Chapter 5 Network Layer: Control Plane

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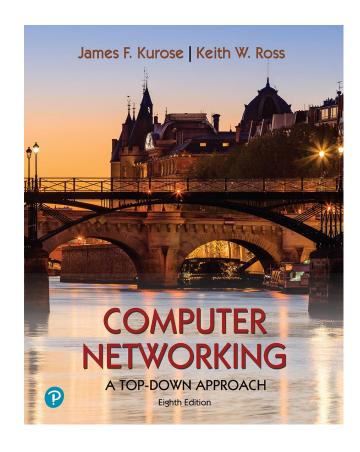
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Computer Networking: A Top-Down Approach

8th edition Jim Kurose, Keith Ross Pearson, 2020

Network layer: "control plane" roadmap

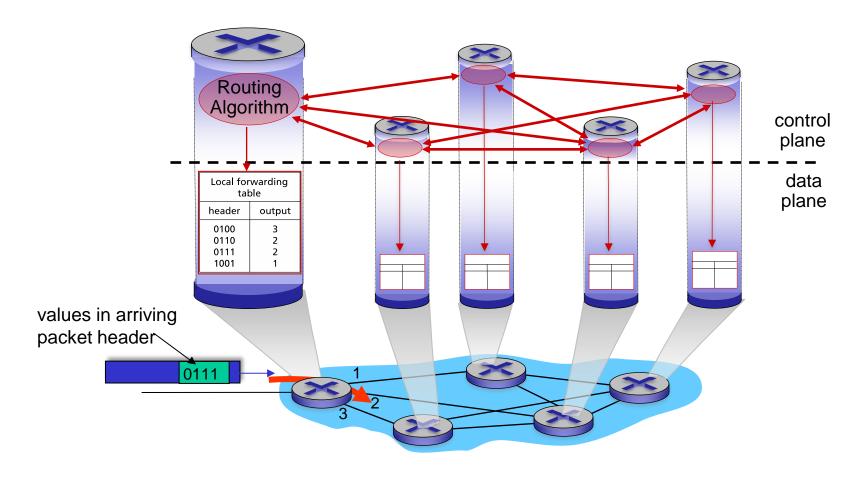
- introduction
- routing protocols
- intra-ISP routing: OSPF
- routing among ISPs: BGP
- SDN control plane
- Internet Control MessageProtocol



- Internet network layer: historically implemented via distributed, per-router control approach:
 - monolithic router contains switching hardware, runs proprietary implementation of Internet standard protocols (IP, RIP, IS-IS, OSPF, BGP) in proprietary router OS (e.g., Cisco IOS)
 - different "middleboxes" for different network layer functions: firewalls, load balancers, NAT boxes, ..
- ~2005: renewed interest in rethinking network control plane

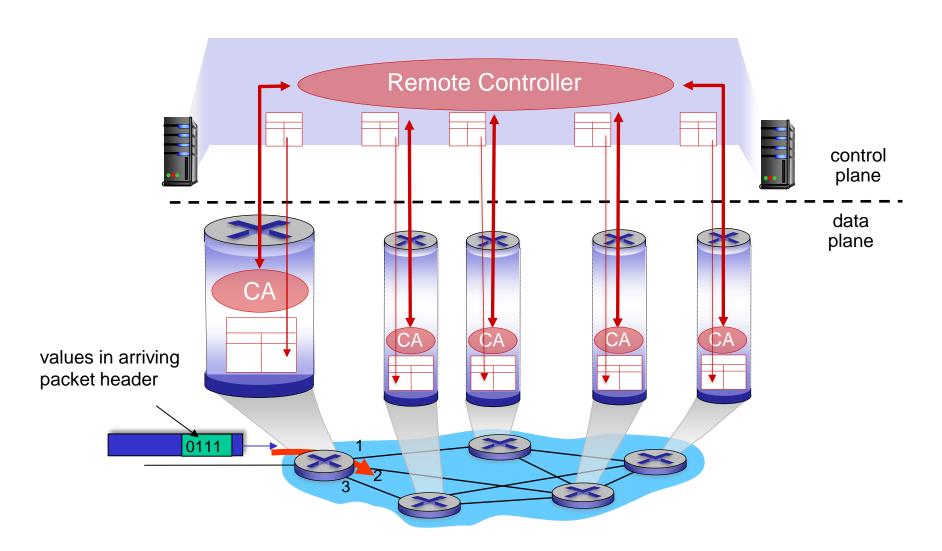
Per-router control plane

Individual routing algorithm components in each and every router interact in the control plane to computer forwarding tables



Software-Defined Networking (SDN) control plane

Remote controller computes, installs forwarding tables in routers



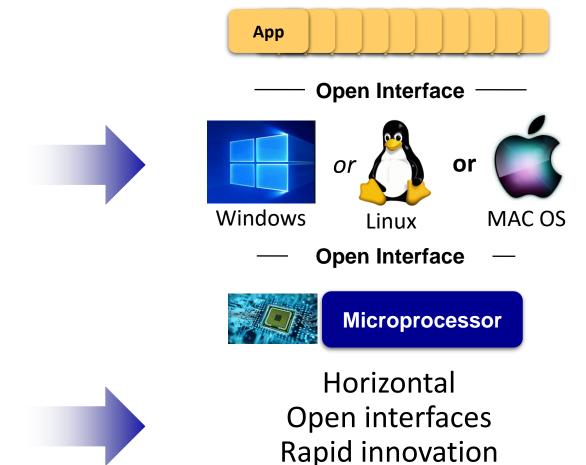
Why a logically centralized control plane?

- easier network management: avoid router misconfigurations, greater flexibility of traffic flows
- table-based forwarding (recall OpenFlow API) allows "programming" routers
 - centralized "programming" easier: compute tables centrally and distribute
 - distributed "programming" more difficult: compute tables as result of distributed algorithm (protocol) implemented in each-and-every router
- open (non-proprietary) implementation of control plane
 - foster innovation: let 1000 flowers bloom

SDN analogy: mainframe to PC revolution



Vertically integrated Closed, proprietary Slow innovation Small industry

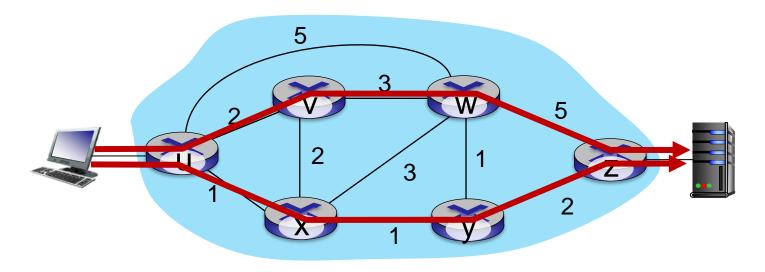


Huge industry

* Slide courtesy: N. McKeown

Network Layer: 5-7

Traffic engineering: difficult with traditional routing

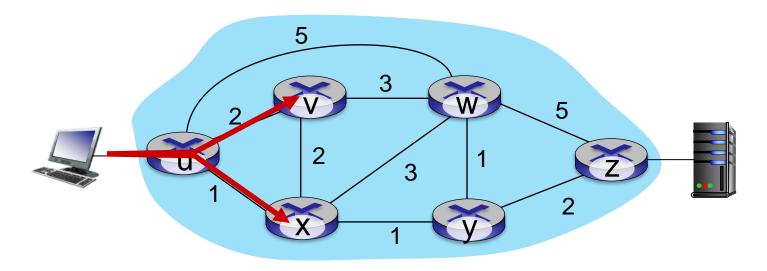


<u>Q:</u> what if network operator wants u-to-z traffic to flow along uvwz, rather than uxyz?

<u>A:</u> need to re-define link weights so traffic routing algorithm computes routes accordingly (or need a new routing algorithm)!

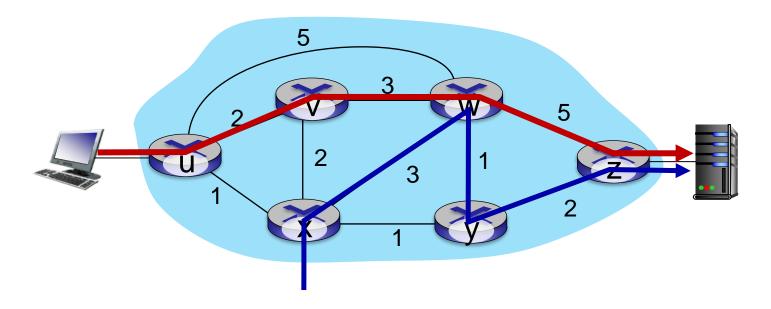
link weights are only control "knobs": not much control!

Traffic engineering: difficult with traditional routing



<u>Q:</u> what if network operator wants to split u-to-z traffic along uvwz <u>and</u> uxyz (load balancing)? <u>A:</u> can't do it (or need a new routing algorithm)

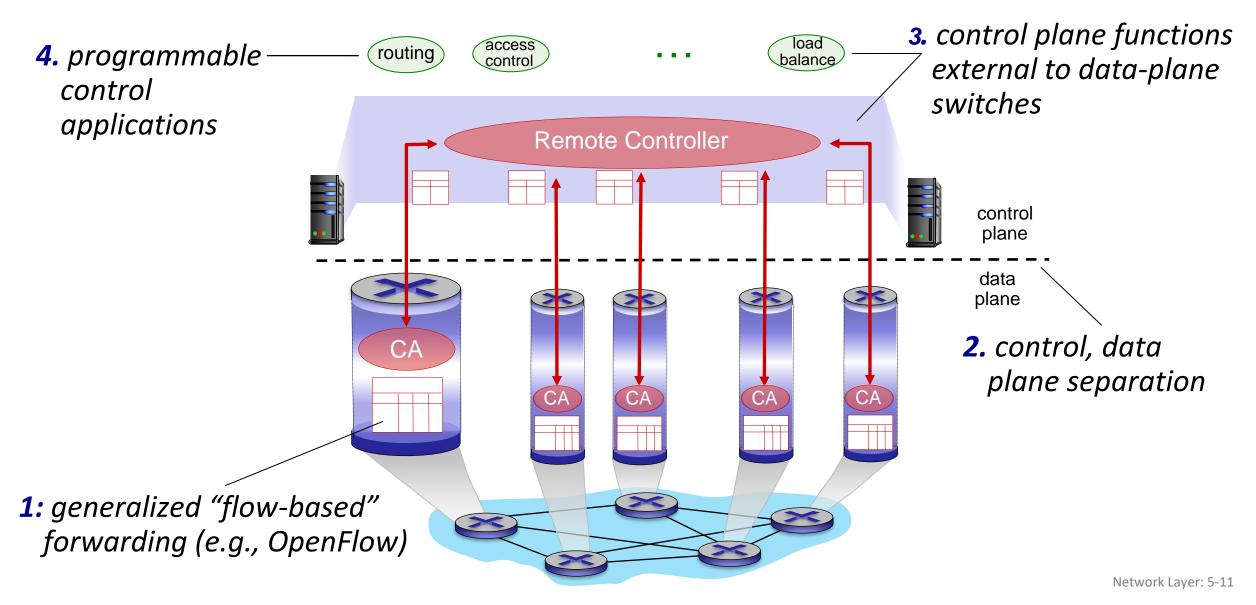
Traffic engineering: difficult with traditional routing



<u>Q:</u> what if w wants to route blue and red traffic differently from w to z?

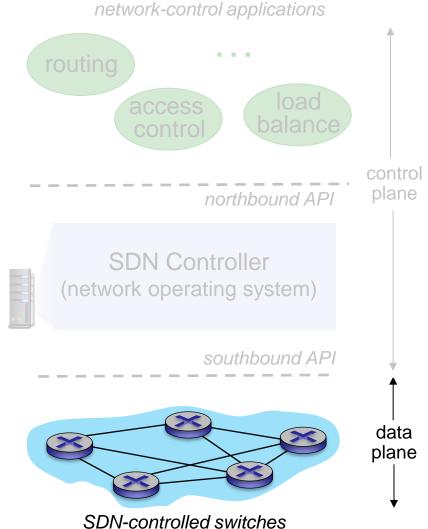
A: can't do it (with destination-based forwarding, and LS, DV routing)

We learned in Chapter 4 that generalized forwarding and SDN can be used to achieve *any* routing desired



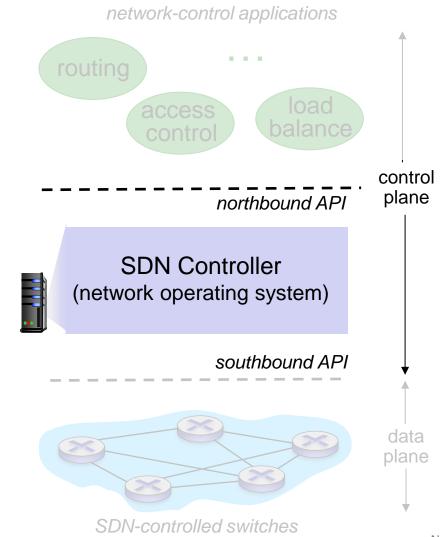
Data-plane switches:

- fast, simple, commodity switches implementing generalized data-plane forwarding (Section 4.4) in hardware
- flow (forwarding) table computed, installed under controller supervision
- API for table-based switch control (e.g., OpenFlow)
 - defines what is controllable, what is not
- protocol for communicating with controller (e.g., OpenFlow)



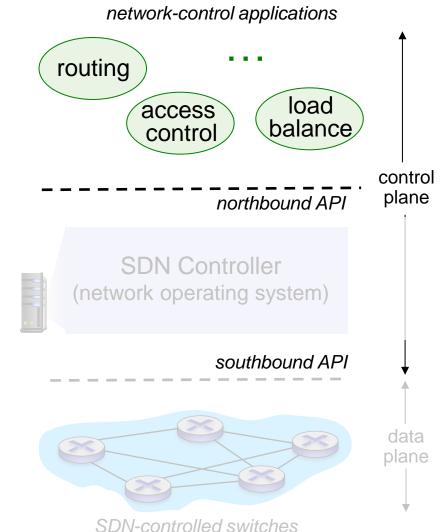
SDN controller (network OS):

- maintain network state information
- interacts with network control applications "above" via northbound API
- interacts with network switches "below" via southbound API
- implemented as distributed system for performance, scalability, faulttolerance, robustness



network-control apps:

- "brains" of control: implement control functions using lower-level services, API provided by SDN controller
- unbundled: can be provided by 3rd party: distinct from routing vendor, or SDN controller



Components of SDN controller

interface layer to network control apps: abstractions API

network-wide state

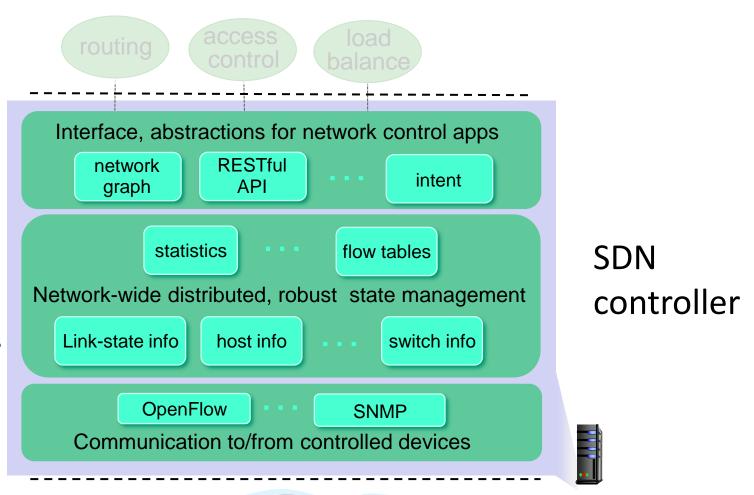
management: state of

networks links, switches,

services: a distributed database

communication: communicate between SDN controller and

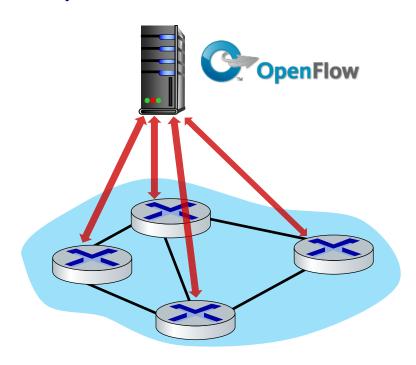
controlled switches



OpenFlow protocol

- operates between controller, switch
- TCP used to exchange messages
 - optional encryption
- three classes of OpenFlow messages:
 - controller-to-switch
 - asynchronous (switch to controller)
 - symmetric (misc.)
- distinct from OpenFlow API
 - API used to specify generalized forwarding actions

OpenFlow Controller

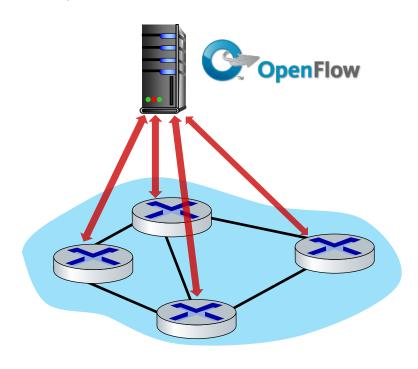


OpenFlow: controller-to-switch messages

Key controller-to-switch messages

- *features:* controller queries switch features, switch replies
- configure: controller queries/sets switch configuration parameters
- modify-state: add, delete, modify flow entries in the OpenFlow tables
- packet-out: controller can send this packet out of specific switch port

OpenFlow Controller

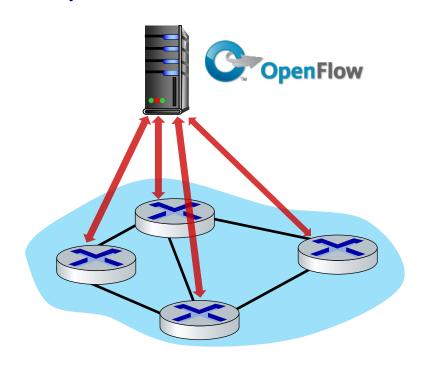


OpenFlow: switch-to-controller messages

Key switch-to-controller messages

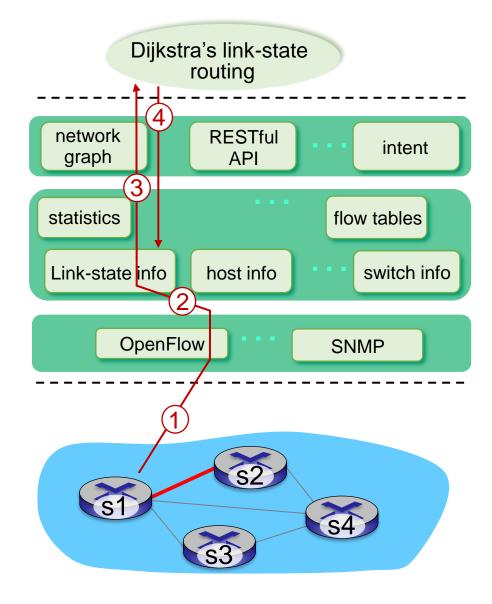
- packet-in: transfer packet (and its control) to controller. See packet-out message from controller
- flow-removed: flow table entry deleted at switch
- port status: inform controller of a change on a port.

OpenFlow Controller



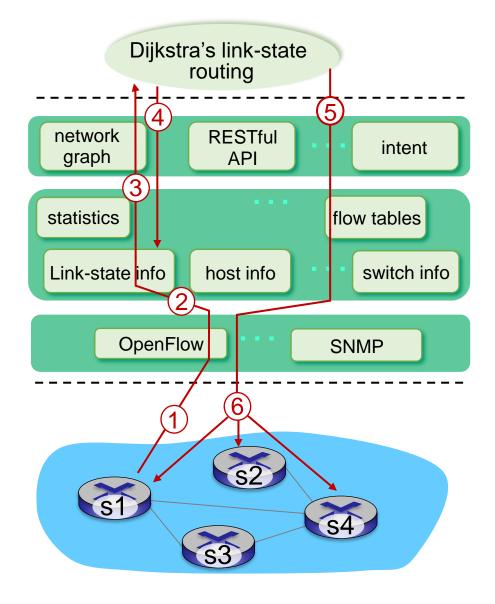
Fortunately, network operators don't "program" switches by creating/sending OpenFlow messages directly. Instead use higher-level abstraction at controller

SDN: control/data plane interaction example



- 1 S1, experiencing link failure uses OpenFlow port status message to notify controller
- 2 SDN controller receives OpenFlow message, updates link status info
- 3 Dijkstra's routing algorithm application has previously registered to be called when ever link status changes. It is called.
- Dijkstra's routing algorithm access network graph info, link state info in controller, computes new routes

SDN: control/data plane interaction example



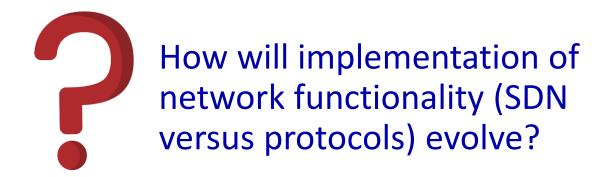
- 5 link state routing app interacts with flow-table-computation component in SDN controller, which computes new flow tables needed
- 6 controller uses OpenFlow to install new tables in switches that need updating

SDN: selected challenges

- hardening the control plane: dependable, reliable, performancescalable, secure distributed system
 - robustness to failures: leverage strong theory of reliable distributed system for control plane
 - dependability, security: "baked in" from day one?
- networks, protocols meeting mission-specific requirements
 - e.g., real-time, ultra-reliable, ultra-secure
- Internet-scaling: beyond a single AS
- SDN critical in 5G cellular networks

SDN and the future of traditional network protocols

- SDN-computed versus router-computer forwarding tables:
 - just one example of logically-centralized-computed versus protocol computed
- one could imagine SDN-computed congestion control:
 - controller sets sender rates based on router-reported (to controller) congestion levels





Network layer: "control plane" roadmap

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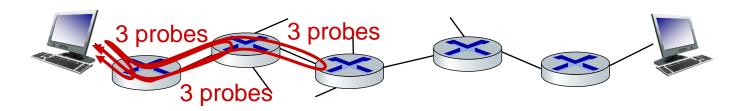


ICMP: internet control message protocol

- used by hosts and routers to communicate network-level information
 - error reporting: unreachable host, network, port, protocol
 - echo request/reply (used by ping)
- network-layer "above" IP:
 - ICMP messages carried in IP datagrams
- ICMP message: type, code plus first 8 bytes of IP datagram causing error

<u>Type</u>	<u>Code</u>	<u>description</u>
0	0	echo reply (ping)
3	0	dest. network unreachable
3	1	dest host unreachable
3	2	dest protocol unreachable
3	3	dest port unreachable
3	6	dest network unknown
3	7	dest host unknown
4	0	source quench (congestion
		control - not used)
8	0	echo request (ping)
9	0	route advertisement
10	0	router discovery
11	0	TTL expired
12	0	bad IP header

Traceroute and ICMP



- source sends sets of UDP segments to destination
 - 1st set has TTL =1, 2nd set has TTL=2, etc.
- datagram in *n*th set arrives to nth router:
 - router discards datagram and sends source ICMP message (type 11, code 0)
 - ICMP message possibly includes name of router & IP address
- when ICMP message arrives at source: record RTTs

stopping criteria:

- UDP segment eventually arrives at destination host
- destination returns ICMP "port unreachable" message (type 3, code 3)
- source stops

Network layer: Summary

we've learned a lot!

- approaches to network control plane
 - per-router control (traditional)
 - logically centralized control (software defined networking)
- traditional routing algorithms
 - implementation in Internet: OSPF, BGP
- SDN controllers
- Internet Control Message Protocol

next stop: link layer!