

PART-B

1. Back Patching:

- info is of ~~the~~ labels
- it indicates the address of the label in goto state. while producing TACs for an exp.
- 2 passes are used becoz assigning the position of these labels are challenging
- It can leave these addresses unidentified in 1st pass then populate in 2nd round.
- the process of filling up gaps in incomplete transformation 2 info.

need for BP:

- Boolean exp
- flow of Control Statement
- labels 2 goto.
 - flow of Prog in L 2 A.
 - Atleast one statement with L in scope of this g statement.

2. triples of exp $a * (b + c)$

1: $t1 := b + c$ → 1st triple, 2nd, 3rd.

2: $t2 := -t1$

3: $t3 := a * t2$

↓
temp variables to
store intermediate results
used to build up
the final result of exp

3. $x = *y; a = 2x;$

1: $t1 := *y$

2: $x := t1$

3: $a := 2x$

4. idea behind generating 3-A-C during compilation.

→ to create I-C representation of S-C.
easily translated into Machine code

→ low-level rep. uses small no. of instructions

Benefits

Simplification: → break into small pieces.

→ simple to perform

Portability → platform independent representation
→ same code in diff platform.

Debugging → structured rep.

5. Quadruples preferred over triples in an optimizing compiler?

→ When the statements are often moved around.

→ introduces an extra degree of indirection bet the computation of a value & its use

In triples notation,

Moving statement that defines a temporary value requires to change all pointers to that statement in Arg1 & 2.

→ this makes difficult use in optimizing compiler.

6. Notations to represent an intermediate lang.

TAC: → low-level representation of code.

Quadruples: → 4 field than TAC

→ Source 1, 2, Destination.

Syntax tree → hierarchical representation

→ relation bet exp & statements

Abstract Syntax tree:

→ Similar to Syntax tree

→ more abstract, omitting nodes that not affect the meaning of code

Byte code:

→ virtual mach. (low-level rep of code)
Directly by processor.

7. how to solve ~~the~~ the issues in the Design of Code generator.

* I/p to code generator.

→ intermediate code generated from front end.

Linear rep. → postfix

TAC → R, T, I

graphical rep. → syntax tree
→ DAG.

* target program

→ o/p of code generator

Absolute ML, Relocatable ML, Assembly lang.

* Memory management.

Mapping name to address of Data obj.

* instruction selection.

→ Best instrumt. improves efficiency.

* Register allocation issues

→ computation faster

→ Register Allocation & assignment.

* choice of evaluation order.

→ order in which the instructions are executed.

8. four principle use of registers in C. Generation

Registers → small, high-speed memory locations

temporay store data & perform logic &

Arith. exp/op.

Variable storage:

→ store values.

Parameter passing.

→ pass function arg bet func.

Return values.

Retur. Val. from func.

temporay storage: