

Name: Bhavesh Kewalramani
Roll No.: A-25

Practical No. 3

Theory

LL(1)

Parsing:

Here the 1st L represents that the scanning of the Input will be done from Left to Right manner and the second L shows that in this parsing technique we are going to use Left most Derivation Tree. And finally, the 1 represents the number of look-ahead, which means how many symbols are you going to see when you want to make a decision.

Algorithm to construct LL(1) Parsing Table:

Step 1: First check for left recursion in the grammar, if there is left recursion in the grammar remove that and go to step 2.

Step 2: Calculate First() and Follow() for all non-terminals.

1. **First():** If there is a variable, and from that variable, if we try to drive all the strings then the beginning Terminal Symbol is called the First.
2. **Follow():** What is the Terminal Symbol which follows a variable in the process of derivation.

Step 3: For each production $A \rightarrow \alpha$. (A tends to alpha)

1. Find First(α) and for each terminal in First(α), make entry $A \rightarrow \alpha$ in the table.
 2. If First(α) contains ϵ (epsilon) as terminal than, find the Follow(A) and for each terminal in Follow(A), make entry $A \rightarrow \alpha$ in the table.
 3. If the First(α) contains ϵ and Follow(A) contains \$ as terminal, then make entry $A \rightarrow \alpha$ in the table for the \$.
- To construct the parsing table, we have two functions:

In the table, rows will contain the Non-Terminals and the column will contain the Terminal Symbols. All the **Null Productions** of the Grammars will go under the Follow elements and the remaining productions will lie under the elements of the First set.

Practicals

Aim:

- A. Write a program to find FIRST for any grammar. All the following rules of FIRST must be implemented.
- B. Further, write a program to find Follow for the given grammar.
- C. Construct the LL(1) parsing table using the FIRST and FOLLOW values computed above.

Program:

```

from collections import OrderedDict

# ` symbol represents EPSILON

def EnterGrammar():
    d = OrderedDict()
    f = open("grammar.txt", "r")

    print('Grammar\n'+f.read())

    f.seek(0)
    for line in f:
        k = ""
        for c in line:
            if c != "~" and k == "":
                d[c] = []
                k = c
            elif c != "~" and c != "\n":
                d[k].append(c)

    f.seek(0)
    nonterminal = []
    terminal = []
    for line in f:
        for c in line:
            if c not in d.keys() and c != "~" and c != "\n" and c != "" and c != "/" and c not in terminal:
                terminal.append(c)
            if c >= 'A' and c <= 'Z' and c not in nonterminal:
                nonterminal.append(c)

    return d, nonterminal, terminal

def Calculate_First(rule, index):
    first = ""
    v = rule[index]
    j = 1
    for i in range(len(v)):
        if v[i] == "/":
            j = 1
        elif j == 1:
            if v[i] not in rule.keys():
                if v[i] not in first and v[i] != "/":
                    first = first + v[i]

```

```

        j = 0
    else:
        a = list(Calculate_First(rule, v[i]))
        while "" in a and i+1 < len(v) and v[i+1] != "/":
            a.remove("")
            if v[i+1] not in rule.keys():
                a.append(v[i+1])
            else:
                a = list(set().union(a, Calculate_First(rule, v[i+1])))
                i += 1
        a.extend(first)
        first = "".join(list(set(a)))
        j = 0
    return first

```

```

def Calculate_Follow(rule, n, start_symbol):
    follow = ""
    if n == start_symbol:
        follow += "$"
    for k, v in rule.items():
        for i in range(len(v)):
            if v[i] == n:
                if i == len(v) - 1:
                    follow += Calculate_Follow(rule, k, start_symbol)
                elif i + 1 < len(v) and v[i + 1] not in rule.keys() and v[i+1] != "/" and v[i+1] not in follow:
                    follow += v[i + 1]
            elif i + 1 < len(v) and v[i+1] != "/" and v[i+1] not in follow:
                a = []
                for j in first[v[i + 1]]:
                    a.append(j)
                if "" in a:
                    a.remove("")
                    a.append(Calculate_Follow(d, v[i+1], start_symbol))
                follow += "".join(list(set("".join(a))))
            elif k == start_symbol:
                follow += "$"
    return follow

```

```

def parsingTable(rule, nonterminals, terminals, first, follow):
    terminals.append('$')
    #make table
    parse_table = [ ["_"]*(len(terminals) + 1) for i in range(len(nonterminals) + 1) ]
    for i in range(len(parse_table)):

```

```

for j in range(len(parse_table[0])):
    if i == 0 and j != 0:
        parse_table[i][j] = terminals[j-1]
    if i != 0 and j == 0:
        parse_table[i][j] = nonterminals[i-1]

#fill table
for i in range(1,len(parse_table)):
    for j in range(1,len(parse_table[0])):
        if parse_table[0][j] in first[parse_table[i][0]]:
            key = parse_table[i][0]
            for k,v in rule.items():
                if k == key:
                    val = v

            c = 0
            flag = 0
            while c < len(val) and flag == 0:
                if val[c] in nonterminals:
                    if parse_table[0][j] in first[val[c]]:
                        rhs = ""
                        for k in range(c, len(val)):
                            if val[k] == '/':
                                break
                            rhs += val[k]

                        ans = parse_table[i][0] + '~' + rhs
                        parse_table[i][j] = ans
                        flag = 1

                elif val[c] == parse_table[0][j]:
                    rhs = ""
                    for k in range(c,len(val)):
                        if val[k] == '/':
                            break
                        rhs += val[k]

                    ans = parse_table[i][0] + '~' + rhs
                    parse_table[i][j] = ans
                    flag = 1

            else:
                while val[c] != '/' and c < len(val):
                    c += 1
                if c < len(val):
                    c += 1

```

```

        elif parse_table[0][j] in follow[parse_table[i][0]] and '' in first[parse_table[i][0]]:
            parse_table[i][j] = parse_table[i][0] + '~'

        elif parse_table[0][j] in follow[parse_table[i][0]]:
            parse_table[i][j] = '_'

        else:
            pass

    return parse_table

rule, nonterminals, terminals = EnterGrammar()
start_symbol = input("Enter the start symbol : ")
print()
print("Dictionary:", rule)
print()


first = OrderedDict()
for k, v in rule.items():
    first[k] = []
    first[k].extend(Calculate_First(rule, k))
print("First:", first)
print()

follow = OrderedDict()
for k, v in rule.items():
    follow[k] = []
    follow[k].extend(Calculate_Follow(rule, k, start_symbol))
    follow[k] = list(set(follow[k]))
print("Follow:", follow)
print()

parse_table = parsingTable(rule, nonterminals, terminals, first, follow)
print("Parsing Table")
for i in range(len(parse_table)):
    for j in range(len(parse_table[0])):
        print(parse_table[i][j], end="\t\t")
    print()

```

Input:

 jupyter grammar.txt ✓ a few seconds ago

FileEditViewLanguage

1 A~SB/B

2 S~a/Bc/`

3 B~b/d

Output:

```
print()
```

Grammar
A~SB/B
S~a/Bc/`
B~b/d
Enter the start symbol : A

Dictionary: OrderedDict([(('A', ['S', 'B', '/', 'B']), ('S', ['a', '/', 'B', 'c', '/', '']), ('B', ['b', '/', 'd']))])

First: OrderedDict([(('A', ['a', 'd', 'b']), ('S', ['a', 'd', 'b', '']), ('B', ['b', 'd']))])

Follow: OrderedDict([(('A', ['\$']), ('S', ['d', 'b']), ('B', ['\$', 'c']))])

Parsing Table

	a	c	b	d	\$
A	A~SB	—	A~SB	A~SB	—
S	S~a	—	S~Bc	S~Bc	—
B	—	—	B~b	B~d	—

In []:

Name - Bhavesh Kewalramani

Roll No. - A-25

Section - A

Semester - VI

Shift - I

Batch - A1

$A \rightarrow SB|B$

$S \rightarrow a|Bc|E$

$B \rightarrow b|d$

First -

$\text{First}(a) = \{a\}$

$\text{first}(b) = \{b\}$

$\text{first}(c) = \{c\}$

$\text{first}(d) = \{d\}$

$\text{first}(E) = \{E\}$

$\text{first}(B) = \text{first}(b) \cup \text{first}(d) = \{b\} \cup \{d\} = \{b, d\}$

$$\begin{aligned} \text{first}(S) &= \text{first}(a) \cup \text{first}(Bc) \cup \text{first}(E) \\ &= \{a\} \cup \text{first}(B) \cup \{E\} \\ &= \{a\} \cup \{b, d\} \cup \{E\} \\ &= \{a, b, d, E\} \end{aligned}$$

$$\begin{aligned} \text{first}(A) &= \text{first}(SB) \cup \text{first}(B) \\ &= \text{first}(S) - \{E\} \cup \text{first}(B) \cup \text{first}(B) \\ &= \{a, b, d, E\} - \{E\} \cup \{b, d\} \cup \{b, d\} \\ &= \{a, b, d\} \cup \{b, d\} \cup \{b, d\} \\ &= \{a, b, d\} \end{aligned}$$

Follow -

$$\text{follow}(A) = \{\$ \}$$

$$\begin{aligned}\text{follow}(S) &= \text{first}(B) \\ &= \{b, d\}\end{aligned}$$

$$\begin{aligned}\text{follow}(B) &= \text{first}(\epsilon) - \{ \epsilon \} \cup \text{follow}(A) \cup \text{first}(C) \\ &= \{ \epsilon \} - \{ \epsilon \} \cup \{ \$ \} \cup \{ c \} \\ &= \{ c, \$ \}\end{aligned}$$

Parsing Table -

	a	b	c	d	\$
A	$A \rightarrow SB$	$A \rightarrow SB B$		$A \rightarrow SB B$	
S	$S \rightarrow a$	$S \rightarrow Bc E$		$S \rightarrow Bc E$	
B		$B \rightarrow b$		$B \rightarrow d$	