

Trees 2

Diameter

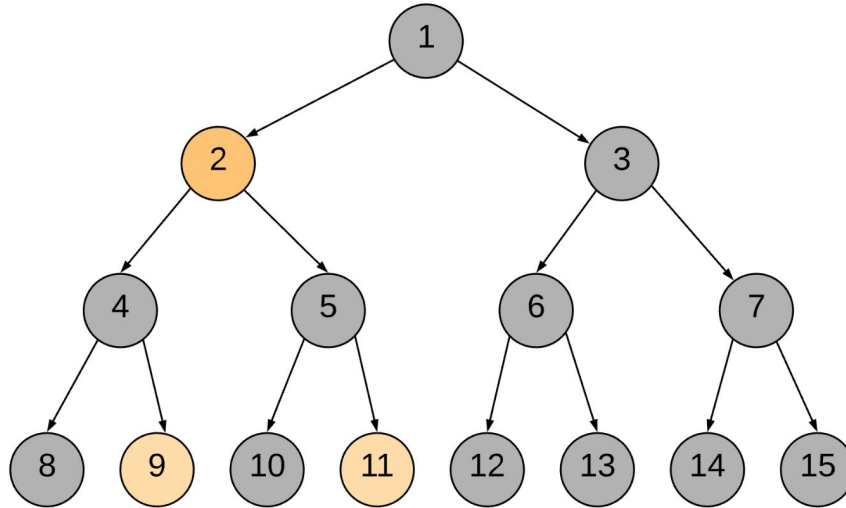
- Longest path present in the tree is the diameter
- Useful technique to think about problems involving trees

LCA

- Lowest Common Ancestor
- For a rooted tree LCA of 2 nodes is the farthest common ancestry those nodes
- Exists only for a certain root of the tree
- Can change for different roots

LCA results

- $\text{dist}(A, B) = \text{level}_A + \text{Level}_B - 2 * \text{Level}_{\text{LCA}}$



Lowest Common Ancestor for **Node 9** and **Node 11** is **Node 2**

How to find LCA?

- A powerful technique called binary lifting is used
- Crux of the idea is the base 2 representation of any number
- Move in incremental steps of powers of 2
- Time Complexity : $O(\log n)$
- Space Complexity : $O(n \log n)$

```

27 const int N = 4e4 + 2;
26 const int LN = 20;
25
24 vector< int > g[N];
23 int tin[N];
22 int tout[N];
21 int timer;
20 int up[N][LN];
19 int m;
18
17 void dfs (int u, int p) {
16     tin[u] = ++timer;
15     up[u][0] = p;
14     for (int i = 1; i < LN; ++i) {
13         up[u][i] = up[up[u][i - 1]][i - 1];
12     }
11     for (auto &to : g[u]) {
10         if (to != p) {
9             dfs(to, u);
8         }
7     }
6     tout[u] = ++timer;
5 }
4
3 bool is_ancestor (int u, int v) {
2     return tin[u] ≤ tin[v] && tout[v] ≤ tout[u];
1 }
28
1 int lca (int u, int v) {
2     if (is_ancestor(u, v)) return u;
3     if (is_ancestor(v, u)) return v;
4     for (int i = LN - 1; i ≥ 0; --i) {
5         if (!is_ancestor(up[u][i], v)) {
6             u = up[u][i];
7         }
8     }
9     return up[u][0];
10 }

```

LCA - I

- Find the maximum in path between 2 nodes for a rooted tree

K^{th} ancestor

- Main idea : Represent the number k in binary
- Move in powers of 2 using the table built during binary lifting


```

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20 void dfs (int u, int p) {
19     tin[u] = ++timer;
18     up[u][0] = p;
17     for (int i = 1; i < LN; ++i) {
16         up[u][i] = up[up[u][i - 1]][i - 1];
15     }
14     for (auto &to : g[u]) {
13         if (to != p) {
12             dfs(to, u);
11         }
10     }
9     tout[u] = ++timer;
8 }
7
6 int kth_ancestor (int u, int k) {
5     for (int i = 0; i < LN; ++i) {
4         if (k >> i & 1) {
3             u = up[u][i];
2         }
1     }
31 return u;
1 }

```

Euler Tour

- A way to make an array representation of the tree
- Useful for answering queries regarding subtrees
- Many ways to represent Euler Tour
- Will study the one derived from what we learnt with the in-time/out-time concept

Training Tasks

1. <https://codeforces.com/contest/208/problem/E>
2. <https://www.codechef.com/submit/LGSEG?tab=statement>