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Practical no. 8

Aim: Design of a Predictive parser of a given language.

Theory :

Predictive Parser is also another method that implements the technique of Top-Down parsing without Backtracking. A predictive parser is an effective technique of executing recursive-descent parsing by managing the stack of activation records, particularly.

Predictive Parsers has the following components –

- **Input Buffer** – The input buffer includes the string to be parsed followed by an end

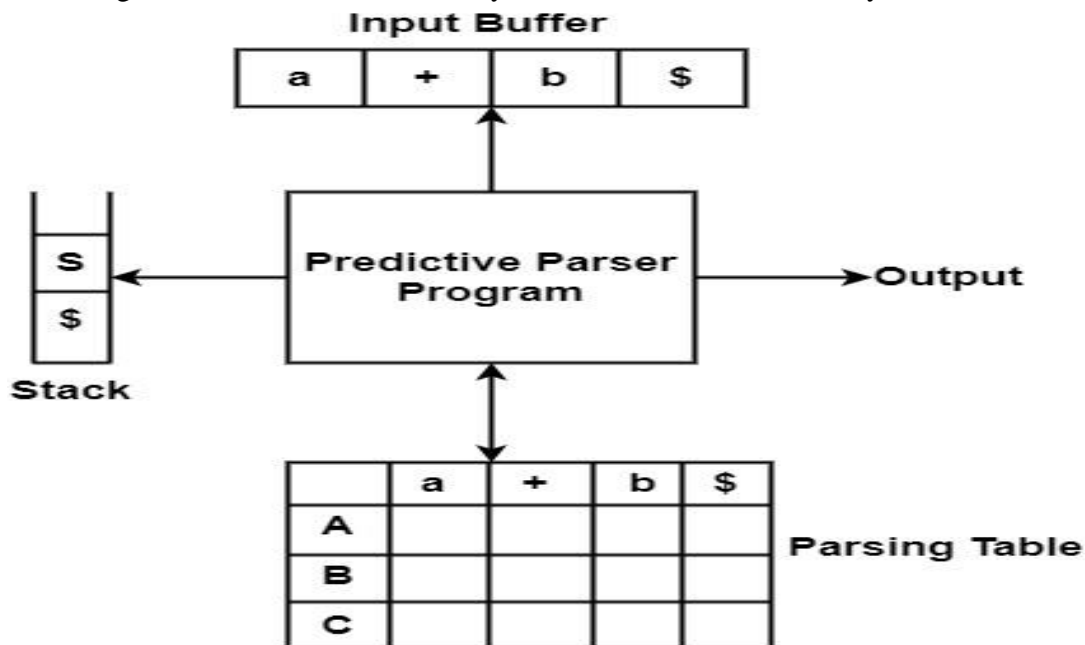
a	+	b	\$
---	---	---	----

Input String

marker \$ to denote the end of the string.

Here a, +, b are terminal symbols.

Stack – It contains a combination of grammar symbols with \$ on the bottom of the stack. At the start of Parsing, the stack contains the start symbol of Grammar followed by \$.

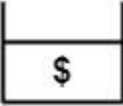
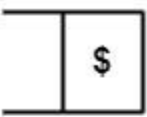

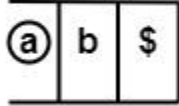

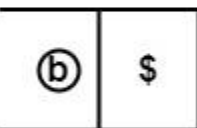
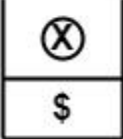



Predictive Parser

- **Parsing Table** – It is a two-dimensional array or Matrix $M[A, a]$ where A is nonterminal and 'a' is a terminal symbol.

All the terminals are written column-wise, and all the Non-terminals are written rowwise.

- **Parsing Program** – The parsing program performs some action by comparing the symbol on top of the stack and the current input symbol to be read on the input buffer.
- **Actions** – Parsing program takes various actions depending upon the symbol on the top of the stack and the current input symbol. Various Actions taken are given below –

Description	Top of Stack	Current Input Symbol	Action
1. If stack is empty, i.e., it only contains \$ and current Input symbol is also \$.			Parsing will be successful and will be halted.
2. If symbol at top of stack and the current input symbol to be read are both terminals and are same.			Pop a from stack & advance to next input symbol.
3. If both top of stack & current input symbol are terminals and top of stack \neq current input symbol e.g. $a \neq b$.			Error
4. If top of stack is non-terminal & input symbol is terminal.			Refer to entry $M[X, a]$ in Parsing Table. If $M[X, a] = X \rightarrow ABC$ then Pop X from Stack Push C, B, A onto stack.

Algorithm to construct Predictive Parsing Table

Input – Context-Free Grammar G

Output – Predictive Parsing Table M

Method – For the production $A \rightarrow \alpha$ of Grammar G.

- For each terminal, a in FIRST (α) add $A \rightarrow \alpha$ to $M[A, a]$.
- If ϵ is in FIRST (α), and b is in FOLLOW (A), then add $A \rightarrow \alpha$ to $M[A, b]$.
- If ϵ is in FIRST (α), and \$ is in FOLLOW (A), then add $A \rightarrow \alpha$ to $M[A, \$]$.

- All remaining entries in Table M are errors.

		Terminal Symbols			
		a	b	+	\$
Non-Terminals Symbols	A				
	B		←		
	C			↑	
	D				

M [B, b]

M[C, +]

Following are the steps to perform Predictive Parsing

- Elimination of Left Recursion
- Left Factoring
- Computation of FIRST & FOLLOW
- Construction of Predictive Parsing Table
- Parse the Input String

Program Code:

```
#include<stdio.h>
#include<conio.h>
#include<string.h>
char prol[7][10]={"S","A","A","B","B","C","C"};
char pror[7][10]={"A","Bb","Cd","aB","@","Cc","@"};
char prod[7][10]={"S->A","A->Bb","A->Cd","B->aB","B->@","C->Cc","C->@"};
char first[7][10]={"abcd","ab","cd","a@","@","c@","@"};
char follow[7][10]={"$","$","$","a$","b$","c$","d$"};
char table[5][6][10];
numr(char c)
{
switch(c)
{
case 'S': return 0;
case 'A': return 1;
case 'B': return 2;
case 'C': return 3;
case 'a': return 0;
case 'b': return 1;
case 'c': return 2;
```

```

case 'd': return 3;
case '$': return 4;
}
return(2);
}
void main()
{
int i,j,k;

for(i=0;i<5;i++)
for(j=0;j<6;j++)
strcpy(table[i][j]," ");
printf("\nThe following is the predictive parsing table for the following
grammar:\n");
for(i=0;i<7;i++)
printf("%s\n",prod[i]);
printf("\nPredictive parsing table is\n");
fflush(stdin);
for(i=0;i<7;i++)
{
k=strlen(first[i]);
for(j=0;j<10;j++)
if(first[i][j]!='@')
strcpy(table[numr(prol[i][0])+1][numr(first[i][j])+1],prod[i]);
}
for(i=0;i<7;i++)
{
if(strlen(pror[i])==1)
{
if(pror[i][0]=='@')
{

k=strlen(follow[i]);
for(j=0;j<k;j++)
strcpy(table[numr(prol[i][0])+1][numr(follow[i][j])+1],prod[i]);

}
}
}
strcpy(table[0][0]," ");

```

```

strcpy(table[0][1],"a");
strcpy(table[0][2],"b");
strcpy(table[0][3],"c");
strcpy(table[0][4],"d");
strcpy(table[0][5],"$");
strcpy(table[1][0],"S");
strcpy(table[2][0],"A");
strcpy(table[3][0],"B");
strcpy(table[4][0],"C");
printf("\n-----\n");
for(i=0;i<5;i++)
for(j=0;j<6;j++)
{
printf("%-10s",table[i][j]);
if(j==5)
printf("\n-----\n");
}
getch();
}

```

Output:

```

The following is the predictive parsing table for the following grammar:
S->A
A->Bb
A->Cd
B->aB
B->@
C->Cc
C->@

Predictive parsing table is
-----
      a      b      c      d      $
-----
S      S->A    S->A    S->A    S->A
-----
A      A->Bb    A->Bb    A->Cd    A->Cd
-----
B      B->aB    B->@      B->@
-----
C      C->@      C->@      C->@
-----

```

Conclusion : In this practical we implemented predictive parser.