

DESIGN AND ANALYSIS OF ALGORITHMS

LAB WORKBOOK WEEK – 8

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ROLL NUMBER: CH.SC.U4CSE24120

CLASS: CSE-B

Huffman Coding:

DATA ANALYTICS AND INTELLIGENCE LABORATORY

Code:

```
//CH.SC.U4CSE24120
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#define MAX 100
struct Node {
    char data;
    int freq;
    struct Node *left, *right;
};
struct Node* createNode(char data, int freq) {
    struct Node* node = (struct Node*)malloc(sizeof(struct Node));
    node->data = data;
    node->freq = freq;
    node->left = node->right = NULL;
    return node;
}
void sort(struct Node* arr[], int n) {
    for(int i = 0; i < n-1; i++) {
        for(int j = i+1; j < n; j++) {
            if(arr[i]->freq > arr[j]->freq) {
                struct Node* temp = arr[i];
                arr[i] = arr[j];
                arr[j] = temp;
            }
        }
    }
}
void printCodes(struct Node* root, int code[], int top,
               int *totalBits, int *totalFreq) {

    if(root->left) {
        code[top] = 0;
        printCodes(root->left, code, top+1, totalBits, totalFreq);
    }
    if(root->right) {
        code[top] = 1;
        printCodes(root->right, code, top+1, totalBits, totalFreq);
    }
}
```

```

if(!root->left && !root->right) {
    printf("%c : ", root->data);
    for(int i = 0; i < top; i++)
        printf("%d", code[i]);
    printf(" (freq=%d, length=%d)\n", root->freq, top);
    *totalBits += root->freq * top;
    *totalFreq += root->freq;
}
}

int main() {
    char text[] = "DATA ANALYTICS AND INTELLIGENCE LABORATORY";
    int freq[256] = {0};
    for(int i = 0; text[i]; i++) {
        if(text[i] != ' ')
            freq[(int)text[i]]++;
    }
    struct Node* nodes[MAX];
    int n = 0;
    for(int i = 0; i < 256; i++) {
        if(freq[i] > 0) {
            nodes[n++] = createNode((char)i, freq[i]);
        }
    }
    while(n > 1) {
        sort(nodes, n);
        struct Node* left = nodes[0];
        struct Node* right = nodes[1];
        struct Node* newNode = createNode('$',
                                         left->freq + right->freq);
        newNode->left = left;
        newNode->right = right;
        nodes[0] = newNode;
        nodes[1] = nodes[n-1];
        n--;
    }
    struct Node* root = nodes[0];
    int code[100], totalBits = 0, totalFreq = 0;
    printf("Huffman Codes:\n\n");
    printCodes(root, code, 0, &totalBits, &totalFreq);
    printf("\nTotal Bits used: %d\n", totalBits);
}

```

```
    printf("\nTotal Compressed Bits = %d\n", totalBits);
    float avg = (float)totalBits / totalFreq;
    printf("Average Code Length = %.2f bits\n", avg);
    return 0;
}
```

Output:

```
Huffman Codes:
R : 0000 (freq=2, length=4)
D : 0001 (freq=2, length=4)
C : 0010 (freq=2, length=4)
O : 0011 (freq=2, length=4)
L : 010 (freq=4, length=3)
T : 011 (freq=4, length=3)
N : 100 (freq=4, length=3)
Y : 1010 (freq=2, length=4)
S : 10110 (freq=1, length=5)
B : 101110 (freq=1, length=6)
G : 101111 (freq=1, length=6)
E : 1100 (freq=3, length=4)
I : 1101 (freq=3, length=4)
A : 111 (freq=7, length=3)

Total Compressed Bits = 138
Average Code Length = 3.63 bits
```

Working:

Huffman coding :-

DATA ANALYTICS AND INTELLIGENCE LABORATORY

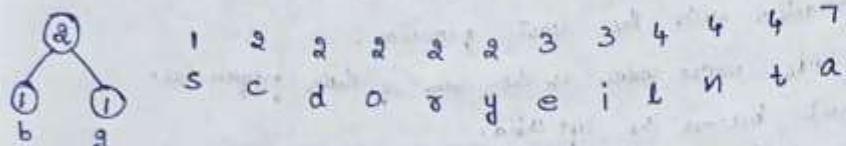
Algorithm:-

- 1) Write all the characters along with their frequencies in a table.
- 2) Arrange the characters in ascending order of frequency.
- 3) Select the two characters with the least frequency.
Create a root node whose value is the sum of their frequencies.
→ The first element becomes the left child.
→ The second element becomes the right child.
- 4) If two elements have the same frequencies, follow these rules:
→ character vs character → arrange in alphabetical order
→ character vs Tree → character comes first.
→ Tree vs Tree → choose the tree formed earlier.
- 5) Insert the new node back into the list and again arrange the list in ascending order.
- 6) Repeat this process until only one final tree is obtained.
- 7) Assign codes to the branches:
Left child → 0
Right child → 1
- 8) Traverse the tree and write the Binary code for each character.
- 9) Compute the no. of compressed bits using
$$\text{compressed bits} = \sum (\text{code length} \times \text{frequency})$$
- 10) Average HC length per character $= \frac{\sum (\text{code length} \times \text{frequency}_i)}{\sum (\text{frequency}_i)}$

Character	d	a	t	n	l	y	i	c	s	e	g	b	o	'
frequency	2	7	4	4	4	2	3	2	1	3	1	1	2	2

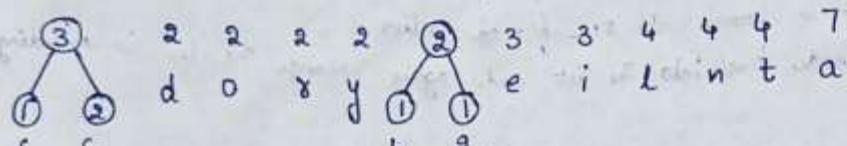
→ Ascending Order

1	1	1	2	2	2	2	3	3	4	4	4	7
b	g	s	c	d	o	y	e	i	l	n	t	a



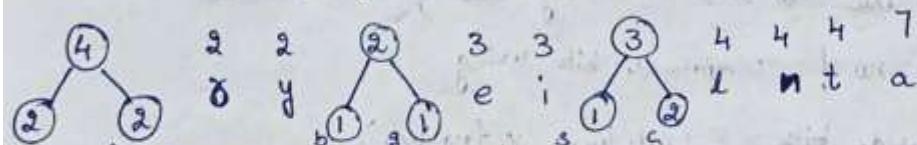
Arrange in Ascending order character first then tree.

1	2	2	2	2	2	2	3	3	4	4	4	7	
s	c	d	o	y	r	b	g	e	i	l	n	t	a



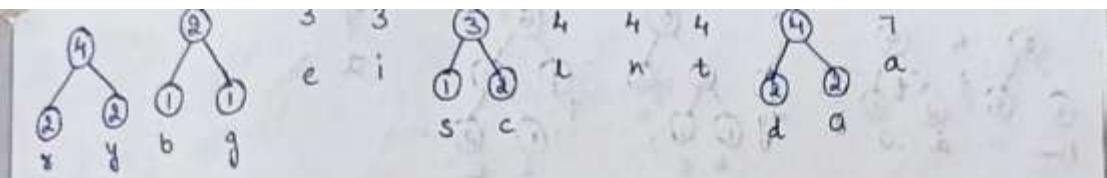
Arrange in Ascending order

2	2	2	2	2	2	3	3	3	4	4	4	7
d	o	r	y	b	g	e	i	s	l	n	t	a

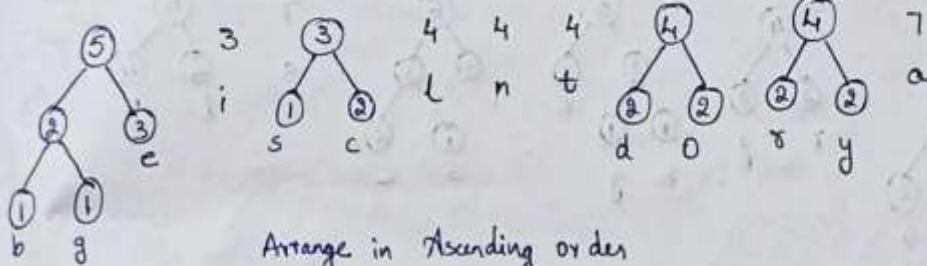
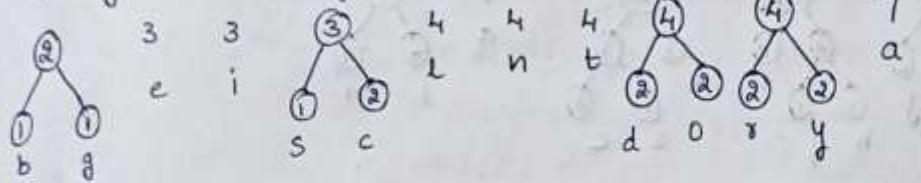


Arrange in Ascending order

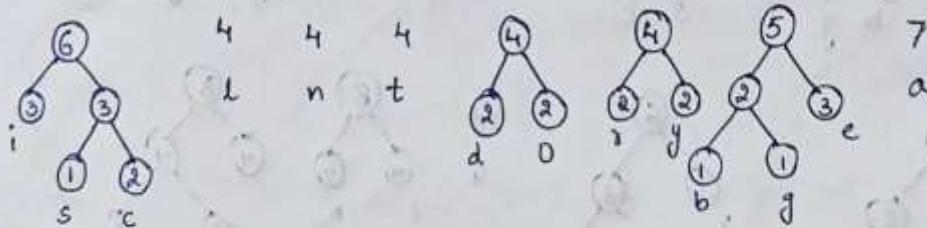
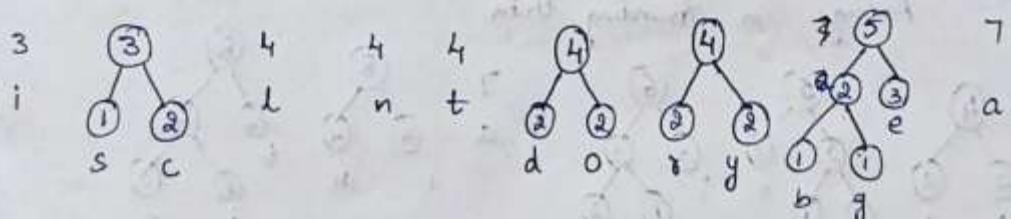
2	2	2	2	2	2	3	3	3	4	4	4	7
r	y	b	g	e	i	s	c	l	n	t	d	o



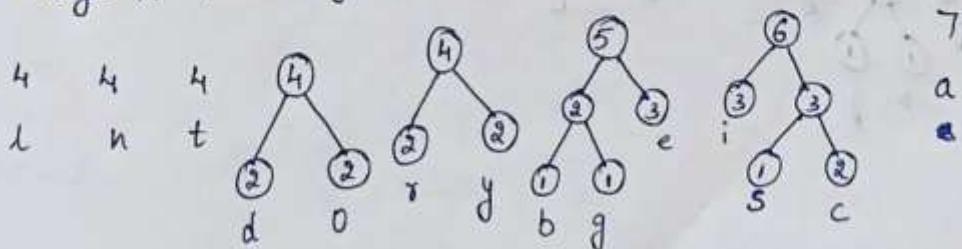
Arrange in Ascending order

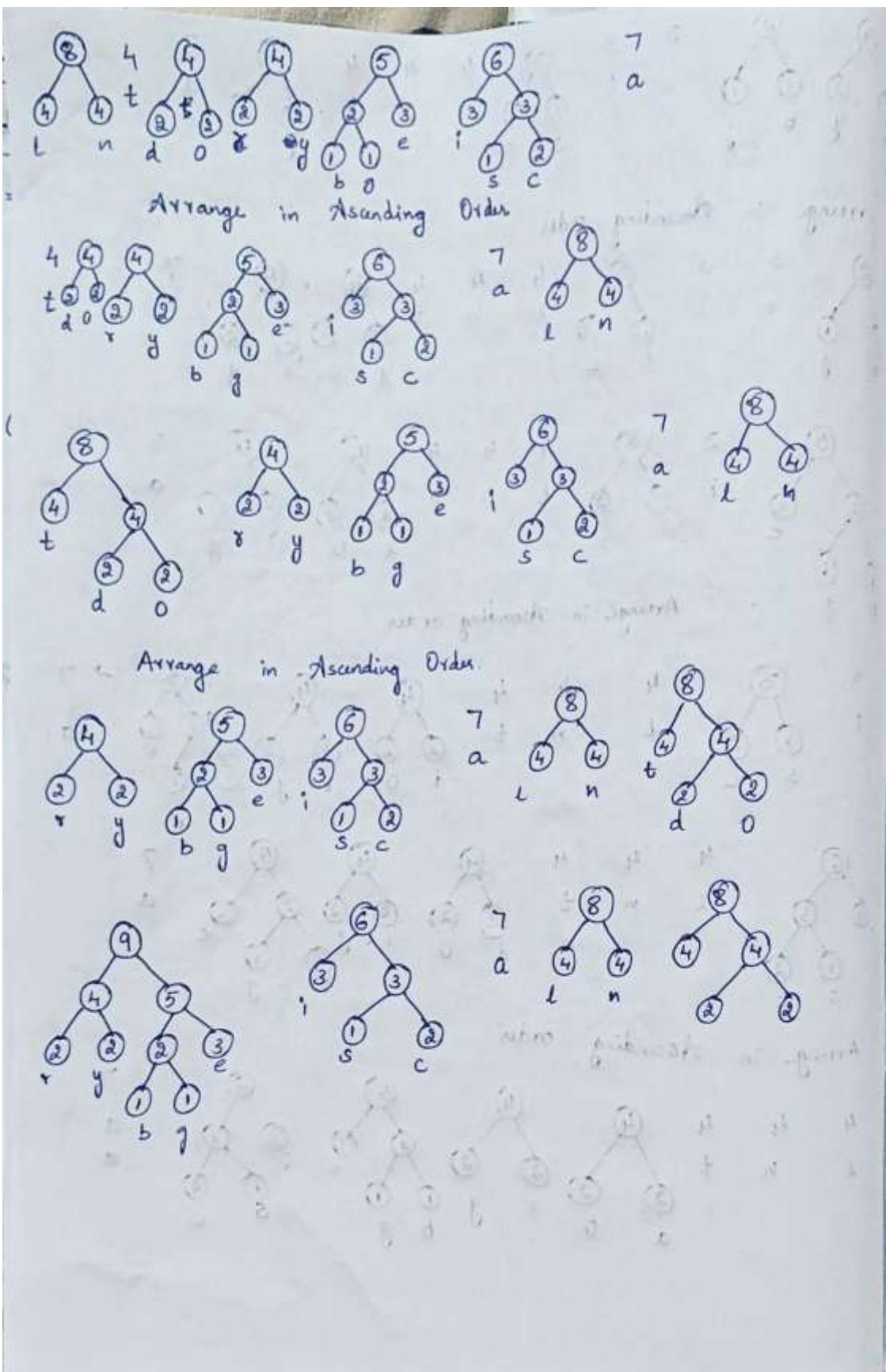


Arrange in Ascending order



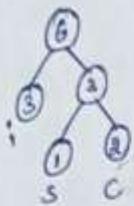
Arrange in Ascending order





Arrange

in descending order



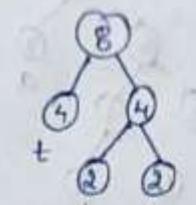
7

a



L

n

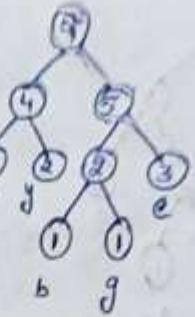


t

l

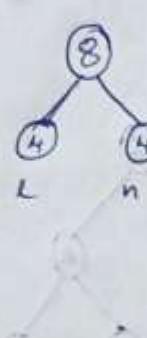
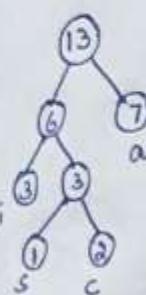
d

o



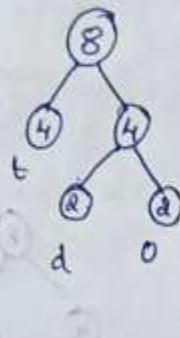
b

j



L

n

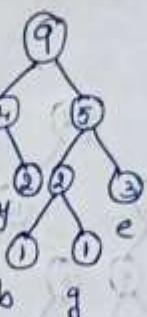


t

l

d

o

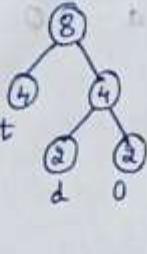
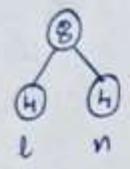


b

j

Arrange

in Ascending order

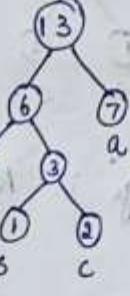


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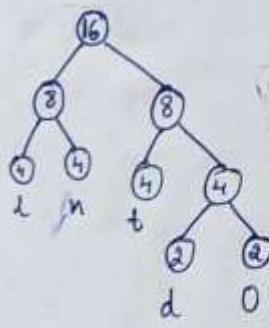
d

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s

c

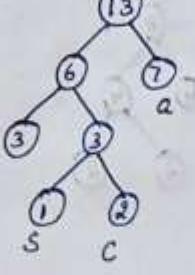


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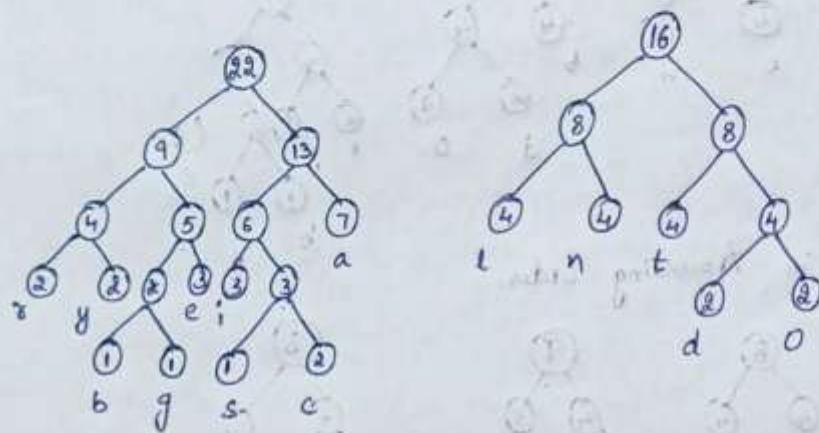
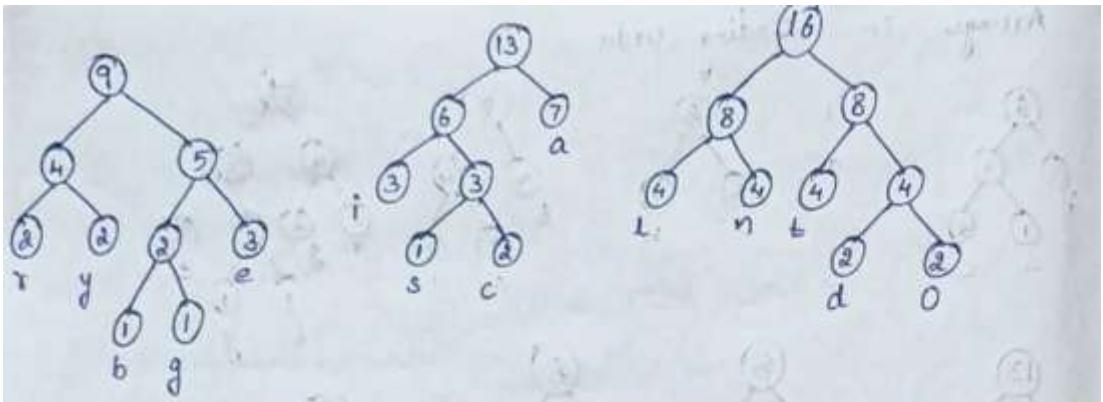


s

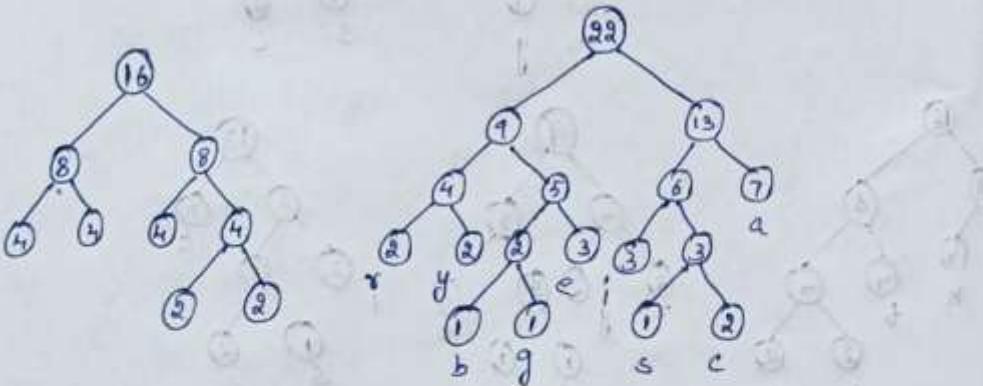
c

Arrange

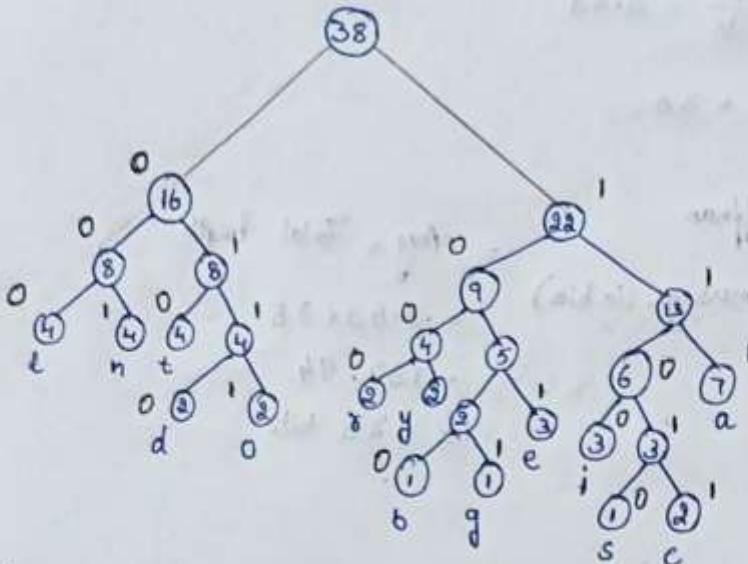
in Ascending order.



Arrange in Ascending order:



Final Tree



$$b - 10100 = 5 \times 1 = 5$$

$$g - 10101 = 5 \times 1 = 5$$

$$s - 11010 = 5 \times 1 = 5$$

$$c - 11011 = 5 \times 2 = 10$$

$$d - 0110 = 4 \times 2 = 8$$

$$o - 0111 = 4 \times 2 = 8$$

$$y - 1000 = 4 \times 2 = 8$$

$$j - 1001 = 4 \times 2 = 8$$

$$e - 1011 = 4 \times 3 = 12$$

$$i - 1100 = 4 \times 3 = 12$$

$$l - 000 = 3 \times 4 = 12$$

$$n - 001 = 3 \times 4 = 12$$

$$t - 010 = 3 \times 4 = 12$$

$$a - 111 = 3 \times 7 = 21$$

$$\text{Avg} = \frac{5+5+5+10+8+8+8+8+12+12+12+12+12+12+21}{1+1+1+2+2+2+2+3+3+4+4+4+7}$$

$$= \frac{138}{38} = 3.63$$

Total length = 38

Length of Hoffman

encoded message (in bits)

= Avg x Total length

$$= 3.63 \times 38$$

$$= 137.94$$

$$\approx 138 \text{ bits}$$

C = 1x0 = 00101 - d

E = 1x0 = 10101 - l

G = 1x0 = 01011 - s

H = 1x0 = 11011 - z

B = 2x0 = 01100 - r

S = 2x0 = 11100 - v

A = 3x0 = 00011 - n

F = 3x0 = 10011 - t

D = 3x0 = 11011 - g

I = 4x0 = 00111 - u

M = 4x0 = 00001 - p

L = 5x0 = 10001 - q

K = 5x0 = 01101 - w

J = 5x0 = 11111 - x



Time Complexity:

The algorithm repeatedly sorts the nodes in ascending order and merges the two smallest nodes.

Since Bubble Sort is used inside a loop, sorting is done multiple times.

- Best / Average Case = $O(n^3)$
- Worst Case = $O(n^3)$

Space Complexity:

Space is required for storing the Huffman tree and node list.

Recursion is used to generate codes.

- Average Case = $O(n)$
- Worst Case = $O(n)$