

Todays Content:

- IsPrime
- Prime Sieve
- Count factors
- spf[]
- Segmented Sieve [Optimal Content]

IsPrime(): #factors = 2

$N=10$: Not prime

$N=7$: prime

IsPrime(): Calculate no: of factors

↳ TC1: $O(N)$: $1..N$

↳ TC2: $O(\sqrt{N})$: $1..\sqrt{N}$

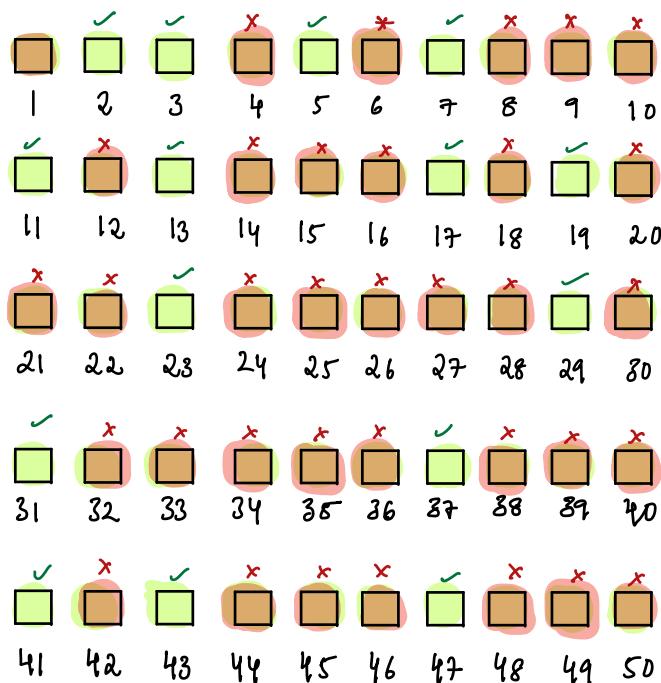
Print all Primes from 1-N:

$N=10$: Output: 2 3 5 7

Ideal1: Iterate from $1-N$ & check if every number prime or not

TC: $N * \sqrt{N} \approx O(N\sqrt{N})$ ↳ $O(\sqrt{N})$
↳ TC for isprime()

Ideal2: Say we need all prime 1-50



Pseudocode: Sieve of Eratosthenes or Prime Sieve

```
void allPrime(int n) {
```

```
    bool p[n+1] = {True} // 1-based indexing
```

```
    p[1] = false
```

```
    i = 2; j; i <= N; i++) {
```

```
        // iterate on multiples of i, if i is a prime number
```

```
        if (p[i] == True) {
```

```
            // multiple of i = 2i  $\rightarrow$  3i  $\rightarrow$  4i ....
```

```
j = 2i; j <= n; j = j + i) {
```

```
        // j is not prime
```

```
p[j] = false
```

```
}
```

```
// iterate on p[] & get indices with True
```

```
}
```

TC: Given N :

i	j	Iteration
2	mul of 2 till N	$N/2$
3	mul of 3 till N	$N/3$
4	x	
5	mul of 5 till N	$N/5$
6	*	
7	mul of 7 till N	$N/7$
8		
11	mul of 11 till N	$N/11$

Total Iterations:

$$= \frac{N}{2} + \frac{N}{3} + \frac{N}{5} + \frac{N}{7} + \frac{N}{11} + \dots$$

$$N \left[\frac{1}{2} + \frac{1}{3} + \frac{1}{5} + \frac{1}{7} + \frac{1}{11} + \dots \right]$$

[Sum of inverse of primes till N]
 $\log(\log N)$ mathts: ref

TC: $N(\log(\log N))$ SC: $O(N)$

Optimization to above idea :

$$\underline{N=25}$$

1 mut of i

1 :

2 : $2 \times 2 \quad 2 \times 3 \quad 2 \times 4 \quad 2 \times 5 \dots$

3 : $\boxed{3 \times 2} \quad \boxed{3 \times 3} \quad \cancel{3 \times 4} \quad \cancel{3 \times 5} \dots$

4 :

5 : $\boxed{5 \times 2} \quad \boxed{5 \times 3} \quad \boxed{5 \times 4} \quad \cancel{5 \times 5} \dots$

6 :

7 : $\boxed{7 \times 2} \quad \boxed{7 \times 3} \quad \boxed{7 \times 4} \quad \boxed{7 \times 5} \quad \boxed{7 \times 6} \quad \cancel{7 \times 7} \dots$

void allPrimeC(pnt n) { TC: $N \log(\log N)$ SC: $O(N)$

 bool p[N+1] = {True} // 1-based indexing

 p[1] = False

 i = 2; j = i * i = N; i++ {

 // iterate in multiples of i, if i is a prime number

 if (p[i] == True) {

 // multiple of i = $2i \rightarrow 3i \rightarrow 4i \dots$

 j = i^2 ; j = N; j = j + i {

 // j is not prime

 p[j] = False

↳ To Do:

Construct Table you
will understand

if Enter 2nd loop

i = : $j = (i^2)$

\sqrt{N} N ✓

$\sqrt{N+1}$ $\frac{[\sqrt{N+1}][\sqrt{N+1}]}{2} *$

$\frac{N+2\sqrt{N+1}}{2} > N$

*

3 // iterate in p[] & get indices with True

Find no: of factors for all [1-N]

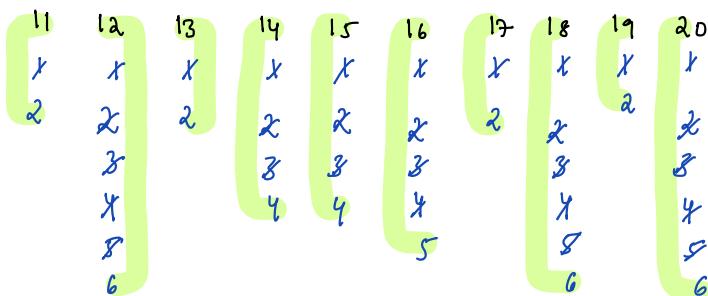
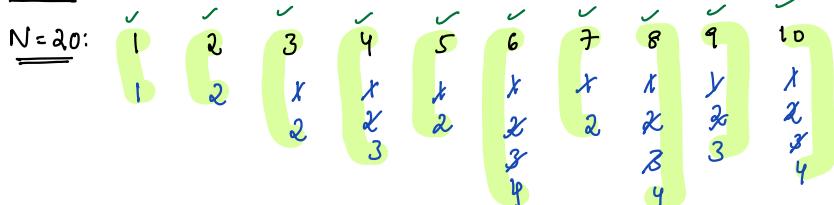
N=10: 1 2 3 4 5 6 7 8 9 10

factors: 1 2 2 3 2 4 2 4 3 4

Ideas: for every number from [1-N] find no: of factors

$$\underline{\underline{TC}}: N * O(\sqrt{N}) = O(N\sqrt{N})$$

Ideas:



int[] All factors (int n) { Note: count of factors of N = f[N] }

int f[n+1] = {1} // Because for all numbers 1 is factor

f[0] = // No: of factors for 0, depends on question

i=2; j=i; i++; } i is a factor to itself

// multiply of i = i, 2i, 3i, 4i, ... -

j = i; j++; j = j + i; }

// i is a factor of j, inc count by 1

f[j] = f[j] + 1

return f

Table Given N

<u>i</u>	<u>j</u>	<u>iter</u>
2	Itc mul of 2	$N/2$
3	Itc mul of 3	$N/3$
4	Itc mul of 4	$N/4$
5	Itc mul of 5	$N/5$
6	Itc mul of 6	$N/6$
.	.	
.	.	
N	Itc mul of N	N/N

Total Iter =

$$= N/2 + N/3 + N/4 + \dots + N/N$$

$$= N \left[\frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \frac{1}{5} + \dots + \frac{1}{N} \right]$$

= sum of all reciprocals from [2-N]

$$= \underline{\underline{\log N}}$$

Tc: $N \log N$ SC: $O(N)$

Spf: Find smallest prime factor for all numbers from 1-N

$N=10$: 1 2 3 4 5 6 7 8 9 10

S_{PFTJ} : 0 2 3 2 5 2 7 2 3 2

Idea: Say we need spf for all numbers from 1-50

1	2	3	4	5	6	7	8	9	10
11	2	13	2	3	2	17	2	19	2
11	12	13	14	15	16	17	18	19	20
21	22	23	24	25	26	27	28	29	2
31	2	3	2	5	2	37	2	8	2
31	32	33	34	35	36	37	38	39	40
41	2	43	2	3	2	47	2	7	2
41	42	43	44	45	46	47	48	49	50

`int[] AllSpf(int N) {` $TC: O(N \cdot N \log(\log N))$ $SC: O(N)$

`int SPF[N+1] = {0};`

`i = 0; j <= N; i++) {` $\rightarrow TC: O(N)$

`// Initialize SPF[i] = i`
 `j = 2;`

`while (j <= N; j++) {` $\rightarrow TC: O(N \log(\log N))$

`// Is i a prime number? i can be SPF to all mul of i`

`if (SPF[i] == i) {` \rightarrow // goto mul of i

`j = i * i;` $j <= N; j = j + i;$

`SPF[j] = min(i, SPF[j]);`

`}`
 `return SPF;`

i mul of i

$\#SPF = 2$

`2 :` $\overbrace{2 \times 2 \quad 2 \times 3 \quad 2 \times 4 \quad 2 \times 5 \quad 2 \times 6 \quad 2 \times 7 \dots}$

`3 :` $\overbrace{\begin{matrix} 3 \times 2 \\ 3 \times 3 \end{matrix} \quad 3 \times 4 \quad 3 \times 5 \quad 3 \times 6 \quad 3 \times 7 \dots}$

`4 :`

`5 :` $\overbrace{\begin{matrix} 5 \times 2 \\ 5 \times 3 \\ 5 \times 4 \end{matrix} \quad \begin{matrix} 5 \times 5 \\ 5 \times 6 \end{matrix} \quad 5 \times 7 \dots}$

`6 :`

`7 :` $\overbrace{\begin{matrix} 7 \times 2 \\ 7 \times 3 \\ 7 \times 4 \end{matrix} \quad \begin{matrix} 7 \times 5 \\ 7 \times 6 \end{matrix} \quad 7 \times 7 \dots}$

Properties:

If N is a prime it will have no factor $[2 \sqrt{N}]$

for N : To get no: of factors

$i = 1; i^*i <= N; i++ \}$

if ($N \% i == 0$) {

 if ($i == N/i$) {

$c = c + 1$

 else { $c = c + 2$ }

}

$N = 7$

<u>i</u>	<u>$i^*i <= 7$</u>	<u>i</u>	<u>N/i</u>	<u>$c = 0$</u>
1	✓	1	7	$c = c + 2$
2	✓	2		
3	*	return 2		

$N = 23$

<u>i</u>	<u>$i^*i <= 23$</u>	<u>i</u>	<u>N/i</u>	<u>$c = 0$</u>
1	✓	1	23	$c = c + 2$
2	✓	2		not factor
3	✓	3		not factor
4	✓	4		not factor
5	*	return 2		

Claims:

If N is a prime it will have no factor $[2 \sqrt{N}]$

→ If N is not prime it will have at least 1 factor $[2 \sqrt{N}]$

1 prime

Segmented Sieve:

Given $[a, b]$ find no: of primes in range $[a, b]$

Constraints:

$$1 \leq a, b \leq 10^9$$

$$1 \leq b - a \leq 10^5$$

Note: In general max array size
we can declare is $10^5 - 10^6$

Idea: Get all primes from $[1 - b]$ using sieve

TC: $b \log(\log(b))$ SC: $\mathcal{O}(b)$ $b = 10^9 \rightarrow p[10^9] +$

Ex: $a = 100, b = 130, \sqrt{130} = 11 = \{2, 3, 5, 7, 11\}$ bool $p[31]$

100 * 0

101 1

102 2

103 3

104 4

105 5

106 6

107 7

108 8

109 9

110 10

111 11

112 12

113 13

114 14

115 15

116 16

117 17

118 18

119 19

120 20

121 21

122 22

123 23

124 24

125 25

126 26

127 27

128 28

129 29

130 30

$\text{prmul of } i : [a = 100, b = 120]$

$\text{prmul of } a : 100$

$$\text{prmul of } 3 : \left[\frac{100}{3} \right] * 3 + 3 = 33 * 3 = 99 + 3 = 102$$

$\text{prmul of } 5 : 100$

$$\text{prmul of } 7 : \left[\frac{100}{7} \right] * 7 + 7 = 98 + 7 = 105$$

$$\text{prmul of } 11 : \left[\frac{100}{11} \right] * 11 + 11 = 99 + 11 = 110$$

$\text{prmul of } i : \text{if } (a \% i == 0) \{ \text{prmul} = a \}$

$$\text{else } \{ \text{prmul} = \left[\frac{a}{i} \right] * i + i \}$$

```
int SegPrime(int a, int b) {
```

// Step: 1 → Get all prime from $[1, \sqrt{b}]$ & store in list

↳ Solve this using seive

↳ // Say we are storing only prime in list prime
list<int> primes // All prime are $[1, \sqrt{b}]$

// Step: 2 → ch[b-a+1] = True

i = 0; j < primes.size(); i++) {

p = prime[i]

// Get 1st mul of p in [a b]

if (a % p == 0) { mul = a; }

else { mul = [a/p]*p + p; }

if (mul == p) { mul = mul + p; }

j = mul; j <= b; j = j + p) {

ch[j - a] = false;

}

}

All primes $[1, \sqrt{50}] = 1, 2, 3, 5, 7, 11$

Edge Case: no. of primes $[1, 50] = 12, 3, 5, 7, 11$

$\begin{bmatrix} a & b \\ 1 & 50 \end{bmatrix}$

1st mul of 2: $[a/2]*2 + 2 = \boxed{[1/2]} * 2 + 2 = 0 + 2 = \underline{\underline{2}} \rightarrow 4$

1st mul of 3: $[a/3]*3 + 3 = \boxed{[1/3]} * 3 + 3 = 0 + 3 = \underline{\underline{3}} \rightarrow 6$

1st mul of 5: $[a/5]*5 + 5 = \boxed{[1/5]} * 5 + 5 = 0 + 5 = \underline{\underline{5}} \rightarrow 10$

1st mul of 7: $[a/7]*7 + 7 = \boxed{[1/7]} * 7 + 7 = 0 + 7 = \underline{\underline{7}} \rightarrow 14$

$$1 \leq a \leq b = 10^9$$

$$1 \leq b - a \leq 10^5$$

Step 1: List of all primes $[1 - \sqrt{10^9}] = \checkmark$

Step 2: $\text{ch}[b - a + 1] = \text{ch}[10^5] = \checkmark$