O Vieng emp name as a clustered anden is possible only when every employer will have a tinique name. It this is enued, the & anishes will be organized according emprome alphatitically Ouring empired as a churted index is definitely houdele consideing everyone valready has a anique id arigned to them the by implus arill le organize according to empire. Oung both emp name fempid en dufered einden may no of be houble but it is posible two name on clustered render and one nonchentined under. . Then & payer of to intolline region with a start of the state of the s And the property is to other the west of the second The state of the s

(2) * DDL is important in sepresenting derformation in DBMS wood to because is seved to describe external and logical * DMZ in ared to applate and a west date it in not important for representatives data

* DBMS*

L. Bharath 19BCSO62 CSE-Asection.

B) True, A DBMS as typically shared among many cuers. Transactions from there wers can be interleaved to compoure the execution time Juver's queries. By interheaving queries, wers do not have to wait for other over's transactions to the Complete fully before their own transaction is to complete begins without unterteaving if cuers. A begin a braneaction that will take 10 seconds ito complete; and were B wanter the begins a braneation overs & would have to exist an additional to records for arein A's transaction for complete fefore the database would beign processing over's B's request.

4. Bhoreath 19BCSO62

(4) a) A were must guarantee that his her transaction does not correct data or inset non une in the database, for example, in a banking datobare; a over mut guarantee that a cach withdraw transaction accurately models the amount a hereon removes from his a her account. A database application would be worthless if a hereon summed 20 dollar from an ATM but the transaction set their blalance to zero. b) & DBMS must guarantee that transactions are executed fully and condependently of other transactions. Is eventual property of a DEMS is that a transaction should execute atomically of as if it the only transaction running. Also, transaction will either complety fully or will be aborted and the database returned to its cintial that This ensures that the database seamains Conjustent

By ter, we can obtermine the key of suchation with the help of instance eg: Ing one to many relation we can consider the column lattribule with unique values as a frimwry key.

P(R,, Catalog)

P(R2, Cotalog)

TR, histor, pid = R2. hist 1R, - Sid6 = R2. sid(R, x R2)

* SQL

* SQL

SELECT C. ind

FROM Catalog C

WHERE EXISTS (SELECT G. ind

FROM Catalog C, WHERE C. pid = C. pid AND

C. sid 6 = C. sid)

(a) The some (This ((oca) or = Oredo (Parts)) x

(= Cost < 100 (Catalog)) x . Suppliers))

I malis query.

Their relational algebra ilst ment does not seturn anything because of the requesses of projection fresh in the sid in projected; it is the only aill not return anything.

A. Bhoroth 1) The following view on Emp can be updated automatically by updating Emp: CREATE VIEW Series Emp (end oname age, ralary) AS SELECT E. eid, E. ename, E. age, F. rolary. FROM EMPE. WHERE E.age \$50. Scanned with CamScanner