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Practical No. 4

Theory

LL(1)

Parsing:

Here the 1st **L** represents that the scanning of the Input will be done from Left to Right manner and the second **L** shows that in this parsing technique we are going to use Left most Derivation Tree. And finally, the **1** represents the number of look-ahead, which means how many symbols are you going to see when you want to make a decision.

Algorithm to construct LL(1) Parsing Table:

Step 1: First check for left recursion in the grammar, if there is left recursion in the grammar remove that and go to step 2.

Step 2: Calculate First() and Follow() for all non-terminals.

1. **First():** If there is a variable, and from that variable, if we try to drive all the strings then the beginning Terminal Symbol is called the First.
2. **Follow():** What is the Terminal Symbol which follows a variable in the process of derivation.

Step 3: For each production $A \rightarrow \alpha$. (A tends to alpha)

1. Find First(α) and for each terminal in First(α), make entry $A \rightarrow \alpha$ in the table.
 2. If First(α) contains ϵ (epsilon) as terminal than, find the Follow(A) and for each terminal in Follow(A), make entry $A \rightarrow \alpha$ in the table.
 3. If the First(α) contains ϵ and Follow(A) contains \$ as terminal, then make entry $A \rightarrow \alpha$ in the table for the \$.
- To construct the parsing table, we have two functions:

In the table, rows will contain the Non-Terminals and the column will contain the Terminal Symbols. All the **Null Productions** of the Grammars will go under the Follow elements and the remaining productions will lie under the elements of the First set.

Practicals

Aim:

- (A) Write a program to validate a natural language sentence. Design a natural language grammar, compute and input the LL(1) table. Validate if the given sentence is valid based on the grammar or not.
- (B) Use Virtual Lab on LL1 parser to validate the string and verify your string validation using simulation.

Program:

```
from collections import OrderedDict
import re
```

```
# ` represents EPSILON
```

```
def getGrammar():
    terminal=[]
    nonterminal=[]
    start=""
    rule=dict()
    flag=0
    print("Enter the Production Rules : ")
    while(True):
        inp=input("====> ")
        if(inp==""):
            break
        s1,s2=inp.split("~")
        if flag==0:
            start=s1
            flag=1
        rule[s1]=[]
        s2=list(s2.split("/"))
        for i in s2:
            rule[s1].append(i)

    r=[]
    k=rule.keys()
    for i in rule.values():
        for j in i:
            for a in list(j.split(" ")):
                if a not in k:
                    if a not in r:
                        r.append(a)

    r.append("$")
    t=dict()
    for i in rule.keys():
        t[i]=dict()
        for j in r:
            t[i].update( {j:set()} )

    return rule,start,r,t

def Calculate_First(s,rule,v):
    if s[0] in v:
        return set([s[0]])
```

```

elif s[0] == '':
    return set([""])
else:
    res = set()
    for j in rule[s[0]]:
        h = Calculate_First(list(j.split(' ')),rule,v)
        res.update(set(h))
    if len(s) == 1:
        return res
    else:
        if '' in res:
            res.remove('')
        return res.union(Calculate_First(list(s.split(' '))[1:],rule,v))
    return res

```

```

def getValue(v,rule):
    for key, value in rule.items():
        if value == v:
            return key

```

```

def Calculate_Follow(s,rule,v,start):
    res = set()
    if s == start:
        res = set(['$'])
    for i in rule.values():
        for j in range(len(i)):
            l = list(i[j].split(' '))
            for k in range(len(l)):
                if l[k] == s:
                    if k == len(l) - 1:
                        if getValue(i,rule) == s:
                            continue
                    else:
                        res.update(set(Calculate_Follow(getValue(i,rule),rule,v,start)))
            else:
                c = set(Calculate_First(l[k + 1:],rule,v))
                if '' in c:
                    c.remove('')
                fol = Calculate_Follow(getValue(i,rule),rule,v,start)
                c.update(fol)
                res.update(c)
            else:
                res.update(set(c))
    return res

```

```

def parseString(string,t):
    print("Given String is : ", string)
    string = list(string.split(' '))
    string.append('$')
    stk = ['$', start]
    print("Input                                     Stack")
    print(string[::-1] , "\t\t\t\t\t", stk[::-1])
    while not len(stk) == 0:
        top = stk[-1]
        stk.pop()
        if string[0] == '$' and top == '$' and len(stk) == 0:
            print("***** String Accepted *****")
        elif (string[0] == '$' and len(stk) != 0) or (string[0] != '$' and len(stk) == 0):
            print("***** String Rejected *****")
        elif top == string[0]:
            string = string[1:]
        else:
            for i in t[top][string[0]]:
                l = list(i.split(' '))
                for j in l[::-1]:
                    stk.append(j)
            print(stk[::-1] , "\t\t\t\t\t", stk[::-1])

if __name__=="__main__":

    d,start,r,t=getGrammar()

    print("===== FIRST =====")
    for i in d.keys():
        print("first(", i, ") : ", Calculate_First([i],d,r))
    print("===== FOLLOW =====")
    for i in d.keys():
        print("follow(", i, ") : ", Calculate_Follow(i,d,r,start))
    print("===== Parsing Table =====")

    for i in d.keys():
        for rule in d[i]:
            f = Calculate_First(list(rule.split(' ')),d,r)
            if '' in f:
                fol = Calculate_Follow(i,d,r,start)
                for j in fol:
                    if j != '$':
                        t[i][j].add(rule)
                if '$' in fol:

```

```

        t[i]['$'].add(rule)
    else:
        for j in f:
            t[i][j].add(rule)
print("\t\t", end=" ")
for i in r:
    print(i, end="\t\t\t\t")
print()
print("===== String Parsing =====")

for j in r:
    print("Terminal : ", j)
    for i in d.keys():
        if len(t[i][j]) != 0:
            print(i, "->", t[i][j])

print()

string = "India won the championship"
parseString(string,t)

```

Input:

```

Enter the Production Rules :
===> S~NP VP
===> NP~P/PN/D N
===> VP~V NP
===> N~championship/ball/toss
===> V~is/want/won/played
===> P~me/I/you
===> PN~India/Australia/Steve/John
===> D~a/an/the
===>
-----

```

Output:

```

===== FIRST =====
first( S ) : {'the', 'you', 'John', 'Steve', 'me', 'a', 'an', 'India', 'Australia', 'I'}
first( NP ) : {'the', 'you', 'John', 'Steve', 'me', 'a', 'an', 'India', 'Australia', 'I'}
first( VP ) : {'want', 'is', 'won', 'played'}
first( N ) : {'championship', 'toss', 'ball'}
first( V ) : {'want', 'is', 'won', 'played'}
first( P ) : {'me', 'I', 'you'}
first( PN ) : {'Steve', 'John', 'India', 'Australia'}
first( D ) : {'a', 'an', 'the'}
===== FOLLOW =====
follow( S ) : {'$'}
follow( NP ) : {'want', '$', 'won', 'is', 'played'}
follow( VP ) : {'$'}
follow( N ) : {'want', '$', 'won', 'is', 'played'}
follow( V ) : {'the', 'you', 'John', 'Steve', 'me', 'a', 'an', 'India', 'Australia', 'I'}
follow( P ) : {'want', '$', 'won', 'is', 'played'}
follow( PN ) : {'want', '$', 'won', 'is', 'played'}
follow( D ) : {'championship', 'toss', 'ball'}

follow( V ) : {'championship', 'toss', 'ball'}
===== Parsing Table =====
                                championship          ball          toss          is
want                                won                    played          me
I                                  you                    India          Australia
Steve                             John                    a              an
the                                $

===== String Parsing =====
Terminal : championship
N -> {'championship'}

Terminal : ball
N -> {'ball'}

Terminal : toss
N -> {'toss'}

```

```

Terminal : is
VP -> {'V NP'}
V -> {'is'}

Terminal : want
VP -> {'V NP'}
V -> {'want'}

Terminal : won
VP -> {'V NP'}
V -> {'won'}

Terminal : played
VP -> {'V NP'}
V -> {'played'}

```

```

Terminal : me
S -> {'NP VP'}
NP -> {'P'}
P -> {'me'}

Terminal : I
S -> {'NP VP'}
NP -> {'P'}
P -> {'I'}

Terminal : you
S -> {'NP VP'}
NP -> {'P'}
P -> {'you'}

```

Terminal : India
S -> {'NP VP'}
NP -> {'PN'}
PN -> {'India'}

Terminal : Australia
S -> {'NP VP'}
NP -> {'PN'}
PN -> {'Australia'}

Terminal : Steve
S -> {'NP VP'}
NP -> {'PN'}
PN -> {'Steve'}

Terminal : John
S -> {'NP VP'}
NP -> {'PN'}
PN -> {'John'}

Terminal : a
S -> {'NP VP'}
NP -> {'D N'}
D -> {'a'}

Terminal : an
S -> {'NP VP'}
NP -> {'D N'}
D -> {'an'}

D -> {'an'}

Terminal : the
S -> {'NP VP'}
NP -> {'D N'}
D -> {'the'}

Terminal : \$

Terminal : \$

Given String is : India won the championship
Input
['\$', 'championship', 'the', 'won', 'India']
['NP', 'VP', '\$']
['PN', 'VP', '\$']
['India', 'VP', '\$']
['VP', '\$']
['V', 'NP', '\$']
['won', 'NP', '\$']
['NP', '\$']
['D', 'N', '\$']
['the', 'N', '\$']
['N', '\$']
['championship', '\$']
['\$']
***** String Accepted *****
[]

Stack
['NP', 'VP', '\$']
['PN', 'VP', '\$']
['India', 'VP', '\$']
['VP', '\$']
['V', 'NP', '\$']
['won', 'NP', '\$']
['NP', '\$']
['D', 'N', '\$']
['the', 'N', '\$']
['N', '\$']
['championship', '\$']
['\$']
[]

['S', '\$']

Virtual Lab Simulation:

1. Write your LL(1) grammar (empty string "" represents ϵ):

S ::= NP VP
NP ::= P
NP ::= PN
NP ::= D N
VP ::= V NP
N ::= championship
N ::= ball
N ::= toss
V ::= is
V ::= want
V ::= won
V ::= played
P ::= me
P ::= I
P ::= you

1. Write your LL(1) grammar (empty string "" represents ϵ):

N ::= toss
V ::= is
V ::= want
V ::= won
V ::= played
P ::= me
P ::= I
P ::= you
PN ::= India
PN ::= Australia
PN ::= Steve
PN ::= John
D ::= a
D ::= the
D ::= an

2. Nullable/First/Follow Table and Transition Table

Nonterminal	Nullable?	First	Follow
S	×	me, I, you, India, Australia, Steve, John, a, the, an	\$
NP	×	me, I, you, India, Australia, Steve, John, a, the, an	is, want, won, played, \$
VP	×	is, want, won, played	\$
N	×	championship, ball, toss	is, want, won, played, \$
V	×	is, want, won, played	me, I, you, India, Australia, Steve, John, a, the, an
P	×	me, I, you	is, want, won, played, \$
PN	×	India, Australia, Steve, John	is, want, won, played, \$
D	×	a, the, an	championship, ball, toss

	\$	championship	ball	toss	is	want	won	played	me	I	you	India	Australia	Steve	John	a	the	an
S									S ::= NP VP S ::= S \$	S ::= NP VP S ::= S \$	S ::= NP VP S ::= S \$	S ::= NP VP S ::= S \$	S ::= NP VP S ::= S \$	S ::= NP VP S ::= S \$	S ::= NP VP S ::= S \$	S ::= NP VP S ::= S \$	S ::= NP VP S ::= S \$	S ::= NP VP S ::= S \$
NP									NP ::= P	NP ::= P	NP ::= P	NP ::= PN	NP ::= PN	NP ::= PN	NP ::= PN	NP ::= D N	NP ::= D N	NP ::= D N
VP					VP ::= V NP	VP ::= V NP	VP ::= V NP	VP ::= V NP										
N		N ::= championship	N ::= ball	N ::= toss														
V					V ::= is	V ::= want	V ::= won	V ::= played										
P									P ::= me	P ::= I	P ::= you							
PN												PN ::= India	PN ::= Australia	PN ::= Steve	PN ::= John			
D																D ::= a	D ::= the	D ::= an

3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ S

Remaining Input

India won the championship \$

Rule

Partial Parse Tree

S

3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ VP NP

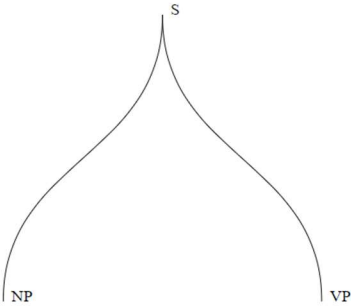
Remaining Input

India won the championship \$

Rule

S ::= NP VP

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ VP PN

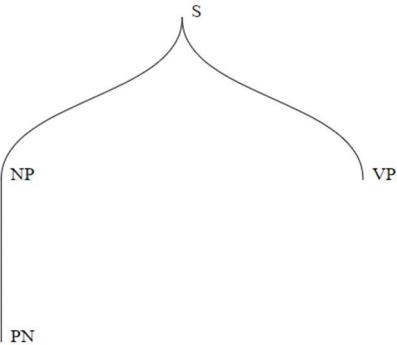
Remaining Input

India won the championship \$

Rule

NP ::= PN

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ VP India

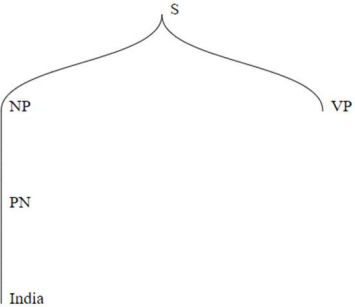
Remaining Input

India won the championship \$

Rule

PN ::= India

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ VP

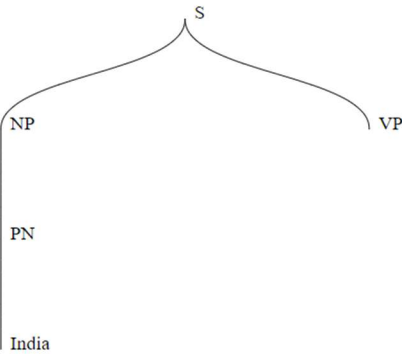
Remaining Input

won the championship \$

Rule

Match India

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ NP V

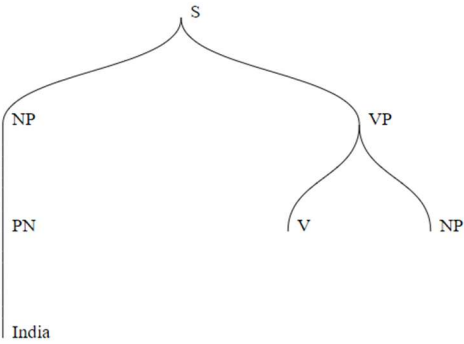
Remaining Input

won the championship \$

Rule

VP ::= V NP

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ NP won

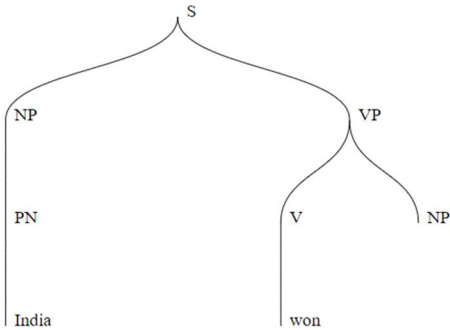
Remaining Input

won the championship \$

Rule

V ::= won

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ NP

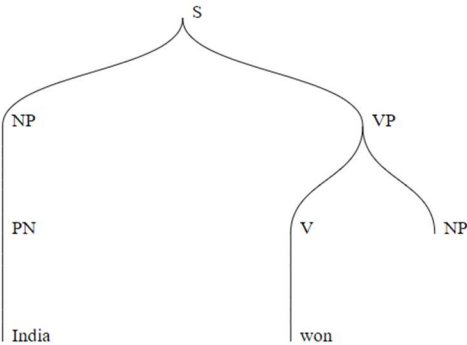
Remaining Input

the championship \$

Rule

Match won

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset Step Forward

Stack

\$ N D

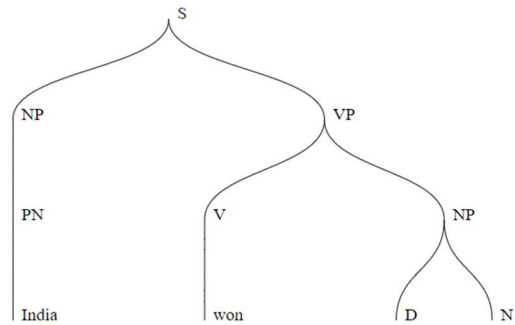
Remaining Input

the championship \$

Rule

NP ::= D N

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset Step Forward

Stack

\$ N the

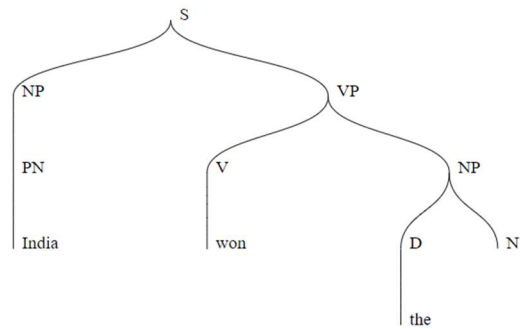
Remaining Input

the championship \$

Rule

D ::= the

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ N

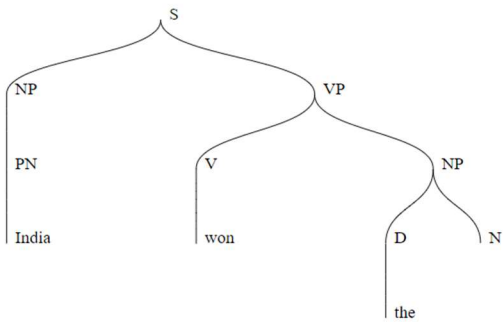
Remaining Input

championship \$

Rule

Match the

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

India won the championship

Start/Reset

Step Forward

Stack

\$ championship

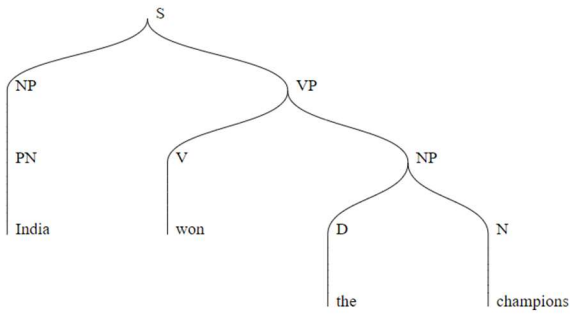
Remaining Input

championship \$

Rule

N ::= championship

Partial Parse Tree



3. Parsing

Token stream separated by spaces:

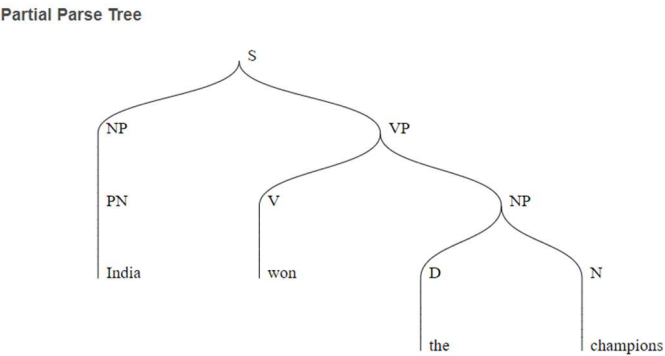
Start/Reset

Step Forward

Stack

Remaining Input

Rule



3. Parsing

Token stream separated by spaces:

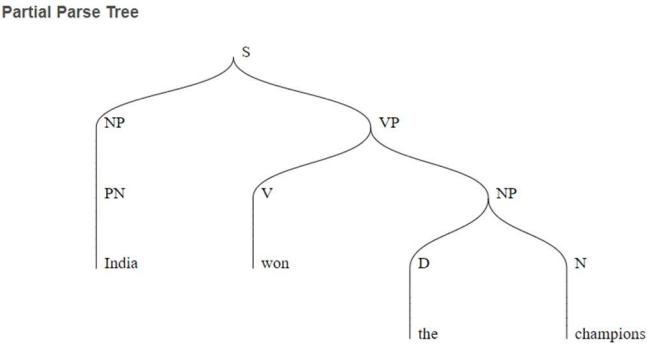
Start/Reset

Step Forward

Stack

Remaining Input

Rule



www.cs.princeton.edu says

Parsing complete! Press reset to see it again.

OK