Detecting Selfish Node in MANET Using Collaborative Contact Based Watchdog

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Abstract:- Wireless mobile ad hoc networks are dynamic networks, self-configuring in that nodes are free to move .Mobile ad-hoc networks (MANETs) assume that mobile nodes voluntary cooperate in order to works properly. So, this cooperation is a cost-intensive activity and that nodes can refuse to cooperate, leading to selfish node behaviour. So in this way, the overall network performance could be affected. The watchdog is used to well-known mechanism to detect a selfish node. The detection process is performed by watchdog can fail, generating false positives and false negatives that can induce to wrong operations. Moreover, the relying on local watchdogs alone can lead to poor performance when detecting selfish nodes, in term of precision and speed. The collaborative contactbased watchdog (CoCoWa) is a collaborative approach based on the diffusion of local selfish nodes awareness when a contact occurs, that information about selfish nodes is quickly propagated. The collaborative approach reduces the time and increases the precision when detecting selfish nodes.

Keywords—Wireless networks, opportunistic and delay tolerant networks, selfish nodes, MANET, DTN.

I. INTRODUCTION

Cooperative networking is a currently receiving significant attention as an emerging network design strategy for future mobile wireless networks. The successful cooperative networking can prompt the development of advanced wireless networks to cost-effectively provides services and applications in a contexts such as vehicular ad hoc networks (VANETs) or mobile social networks. Two basic technologies that are considered as the core for these types of networks are mobile ad-hoc networks and opportunistic and delay tolerant networks. The literature review provides two main strategies to deal with selfish behavior: a) motivation or incentive based approaches, b) detection and exclusion. First approach, tries to motivate nodes to actively participate in the forwarding activities.

The impact of node selfishness in MANETs has been studied in [7], [8], [9]. A more detailed study [7] shows that a moderate concentration of node selfishness has a huge impact on the overall performance of Mobile Adhoc Networks, such as the average hop count, the number of packets dropped, the throughput, and the probability of reachability.

This paper introduces a Collaborative Contact-based Watchdog (CoCoWa) as a new scheme for detecting selfish nodes that combines the local watchdog detections and the dissemination of this information on the network. If one node has previously detected a selfish node and it can transmit this information to other nodes when a contact occurs. This way, nodes have second hand information

about selfish nodes in a network. The goal is to reduce the detection time and to improve the precision by reducing effect of both false negatives and false positives.

II. LITERATURE SURVEY

Sybil attacker can create the more than one identity on a single physical device in order to launch a coordinated attack on the network or can switch the identities in order to weaken the detection process, thereby promoting lack of accountability in the network. In this paper, we propose a light weight scheme to detect the new identities of Sybil attackers without using a centralized trusted third party or any extra hardware, such as directional antennae or geographical positioning system [1].

The Ad hoc networks rely on the cooperation of the nodes participating in the network to forward the packets for each other. A node may decide not to be cooperate to save its resources while still using the network to relay its traffic. [2].

In this paper, describe the use of a self-policing mechanism based on reputation to enable mobile ad hoc networks to keep functioning despite the presence of misbehaving nodes. [3].

We can see, the problem of service availability in mobile ad-hoc WANs. We present a secure mechanism to stimulate end users to keep their devices turned on, to refrain from overloading the network, and to oppose the tampering aimed at converting the device into the "selfish"[4].

In this paper, we can see that each node have its own authority and tries to maximize the benefits it gets from the network. So, we assume that the nodes are not willing to forward packets for the benefit of the other nodes. So, in order to stimulate the nodes for the packet forwarding, we propose a simple mechanism based on a counter in each node. [5].

III. ARCHITECTURE OVERVIEW

The selfish node usually denies packet forwarding in order to save its own resources. Such type of behaviour implies that the selfish node is neither participates in routing nor relays data packets. A common technique to detect this selfish behaviour is network observe using local watchdogs. A node's watchdog consists on overhearing the packets transmitted and received by its neighbors in order to detect the anomalies, such as ratio between packets received to packets being re-transmitted. So, Byusing this technique, the local watchdog can generate a positive (or negative) detection in case the node is acting selfish (or not).

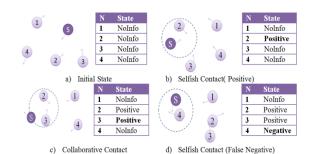


Fig. 1:An example of how CoCoWa works. a) Initially all nodes have no information about the selfish node. b) Node 2 detects the selfish node using its own watchdog. c) Node 2 contacts with node 3 and it transmit the positive about the selfish node. d) The local watchdog of Node 4 fails to detect the selfish node and it generates a negative detection (a false negative)

The example of how CoCoWa works is outlined in Fig. 1. It is based on the combination of a local watchdog and the diffusion of information when a contact between pairs of nodes occurs. A contact is defined as an opportunity of transmission between a pair of nodes. Assuming that there is only one selfish node, above figure shows how initially no node has information about the selfish node. When node detects a selfish node by using its watchdog, it is marked as a positive, and if it detected as a non-selfish node, it marked as a negative.

This scheme is the uncontrolled diffusion of positive and negative detections can produce the fast diffusion of wrong information, therefore, a poor network performance. For example, in figure 1, on the last state d), node two and three have positive detection and node four has negative detection (a false negative). Now, node one, which has the no information about selfish node, has several possibilities: if it contacts the selfish node it may be able to detect it; if it contacts a node two or three it can get positive detection; but if it contacts node four, it can get a false negative.

Figure 2 shows the functional structure of CoCoWa



The Local Watchdog having two functions: the detection of selfish nodes and the detection of new contacts. Local watchdog can generate following events about the neighbor nodes: PosEvt (positive event) when the watchdog detects selfish node, NegEvt (negative event) when watchdog detects that node is not selfish, and NoDetEvt (no detection event) when watchdog does not have the enough information about a node.

A Diffusion module has two functions: the transmission as well as the reception of positive and negative

detections. The key issue of our approach is the diffusion of information. The number of selfish nodes is low compared to the total number of nodes; the positive detections can always be transmitted with a low overhead. So, transmitting only positive detections is a drawback: false positives can be a spread over the network very fast.

The updating or consolidating the information is the another key issue. This is a function of the Information Update module. Anodes have the following internal information about other nodes: No Info state, Positive state and Negative state. The node have direct information (from local watchdog) and indirect information (from the neighbor nodes). CoCoWa is event driven, so the state of a node is updated when the PosEvt or NegEvt events are received from local watchdog and diffusion modules. In particular, these events update a reputation value ρ using following expression:

$$P = p + \Delta \Delta = \begin{cases} +\delta & (\text{PosEvt, Local}) \\ +1 & (\text{Posevt, Indirect}) \\ -\delta & (\text{NegEvt, Local}) \end{cases} \delta \ge 1$$
 (1)

So, a PosEvt event increments reputation value while NegEvt event decrements it. Defining θ as a threshold and using reputation value ρ , and the state of the node changes to Positive if $\rho \geq \theta$, and to Negative if $\rho \leq -\theta$. Otherwise, the state is a NoInfo. The combination of δ and θ parameters allows is a very flexible and dynamic behaviour. First of all, if $\theta > 1$ and $\delta < \theta$ we need a several events in order to change the state. For example, starting from the NoInfo state, and if $\theta = 2$ and $\delta = 1$, at least a local and an indirect event is needed to change the state, but if $\theta = 1$, only one event is a needed. Second, we can give a more trust to the local watchdog or to a indirect information. For example, a value of $\delta = 2$ and $\theta = 3$, means that we needed one local event and one indirect event, or three indirect events, to change state.

IV. PROPOSED SYSTEM

The network is modeled as set of N wireless mobile nodes, with C collaborative nodes, M is malicious nodes and S is selfish nodes (N = C + M + S). The goal is to obtain time and overhead that a set of D \leq C nodes need to detect the selfish nodes in the network. In case of several selfish nodes (S > 1) on a network with N nodes, we can assume that C = N - S are cooperative nodes.

4.1 The Model for the CoCoWa Architecture

The goal of this model is the behaviour of the different modules of our architecture (see figure 2). The local watchdog is modeled using three parameters: the probability of detection p_d , the ratio of false positives p_{fp} , and the ratio of false negatives p_{fn} . The first parameter, the probability of detection (p_d) , reflects the probability that, when a node contacts to the another node, the watchdog has enough information to generate a PosEvt or NegEvt event. So, this value depends on effectiveness of the watchdog, the traffic load, and the mobility pattern of nodes. Moreover,

the watchdog can a generate false positives and false negatives. A false positive is when watchdog generates positive detection for a node that is not selfish node. A false negative is generated when selfish node is marked as negative detection. In order to measure the performance of a watchdog, these values can be expressed as a ratio or probability: p_{fp} is the ratio (or probability) of false positives generated when a node contacts a non-selfish node, and pfn is the ratio (or probability) of false negatives generated when a node contacts a selfish node. By using the previous parameters we have model the probability of generating local PosEvt and NegEvt events when a contact occurs:

- PosEvt event: the node contacts with a selfish node and the watchdog detects it, with probability $p_{d}(1-p_{fn}).$ Note down the false positive can also be generated with probability $p_d \cdot p_{fp}$.
- NegEvt event: the node contacts with a non-selfish node and detect it with probability $p_d(1-p_{fp})$. A false negative can also be generated when it contacts with the selfish node with probability $p_d \cdot p_{fn}$.

So, the diffusion module can be generate indirect events when a contact with neighbor nodes occurs. Nevertheless, a contact does not always imply collaboration, so we have model this probability of collaboration as pc. In a real networks, full collaboration (p_c = 1) is almost impossible. Finally, the probabilities of generating the indirect events are the following:

- PosEvt event: When a contact with another node that has a Positive state of the selfish node with probability pc.
- NegEvt event: When a contact with another node that has a Negative state, being the probability γ · p_c . Note down a not all Negative states are to be transmitted, it depends on a diffusion factor γ.

The information update module is driven byprevious local and indirect events. These event are update the reputation ρ about a node, and it is used to finally decide if node is a selfish or not by using the threshold θ .

4.2 Malicious Nodes and Attacker Model

The malicious nodes attempt to attack the CoCoWa system by generating a wrong information about the nodes. Behaviour of malicious nodes is modeled from the receiver perspective, which is a based on probability of receiving wrong information about given node when a contact with a malicious node occurs (that is, it receives a Negative about selfish node, and a Positive about other nodes). We denote this behaviour as the maliciousness probability p_m . Following give the details of several aspects that can affect this probability:

- 1) The reception of the information, considering that not all contacts produce this reception. This aspect is similar to collaboration degree, but an increase communication range of the malicious nodes will increase information reception.
- 2) The malicious node does not have information about all nodes; so, in order to send a positive/negative about a node, they must have contacted to this node previously or have received a message from other nodes.

3) Another issue is to consider the proper generation of wrong information, for example when the receiving a positive of a node that is not the selfish node.

Summing up, above parameter reflects the average intensity or the effectiveness of the attack of the malicious nodes.

4.3 The Model for the Detection of Selfish Node

In this section we introduce an analytical model for evaluating the performance of the CoCoWa. The goal is to obtain the detection time of selfish node in a network. This model takes into account the effect of false negatives. The false positives does not affect on the detection time of the selfish node, so p_{fp} is not introduced in this model.

Using λ as the contact rate between the nodes, we can model the network by using a 4D Continuous Time Markov chain (4DCTMC). For modeling purposes, the collaborative nodes are divided into two sets: a set with D destination nodes and a set of E = C - D as intermediate nodes. Thus, the 4D-CTMC states are: $(d_p(t), d_n(t), e_p(t), e_n(t))$, where e_n(t) represents number of intermediate nodes that have a Positive state, $e_n(t)$ the intermediate nodes with a Negative state, $d_p(t)$ the destination nodes with the Positive state and d_n(t) the destination nodes with the Negative state. Note down, in this model, a Negative is a false negative. The states must beverify the following conditions are:

 $d_p(t) + d_n(t) \le D$ and $e_p(t) + e_n(t) \le E$. Our 4D-CTMC model has an initial state (0, 0, 0, 0) (that is, all the nodes have no information). The final states are when $d_p(t) = D$. We define υ as the number absorbing states, that are all the possible permutations of states ($\{(D, 0, *, *)\}$) that sum E. It is easy to derive that $v = p^s(E) = 0.5(E + 1)(E + 2)$. The number of transient states are tis obtained in similar way: $\tau = (p^s(D) - 1)p^s(E)$. So this model can be expressed by using the following generator matrix is Q:

$$Q = \begin{pmatrix} T & R \\ 0 & 0 \end{pmatrix}, \tag{2}$$

 $Q=\begin{pmatrix} T & R \\ 0 & 0 \end{pmatrix}, \qquad (2)$ Where T is a $\tau \times \tau$ matrix with elements q_{ij} denoting the transition rate from transient state s_i to transient state s_j , R is a $\tau \times v$ matrix with elements q_{ij} denoting the transition rate from the transient state is s_i to the absorbing state s_i , the left 0 is a $\upsilon \times \tau$ zero matrix, and the right 0 is a $\upsilon \times \upsilon$ zero matrix. Now, we have derive the transition rates are q_{ij} . Given the state $s_i = (e_p, e_n, d_p, d_n)$, we have:

$$q_{ij} = \begin{cases} R_{p}(E - e_{p} - e_{n})e_{p} + \\ R_{fn}(E - e_{p} - e_{n})e_{n} + \\ R_{fn}e_{p}e_{p} - \\ R_{p}e_{n}e_{n} - \\ R_{p}(D - d_{p} - d_{n})d_{p} + \\ R_{fn}(D - d_{p} - d_{n})d_{n} + \\ R_{fn}d_{p}d_{p} - \\ R_{p}d_{n}d_{n} - \end{cases}$$
(3)

Where x+are represents the transition from a state(\cdots , x, \cdot $\cdot \cdot \cdot$) to $(\cdot \cdot \cdot \cdot, x + 1, \cdot \cdot \cdot)$, and x-represents transition from state $(\cdots, x + 1, \cdots)$ to (\cdots, x, \cdots) . Finally, $q_{ij} = -\sum_{i \neq j} q_{ij}$.

The first transition e_p + is when a intermediate collaborative node changes from NoInfo state to a Positive state $((d_p,d_n,e_p,e_n)$ to (d_p,d_n,e_p+l,e_n)). The rate of change depends on the updating of ρ , and on the δ and θ parameters. The reputation value is ρ increments according to the expression 1. First, the local watchdog can be generate a local PosEvt with rate $\lambda p_d(1 - p_{fn})$ reputation is incremented by δ . Then, the rate of increment due to the local events is $\lambda \delta p_d$ (1 – p_{fn}). Second, updating from an indirect event depends on a number of nodes with Positive and Negative states and the probability of collaboration: $\lambda p_c(c_p - \gamma c_n)$ where $c_p = e_p + d_p$ and $c_n =$ e_n+d_n. Malicious nodes affect this updating by generating indirect NegEvt with a rate $\lambda M p_m$. Since we are evaluating the increment, this term must be positive. So, the final rate due to indirect events is $\lambda \max(p_c(c_p - c_n) - Mp_m)$. All the previous terms are divided by threshold θ in order to obtain the rate of changing when a node contacts with a collaborative node:

$$R_{p} = \lambda \left(\delta p_{d}(1 - p_{fn}) + \max(p_{c}(c_{p} - \gamma c_{n}) - Mp_{m}, 0)\right) / \theta \tag{4}$$

Finally, there are $(E - e_p - e_n)$ nodes with the NoInfo state so the final transition rate is $R_p(E - e_p - e_n)$.

The second transition, e_n+ , is when a intermediate collaborative node changes from $((d_p,d_n,\,e_p,\,e_n)$ to (d_p,d_n,e_p,e_n+1) . This means that a intermediate collaborative node is changes to the Negative state (a false negative). We can derive a similar expression for the rate of change to a (false) Negative state R_{fN} . In this case, when a node contacts with the selfish node, the reputation is decreased with rate $\lambda \delta p_d p_{fn}$, and also by indirect events with rate $\lambda (p_c(\gamma c_n-c_p)+Mp_m)$. Finally, we have:

$$\begin{split} R_{fn} = & \lambda (\delta d_p p_{fn} + \text{max}(p_c(\gamma c_n - c_p) + \text{Mp}_m, 0)) / \theta \\ \text{and the transition is } & R_{fn}(E - e_p - e_n). \end{split} \tag{5}$$

The transition e_p^- is when the intermediate collaborative node that has a Positive state changes to the NoInfo. So, this event is similar to e_n^+ and the transition rate is similar: $R_{fn}e_p^-$. Note that in this case we can multiply by the number of nodes that have the Positive state instead of a number of pending nodes. In similar way, the transition e_n^+ occurs when a intermediate collaborative node that has the Negative state changes to NoInfo. So, transition rate is $R_pe_p^-$.

By using the generator matrix Q we can derive two different expressions: one for a detection time T_d and another for a overall overhead (or cost) O_d . Starting with the detectison time, from the 4D-CTMC we can be obtain how long it will betake for the process to be absorbed. Using the fundamental matrix $N = -T^{-1}$, so ,we can obtain a vector t of the expected time to absorption as $t = N_v$, where v is column vector of ones $(v = [1,1,...1]^T)$. Each entry t_i of t is represents the expected time to absorption

from the state si. Since we only need the expected time from $states_1 = (0, 0, 0, 0)$ to absorption (that is, the expected time for all the destination nodes to have a Positive state), the detection time T_d , is:

$$T_{d} = E[T] = v_1 N_v \tag{6}$$

Where T is a random variable denoting the detection time for all nodes and $v_1 = [1, 0, ..., 0]$. Concerning the overhead we have needed to obtain the number of transmitted messages for each state is s_i . First, the duration of each state is s_i can be obtained using the fundamental matrix N. By definition, the elements of first row of N are to be the expected times in each state starting from state 0. Then, the duration of state s_i is $f_i = N(1, i)$. Now, we calculate the expected number of message sm_i. The number of messages depends on a diffusion model. So the easier exposition, we can start with $\gamma = 0$, that is, only the positive detections are transmitted. From state $s_1 = (0, 0, 0, 0)$ to $s_{E+1}=(0, 0, 0, E)$ no node has aPositive state, so no messages are transmitted and m1 = 0. From states $s_{E+2} = (0, 0, 1, 0)$ to $s_{2E+1} = (0, 0, 1, E - 1)$, one node has a Positive state. In such cases, the Positive can be transmitted to the all nodes (except itself) for the duration of each state i (N(1, i)) with the rate $\boldsymbol{\lambda}$ and the probability $\boldsymbol{p}_c.$ Then, the expected number of messages can be obtained as $m_i = N(1, i)\lambda(C - 1)p_c$. From states $s_{2E+2} = (0, 0, 2, 0)$ to $s_{3E+1} = (0, 0, 2, E - 2)$, we have two possible senders and $m_i = 2N(1, i)\lambda(C - 1)p_c$. Considering both types of nodes (destination and intermediate), the number of nodes with a Positive for state s_i is $\Phi(s_i)=d_p+e_p$. Summarizing, the overhead of transmission (number of messages) is:

$$O_d = E[Msg] = \lambda(C-1)p_c \sum_{i=1}^{\tau} \Phi(s_i)N(1,i).$$
 (7)

Finally, for $\gamma > 0$, the ratio of nodes can that will be transmit a Negative is precisely γ , so $\Phi(s_i) = d_p + e_p + \gamma(d_n + e_n)$.

By using the previous model, we can evaluate the time when the destination nodes D have a "false negative" about the selfish node. In this case absorbing states are $\{0, D, *, *\}$, that is, when $d_n = D$. A high rate of the false negatives and the malicious nodes may be cause a false negative state to reached in less time than a true positive detection.

4.4 The Model for False Positives

This model describe, evaluating the effect of false positives. The diffusion time is similar to detection time of true positives described in the previous subsection, and it canobtained in a similar way. Following process is same that in the previous model for false negatives, we have a 4D-CMTC with a same states $(d_p,d_n,\,e_p,\,e_n)$, but in this case $c_p=d_p+e_p$ represents number of nodes with a false positive, and $c_n=d_n+e_n$ the number of nodes with a negative detection. We can derive expressions similar to 4 and 5, for the case of the false positives. So in this case, R_{fp}

represents the rate of false positive and it is derived in a similar way:

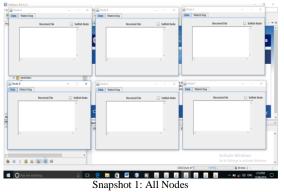
$$R_{fp} = \lambda(\delta p_d p_{fp} + \max(p_c(c_p - \gamma c_n) + Mp_m, 0))/\theta(8)$$

and R_n represents the rate of negative detection:

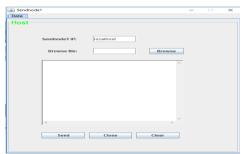
$$Rn = \lambda(\delta p_d (1-p_{fp}) + \max(p_c (\gamma c_n - c_p) - Mp_m, 0))/\theta(9)$$

Using these expressions, the transition rates (qij) of the generator matrix Q are similar to expression 3, substituting R_{p} and R_{fn} by $R_{fp} and R_{n},$ respectively. Finally, byusing equations 6 and 7 described in previous model.So, we can obtain the diffusion time and the overhead.

RESULT AND EVALUATION V.

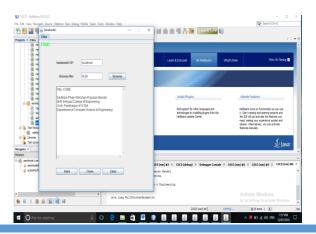


Above figure shows there is 6 node i.e. Node A, Node B, Node C, Node D, Node E, Node V.



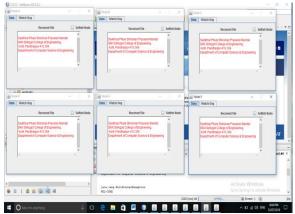
Snapshot 2: Server Frame

In above figure, the Send node work as host. It will browse the data and send to the all node.



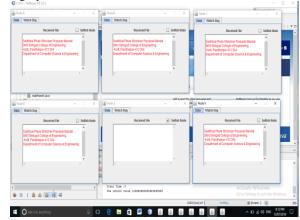
Snapshot 3: Server frame browse the file

In that figure, it will browse the file & open it. After that it will send to all nodes.



Snapshot 4: Without Selfish Node

Above figure shows that all nodes have received file sending by server because there is no Selfish Node.



Snapshot 5: With Selfish Node

Above figure shows that Node E is Selfish Node. So Node E not receive the file. Another all node receives the file.

CONCLUSION VI.

The CoCoWa as a collaborative contact-based watchdog to reduce the time and improve the effectiveness of detectingselfish nodes, reducing the harmful effect of false positives, false negatives and a malicious node. CoCoWa is a based on diffusion of the known positive and the negative detections. When a contact is occurs between two collaborative nodes, the diffusion module is transmits and processes the positive (and negative) detections.

CoCoWa can be reduce the overall detection time with respect to original detection time when the collaboration scheme is not used, with reduced overhead (message cost).

The combined effect of the collaboration and reputation of this approach can be reducing the detection time while increasing the global accuracy by using a moderate local precision watchdog.

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