

18CSC302J-Computer Networks





Syllabus - Unit I

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- IP Fragmentation
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- > RARP
- > ICMP
 - > Introduction
 - Messages
 - Debugging Tools
 - ICMP Package

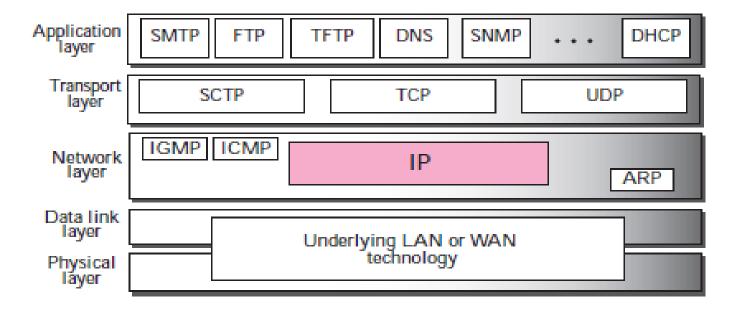
- UDP Datagram
 - Characteristics
- TCP Header
 - > TCP Connection Establishment Process
 - Error Control
 - Congestion Control
 - Flow Control
- Multicasting & Multicast Routing Protocols
- Stream Control Transmission Protocol





IP- An Introduction

➤ The **Internet Protocol (IP)** is the transmission mechanism used by the TCP/IP protocols at the network layer Operates at higher level



Position of IP in TCP/IP protocol suite



IP- An Introduction

- ➤ IP is an unreliable and connectionless datagram protocol—a **best-effort delivery** service.
- ➤ The term *best-effort* means that IP packets can be corrupted, lost, arrive out of order, or delayed and may create congestion for the network.
- > If reliability is important, IP must be paired with a reliable protocol such as TCP.



IP- An Introduction

> Example

- ✓ The post office does its best to deliver the mail but does not always succeed. If an unregistered letter is lost, it is up to the sender or would-be recipient to discover the loss and rectify the problem.
- ✓ The post office itself does not keep track of every letter and cannot notify a sender of loss or damage.
- ➤ IP is also a connectionless protocol for a packet switching network that uses the datagram approach
- This means that each datagram is handled independently, and each datagram can follow a different route to the destination.



IP Datagram

Bit 0

Bit 31

version (4)	Har Len (4)	TOS (8)	Total Length in bytes (16)		
Identification (16 bits)		Flags (3) Fragment Offset (13)			
Time to	Live (8)	Protocol (8)	Header Checksum (16)		
		Source IF	P Address		
	Destination IP Address				
	Options (if any)				

Header

Data

Data (variable length)



Version

- Version number of IP protocol
 IP Packet Header
- Current version is Version 4
- Version 6 has different header format

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)		
Identification (16 bits)		Flags (3) Fragment Offset (13			
Time to	Live (8)	Protocol (8)	Hea	der Checksum (16)	
	Source IP Address				
	Destination IP Address				
Options (if any)					



- Header Length (in 32 bit words)
 - Indicates end of header and beginning of payload
 - If no options, Header length = 5

• Bit 0 Bit 31

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)		
Identification (16 bits)		Flags (3)	Fragment Offset (13)		
Time to	Time to Live (8) Protocol (8)		Header Checksum (16)		
	Source IP Address				
	Destination IP Address				
Options (if any)					



- Type of Service (TOS)
 - Allows different types of service to be requested
 - Initially, meaning was not well defined
 - Currently being defined (diffserv)

Bit 0 Bit 31

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)	
Identification (16 bits)		Flags (3)	Fragment Offset (13)	
Time to	Live (8)	Protocol (8)	Hea	der Checksum (16)
		Source IF	P Address	
	Destination IP Address			
Options (if any)				



- Packet Length (in Bytes)
 - Unambiguously specify end of packet
 - Max packet size = 2^{16} = 65,535 Bytes

Bit 0

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)	
Identification (16 bits)		Flags (3)	Fragment Offset (13)	
Time to	Live (8)	Protocol (8)	(8) Header Checksum (16)	
		Source IF	P Address	
	Destination IP Address			
Options (if any)				



• These three fields for Fragmentation Control

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)		
Identification (16 bits)		Flags (3)	Fragment Offset (13)		
Time to	Live (8)	Protocol (8)	Header Checksum (16)		
	Source IP Address				
Destination IP Address					
Options (if any)					



Time to Live

- Initially set by sender (up to 255)
- Decremented by each router
- Discard when TTL = 0 to avoid infinite routing loops

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)	
Identification (16 bits)		Flags (3)	Fragment Offset (13)	
Time to	Live (8)	Protocol (8)	Hea	der Checksum (16)
		Source IF	P Address	
	Destination IP Address			
Options (if any)				



Protocol

Value indicates what is in the data field

Bit 0

Example: TCP or UDP

Version Hdr Len (4) (4)	TOS (8)	Total Length in bytes (16)			
Identification (16 bits)		Flags (3)	Fragment Offset (13)		
Time to Live (8)	Protocol (8)	Header Checksum (16)			
	Source IP Address				
Destination IP Address					
Options (if any)					



Header Checksum

- Checks for error in the header only
- Bad headers can harm the network
- If error found, packet is simply discarded

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)		
Identification (16 bits)		Flags (3) Fragment Offset (13)			
Time to	Time to Live (8) Protocol (8) Header Checksum (16)		nder Checksum (16)		
	Source IP Address				
Destination IP Address					
Options (if any)					



- Source and Destination IP Addresses
 - Strings of 32 ones and zeros

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)	
Identification (16 bits)		Flags (3) Fragment Offset (13)		
Time to	Live (8)	Protocol (8)	Hea	der Checksum (16)
		Source IF	P Address	
	Destination IP Address			
Options (if any)				



Options

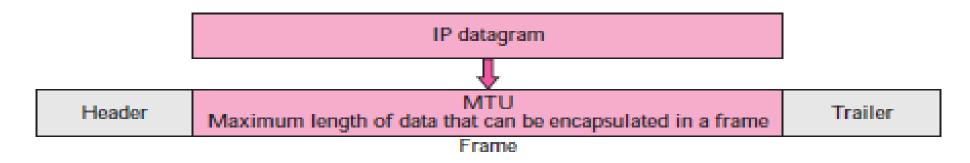
- Example: timestamp, record route, source route

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)		
Identification (16 bits)		Flags (3)	Fragment Offset (13)		
Time to	Live (8)	ive (8) Protocol (8) Header Checksum (16)		der Checksum (16)	
	Source IP Address				
	Destination IP Address				
Options (if any)					



IP Fragmentation & Reassembly

- Maximum Transmission Unit (MTU)
 - Largest IP packet a network will accept
 - Arriving IP packet may be larger (max IP packet size = 65,535 bytes)
- > Sender or router will split the packet into multiple fragments
- Destination will reassemble the packet
- > IP header fields used to identify and order related fragments





IP Fragmentation & Reassembly

- ➤ Divide the datagram to make it possible to pass through these networks called **fragmentation.**
- ➤ A fragmented datagram may itself be fragmented if it encounters a network with an even smaller MTU.
- ➤ A datagram can be fragmented by the source host or any router in the path
- the reassembly of the datagram, however, is done only by the destination host



- Identification
 - - All fragments of a single datagram have the same identification number

• Bit 0 Bit 31

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)	
Identification (16 bits)			Flags (3)	Fragment Offset (13)
Time to Live (8) Protocol (8)		Header Checksum (16)		
		Source IF	P Address	
Destination IP Address				
Options (if any)				



Flags:

- 1st bit: reserved, must be zero
- 2nd bit: DF -- Do NotFragment
- 3rd bit: MF -- More Fragments

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)				
Identification (16 bits)			Flags (3)	Fragment Offset (13)			
Time to	Live (8)	Protocol (8)	Header Checksum (16)				
Source IP Address							
Destination IP Address							
Options (if any)							



- Fragment Offset (in units of 8 bytes)
 - Used for reassembly of packet
 - 1st fragment has offset = 0

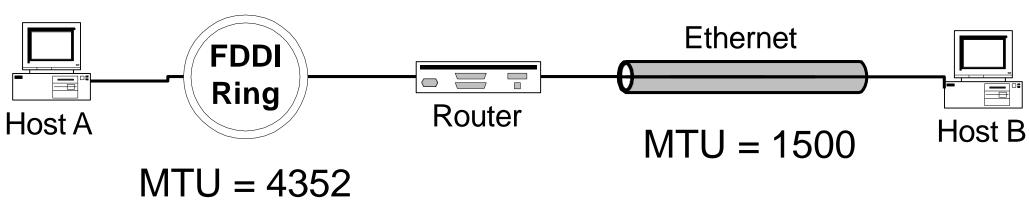
Bit 0 Bit 31

Version (4)	Hdr Len (4)	TOS (8)	Total Length in bytes (16)				
lde	Identification (16 bits)			Fragment Offset (13)			
Time to	Live (8)	Protocol (8)	Header Checksum (16)				
Source IP Address							
Destination IP Address							
Options (if any)							



IP Fragmentation Example

Host A wants to send to Host B an IP datagram of size = 4000
 Bytes

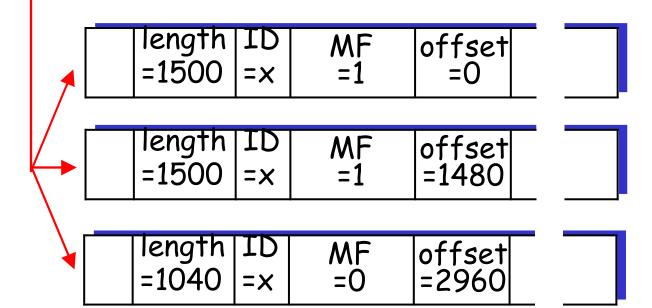




IP Fragmentation Example

length ID	MF =0	offset =0	
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One large datagram becomes several smaller datagrams





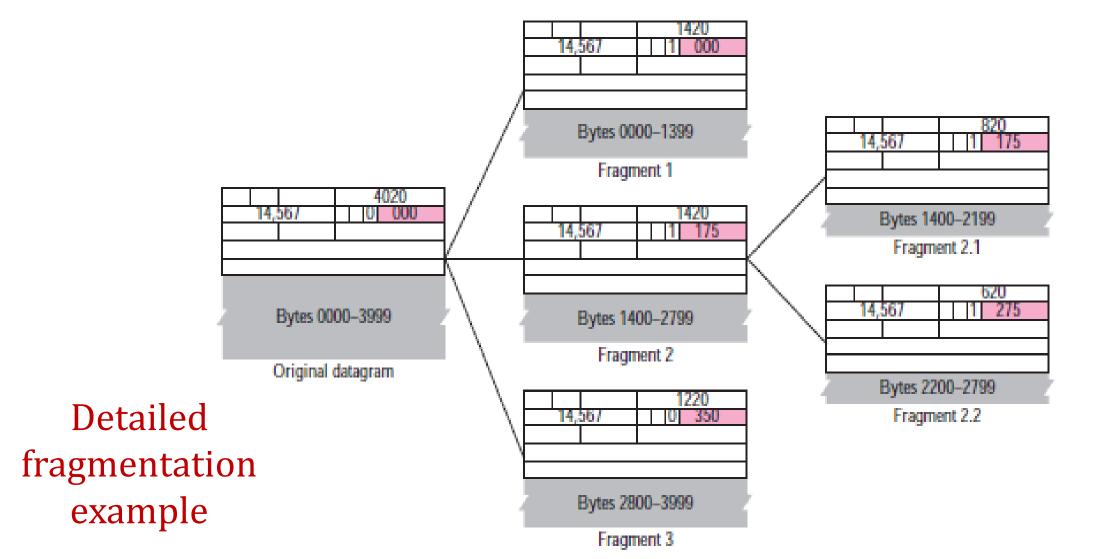
Multiple Fragmenting Points

Let MTUs along internet path be

- 1500
- 1500
- 1000
- 1500
- 576
- 1500

Result: fragmentation can occur twice

Multiple Fragmenting Points





- The figure shows what happens if a fragment itself is fragmented.
 - In this case the value of the offset field is always relative to the original datagram.
 - a. The first fragment has an offset field value of zero.
 - **b.** Divide the length of the first fragment by 8. The second fragment has an offset value equal to that result.
 - **c.** Divide the total length of the first and second fragment by 8. The third fragment has an offset value equal to that result.
 - **d.** Continue the process. The last fragment has a *more* bit value of 0.