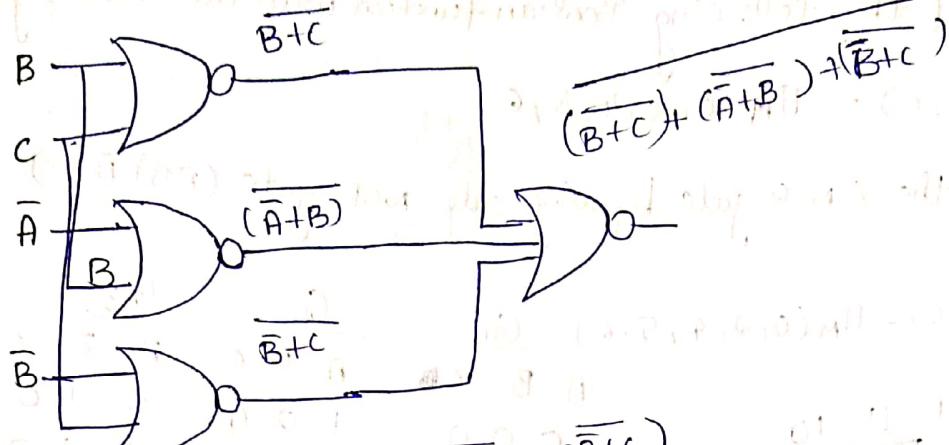


### (ii) NOR-Bubbled AND



(iii) NOR-NOR

$$Y = (B+C) \cdot (\bar{A}+B) \cdot (\bar{B}+C)$$



$$y = (\overline{B+C}) + (\overline{A+B}) + (\overline{B+C}) \\ = (\overline{\overline{B+C}}), (\overline{\overline{A+B}}), (\overline{\overline{B+C}})$$

$$Y = (B+C) \cdot (\bar{A}+B) \cdot (\bar{B}+C)$$

② Implement the X-NOR gate by using only NOR gate.

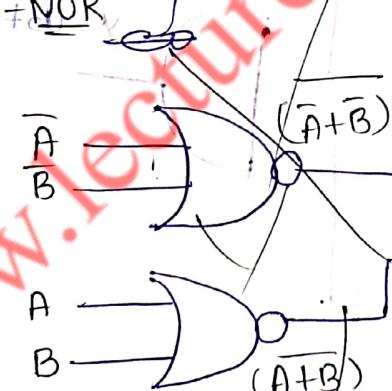
$$y = AB \oplus \bar{A} \bar{B}$$

NOR = NOR

$$\bar{Y} = \overline{A/B + \bar{A}\bar{B}}$$

$$\frac{A}{B} = (\bar{A} + \bar{B}) \cdot (A + B)$$

NOR + NOR



$$\bar{x} = \frac{A+B}{A+B}$$

$$\cdot (\overline{8+i}) \cdot (\overline{5-i}) = \overline{A+B} \quad A+B$$

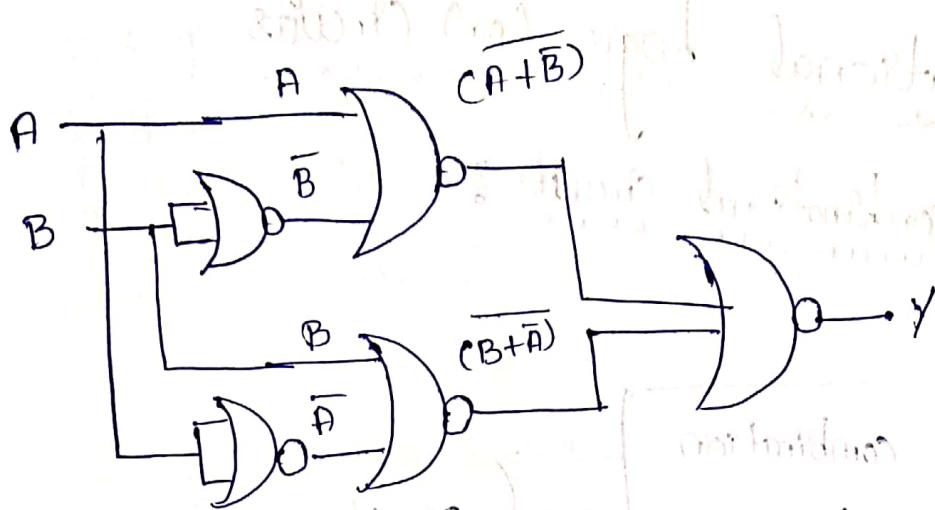
$$(\overline{5+i}) = -1 \cdot (\overline{A+B}) \cdot (A+B)$$

$$Y = AB + \bar{A} \bar{B}$$

11

Q1: Implement the  $\bar{A} \oplus B$

A-TING



$$Y = (A \cdot B) + (\bar{A} \cdot B)$$

$$Y = (\overline{A \cdot B}) \cdot (\overline{\bar{A} \cdot B})$$

$$Y = (A + \bar{B})(B + \bar{A})$$

$$= AB + \bar{B}\bar{B} + A\bar{A} + \bar{A}\bar{B}$$

$$= AB + \bar{A}\bar{B}$$

$(\text{as } B\bar{B} = 0$   
 $\text{and } A\bar{A} = 0)$

