

Computer based Optimization Technique Lecture-II & III

LPP

Graphical Method

Types of Solution

Conversion of Normal LPP to standard LPP

Simplex method

Brief History

- It is both Arts and Science
- It has started in World War-II (since 1930s) and then soon become very popular in mathematics, Industry, economics, management, and all most in any fields.
- Primary application areas include forecasting, production, scheduling, inventory control and project management
- It is also known as “Operation Research”, “Decision Science”, and Optimization Technique

Optimization technique is the scientific study of operations for the purpose of making better decisions

Objective

- Decision making and improve its quality.
- Identify optimum solution.
- Integrating the systems.
- Improve the objectivity of analysis.
- Minimize the cost and maximize the profit.
- Improve the productivity.
- Success in competition and market leadership

Different Methods

1. Simulation methods

- It gives ability to conduct sensitive study to -
 - (a) search for improvements and
 - (b) test the improvement ideas that are being made.

2. *Optimization methods.*

- Here goal is to enable the decision makers to identify and locate the very best choice, where innumerable feasible choices are available and comparing them is difficult.

- ***3. Data-analysis methods***

- The goal is to aid the decision-maker in detecting actual patterns and inter-connections in the data set and
- Use of this analysis for making solutions.
- This method is very useful in Public Health.

Process

1. Identification of program problem.

- Most critical step in the process.
- Unless problem is clearly defined it is impossible to develop good solutions.

2. Identification of possible reasons and solutions .

- Once the problem has been identified , it is the job of the program implementer and researcher to determine the reasons for the problem and generate possible solutions

3. Testing of potential solution

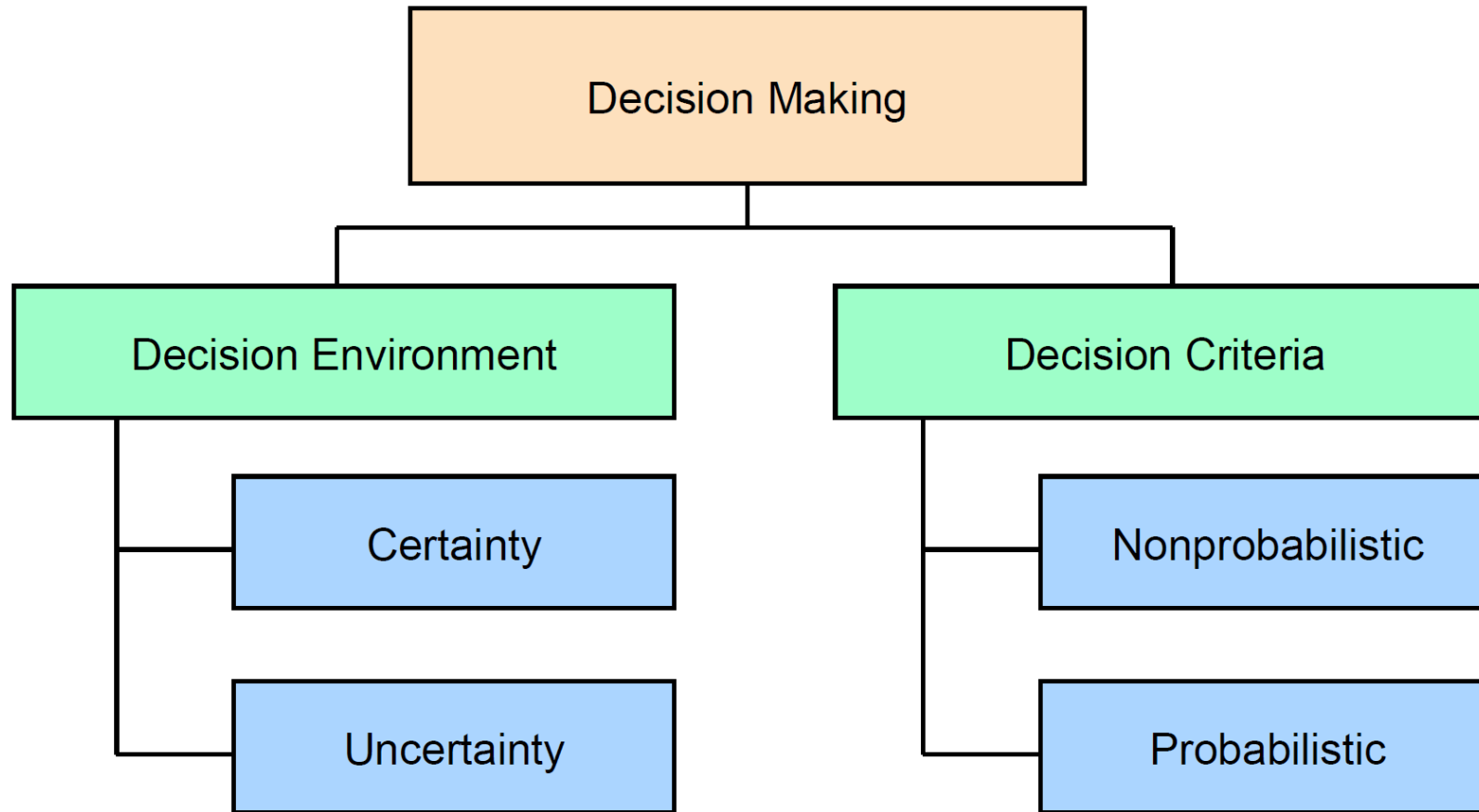
4. Result utilization

- It is necessary to decide how its results are meant to be used.

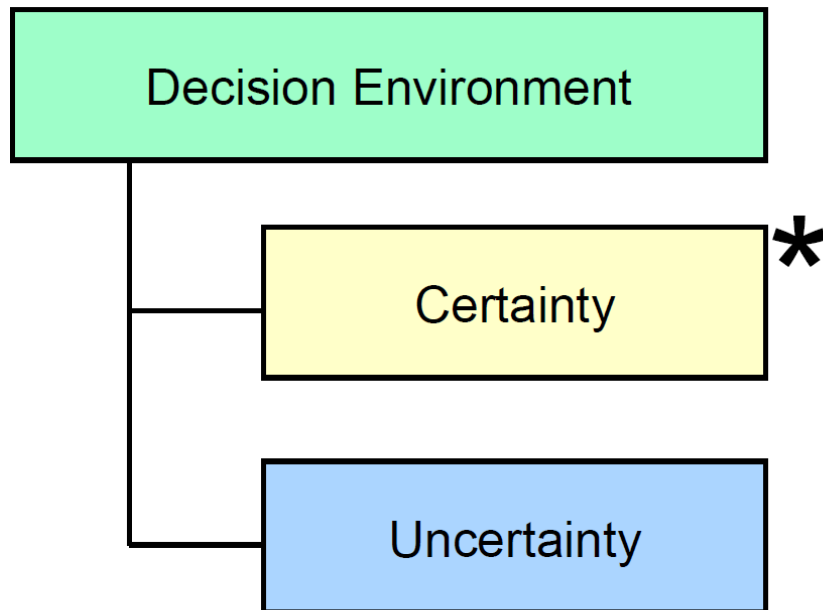
5. Results dissemination

- Results dissemination are done in the form of *seminars* or by *meeting* with decision makers.

Decision Making Overview



Decision Environment



Certainty: The results of decision alternatives are known

Decision Making Under Certainty When it is known for certain which is of the possible future conditions will happen, just choose the alternative that has the best payoff under the state of nature

Example:

Must print 10,000 color brochures

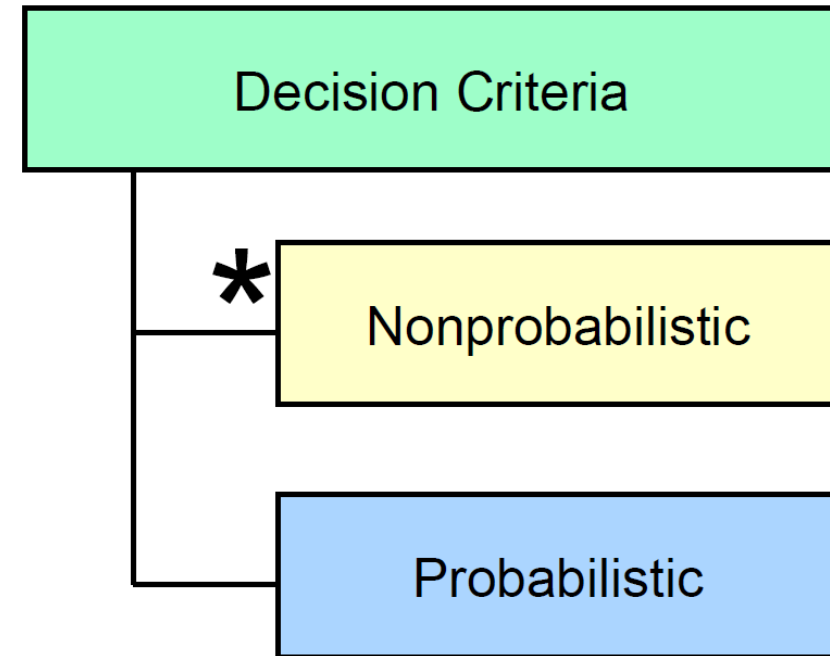
Offset press A: 2,000 fixed cost
+ 0.24 per page

Offset press B: 3,000 fixed cost
+ 0.12 per page

Decision Criteria

Nonprobabilistic Decision Criteria: Decision rules that can be applied if the probabilities of uncertain events are not known.

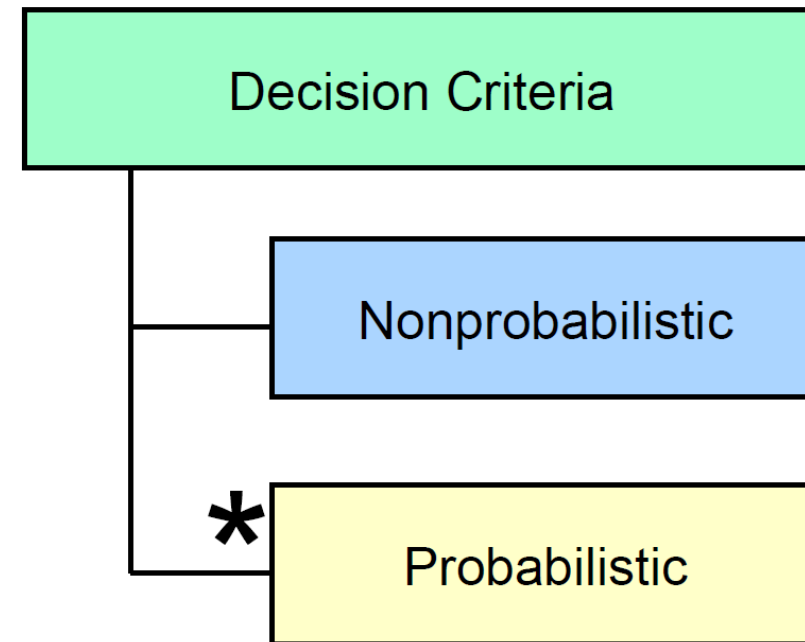
- maximax criterion
- maximin criterion
- minimax regret criterion



Decision Criteria

Probabilistic Decision Criteria: Consider the probabilities of uncertain events and select an alternative to maximize the expected payoff or minimize the expected loss

- maximize expected value
- minimize expected opportunity loss



Pay-off

- The **payoffs** of a decision analysis problem are ***the benefits or rewards that result from selecting a particular decision alternative***

	Strong Economy	Stable Economy	Weak Economy
Large Store	200	50	-100
Average Store	100	100	-50
Small Store	50	30	25

- *Maximax, Minimax and Minimax regret (opportunity loss)*

Linear Programming Problem (LPP)

- Objective function
- Constraints
 - Unequal
 - Greater than equal
 - Less than equal
 - Equal
- Solution Approach
 - Graphical
 - Simplex Method

Linear Problem

- Objective Function
 - Optimization Criteria
 - Maximization
 - Minimization
- Constraints
 - Inequal
 - Greater than equal to
 - Less than equal to
 - Equal

Two-Dimensional LP

- There are only two decision variable

2-dim Linear Programming Problem

$$\text{Min/ max } c_1x_1 + c_2x_2$$

Subject to

$$a_{11}x_1 + a_{12}x_2 \leq b_1$$

$$a_{21}x_1 + a_{22}x_2 \leq b_2$$

$$a_{m1}x_1 + a_{m2}x_2 \leq b_m$$

$$x_1, x_2, b_1, b_2 \geq 0$$

Four possible Solution cases

- Unique Solution
 - There is only one solution to the problem
- Multiple Solutions
 - There are multiple solution to the problem which gives the same answer
- Unbounded Solution
 - The solution is not bounded i.e almost infinite number of solutions
- Infeasible Solution
 - No solution exists

Unique Solution

1. Unique solution

$$\text{Maximize } 5x_1 + 3x_2$$

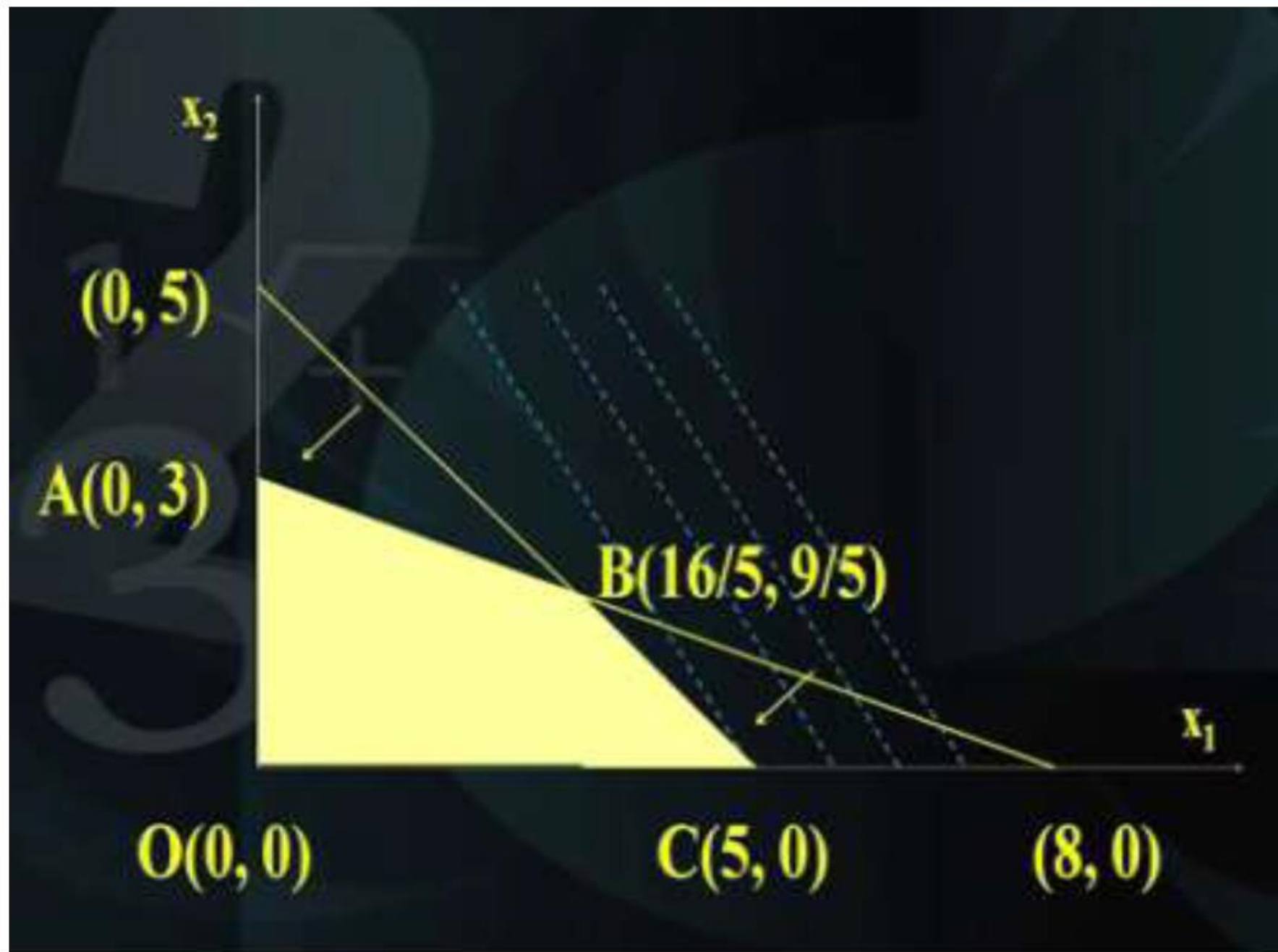
$$\text{Subject to } x_1 + x_2 \leq 5$$

$$3x_1 + 8x_2 \leq 24$$

$$x_1 \geq 0, x_2 \geq 0$$

Solution steps

- Plot the x-axis on the horizontal axis, that is, the x_1 variable will be represented
- Plot the y-axis on the vertical axis, that is, the x_2 variable will be represented
 $O(0,0)$ is origin
- Plot the first constraint that is $x_1 + x_2 < 5$,
 - This is the straight line passing through the two points $(0, 5)$ and $(5, 0)$.
 - A line uniquely passes through two points and this can be obtained by putting $x_1 = 0$ and getting the corresponding value of x_2 similarly, putting $x_2 = 0$ and getting the corresponding value of x_1 .
- Similarly plot for the second constraint
- Find out the feasible region
- Calculate the objective function value (At corner points)
- Find the Maximum



The Best Solutions

Vertices of feasible region	Objective function value $5x_1 + 3x_2$
O(0, 0)	0
A(0, 3)	9
B(16/5, 9/5)	21.4
C(5, 0)	25 Max

Multiple Solutions

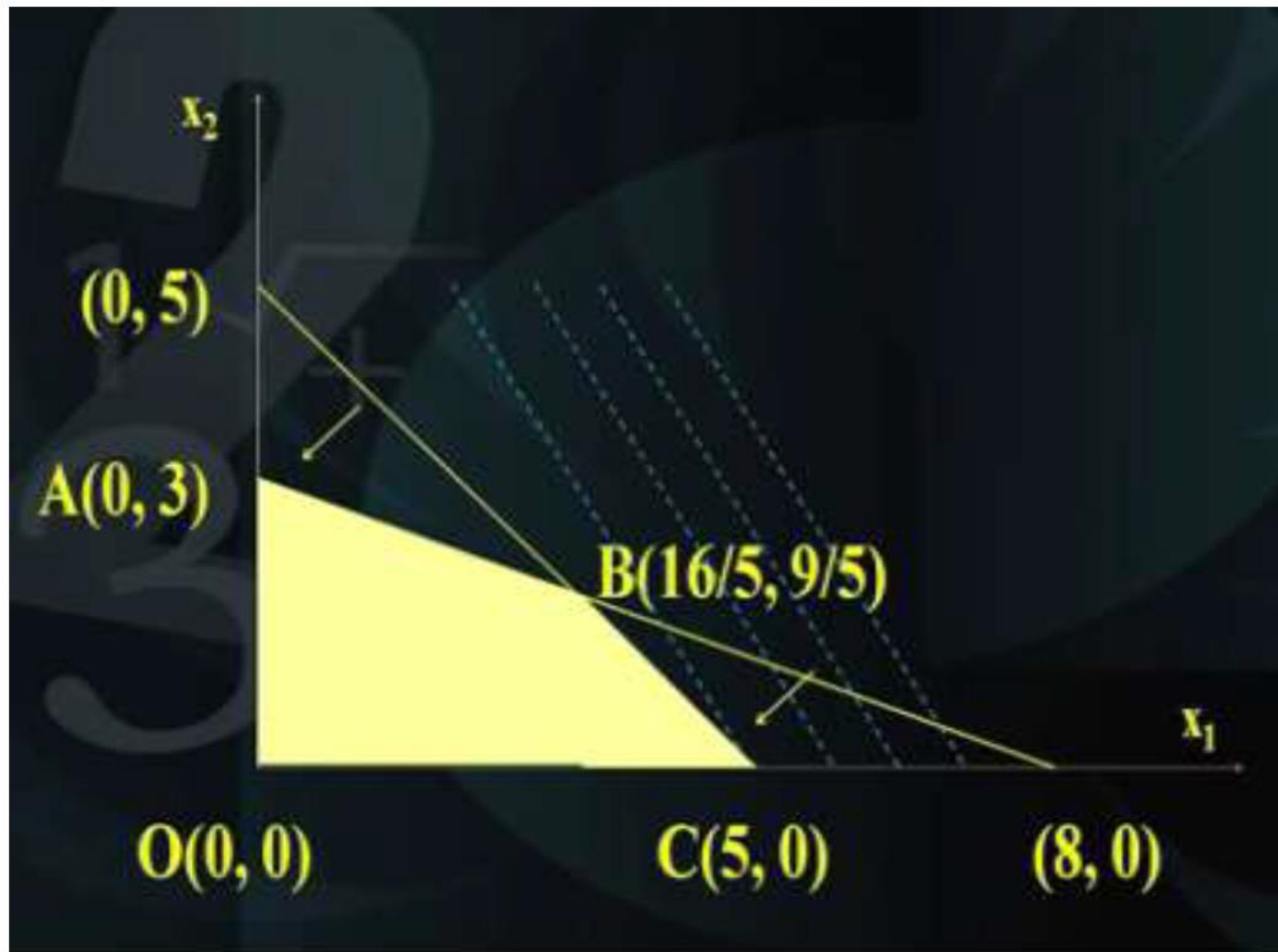
2. Multiple solutions

Maximize $3x_1 + 8x_2$

Subject to $x_1 + x_2 \leq 5$

$3x_1 + 8x_2 \leq 24$

$x_1 \geq 0, x_2 \geq 0$



Multiple best solutions

Vertices of feasible region	Objective function value $3x_1 + 8x_2$
O(0, 0)	0
A(0, 3)	24 max
B(16/5, 9/5)	24 Max
C(5, 0)	15

Unbounded Solution

3. Unbounded solutions

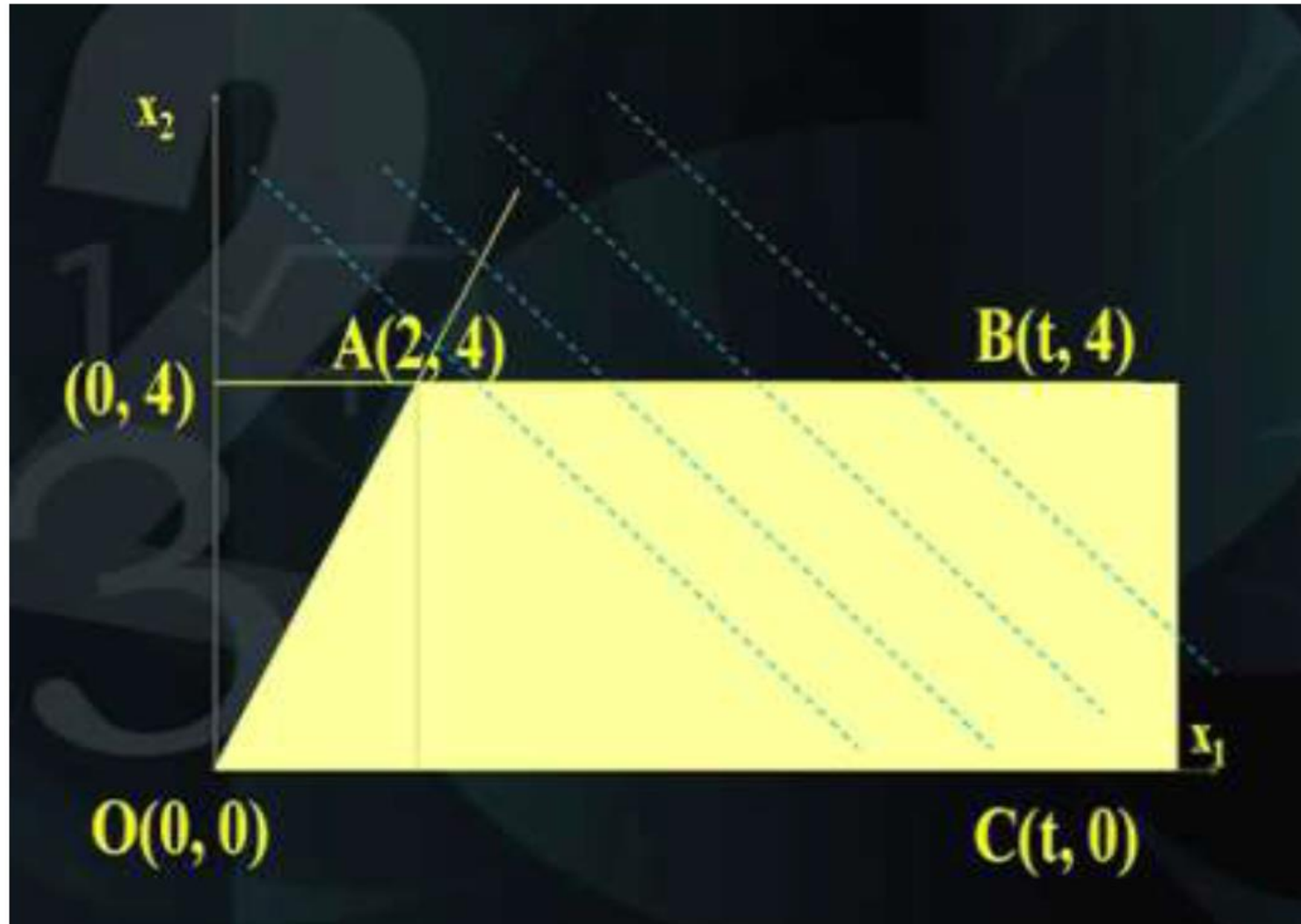
Maximize $3x_1 + 8x_2$

Subject to $2x_1 - x_2 \leq 0$

$x_2 \leq 4$

$x_1 \geq 0, x_2 \geq 0$

The feasible solution Region



Unbounded number of Solutions

Vertices of feasible region	Objective function value $3x_1 + 8x_2$
O(0, 0)	0
A(2, 4)	38
B(t, 4)	$6 + 8t$
C(t, 0)	$8t$

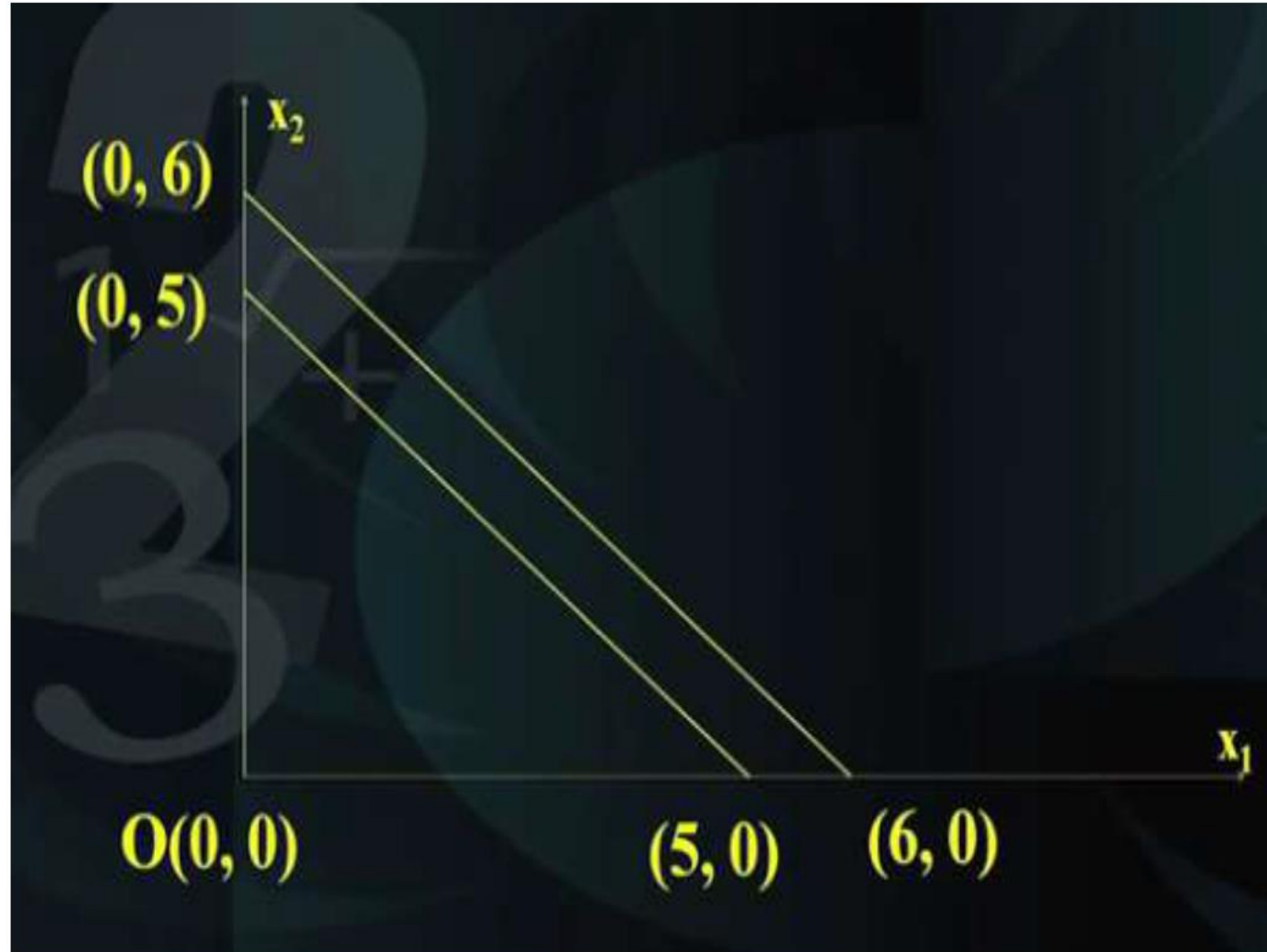
As t increases $6 + 8t$ also increases

Infeasible Solution

4. Infeasible solution

$$\begin{array}{ll}\text{Maximize} & 3x_1 - 2x_2 \\ \text{subject to} & x_1 + x_2 \leq 5 \\ & x_1 + x_2 \geq 6 \\ & x_1, x_2 \geq 0\end{array}$$

No intersection of the constraint plots and feasible region is not bounded



Simplex Method

- Definition of general LPP
- Conversion of general LPP to standard LPP
 - Basic variable
 - Slack and surplus variable
 - To convert the inequality constraints to equality constraints
- Find out the Incoming and outgoing variable
 - Identify the Pivot Column, Pivot Row and Pivot Element

Definition of a general LPP

DEF OF A GENERAL LPP

Min / Max $c_1x_1 + c_2x_2 + \dots c_nx_n$

Subject to

$$a_{11}x_1 + a_{12}x_2 + \dots a_{1n}x_n \leq b_1$$

$$a_{21}x_1 + a_{22}x_2 + \dots a_{2n}x_n \leq b_2$$

.....

$$a_{m1}x_1 + a_{m2}x_2 + \dots a_{mn}x_n \leq b_m$$

$$x_i \geq 0, i = 1, 2, \dots n$$

$$m \neq n$$

Definition of standard LPP

LPP IN STANDARD FORM

$$\text{Max } c_1x_1 + c_2x_2 + \dots c_nx_n$$

Subject to

$$a_{11}x_1 + a_{12}x_2 + \dots a_{1n}x_n = b_1$$

$$a_{21}x_1 + a_{22}x_2 + \dots a_{2n}x_n = b_2$$

$$\dots \dots \dots$$
$$a_{m1}x_1 + a_{m2}x_2 + \dots a_{mn}x_n = b_m$$

$$x_i \geq 0, i = 1, 2, \dots n$$

$$b_j \geq 0, j = 1, 2, \dots m$$

$$m \neq n$$

Use of Slack and Surplus Variable

- Slack variable is used to convert \leq to equal
 - $X_1 \leq 5$ can be written as $X_1 + s_1 = 5$, where $s_1 \geq 0$
- Surplus Variable is used to convert \geq to equal
 - $X_1 \geq 5$ can be written as $X_1 - s_1 = 5$, where $s_1 \geq 0$
- For unrestricted constraint (X_1 can be positive or negative)
 - $X_1 + s_1 - s_2 = 0$ (where $s_1, s_2 \geq 0$)
- Conversion of RHS to +ve
 - Multiply the row by -1 to make RHS +ve if it is -ve

Example

EXAMPLE

Maximize $4x_1 + 2x_2$

Subject to

$$7x_1 - 3x_2 \leq 15$$

$$7x_1 - 3x_2 + x_3 = 15$$

$$3x_1 + 4x_2 \geq 20$$

$$3x_1 + 4x_2 - x_4 = 20$$

$$5x_1 - 4x_2 = -60$$

$$-5x_1 + 4x_2 = 60$$

$$\text{All } x_i \geq 0$$

Elementary Row Operations

- Multiply Row by a scalar
- Add/Subtract Row to/from another row
- Interchange two rows

Terminology

- Basic variable: A variable appears with coefficient 1 in equation 'I' and with 0 in all other
- Non Basic Variable: All variable which are not basic
- Pivot Operation: By using elementary row operation a non-basic variable can be converted to basic variable
- Initial basic feasible solution obtained by setting all non-basic variables to zero and solving for basic variables where the basic variables are non –ve
-

Simplex Method Explanation

PROCEDURE OF SIMPLEX METHOD

Example:

$$\text{Maximize } z = 5x_1 + 2x_2 + 3x_3 - x_4 + x_5$$

$$\text{S. t. } x_1 + 2x_2 + 2x_3 + x_4 = 8$$

$$3x_1 + 4x_2 + x_3 + x_5 = 7$$

$$x_i \geq 0, i = 1, 2, 3, 4, 5.$$

Initial b.f.s. is

$$x_4 = 8, x_5 = 7, x_1 = x_2 = x_3 = 0.$$

Procedure for Simplex method

- Construct the initial simplex table

INITIAL TABLE

.		5	2	3	-1	1	
C_0	Basis	x_1	x_2	x_3	x_4	x_5	RHS
-1	x_4	1	2	2	1	0	8
1	x_5	3	4	1	0	1	7
	Dev.	3	0	4	0	0	$Z=-1$

Identifying Incoming variable

ENTERING VARIABLE IN BASIS

Deviations

$$5 - (-1 \ 1) (1 \ 3)^t = 3$$

$$2 - (-1 \ 1) (2 \ 4)^t = 0$$

$$3 - (-1 \ 1) (2 \ 1)^t = 4 \quad \text{largest}$$

$$5 - (-1 \ 1) (1 \ 3)^t = 0$$

$$-1 - (-1 \ 1) (1 \ 0)^t = 0$$

$$1 - (-1 \ 1) (0 \ 1)^t = 0$$

Entering variable = x_3

Outgoing Variable using Minimum Ratio Test

LEAVING VARIABLE FROM BASIS

Variable corresponding to
minimum Ratio Test

= mini (RHS and +ive pivot column)

= min ($8/2$, $7/1$)

= $8/2$

Leaving variable = x_4

Second Iteration

SECOND TABLE

		5	2	3	-1	1	
C_0	Basis	x_1	x_2	x_3	x_4	x_5	RHS
3	x_3	1/2	1	1	1/2	0	4
1	x_5	5/2	3	0	-1/2	1	3
	Dev.	1	-4	0	-2	0	Z=15

Third Iteration

THIRD TABLE

		5	2	3	-1	1	
C_0	Basis	x_1	x_2	x_3	x_4	x_5	RHS
3	x_3	0	$2/5$	1	$3/5$	$-1/5$	$17/5$
5	x_1	1	$6/5$	0	$-1/5$	$2/5$	$6/5$
	Dev.	0	$-26/5$	0	$-9/5$	$-2/5$	$81/5$

Computational Method

1. Convert all right side constant of constraints are positive
2. Convert unequal constraints to equations by introducing slack/surplus variable
3. Present the constraints equations in matrix form
4. Construct the starting simplex table (as discussed)
5. Test basic feasible solution for optimality

$$\Delta_j = c_j - x_j = c_j - C_B X_j$$

Optimality test

- (i) If all $\Delta_j \leq 0$ the solution is optimal (alternate solution exists if any non-basic is also zero)
 - (ii) At least one Δ_j is positive, The solution is non-optimal, proceed to next iteration
 - (iii) If corresponding to most positive Δ_j , all elements of the column is negative or zero, the solution is unbounded
6. Find out the incoming and outgoing vector

Contd...

7. Make the key element equal to 1 (unity), perform required matrix operations

8. Repeat steps 5 thru 7 until optimality condition is satisfied

Quick solution

1. Mark the incoming vector (max of Δ_j), let it be column 'k'

2. Key element is found at the intersection of incoming and outgoing vector (Min X_B / X_k)

3. ***Perform Row Operations***

Contd..

4. Δ_j corresponding to unit matrix col. Is always zero
5. Δ_j Will be equal to C_j in the initial solution with only slack or surplus variable in the initial basic solution matrix
6. While transforming the table, the value of 'z' and corresponding is also calculated Δ_j

Two phase simplex method

- Artificial Variable
 - To deal with \geq and $=$ constraints
 - Basic matrix can be obtained as an identity matrix in the simplex table
 - With slack/surplus variable, basic matrix can not be obtained
 - Use artificial variable (Fictitious/ no meaning)

Phase-I

- Remove the artificial variables
 - Start with artificial variable in the basic solution
 - Once the artificial variable is outgoing (drop the entire column)
 - Obtain the standard simplex table

- Phase-II
 - Proceed as discussed earlier

Big-M (Penalty) method

- $X_1, X_2, \dots, S_1, S_2, \dots, A_1, A_2, \dots$
- Assign $-M$ (M is very large +ve integer) as cost associated with Artificial Variable
- In the Δ_j row, choose the minimum negative M value (ignore other constants) as outgoing row

Difficult for computation

Comparison between Two-phase and Big-M method

Two-Phase	Big-M
Suitable for Computation	Not Suitable for Computation
No. of Iteration is large if there are more Artificial Variables	No. of iterations do not depend on the number of artificial variable
Preferable if less number of artificial variable	Preferable if less number of artificial variable
No choice of M is required	Choice of M is crucial
It is a two stage process	Single stage process

Other Mathematical Solutions

- Solution of Simultaneous linear equations using simplex method
 - Formulate the objective function with associated cost as zero
 - Simultaneous linear equations are written as equality constraint
 - Use two phase simplex method for solving the equations
 - The solution will give the solution of variables
- Finding matrix inversion problem using simplex method
 - Similar to solving the simultaneous linear equations
 - The values corresponding to artificial variable columns is the matrix inversion