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Practical no. 5

Aim: Simulate First and Follow of a Grammar.

Theory :

FIRST and FOLLOW are two functions associated with grammar that help us fill in the entries of an M-table.

FIRST ()– It is a function that gives the set of terminals that begin the strings derived from the production rule.

A symbol c is in FIRST (α) if and only if $\alpha \Rightarrow c\beta$ for some sequence β of grammar symbols.

A terminal symbol a is in FOLLOW (N) if and only if there is a derivation from the start symbol S of the grammar such that $S \Rightarrow \alpha N \alpha \beta$, where α and β are a (possible empty) sequence of grammar symbols. In other words, a terminal c is in FOLLOW (N) if c can follow N at some point in a derivation.

Benefit of FIRST () and FOLLOW ()

- It can be used to prove the LL (K) characteristic of grammar.
- It can be used to promote in the construction of predictive parsing tables.
- It provides selection information for recursive descent parsers.

Computation of FIRST

FIRST (α) is defined as the collection of terminal symbols which are the first letters of strings derived from α .

$$\text{FIRST}(\alpha) = \{a \mid \alpha \rightarrow^* a\beta \text{ for some string } \beta\}$$

If X is Grammar Symbol, then First (X) will be –

If X is a terminal symbol, then $\text{FIRST}(X) = \{X\}$

If $X \rightarrow \epsilon$, then $\text{FIRST}(X) = \{\epsilon\}$

If X is non-terminal & $X \rightarrow a\alpha$, then $\text{FIRST}(X) = \{a\}$

If $X \rightarrow Y_1, Y_2, Y_3$, then $FIRST(X)$ will be

(a) If Y is terminal, then

$$FIRST(X) = FIRST(Y_1, Y_2, Y_3) = \{Y_1\}$$

(b) If Y_1 is Non-terminal and

If Y_1 does not derive to an empty string i.e., If $FIRST(Y_1)$ does not contain ϵ then, $FIRST(X) = FIRST(Y_1, Y_2, Y_3) = FIRST(Y_1)$

(c) If $FIRST(Y_1)$ contains ϵ , then.

$$FIRST(X) = FIRST(Y_1, Y_2, Y_3) = FIRST(Y_1) - \{\epsilon\} \cup FIRST(Y_2, Y_3)$$

Similarly, $FIRST(Y_2, Y_3) = \{Y_2\}$, If Y_2 is terminal otherwise if Y_2 is Non-terminal then

$FIRST(Y_2, Y_3) = FIRST(Y_2)$, if $FIRST(Y_2)$ does not contain ϵ .

If $FIRST(Y_2)$ contain ϵ , then

$$FIRST(Y_2, Y_3) = FIRST(Y_2) - \{\epsilon\} \cup FIRST(Y_3)$$

Similarly, this method will be repeated for further Grammar symbols, i.e., for $Y_4, Y_5, Y_6 \dots Y_K$.

Computation of FOLLOW

Follow (A) is defined as the collection of terminal symbols that occur directly to the right of A.

$$FOLLOW(A) = \{a | S \Rightarrow^* \alpha A a \beta \text{ where } \alpha, \beta \text{ can be any strings}\}$$

Rules to find FOLLOW

If S is the start symbol, $FOLLOW(S) = \{\$ \}$

If production is of form $A \rightarrow \alpha B \beta$, $\beta \neq \epsilon$.

(a) If $FIRST(\beta)$ does not contain ϵ then, $FOLLOW(B) = \{FIRST(\beta)\}$

Or

(b) If $\text{FIRST}(\beta)$ contains ϵ (i. e. , $\beta \Rightarrow^* \epsilon$), then

$$\text{FOLLOW}(B) = \text{FIRST}(\beta) - \{\epsilon\} \cup \text{FOLLOW}(A)$$

\therefore when β derives ϵ , then terminal after A will follow B.

If production is of form $A \rightarrow \alpha B$, then $\text{Follow}(B) = \{\text{FOLLOW}(A)\}$.

Program Code:

```
#include<bits/stdc++.h>
using namespace std;

set<char> ss;
bool dfs(char i, char org, char last, map<char,vector<vector<char>>> &mp){
    bool rtake = false;
    for(auto r : mp[i]){
        bool take = true;
        for(auto s : r){
            if(s == i) break;
            if(!take) break;
            if(!(s>='A'&&s<='Z')&&s!='e'){
                ss.insert(s);
                break;
            }
            else if(s == 'e'){
                if(org == i||i == last)
                    ss.insert(s);
                rtake = true;
                break;
            }
            else{
                take = dfs(s,org,r[r.size()-1],mp);
                rtake |= take;
            }
        }
    }
    return rtake;
}
```

```
}
```

```
int main(){
    int i,j;
    ifstream fin("inputfirstfollow.txt");
    string num;
    vector<int> fs;
    vector<vector<int>> a;
    map<char,vector<vector<char>>> mp;
    char start;
    bool flag = 0;
    cout<<"Grammar: "<<"\n";
    while(getline(fin,num)){
        if(flag == 0) start = num[0],flag = 1;
        cout<<num<<"\n";
        vector<char> temp;
        char s = num[0];
        for(i=3;i<num.size();i++){
            if(num[i] == "|"){
                mp[s].push_back(temp);
                temp.clear();
            }
            else temp.push_back(num[i]);
        }
        mp[s].push_back(temp);
    }
    map<char,set<char>> fmp;
    for(auto q : mp){
        ss.clear();
        dfs(q.first,q.first,q.first,mp);
        for(auto g : ss) fmp[q.first].insert(g);
    }

    cout<<"\n";
    cout<<"FIRST: "<<"\n";
    for(auto q : fmp){
        string ans = "";
        ans += q.first;
        ans += " = {";
        for(char r : q.second){
```

```

        ans += r;
        ans += ',';
    }
    ans.pop_back();
    ans+="}";
    cout<<ans<<"\n";
}

```

```

map<char,set<char>> gmp;
gmp[start].insert('$');
int count = 10;
while(count--){
    for(auto q : mp){
        for(auto r : q.second){
            for(i=0;i<r.size()-1;i++){
                if(r[i]>='A'&&r[i]<='Z'){
                    if(!(r[i+1]>='A'&&r[i+1]<='Z')) gmp[r[i]].insert(r[i+1]);
                }
                else {
                    char temp = r[i+1];
                    int j = i+1;
                    while(temp>='A'&&temp<='Z'){
                        if(*fmp[temp].begin()=='e'){
                            for(auto g : fmp[temp]){
                                if(g=='e') continue;
                                gmp[r[i]].insert(g);
                            }
                        }
                        j++;
                        if(j<r.size()){
                            temp = r[j];
                            if(!(temp>='A'&&temp<='Z')){
                                gmp[r[i]].insert(temp);
                                break;
                            }
                        }
                    }
                }
                else{
                    for(auto g : gmp[q.first]) gmp[r[i]].insert(g);
                    break;
                }
            }
        }
    }
    else{

```

```

        for(auto g : fmp[temp]){
            gmp[r[i]].insert(g);
        }
        break;
    }
}
}
}
}
}
if(r[r.size()-1]>='A'&&[r.size()-1]<='Z'){
    for(auto g : gmp[q.first]) gmp[r[i]].insert(g);
}
}
}
}
}

```

```

cout<<"\n";
cout<<"FOLLOW: ";<<"\n";
for(auto q : gmp){
    string ans = "";
    ans += q.first;
    ans += " = {";
    for(char r : q.second){
        ans += r;
        ans += ',';
    }
    ans.pop_back();
    ans+="}";
    cout<<ans<<"\n";
}
return 0;
}

```

Input (inputfirstfollow.txt):

S->aBDh

B->cC

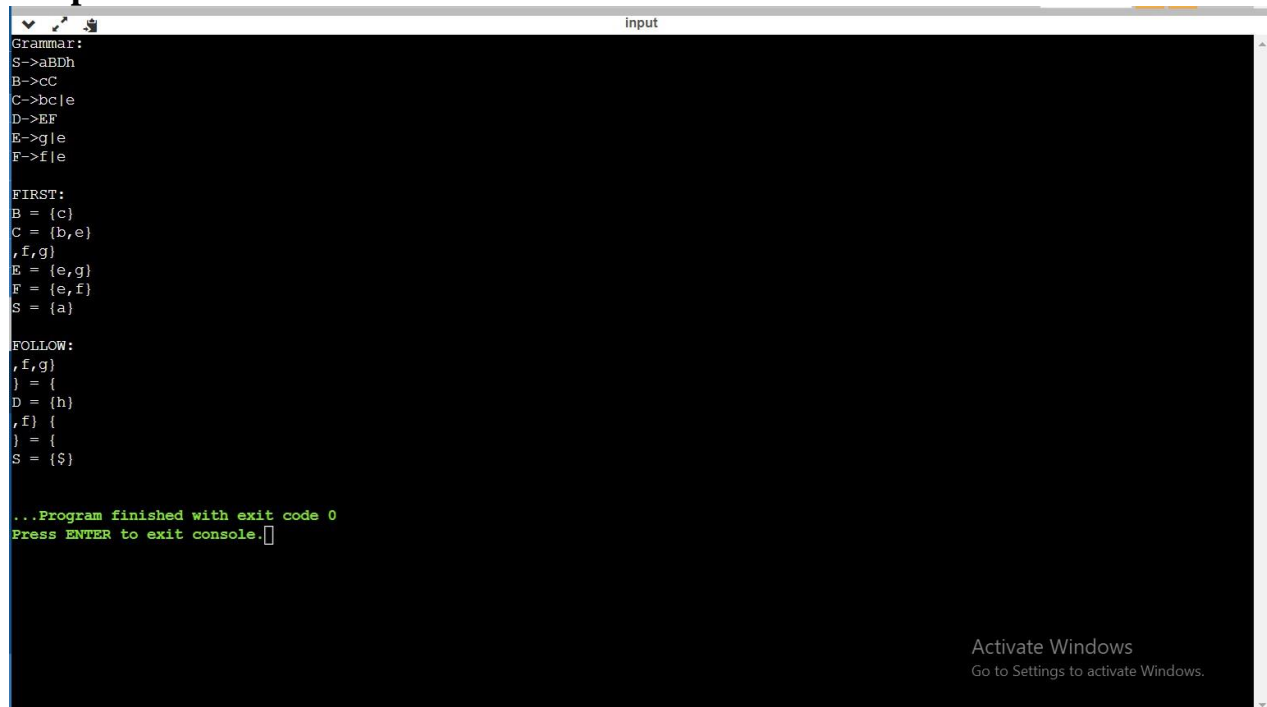
C->bc|e

D->EF

E->g|e

F->f|e

Output:



```
Grammar:
S->aBDh
B->cC
C->bc|e
D->EF
E->g|e
F->f|e

FIRST:
B = {c}
C = {b,e}
, f, g
E = {e,g}
F = {e,f}
S = {a}

FOLLOW:
, f, g
} = {
D = {h}
, f} {
} = {
S = {$}

...Program finished with exit code 0
Press ENTER to exit console.
```

Activate Windows
Go to Settings to activate Windows.

Conclusion : In this practical we implemented a program that simulate the process of finding first and follow of a grammar.