

Todays Content:

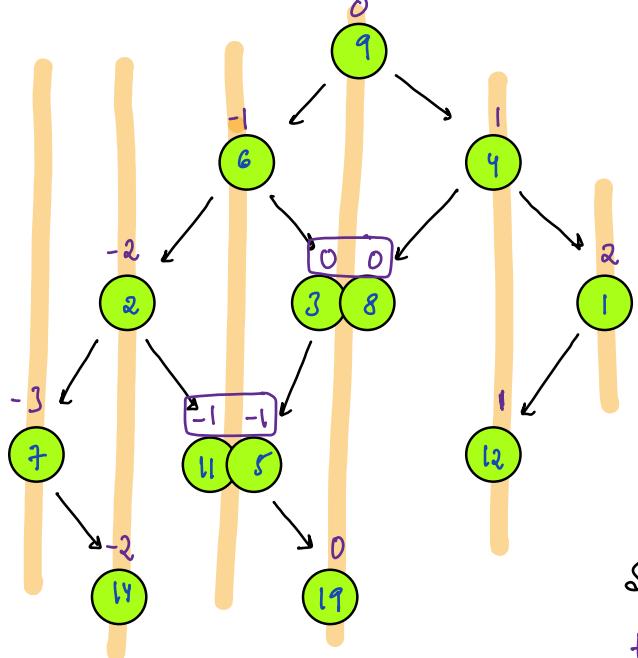
a) Vertical level order traversal

$\left. \begin{array}{l} \rightarrow \text{top view} \\ \rightarrow \text{bottom view} \\ \rightarrow \text{diagonal view} \end{array} \right\}$

b) Construct BT from Inorder & preorder

c) Type of Trees → notes → new class, doubt session

Vertical level order traversal (left → right)



Expected:

```

+
2 14
6 11 5
9 3 8 19
4 12
1

```

Storing data to print

HashMap<int, list<int>>

-3:	+ ↗ level
-2:	2 14 ↗ level
-1:	6 11 5 ↗ level
0:	9 3 8 19 ↗ level
1:	4 12 ↗ level
2:	1 ↗ level

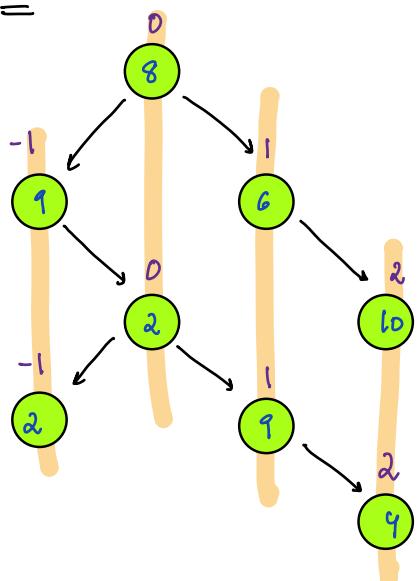
minlevel maxlevel

-3 2

Iterate on all levels from -3 to 2 ↗

print nodes level by level

$\varrho_n:$



hashmap

-1: 9 2
0: 8 2
1: 6 9
2: 10 4

Inorder hashmap *

-1: 9 2
0: 2
1: 9
2: 4

preorder hashmap *

-1: 9 2
0: 8 2
1: 9
2: 4

postorder hashmap *

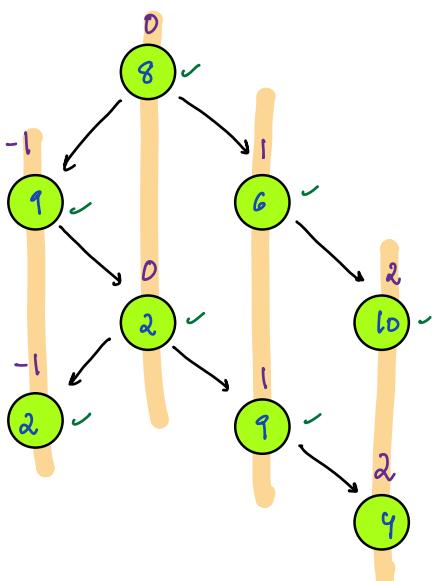
-1:
0:
1:
2:

hashmap level order

8, 07	9, -17	6, 17	2, 07	10, 27
-------	--------	-------	-------	--------

2, -17	9, 17	4, 27
--------	-------	-------

-1: 9 2
0: 8 2
1: 6 9
2: 10 4



```
void Vertical level ( Node root ) { TC: O(N) SC: O(N)
```

```
    hashmap<int, list<int>> hm
```

```
    int minL = 0, maxL = 0
```

```
    queue<pair<node, int>> q
```

node ref ↪ ↪ vertical level of each node

```
    q. enqueue ({root, 0})
```

```
    while ( q.size() > 0 ) {
```

```
        pair<Node, int> data = q.front()
```

```
        q.dequeue()
```

```
        Node t = data.first
```

```
        int l = data.second
```

```
        hm[l].insert(t.data) // insert in hashmap
```

```
        minL = min(l, minL)
```

```
        maxL = max(l, maxL)
```

```
        if ( t.left != NULL ) {
```

```
            q.enqueue ({t.left, l-1})  
        }
```

```
        if ( t.right != NULL ) {
```

```
            q.enqueue ({t.right, l+1})  
        }
```

```
}
```

```
// All nodes at vertical level = i + hm[i]
```

```
// print list hm[i] { TODO }
```

```
// print 1st element hm[i] { Top View }
```

```
} // print last element hm[i] { Bottom View }
```

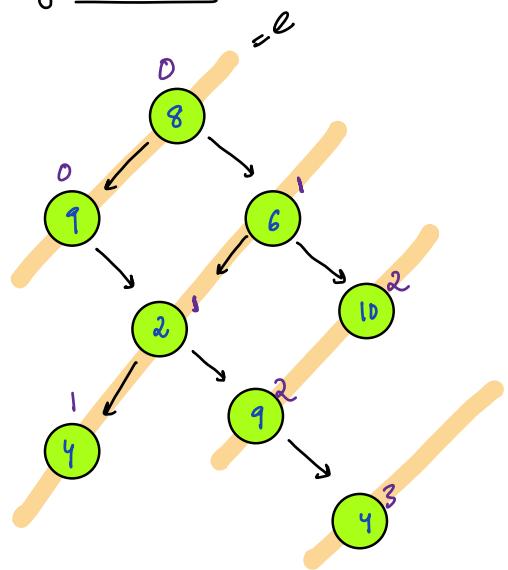
Obs:

Top view: 1st node in every vertical level

Bottom : last node in every vertical level

*better approach
in doubts session*

Diagonal Traversal { R-L }



output :

0 :	8	9
1 :	6	2 4
2 :	10	9
3 :	4	

obs:

if we go to left: level same

if we go to right: level by +1

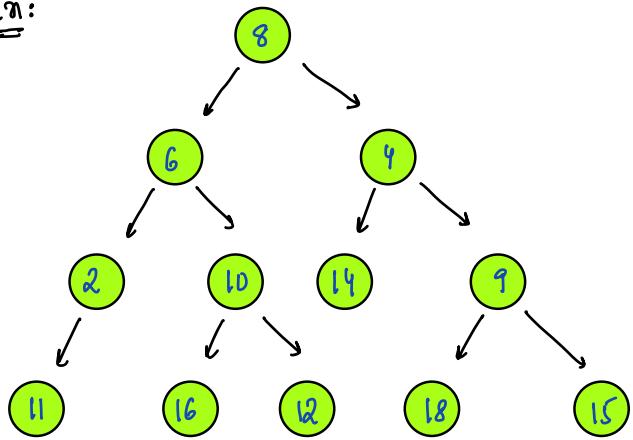
1^n ele on every level = diag of R-L

Diagonal Traversal { L-R } TODO

8:25 → 10 mins → 8:35

Q) Given preorder & inorder of BT of distinct values
print postorder traversal

Eg:



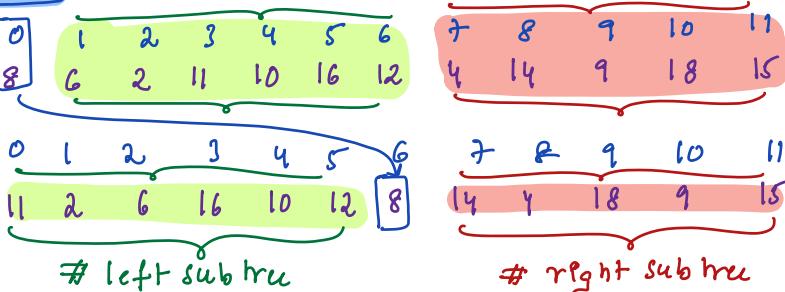
Idea: Construct Tree + Traversal

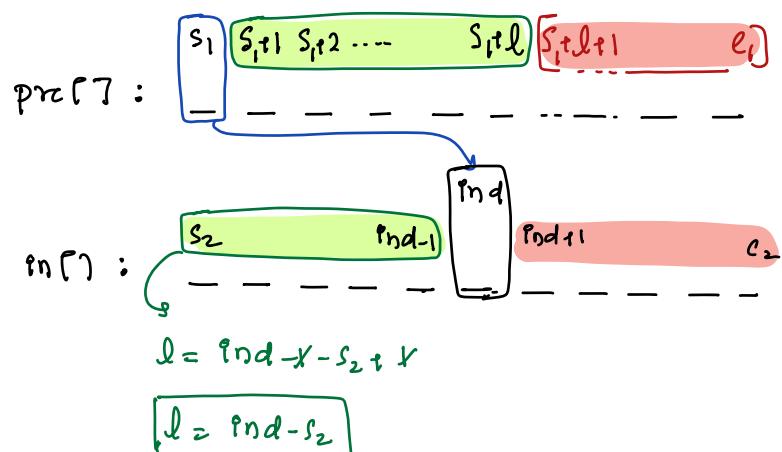
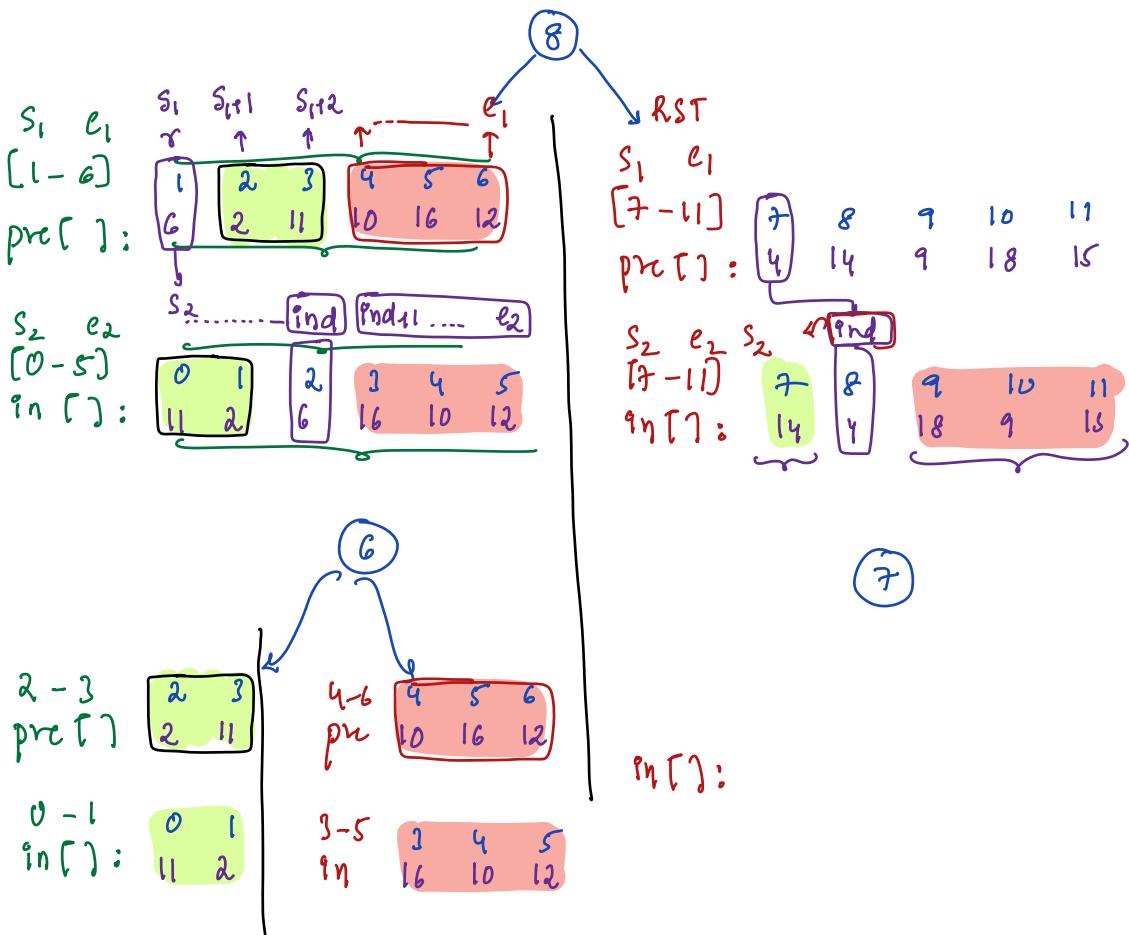
preorder [] :

(DLR)

inorder [] :

(LDR)





Ass: Given pre order & in order construct tree & return root node

Node TreeC int pre[], int s₁, int e₁, int in[], int s₂, int e₂)

// do data
if (s₁ > e₁) { return null }

Step:1 // root node is pre[s₁]

2^{num})

Node root = new Node (pre[s₁])

Step:2 // search for pre[s₁] in in[s₂..e₂]

int ind = —

i = s₂; i <= e₂; i++) {

} if (pre[s₁] == in[i]) { ind = i; break }

}

Step:3 // Splitting

l = ind - s₂ + 1 = ind - s₂ {# no. of in LST}

LST: → pre[s₁+1, s₁+l] → in [s₂, ind-1]

RST: → pre[s₁+l+1, e₁] → in [ind+1, e₂]

} For each search $\Rightarrow O(N)$
Reeach search for every
elem in pre[] $\Rightarrow TC: O(N^2)$

Optimize by storing in[]
elements & puts index
pos in hashmap:

TC: $O(N^2 + N)$ SC: $O(N)$

// Construct LST & return root of LST

root.left = TreeC pre, s₁+1, s₁+l, in, s₂, ind-1)

// Construct RST & return root of RST

root.right = TreeC pre, s₁+l+1, e₁, in, ind+1, e₂)

return root;

// After Tree Construction print postorder

// Add in[] in hashmap in main & pass it as parameter in Tree()

Ques: Given pre order & in order, print postorder

```
void Tree( int pre[], int s1, int e1, int in[], int s2, int e2)
```

```
    // do data  
    if( s1 > e1) { return; }
```

Step:2 // search for pre[s1] in in[s2..e2]

} // using hashmap
for optimization

```
int ind = _____
```

```
i = s2; i <= e2; i++) {
```

```
    if( pre[s1] == in[i]) { ind = i, break; }
```

// print postorder LST

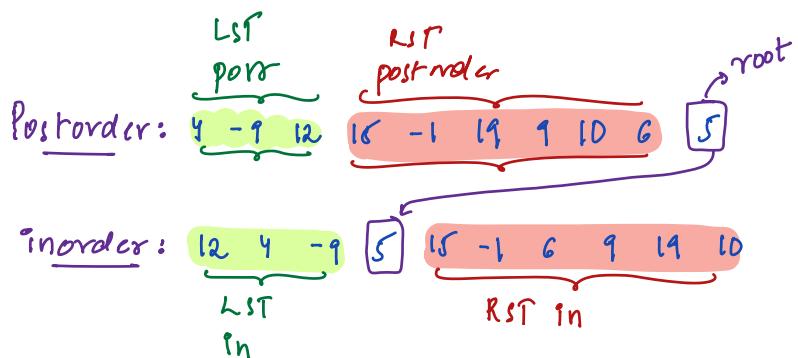
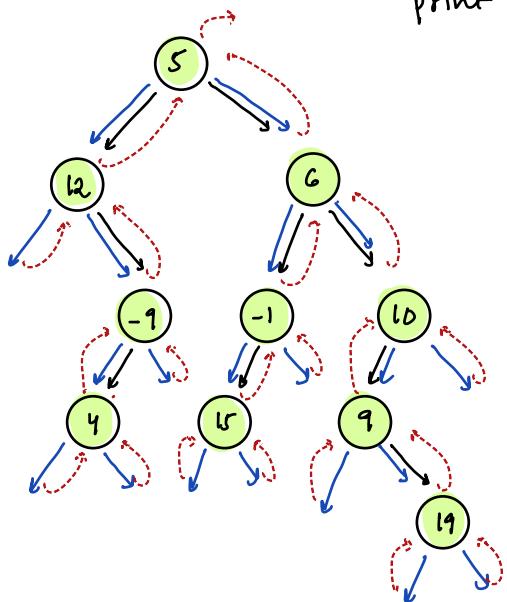
```
Tree( pre, s1+1, s1+l, in, s2, ind-1)
```

// print postorder RST

```
Tree( pre, s1+l+1, e1, in, ind+l, l2)
```

```
print( pre[s1]) // root node
```

Q8) Given `in[]` & `post[]` construct `tree()` print `preorder()` } TODD



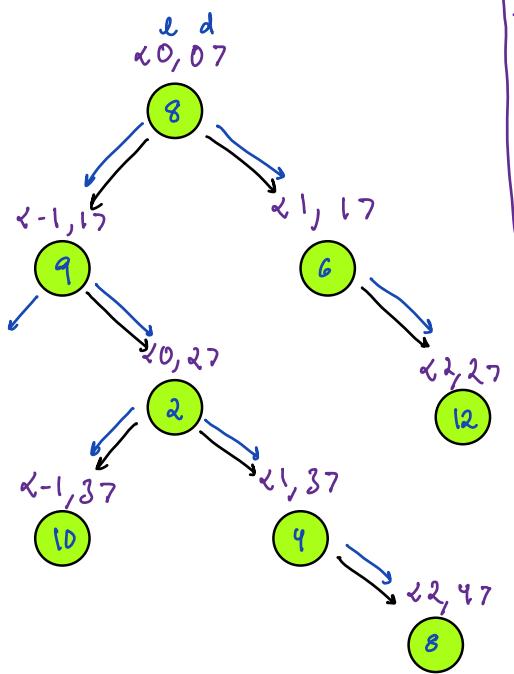
SQ8) Given `pre[]` & `post[]` construct `tree()` ?

Postorder: $4 \text{ LST } -1 \text{ RST } 12 \text{ LST } 16 \text{ LST } -1 \text{ RST } 19 \text{ LST } 9 \text{ RST } 10 \text{ RST } G \text{ post }$ $\text{root } 5$

Preorder: $\text{root } 5 \text{ LST } 12 \text{ LST } -1 \text{ RST } 4 \text{ LST } 6 \text{ LST } -1 \text{ RST } 15 \text{ LST } 10 \text{ LST } 9 \text{ RST } 19$

→ Note: We cannot separate LST & RST hence
Tree construction is not possible

Top view carry approach:



	hash & int, pair<int, int>> hm	root	depth
-1	(9, 1) (10, 3)		
0	(8, 0) (12, 2)		
1	(4, 2) (6, 1)	replace	
2	(8, 4) (12, 2)	replace	

vertical level

hashmap<int, pair<int, int>> hm
minL = 0, maxL = 0

any traversal is fine

```
void preorder(Node root, int vl, int dl) {
    if (root == null) return;
    minL = min(minL, vl); maxL = max(maxL, vl);
    if (hm.search(vl) == false) {
        hm[vl] = {root.data, dl}
    } else {
        pair<int, int> data = hm[vl];
        int cd = data.second;
        if (dl <= cd) { // update hm
            hm[vl] = {root.data, dl}
        } else { // bottom view
            hm[vl] = {root.data, cd}
        }
    }
    preorder(root.left, vl-1, dl+1) // vl is diagonal view L-L
    preorder(root.right, vl+1, dl+1)
}
```

Iterate from minL to maxL on hashmap & print it