Therating System intermediacy b/w the user and 1 what is an operating system? hogram which acts as an the competer hardware. 2) Different operating systems. La Multi-laskung
La Multi-Reogramming Is Real time-operating System Is Beatch operating Lime-Sharingos Listibuted Operating System.

3 basic functions of 08 ? 4 controls and co-ordinates the user of Lardware for various uses. La acts as resource allocators & manager. by controls the use programs to prevent errors and Engropes use of computer. 1) what is kerend? Ly core à essential part of 05 that procedes basec services. Ly is responsible for disk management, memory management, task management etc. Ly just program to load after the bootloads. Ly it is a bridge bloother user and the hardware, it translates user instructions for the hardware and connects them. (3) Micro Kernel: Ly surs services that are minimal foros performance. I all other operations are performed by processor. Macro Kernel: L' combination of minot monolethic kernel. La all 05 code is in single executible file (6) <u>beadlock</u> two processes are waiting for each other 6 condition where ito completel so that they can start. is this results both the processes to hang,

G Process

G program in execution.

Let Two types - Os processes

- were processes

8 States of a process
4 New
4 Rurring
4 wait
4 Ready
4 Terminate

O New: Program going to be priked up by the Os ainto the main memory

(3) Ready The processes which are ready for the execution and reside in main memory.

1 3 week! Running: One of the processes from the Ready state will be chosen by the Os depending on the scheduling algorithm.

The no of nunning process will always be one at a time 'cause only one process will be priched up by the CPU foreneution.

part tour sale thanks

(4) Wait: When a precess waits for a certain sesource to be assigned or input from the uses, then the OS mours this process to the block poait state and assigned the CPV to other processes.

3 Terminale I when a process finishes its execution, it comes to the terminator state.

(2) Staruation

Is it a resource management problem where a process does not get the resources it needs for a long time because the resources are being allocated to other processes.

La technique to avoid staveation un a schederling algorithm.

(10) Semaphore Li variable whose states reports common sesource. GTwo types - benasy and counting semaphora.

Birary semaphores

Lyvalues restricted to 91.

4 wait when O

4 signal when 01

Aduantages

Is allow only one process to be un the critical section.

la follow mutual exclusion.

b nocesource wastage

Disaduantage:

4 complicated.

Is impractical for large scale use.

Counting Semaphoses

- non negative enteger value.

1) corordenate the resource

L' simaphore count = no. of quaitable resources.

Ly incremented when resource added else subtracted.

- (3 Data Structures :
  - @ Search for attarget key in linked list? I apply sequential search.
- I beach mode is transsed and compared with a larger lay.
  Is transsed is continued sentil the key is found/last node is reached.
  - Because algorithm

    Ly a problem is decided wito small subproblem.

    Ly the opp of one recursion after processing one sub-problem becomes the Me to the next recursive process.
    - 3 Ribonacci search

Ly algo applied to a sorted away.
Ly uses a divide- and-conquer approach to reducetime.

- (4) Juffman's algorithm
  Granting extended binary trees that have minimum neighted fath lengths.

  Ly makes we of table that contains the frequency of occurrence for each data element.
- Doubly lined lists:

  Ly transpal can be done in both directions

  Ly each nod has trooling.
- 1) AVITERE
  4 being search tree in a state of partially balanced.
  4 being search tree in a state of partially balanced.
  4 balance is diff blue other houghts of other subtrees from the root.

39/04/2021 . Tower of Hanoi minimum no. of steps required with 'n' disks = [2"-1] Lithrel towers - source, dest and aux. Ly Rules: O only one disk can be moved among the towers cat any quantene. @ Only the "top disk" can be removed. 3 No large disk can sit ouer a small disk. Destination Aux Ly Steps: 1 Nove n-1 disks from source to cur. 2) reque non disks from source to dist. 3) Move n-1 disks from aun to dist. when n = 3, 1 Hove & disks from source to aux 123 1 Marie 3rd desh from source to dest 3 Move 2 distres from our to dest. - Algorithm: S D A Procedure Hanci (disk, source, dest, aux) If disk = = 1, then nove disk from source to dest. Hanvildisk-1, source, aux, dest) 1/ Step1 End frotedure

STOP

Land Protedure

Land Protedure

Land Protedure

End Protedure

Land Protedure

End Protedure

Land Proted noue desk from source to dest 1/ Step 2

```
· Code Walktheorgh
 thindudex wits/std c++h>
 using namespace std;
  class Solution &.
    perblic:
            long long toh ( int N, interom, intero, untaux) ?
            ent count = 0;
  Condition: cif (N==D)
              cout << " move disk! from rod" << from «" to rod" << to
                   endl;
              count tt;
            return count;
                                    destination
   count + = toh(N-1, from, aux, to);
   coute "move disk Nilfrom red" & from & "to red" << to cerd!
    count + = toh( N-1, aux, to, from);
    seturn count +1;
7 6
     main() {
      uit T;
      cins) Ti
      while (T--) ?
       int N;
         cin >> N;
         Solution ob;
        coextex ob toh (N, 1, 3, 2) exends
```

1) Sort an array of Os, 1's and 2s. Given an away AC) consisting of os, Is and 2s. The task is to write a function that sorts the given away. The function should put all 0s first, then all Is and all 2s in last. Injut= 20,1,2,0,1,23. Enangle: Outpet = 20,0,1,1,2,23

> Input = 20,1,1,0,1,2,1,2,0,0,0,13. Output = {0,0,0,0,0,1,1,1,1,1,2,2}

## Approach 1:

Count the number of 0s, I sand 2s in the given array. Store all os in the beginning followed by Is and then 2s.

Algorethm

@ keep 3 counters, c.o for 0, cs for 1 and c2 for 25.

- (2) Keep increasing the count of these variables cowhen o is found, cowhen I is found and cowhen I is found.
- (3) Replace forst co elements with 0, crelements with land or elements with 2's. Print the array.

Time Complemity: O(n) we are only traversing the array once. 3 pace Complexity: 0(1) Li no entraspace is required. Approach 2: The algorithm is also known as Dutch National Plag -of theseway fastitioning. Steps: 1) Dedare these variables downo, high = n-1, mid=0, where n= length of the array. @ Run aloop until mides high, three cases on the basis of value of asscmid), (3) 24 ase[mid] ==2 1 og conid ==0 ewap (ass [mid], sees [high]) swap (ass low], cust mid)) high -- ; mid++5 breaks low++; break; @ 21 aucmid) == 1 11 keep at it is mid++; break; Time Complexity: O(n) Space Complexity: 0(1) no entraspace is required.