

Sorting :

- Understand Sorting
- Few Problems on Sorting
- 1 Sorting Algo ↪
- Comparison _____

Sorting: Arranging data in increasing / decreasing order based
on a parameter

Ex1: 2 4 7 11 15 : Inc, par = value

Ex2: 15 9 6 2 0 : Desc, par = value

Ex3: 1 2 3 7 4 9 6 : Sorted in Inc based on

#fastr: 1 2 2 2 3 3 4 factors

// Inbuilt library:

[↳ sort(), In every language sort() func present
↳ how? logic? ↳ In Adv Content]

Tc: $N \log N$: N represents, # Elements to Sort

Q8] Elements Removal

Given N elements, at every step remove an array element

Cost to remove element = Sum of array elements present in array

Final Min Cost to Remove all Elements?

Note: Add Cost of Removal first then Remove it

$$\underline{\text{Ex1: }} \text{ar}[3] = \{2 \ 1 \ 4\} \quad \text{ar}[3] = \{2 \ 1 \ 4\}$$

$$\begin{array}{l} \text{Cost} \\ \text{remove } 2 : 7 \end{array}$$

$$\text{remove } 1 : 5$$

$$\text{remove } 4 : 4$$

$$\underline{\text{Total Co : 16}}$$

$$\begin{array}{l} \text{Cost} \\ \text{remove } 4 : 7 \end{array}$$

$$\text{remove } 2 : 3$$

$$\text{remove } 1 : 1$$

$$\underline{\text{Total Co : 11}}$$

$$\underline{\text{Ex2: }} \text{ar}[4] = \{3 \ \underline{6} \ 2 \ 4\}$$

$$\begin{array}{l} \text{Cost} \\ \text{remove } 6 : 15 \end{array}$$

$$\text{remove } 4 : 9$$

$$\text{remove } 3 : 5$$

$$\text{remove } 2 : 2$$

$$\underline{\text{Total Cost : 31}}$$

$$\underline{\text{Ex3: }} \text{ar}[3] = \{6 \ -3 \ 4\}$$

$$\begin{array}{l} \text{Cost} \\ \text{remove } 6 : 7 \end{array}$$

$$\text{remove } 4 : 1$$

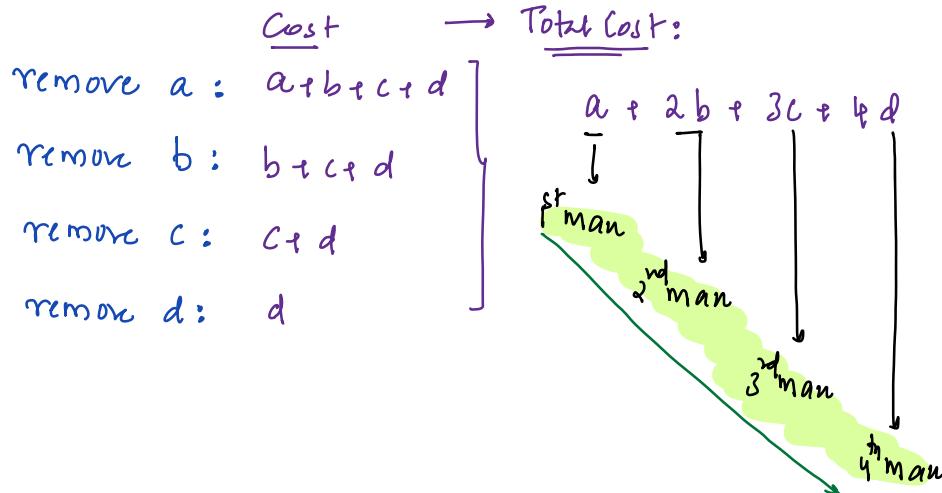
$$\text{remove } -3 : -3$$

$$\underline{\text{Total Cost : 5}}$$

$$\text{ar}[4] = \begin{bmatrix} 0 & 1 & 2 & 3 \\ 6 & 4 & 3 & 2 \\ * & * & * & * \\ 1 & 2 & 3 & 4 \end{bmatrix}$$

// obs: deleting ele by ele in decreasing to get min cost?

Ex: arr[4]: ~~a b c d~~



int MinCost(int arr[]){

 int n = arr.length;

Sort(arr, desc) // sort arr[] in desc order

 int c = 0;
 // Syntax, In your language of choice

 i = 0; i < n; i++) {

 // How many times arr[i] will in TS = $i+1$ times

 c = c + arr[i] * (i+1)

}

 return c;

T_C: $(N \log N + N) \Rightarrow O(N \log N)$

Q8) Noble Integers { Data is Distinct }

Given N array element, Calculate no: of Noble Integers

An element ele in arr[] is said to be Noble if

↳ obs1: Negative { No of elements < ele = ele itself }
 Can never be noble ↳ Count = ele

Ex1: { -1, -5, 3, 5, -10, 4 } // \Rightarrow ans: 3

#less: 2, 1, 3, 5, 0, 4

Ex2: { [-3, 0, 2, 5], [0, 1, 2, 3] } // \Rightarrow ans: 1

#less: 0, 1, 2, 3

Ex3: { [-10, -5, 1, 3, 4, 5, 10], [0, 1, 2, 3, 4, 5, 6] } // \Rightarrow ans: 3

#less:

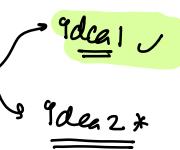
// Ideal:
 for every element in arr[]
 iterate on arr[] & get
 number of elements < than
 & compare with ele itself

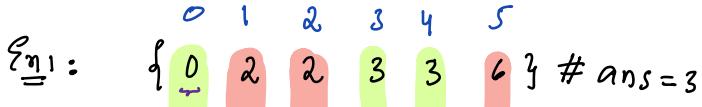
```
int ans=0
for(i=0; i<N; i++) {
    int c=0 // arr[i] need elements <
    for(j=0; j<N; j++) {
        if(arr[i] > arr[j]) {
            c=c+1
        }
    }
    if(arr[i] == c) { ans=ans+1 }
}
```

Tc: O(N^2) Sc: O(1)

Pdear2:
 Sort the array inc
 #elements < arr[i] = i
 int Noble(int arr[]) {
 int n=arr.length
 Sort(arr[], inc) // Sort in inc
 int c=0;
 for(i=0; i<n; i++) {
 if(arr[i] == i) {
 c=c+1
 }
 }
 return c;
 }

Tc: (NlogN + N) = O(NlogN)

// 3Q) Nobie Integers : { Data can repeat } 

E_{n1} : { 0 1 2 3 4 5 } # ans = 3


E_{n2} : { -10 1 1 2 3 4 4 5 6 7 8 } # ans = 7
less: { 0 1 1 1 4 4 4 4 7 8 }

E_{n3} : { -10 1 1 2 4 4 4 8 10 } # ans = 5
less: { 0 1 1 3 4 4 4 7 8 }

E_{n4} : { [0] [-3] [1] [0] [2] [2] [3] [2] [4] [5] [5] [5] [5] [6] [5] [5] [5] [7] [8] [8] [8] [9] [10] [10] [10] [10] [11] [10] [10] [12] [10] [13] [14] } # ans = 7
less: { 0 1 2 2 2 4 4 4 4 8 8 8 10 10 10 10 10 11 10 10 12 10 13 }

Obs1: If element coming first time, if $ar[i] \neq ar[i-1]$

Count of elements, less than that will be, $c = i$

Obs2: If element repeats, if $ar[i] == ar[i-1]$

Count of elements, less than that will remain same

```
int NobleDistinct(int arr[]){
```

```
    int n = arr.length;
```

```
    int c = 0;
```

```
Sort(arr, n) // TODO
```

```
int less = 0 // # no of elements less than arr[i]
```

```
if (arr[0] == 0) { c = 1 } // noble number } Edge case }
```

```
for (int i = 1; i < n; i++) {
```

```
// for arr[i] get no of elements less than that?
```

```
if (arr[i] != arr[i-1]) { // i=0, arr[0] != arr[-1], error
```

```
// arr[i] coming first time, # elements a = i
```

```
    less = i
```

```
}
```

```
else { // skip else
```

```
// arr[i] is repeating, less won't change
```

```
    // Do nothing
```

```
if (less == arr[i]) {
```

```
    c = c + 1;
```

```
}
```

```
return c;
```

$T_C: O(N \log N + N) \approx O(N \log N)$

$$ar[9] = \left\{ \begin{array}{cccccccccc} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ \hline 1 & 1 & 2 & 2 & 4 & 4 & 5 & 7 & 7 \\ \hline l=0 & * & l=0 & l=2 & l=2 & l=4 & l=4 & l=6 & l=7 \\ l=0 & l=0 & l=1 & l=2 & l=3 & l=4 & l=4 & l=5 & l=6 \end{array} \right\} \Rightarrow \underline{\text{ans}} = 2$$

Sorting algo: { Bubble Sort }

0 1 2 3 4
// arr[]: { 3 8 6 2 4 }

Iteration 1: 3 6 2 4 8 : 8 at correct position

Iteration 2: 3 2 4 6 8 : 6, 8 at correct position

Iteration 3: 2 3 4 6 8 : Complete array sorted

Idea: Observe

After 1 iteration → last 1 element in correct pos

After 2 iterations → last 2 elements in correct pos

After 3 iterations → last 3 elements in correct pos

After N-1 iterations → last N-1 elements in correct pos

All N elements in correct pos

In every iteration, → { // keep doing it N-1 times }
iterate from first to last Compare adj elements,
if, they are not in correct order we swap

Bubble Sort (int arr[]) { TC: O(N²) SC: O(1) }

int n = arr.length;

for (int i = 1; i < n; i++) { // N-1 times }

 for (int j = 0; j < n-1; j++) { // if j = N-1

 if (arr[j] > arr[j+1]) {

 Swap arr[j] & arr[j+1]

 arr[N-1] > arr[N]

Error

We can slightly optimize this

// Decreasing:

```
int n = ar.length;
for(int i = 1; i < n; i++) {
    for(int j = 0; j < n - 1; j + i) {
        if(ar[j] < ar[j + 1]) {
            Swap ar[j] & ar[j + 1]
        }
    }
}
```

has power to decide order
of data itself

// If we can, just change the if (condition), it allows our own, sorting order

→ Idea:

```
int n = ar.length;
for(int i = 1; i < n; i++) {
    for(int j = 0; j < n - 1; j + i) {
        if(comp(ar[i], ar[j])) {
            Swap ar[j] & ar[j + 1]
        }
    }
}
```

```
bool comp(int a, int b) {
    // Just implement to get your own sorted order
    // If a should come before b, or b should come before a
}
```

Comparator in Java/ JS

```
Comparator customComparator = (Integer a, Integer b) => {  
    if you want a to come before b : return -1  
    if a & b are same : return 0  
    if you want b to come before a : return 1  
}
```

```
Sort(ar, customComparator)
```

Comparator in Python

```
def compare(a, b):  
    if a should come before b : return -1  
    if a & b same : return 0  
    if b should come before a : return 1  
}
```

```
ar.sort(key = cmp_to_key(compare))
```

Comparator in C++

```
bool comp(a, b){  
    if a should come first : return True  
    else return False  
}
```

```
Sort(ar, comp)
```

1) Given N, Element sort them in increasing order of their
 {No. of factors}

If 2 elements have same No. of factors, elements with
 less value should come first

Note: → {No Extra Space}

ar[]	→	{ 9 3 4 8 16 37 6 13 15 }
#fac		{ 3 2 3 4 5 2 4 2 4 }

Correct order: { 3 13 37 4 9 6 8 15 16 }

Idea:

given ar[]

Sort(ar, comp)

your own order

TC → In Adv

```
int comp (int a, int b){ Java / JS / Python
    int fa = factors(a)           ↗ factors in a
    int fb = factors(b)           ↗ implement factors on
                                  ↗ your own → TODO
    if (fa < fb) {               ↗
        // a to come first
        return -1;                ↗ C++ return True
    }
    else if (fa == fb && a < b) {
        return -1;                ↗ C++ return True
    }
    else return 1;                ↗ C++ return False
}
```

