

Ex:

Convert the PDA $P = (\{p, q\}, \{0, 1\}, \{x, z\}, \delta, q, z_0)$ to EFA.

It δ is given by

$$\begin{aligned} \delta(q, 1, z_0) &= \{(q, x, z_0)\} & \delta(p, \epsilon, z_0) &= \{(q, \epsilon)\} \\ \delta(q, 1, x) &= \{(q, xx)\} & \delta(p, 1, z_0) &= \{(p, \epsilon)\} \\ \delta(q, 0, x) &= \{(p, x)\} & \delta(p, 0, z_0) &= \{(q, z_0)\} \end{aligned}$$

Sol:

$$V = \{s, [q, x, q], [q, x, p], [p, x, q], [p, x, p], [q, z_0, q], [q, z_0, p], [p, z_0, q], [p, z_0, p]\}$$

$$\begin{aligned} \delta &\rightarrow [q, z_0, q] \\ &\quad [q, z_0, p] \end{aligned}$$

$$(i) \delta(q, 1, z_0) = \{(q, xx)\}$$

$$[p, z_0, q] \rightarrow 1 [q, x, q] [q, z_0, q]$$

$$1 [q, x, p] [q, z_0, q]$$

$$[p, z_0, p] \rightarrow 1 [q, x, q] [q, z_0, p]$$

$$1 [q, x, p] [q, z_0, p]$$

$$(ii) \delta(q, 1, x) = \{(q, xxx)\}$$

$$[q, x, q] \rightarrow 1 [q, x, q] [q, x, q]$$

$$1 [q, x, p] [q, x, q]$$

$$[q, x, p] \rightarrow 1 [q, x, q] [q, x, p]$$

$$1 [q, x, p] [q, x, p]$$

$$(iii) \delta(q, 0, x) = \{(p, x)\}$$

$$[q, x, q] \rightarrow 0 [p, x, q]$$

$$[q, x, p] \rightarrow 0 [p, x, p]$$

$$(iv) \delta(q, \epsilon, 20) = \{(q, \epsilon)\}$$

$$(q, 20, q) \rightarrow \epsilon$$

$$(v) \delta(p, 1, x) = \{(p, \epsilon)\}$$

$$(p, x, p) \rightarrow 1$$

$$(vi) \delta(p, 0, 20) = \{(q, 20)\}$$

$$[p, 20, q] \rightarrow 0 [q, 20, q]$$

$$[p, 20, p] \rightarrow 0 [q, 20, p]$$

Final CFA is

$$\delta \rightarrow [q, 20, q]$$

$$[q, 20, p] \rightarrow \{1 [q, x, p] [p, 20, q], \epsilon\}$$

$$[q, x, p] \rightarrow \{1 [q, x, p], [p, x, p], 0 [p, x, p]\}$$

$$[p, x, p] \rightarrow 1$$

$$[p, 20, q] \rightarrow 0 [q, 20, q]$$

Pumping lemma (For CFL) \hookrightarrow CFG $\rightarrow [V, \Sigma, S, P]$ ②

It is used to Prove that a language is not a Context free. d, 2

Assume L is a CFL, there is a Pumping length n such that any string $w \in L$ of length $\geq n$ can be written as

$$|w| \geq n \text{ --- ①}$$

we can break w into 5 strings.

$w = uvxyz$, such that

$$\rightarrow |vxy| \leq n$$

$$\rightarrow |vy| \neq \epsilon$$

\rightarrow For all $k \geq 0$ the string $uv^kxy^kz \in L$

Steps

- ① Assume that L is context free
- ② Pumping length say n
- ③ All strings longer than n can be Pumped
 $|w| \geq n$
- ④ now find a string w in L such that
 $|w| \geq n$
- ⑤ Divide w into $uvxyz$ ⑥ show that $uv^kxy^kz \notin L$
- ⑦ Then consider the ways that w can be divided into $uvxyz$.
- ⑧ show that none of these can satisfy all the 3 Pumping conditions at same time
- ⑨ we cannot be pumped (Contradiction)

Ex:

Find whether $L = \{x^n y^n z^n \mid n \geq 1\}$ is context free or not

(1)

Sol:

Let L is context free

& now we can take a string such that

$$s = x^n y^n z^n$$

& we divide s into 5 parts $uvxyz$

$$\text{Case (1)} \quad n=4 \quad x^4 y^4 z^5$$

v & y each contains only one type of symbols

$$\begin{array}{cccccccccccc} x & x & x & x & y & y & y & y & z & z & z & z & z \\ | & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & | & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} \\ | & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & | & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} \\ u & v & & x & & & & y & & & & & z \end{array} \rightarrow \text{divide into 5 parts } [u, v, x, y, z]$$

$$\Rightarrow uv^k x y^k z \quad [\text{let } k=2]$$

$$\Rightarrow uv^2 x y^2 z$$

$$\Rightarrow \cancel{uv^2} x x x x x y y y y z z z z z$$

$$\Rightarrow x^6 y^4 z^5$$

$$\Rightarrow x^6 y^4 z^5 \notin L$$

$$u = x$$

$$v = xx$$

$$x = xyxyy$$

$$y = z$$

$$z = zz$$

the given language is not context free.

Case (ii)

(5)

Either v or y has more than one kind of symbol

$$\Rightarrow x^n y^n z^n \Rightarrow 5 \text{ Parts}$$

$$uvxy^2$$

$$n=4, \text{ so } x^4 y^4 z^4 \left[uv^k x y^k z^k \notin L \right]$$

$$\Rightarrow xxxxyyyyzzzz \Rightarrow 5 \text{ Parts}$$



$$\begin{aligned} u &= xxx \\ v &= xxyy \\ x &= x \\ y &= y \\ z &= zzzz \end{aligned}$$

$$\Rightarrow \text{let assume } k=2$$

$$\Rightarrow uv^2xy^2z$$

$$\Rightarrow xxxxxxyyyzzzz$$

$$\Rightarrow x^4 y^2 x^2 y^5 z^4$$

$$\Rightarrow \text{It is not belong to } L$$

This is not context free.

Ex: $L = a^i b^j c^i d^j \mid i, j \geq 1$ is not CFL

① Assume L is context free

② now we can find a string such that

$$s = a^i b^j c^i d^j$$

we divide into 5 parts $uvxy^2z$

Case (i) $i = 2, j = 2$

$$= a^2 b^2 c^2 d^2$$

$$\begin{array}{ccccccc} a & a & b & b & c & c & d & d \\ \hline u & & & & & & & \\ & & & & x & y & & z \end{array}$$

$$\begin{array}{l} u = aa \\ v = bb \\ x = c \\ y = c \\ z = dd \end{array}$$

③ $k = 2$

$$\Rightarrow uv^k xy^k z$$

$$\Rightarrow uv^2 xy^2 z$$

$$\Rightarrow aabbbbbbccdd$$

$$\Rightarrow a^2 b^4 c^3 d^2$$

$$\Rightarrow \notin L$$