

* Insertion sort

```

for (i = 2; i < n-1; i++)
{
    curr = arr[i];
    j = i - 1;
    while (arr[j] > curr && j >= 0)
    {
        arr[j+1] = arr[j];
        j--;
    }
    arr[j+1] = curr;
}

```

Insertion time complexity

Best = $O(n)$

avg worst = $O(n^2)$

Worst = $O(n^2)$

* Merge sort

time complexity in all cases $n \log n$

Merge_sort(A, p, r)

If $p < r$ Then

$q = (p+r)/2$

Merge_sort(A, p, q, r) A, p, q

Merge_sort(A, q+1, r)

Merge(A, p, q, r)

Merge (A, p, q, r)
Step 1:- Initialization

$M[r-p+1]$

$i1 \leftarrow p$

$i2 \leftarrow q+1$

$x \leftarrow 0$

Step 2:- Repeat step 3 while $i1 \leq q$ & $i2 \leq r$

Step 3:-

If $A[i1] \leq A[i2]$ Then

$M[x] \leftarrow A[i1];$

$x \leftarrow x+1$

else

$M[x] \leftarrow A[i2];$

$x \leftarrow x+1;$

Step 4:- Repeat step 5 while $i1 \leq q$

Step 5:- ~~merged~~ $M[x] \leftarrow A[i1]$

Step 6:- Repeat step 7 while $i2 \leq r$

Step 7:- $M[x] \leftarrow A[i2]$

Step 8:- [Copy element to original array]

Repeat: while
for $i \leftarrow 1$; $j \leftarrow p$ to r do

$A[j] = \text{Merged}[i]$
 $i \leftarrow i + 1$
 $j \leftarrow j + 1$

Tracing

2, 5, 7, 3, 9, 8

$\text{merge_sort}(A, 1, 6)$
 $q = 3$
 $MS(A, 1, 3)$
 $MS(A, 4, 6)$
 $\text{Merge}(A, 1, 3, 6)$

$MS(A, 1, 3)$
 $q = 2$
 $MS(A, 1, 2)$
 $MS(A, 3, 3)$
 $\text{Merge}(A, 1, 2, 3)$

$MS(A, 1, 2)$
 $q = 1$
 $MS(A, 1, 1) \rightarrow \text{False}$
 $MS(A, 2, 2) \rightarrow \text{False}$
 $\text{Merge}(A, 1, 1, 2)$

$MS(A, 4, 6)$
 $q = 5$
 $MS(A, 4, 5)$
 $MS(A, 6, 6)$
 $\text{Merge}(A, 4, 5, 6)$

$$g = 4$$

$ms(A, 4, 4) \leftarrow \text{False}$

$$ms(A, 5, 5) \leftarrow \text{False}$$

Merge (A, 4, 4, 5)

30, 10, 40, 60, 50

$$p=1, \quad r=5$$
 $q = 3$
$$MS(A, 4, 5)$$

Merge(A, 1, 3) 5

 $q = 2$

MS(A.3.3)

Merge(A, 1, 2,

 $q =$

ANS (A 2 2)

Mammals

Merge CA

30	10	40	60	50
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30	10
----	----

✓
150

30

110

9 - 4

MS (A E S)

MS. CR. 3.5.2

60

→ 50

also

150

Menge (A, 4, 4, 5)