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Aim: Write a C program to implement operator precedence parsing.

Theory :

Operator Precedence Parsing is also a type of Bottom-Up Parsing that can be used to a class of Grammars known as Operator Grammar.

A Grammar G is Operator Grammar if it has the following properties –

- Production should not contain ϵ on its right side.
- There should not be two adjacent non-terminals at the right side of production.

Example1 – Verify whether the following Grammar is operator Grammar or not.

$E \rightarrow E A E \mid (E) \mid id$

$A \rightarrow + \mid - \mid *$

Solution

No, it is not an operator Grammar as it does not satisfy property 2 of operator Grammar.

As it contains two adjacent Non-terminals on R.H.S of production $E \rightarrow E A E$.

We can convert it into the operator Grammar by substituting the value of A in $E \rightarrow E A E$.

$E \rightarrow E + E \mid E - E \mid E * E \mid (E) \mid id$.

Operator Precedence Relations

Three precedence relations exist between the pair of terminals.

Relation	Meaning
$p < . q$	p has less precedence than q.
$p > . q$	p has more precedence than q.
$p = . q$	p has equal precedence than q.

Depending upon these precedence Relations, we can decide which operations will be executed or parsed first.

Association and Precedence Rules

- If operators have different precedence

Since * has higher precedence than +

Example—

In a statement $a + b * c$

$\therefore + < . *$

In statement $a * b + c$

$\therefore * . > +$

- If operators have Equal precedence, then use Association rules.

(a) Example minus; In statement $a + b + c$ here + operators are having equal precedence.

As '+' is left Associative in $a + b + c$

$\therefore (a + b)$ will be computed first, and then it will be added to c.

i.e., $(a + b) + c$

$+ . > +$

Similarly, '*' is left Associative in $a * b * c$

(b) Example — In a statement $a \uparrow b \uparrow c$ here, \uparrow is the Right Associative operator

\therefore It will become $a \uparrow (b \uparrow c)$

$\therefore (b \uparrow c)$ will be computed first.

$\therefore \uparrow < . \uparrow$

- Identifier has more precedence than all operators and symbols.

$\therefore \theta < . id \quad \$ < . id$

$id . > \theta \quad id . > \$$

$id . >)$

$(< . id .$

- \$ has less precedence than all other operators and symbols.

$\$ < . \quad (id . > \$$

$\$ < . + \quad) . > \$$

$\$ < . *$

Example2 – Construct the Precedence Relation table for the Grammar.

$$E \rightarrow E + E \mid E * E / id$$

Solution

Operator-Precedence Relations

	Id	+	*	\$
Id		.>	.>	.>
+	<.	.>	<.	.>
*	<.	.>	.>	.>
\$	<.	<.	<.	

Advantages of Operator Precedence Parsing

- It is accessible to execute.

Disadvantages of Operator Precedence Parsing

- Operator Like minus can be unary or binary. So, this operator can have different precedence's in different statements.
- Operator Precedence Parsing applies to only a small class of Grammars.

Program Code:

```
#include<stdio.h>
#include<stdlib.h>
#include<string.h>
void main()
{
    char stack[20],ip[20],opt[10][10][1],ter[10];
    int i,j,k,n,top=0,col,row;
    for(i=0;i<10;i++)
    {
        stack[i]=NULL;
        ip[i]=NULL;
        for(j=0;j<10;j++)
        {
            opt[i][j][1]=NULL;
        }
    }
    printf("Enter the no.of terminals:\n");
    scanf("%d",&n);
```

```

printf("\nEnter the terminals:\n");
for(i=0;i<n;i++)
{
scanf("%s",&ter[i]);
}
printf("\nEnter the table values:\n");
for(i=0;i<n;i++)
{
for(j=0;j<n;j++)
{
printf("Enter the value for %c %c: ",ter[i],ter[j]);
scanf("%s",opt[i][j]);
}
}
printf("\n** OPERATOR PRECEDENCE TABLE **\n");
for(i=0;i<n;i++)
{
printf("\t%c",ter[i]);
}
printf("\n");
for(i=0;i<n;i++)
{
printf("\n%c",ter[i]);
for(j=0;j<n;j++)
{
printf("\t%c",opt[i][j][0]);
}
}
stack[top]='$';
printf("\nEnter the input string: ");
scanf("%s",ip);
i=0;
printf("\nSTACK\t\t\tINPUT STRING\t\t\tACTION\n");
printf("\n%s\t\t\t%s\t\t\t",stack,ip);
while(i<=strlen(ip))
{
for(k=0;k<n;k++)
{
if(stack[top]==ter[k])
col=k;

```

```

if(ip[i]==ter[k])
row=k;
}
if((stack[top]=='$')&&(ip[i]=='$'))
{
printf("\nString is accepted\n");
break;
}
else if((opt[col][row][0]=='<') ||(opt[col][row][0]=='='))
{
stack[++top]=opt[col][row][0];
stack[++top]=ip[i];
printf("Shift %c",ip[i]);
i++;
}
else
{
if(opt[col][row][0]=='>')
{
while(stack[top]!='<'){--top;}
top=top-1;
printf("Reduce");
}
else
{
printf("\nString is not accepted");
break;
}
}
printf("\n");
for(k=0;k<=top;k++)
{
printf("%c",stack[k]);
}
printf("\t\t");
for(k=i;k<strlen(ip);k++)
{
printf("%c",ip[k]);
}
printf("\t\t");

```

```
}
}
```

Output:

```

Enter the no.of terminals:
4

Enter the terminals:
+
*
i
$

Enter the table values:
Enter the value for + +: >
Enter the value for + *: >
Enter the value for + i: <
Enter the value for + $: >
Enter the value for * +: >
Enter the value for * *: >
Enter the value for * i: <
Enter the value for * $: >
Enter the value for i +: >
Enter the value for i *: >
Enter the value for i i: =
Enter the value for i $: >
Enter the value for $ +: <
Enter the value for $ *: <
Enter the value for $ i: <
Enter the value for $ $: A

** OPERATOR PRECEDENCE TABLE **
      +      *      i      $
+      >      >      <      >
*      >      >      <      >
i      >      >      =      >
$      <      <      <      A
Enter the input string: i+i*i$

STACK      INPUT STRING      ACTION

$          i+i*i$           Shift i
$<i        +i*i$            Reduce
$          +i*i$            Shift +
$<+        1*i$             Shift i
$<+<i      *i$              Reduce
$<+        *i$              Reduce
$          *i$              Shift *
$<*        i$               Shift i
$<+<i      $                Reduce
$<*        $                Reduce
$          $                Reduce
String is accepted

...Program finished with exit code 0
Press ENTER to exit console.

```

Conclusion : In this practical we learnt how operator precedence parser works for parsing operator precedence grammar.