

# A First Look At Prolog

# Outline

- Terms
- Using a Prolog language system
- Rules
- The two faces of Prolog
- Operators
- Lists
- Negation and failure
- What Prolog is good for


# Terms

- Everything in Prolog is built from *terms*:
  - Prolog programs
  - The data manipulated by Prolog programs
- Three kinds of terms:
  - Constants: integers, real numbers, atoms
  - Variables
  - Compound terms

# Constants

- Integer constants: **123**
- Real constants: **1.23**
- Atoms:
  - A lowercase letter followed by any number of additional letters, digits or underscores: **fred**
  - A sequence of non-alphanumeric characters:  
**\*, ., =, @#\$**
  - Plus a few special atoms: **[ ]**

# Atoms Are Not Variables

- An atom can look like an ML or Java variable:
  - **i, size, length** 
- But an atom is not a variable; it is not bound to anything, never equal to anything else
- Think of atoms as being more like string constants: "i", "size", "length"

# Variables

- Any name beginning with an uppercase letter or an underscore, followed by any number of additional letters, digits or underscores: **X**, **Child**, **Fred**, **\_**, **\_123**
- Most of the variables you write will start with an uppercase letter
- Those starting with an underscore, including **\_**, get special treatment

# Compound Terms

- An atom followed by a parenthesized, comma-separated list of one or more terms:

*adam*  $\rightarrow$   $\text{x}(\underline{\text{y}}, \underline{\text{z}}), + (1, 2), . (1, []),$   
 $\text{parent}(\text{adam}, \text{seth}), \text{x}(\text{Y}, \underline{\text{x}(\text{Y}, \text{Z})})$

- A compound term can look like an ML function call:  $\text{f}(\text{x}, \text{y})$  *term*
- Again, this is misleading
- Think of them as structured data

# Terms

$\langle term \rangle ::= \langle constant \rangle \mid \langle variable \rangle \mid \langle compound-term \rangle$   
 $\langle constant \rangle ::= \langle integer \rangle \mid \langle real\ number \rangle \mid \langle atom \rangle$   
 $\langle compound-term \rangle ::= \langle atom \rangle ( \langle termlist \rangle )$   
 $\langle termlist \rangle ::= \langle term \rangle \mid \langle term \rangle , \langle termlist \rangle$

- All Prolog programs and data are built from such terms
- Later, we will see that, for instance, **+ (1 , 2)** is usually written as **1+2**
- But these are not new kinds of terms, just abbreviations



# Unification

- Pattern-matching using Prolog terms
- Two terms *unify* if there is some way of binding their variables that makes them identical
- For instance, parent (adam, Child)<sup>atom</sup> and parent (adam, seth)<sup>variable</sup> unify by binding the variable **Child** to the atom **seth**
- More details later: Chapter 20

# The Prolog Database

- A Prolog language system maintains a collection of facts and rules of inference
- It is like an internal database that changes as the Prolog language system runs
- A Prolog program is just a set of data for this database
- The simplest kind of thing in the database is a *fact*: a term followed by a period

# Example

*Kim is parent of holly. Be consistent.*

```
parent(kim,holly) .  
parent(margaret,kim) .  
parent(margaret,kent) .  
parent(esther,margaret) .  
parent(herbert,margaret) .  
parent(herbert,jean) .
```

- A Prolog program of six facts
- Defining a *predicate* **parent** of *arity* 2
- We would naturally interpret these as facts about families: Kim is the parent of Holly and so on

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# SWI-Prolog

```
Welcome to SWI-Prolog ...
```

```
For help, use ?- help(Topic) . or ?- apropos(Word) .
```

```
?-
```

- Prompting for a query with ?-
- Normally interactive: get query, print result, repeat

# The `consult` Predicate

```
?- consult(relations).  
% relations compiled 0.00 sec, 852 bytes  
true.  
  
?-
```

- Predefined predicate to read a program from a file into the database
- File **relations** (or **relations.pl**) contains our **parent** facts

# Simple Queries

```
?- parent(margaret,kent) .  
true.
```

*unify  
with every fact*

```
?- parent(fred,pebbles) .  
false.
```

*v. if yes → unifiable → true*

```
?-
```

- A query asks the language system to prove something
- Some turn out to be **true**, some **false**
- (Some queries, like **consult**, are executed only for their side-effects)

# Final Period

```
?- parent(margaret,kent)
|      . ← end of question
true.

?-
```

- ❑ Queries can take multiple lines
- ❑ If you forget the final period, Prolog prompts for more input with |



# Queries With Variables

```
?- parent(P,jean) . Who is the parent of Jean?  
P = herbert.  
  
?- parent(P,esther) .  
false.
```

- Any term can appear as a query, including a term with variables
- The Prolog system shows the bindings necessary to prove the query

# Flexibility

- Normally, variables can appear in any or all positions in a query:

– **parent (Parent, jean)**

– **parent (esther, Child)** *Esther is parent of whom?*

– **parent (Parent, Child)**

– **parent (Person, Person)**

*Who has the same name as parent?*

# Multiple Solutions

```
?- parent(Parent,Child) .  
Parent = kim,  
Child = holly .
```

- When the system finds a solution, it prints the binding it found
- If it could continue to search for additional solutions, it then prompts for input
- Hitting Enter makes it stop searching and print the final period...

# Multiple Solutions

- ... entering a semicolon makes it continue the search
- As often as you do this, it will try to find another solution
- In this case, there is one for every fact in the database

```
?- parent(Parent,Child) .  
Parent = kim,  
Child = holly ;  
Parent = margaret,  
Child = kim ;  
Parent = margaret,  
Child = kent ;  
Parent = esther,  
Child = margaret ;  
Parent = herbert,  
Child = margaret ;  
Parent = herbert,  
Child = jean.
```

False (there is no more)

# Conjunctions

```
?- parent(margaret,X) , parent(X,holly) .  
X = kim .
```

- A conjunctive query has a list of query terms separated by commas
- The Prolog system tries prove them all (using a single set of bindings)

and (conjunction).

```
/?- parent(Parent, kim), parent(Grandparent, Parent).  
Parent = margaret,  
Grandparent = esther ;  
Parent = margaret,  
Grandparent = herbert ;  
false
```

there is no more.

```
?- parent(esther, Child),  
|    parent(Child, Grandchild),  
|    parent(Grandchild, GreatGrandchild).  
Child = margaret,  
Grandchild = kim,  
GreatGrandchild = holly .
```

# Outline

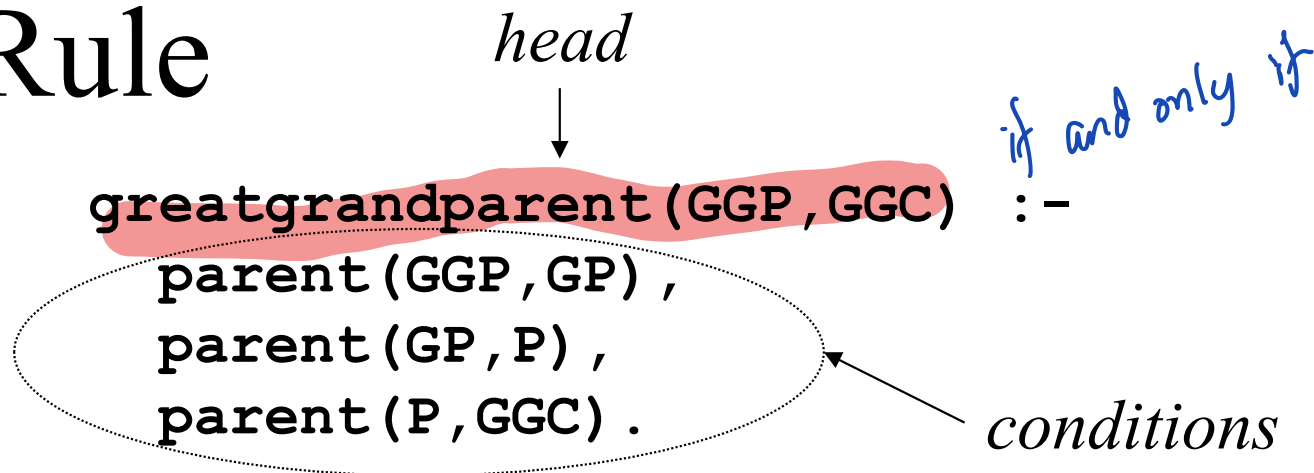
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# The Need For Rules

- Previous example had a lengthy query for great-grandchildren of Esther
- It would be nicer to query directly:  
**greatgrandparent (esther , GGC)**
- But we do not want to add separate facts of that form to the database
- The relation should follow from the **parent** relation already defined



# A Rule



- A rule says how to prove something: to prove the head, prove the conditions
- To prove `greatgrandparent (GGP, GGC)`, find some `GP` and `P` for which you can prove `parent (GGP, GP)`, then `parent (GP, P)` and then finally `parent (P, GGC)`

# A Program With The Rule

```
parent(kim,holly) .  
parent(margaret,kim) .  
parent(margaret,kent) .  
parent(esther,margaret) .  
parent(herbert,margaret) .  
parent(herbert,jean) .  
greatgrandparent(GGP,GGC) :-  
    parent(GGP,GP) , parent(GP,P) , parent(P,GGC) .
```



- A program consists of a list of *clauses*
- A clause is either a fact or a rule, and ends with a period



# Example

```
?- greatgrandparent(esther, GreatGrandchild) .  
GreatGrandchild = holly .
```

- This shows the initial query and final result
- Internally, there are intermediate *goals*:
  - The first goal is the initial query
  - The next is what remains to be proved after transforming the first goal using one of the clauses (in this case, the greatgrandparent rule)
  - And so on, until nothing remains to be proved

- 
- 
1. **parent(kim,holly) .**
  2. **parent(margaret,kim) .**
  3. **parent(margaret,kent) .**
  4. **parent(esther,margaret) .**
  5. **parent(herbert,margaret) .**
  6. **parent(herbert,jean) .**
  7. **greatgrandparent(GGP,GGC) :-**  
    **parent(GGP,GP) , parent(GP,P) , parent(P,GGC) .**

*We will see more  
about Prolog's model  
of execution in  
Chapter 20*

**greatgrandparent(esther,GreatGrandchild)**

↓ Clause 7, binding **GGP** to **esther** and **GGC** to **GreatGrandChild**

**parent(esther,GP) , parent(GP,P) , parent(P,GreatGrandchild)**

↓ Clause 4, binding **GP** to **margaret**

**parent(margaret,P) , parent(P,GreatGrandchild)**

↓ Clause 2, binding **P** to **kim**

**parent(kim,GreatGrandchild)**

↓ Clause 1, binding **GreatGrandchild** to **holly**

*if there any  
other?  
backtrack.*

# Rules Using Other Rules

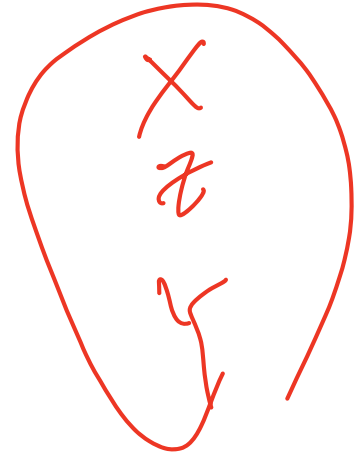
```
grandparent(GP,GC) :-  
    parent(GP,P) , parent(P,GC) .
```

```
greatgrandparent(GGP,GGC) :-  
    grandparent(GGP,P) , parent(P,GGC) .
```

- Same relation, defined indirectly
- Note that both clauses use a variable **P**
- The scope of the definition of a variable is the clause that contains it

# Recursive Rules

```
ancestor(X,Y) :- parent(X,Y) .  
ancestor(X,Y) :-  
    parent(Z,Y) ,  
    ancestor(X,Z) .
```



- **X** is an ancestor of **Y** if:
  - Base case: **X** is a parent of **Y**
  - Recursive case: there is some **Z** such that **Z** is a parent of **Y**, and **X** is an ancestor of **Z**
- Prolog tries rules in the order you give them, so put base-case rules and facts first

```
?- ancestor(jean,jean) .  
false.
```

```
?- ancestor(kim,holly).  
true .
```

```
?- ancestor(A,holly) .  
A = kim ;  
A = margaret ;  
A = esther ;  
A = herbert ;  
false.
```

*Ancestor???*

# Core Syntax Of Prolog

- You have seen the complete core syntax:

$\langle clause \rangle ::= \langle fact \rangle \mid \langle rule \rangle$

$\langle fact \rangle ::= \langle term \rangle .$

$\langle rule \rangle ::= \langle term \rangle :- \langle termlist \rangle .$

$\langle termlist \rangle ::= \langle term \rangle \mid \langle term \rangle , \langle termlist \rangle$

- There is not much more syntax for Prolog than this: it is a very simple language
- Syntactically, that is!



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# The Procedural Side

```
greatgrandparent (GGP,GGC) :-  
    parent (GGP,GP) , parent (GP,P) , parent (P,GGC) .
```

- A rule says how to prove something:
  - To prove `greatgrandparent (GGP,GGC)`, find some `GP` and `P` for which you can prove `parent (GGP,GP)`, then `parent (GP,P)` and then finally `parent (P,GGC)`
- A Prolog program specifies proof procedures for queries

# The Declarative Side

- A rule is a logical assertion:
  - For all bindings of `GGP`, `GP`, `P`, and `GGC`, if  
`parent(GGP, GP)` and `parent(GP, P)` and  
`parent(P, GGC)`, then `greatgrandparent(GGP, GGC)`
- Just a formula – it doesn't say how to *do* anything – it just makes an assertion:

$$\forall GGP, GP, P, GGC . \text{parent}(GGP, GP) \wedge \text{parent}(GP, P) \wedge \text{parent}(P, GGC) \\ \Rightarrow \text{greatgrandparent}(GGP, GGC)$$

# Declarative Languages

- Each piece of the program corresponds to a simple mathematical abstraction
  - Prolog clauses – formulas in first-order logic
  - ML fun definitions – functions
- Many people use *declarative* as the opposite of *imperative*, including both logic languages and functional languages

# Declarative Advantages

- ❑ Imperative languages are doomed to subtle side-effects and interdependencies
- ❑ Simpler declarative semantics makes it easier to develop and maintain correct programs
- ❑ Higher-level, more like *automatic programming*: describe the problem and have the computer write the program

# Prolog Has Both Aspects

- Partly declarative
  - A Prolog program has logical content
- Partly procedural
  - A Prolog program has procedural concerns: clause ordering, condition ordering, side-effecting predicates, etc.
- It is important to be aware of both

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# Operators

- Prolog has some predefined operators (and the ability to define new ones)
- An operator is just a predicate for which a special abbreviated syntax is supported



# The = Predicate

- The goal  $\text{= (X, Y)}$  succeeds if and only if **X** and **Y** can be unified:

```
?- =(parent(adam, seth), parent(adam, X)) .  
X = seth.
```

- Since = is an operator, it can be and usually is written like this:

```
?- parent(adam, seth) = parent(adam, X) .  
X = seth.
```

# Arithmetic Operators

- Predicates  $+$ ,  $-$ ,  $*$  and  $/$  are operators too, with the usual precedence and associativity

```
?- X = +(1, *(2,3)).  
X = 1+2*3.
```

```
?- X = 1+2*3.  
X = 1+2*3.
```

Prolog lets you use operator notation, and prints it out that way, but the underlying term is still  $+(1, *(2, 3))$

# Not Evaluated

```
?- +(X,Y) = 1+2*3.
```

```
X = 1,
```

```
Y = 2*3.
```

```
?- 7 = 1+2*3.
```

```
false.
```

- The term is still  **$+(1, *(2, 3))$**
- It is not evaluated
- There is a way to make Prolog evaluate such terms, but we won't need it yet

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# Lists in Prolog

- A bit like ML lists
- The atom `[]` represents the empty list
- A predicate `.` corresponds to ML's `::` operator

ML expression	Prolog term
<code>[]</code>	<code>[]</code>
<code>1 :: []</code>	<code>. (1, [])</code>
<code>1 :: 2 :: 3 :: []</code>	<code>. (1, . (2, . (3, [])))</code>
No equivalent.	<code>. (1, . (parent(X,Y), []))</code>

# List Notation

List notation	Term denoted
<code>[]</code>	<code>[]</code>
<code>[1]</code>	<code>. (1, [])</code>
<code>[1, 2, 3]</code>	<code>. (1, . (2, . (3, [])))</code>
<code>[1, parent(X, Y)]</code>	<code>. (1, . (parent(X, Y), []))</code>

- ML-style notation for lists
- These are just abbreviations for the underlying term using the `.` Predicate
- Prolog usually displays lists in this notation

# Example

```
?- X = . (1, . (2, . (3, []))) .  
X = [1, 2, 3] .
```

```
?- . (X, Y) = [1, 2, 3] .  
X = 1,  
Y = [2, 3] .
```

# List Notation With Tail

List notation	Term denoted
$[1   \mathbf{x}]$	$. (1, \mathbf{x})$
$[1, 2   \mathbf{x}]$	$. (1, . (2, \mathbf{x}) )$
$[1, 2   [3, 4]]$	same as $[1, 2, 3, 4]$

- Last in a list can be the symbol  $|$  followed by a final term for the tail of the list
- Useful in patterns:  $[1, 2 | \mathbf{x}]$  unifies with any list that starts with  $1, 2$  and binds  $\mathbf{x}$  to the tail

$$\begin{aligned} ?- [1, 2 | X] &= [1, 2, 3, 4, 5] . \\ \mathbf{x} &= [3, 4, 5] . \end{aligned}$$



# The **append** Predicate

```
?- append([1,2],[3,4],Z) .  
Z = [1, 2, 3, 4] .
```

- Predefined **append(X, Y, Z)** succeeds if and only if **Z** is the result of appending the list **Y** onto the end of the list **X**

# Not Just A Function

```
?- append(X, [3,4], [1,2,3,4]) .  
x = [1, 2] .
```

- **append** can be used with any pattern of instantiation (that is, with variables in any positions)

# Not Just A Function

```
?- append(X,Y,[1,2,3]).  
X = [],  
Y = [1, 2, 3] ;  
X = [1],  
Y = [2, 3] ;  
X = [1, 2],  
Y = [3] ;  
X = [1, 2, 3],  
Y = [] ;  
false.
```

# An Implementation

```
append([], B, B) .  
append([Head|TailA], B, [Head|TailC]) :-  
    append(TailA, B, TailC) .
```

# Other Predefined List Predicates

Predicate	Description
<b>member</b> ( <b>X</b> , <b>Y</b> )	Provable if the list <b>Y</b> contains the element <b>X</b> .
<b>select</b> ( <b>X</b> , <b>Y</b> , <b>Z</b> )	Provable if the list <b>Y</b> contains the element <b>X</b> , and <b>Z</b> is the same as <b>Y</b> but with one instance of <b>X</b> removed.
<b>nth0</b> ( <b>X</b> , <b>Y</b> , <b>Z</b> )	Provable if <b>X</b> is an integer, <b>Y</b> is a list, and <b>Z</b> is the <b>X</b> th element of <b>Y</b> , counting from 0.
<b>length</b> ( <b>X</b> , <b>Y</b> )	Provable if <b>X</b> is a list of length <b>Y</b> .

- All flexible, like **append**
- Queries can contain variables anywhere

# Using **select**

```
?- select(2,[1,2,3],Z) .  
Z = [1, 3] ;  
false.
```

```
?- select(2,Y,[1,3]) .  
Y = [2, 1, 3] ;  
Y = [1, 2, 3] ;  
Y = [1, 3, 2] ;  
false.
```

# The **reverse** Predicate

```
?- reverse([1,2,3,4],Y).  
Y = [4, 3, 2, 1].
```

- Predefined **reverse(X,Y)** unifies **Y** with the reverse of the list **X**

# An Implementation

```
reverse([], []).  
reverse([Head|Tail], X) :-  
    reverse(Tail, Y),  
    append(Y, [Head], X).
```

- ❑ Not an efficient way to reverse
- ❑ We'll see why, and a more efficient solution, in Chapter 21



# Non-Terminating Queries

```
?- reverse(X, [1,2,3,4]).
```

```
X = [4, 3, 2, 1] ;
```

```
^CAction (h for help) ? abort
```

```
% Execution Aborted
```

```
?-
```

- Asking for another solution caused an infinite loop
- Hit Control-C to stop it, then *a* for abort
- **reverse** cannot be used as flexibly as **append**

# Flexible and Inflexible

- Ideally, predicates should all be flexible like **append**
- They are more declarative, with fewer procedural quirks to consider
- But inflexible implementations are sometimes used, for efficiency or simplicity
- Another example is **sort**...

# Example

```
?- sort([2,3,1,4],X).
```

```
X = [1, 2, 3, 4].
```

```
?- sort(X,[1,2,3,4]).
```

```
ERROR: Arguments are not sufficiently instantiated
```

- A fully flexible **sort** would also be able to unsort—find all permutations
- But it would not be as efficient for the more common task

# The Anonymous Variable

- The variable `_` is an anonymous variable
- Every occurrence is bound independently of every other occurrence
- In effect, much like ML's `_`: it matches any term without introducing bindings

# Example

**tailof**([\_|A],A) .

- This **tailof**(**X**,**Y**) succeeds when **X** is a non-empty list and **Y** is the tail of that list
- Don't use this, even though it works:

**tailof**([Head|A],A) .

# Dire Warning

```
append([], B, B) .
```

```
append([Head|TailA], B, [Head|TailC]) :-  
    append(TailA, B, TailC) .
```

- Don't ignore warning message about singleton variables
- As in ML, it is bad style to introduce a variable you never use
- More importantly: *if you misspell a variable name, this is the only warning you will see*

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# The **not** Predicate

```
?- member(1,[1,2,3]).  
true .
```

```
?- not(member(4,[1,2,3])).  
false.
```

- For simple applications, it often works quite a bit logical negation
- But it has an important procedural side...



# Negation As Failure

- To prove **not (X)** , Prolog attempts to prove **X**
- **not (X)** succeeds if **X** fails
- The two faces again:
  - Declarative: **not (X)** =  $\neg X$
  - Procedural: **not (X)** succeeds if **X** fails, fails if **X** succeeds, and runs forever if **X** runs forever

# Example

```
sibling(X,Y) :-  
    not(X=Y) ,  
    parent(P,X) ,  
    parent(P,Y) .
```

```
?- sibling(kim,kent) .  
true .
```

```
?- sibling(kim,kim) .  
false.
```

```
?- sibling(X,Y) .  
false.
```

```
sibling(X,Y) :-  
    parent(P,X) ,  
    parent(P,Y) ,  
    not(X=Y) .
```

```
?- sibling(X,Y) .  
X = kim,  
Y = kent ;  
X = kent,  
Y = kim ;  
X = margaret,  
Y = jean ;  
X = jean,  
Y = margaret ;  
false.
```

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# A Classic Riddle

- ❑ A man travels with wolf, goat and cabbage
- ❑ Wants to cross a river from west to east
- ❑ A rowboat is available, but only large enough for the man plus one possession
- ❑ Wolf eats goat if left alone together
- ❑ Goat eats cabbage if left alone together
- ❑ How can the man cross without loss?

# Configurations

- Represent a configuration of this system as a list showing which bank each thing is on in this order: man, wolf, goat, cabbage
- Initial configuration: **[w, w, w, w]**
- If man crosses with wolf, new state is **[e, e, w, w]** – but then goat eats cabbage, so we can't go through that state
- Desired final state: **[e, e, e, e]**

# Moves

- In each move, man crosses with at most one of his possessions
- We will represent these four moves with four atoms: **wolf**, **goat**, **cabbage**, **nothing**
- (Here, **nothing** indicates that the man crosses alone in the boat)

# Moves Transform Configurations

- Each move transforms one configuration to another
- In Prolog, we will write this as a predicate:  
**move (Config, Move, NextConfig)**
  - **Config** is a configuration (like **[w, w, w, w]**)
  - **Move** is a move (like **wolf**)
  - **NextConfig** is the resulting configuration (in this case, **[e, e, w, w]**)

# The **move** Predicate

```
change (e , w) .
```

```
change (w , e) .
```

```
move ( [X , X , Goat , Cabbage] , wolf , [Y , Y , Goat , Cabbage] ) :-  
    change (X , Y) .
```

```
move ( [X , Wolf , X , Cabbage] , goat , [Y , Wolf , Y , Cabbage] ) :-  
    change (X , Y) .
```

```
move ( [X , Wolf , Goat , X] , cabbage , [Y , Wolf , Goat , Y] ) :-  
    change (X , Y) .
```

```
move ( [X , Wolf , Goat , C] , nothing , [Y , Wolf , Goat , C] ) :-  
    change (X , Y) .
```



# Safe Configurations

- A configuration is safe if
  - At least one of the goat or the wolf is on the same side as the man, and
  - At least one of the goat or the cabbage is on the same side as the man

```
oneEq (X, X, _) .
```

```
oneEq (X, _, X) .
```

```
safe ( [Man, Wolf, Goat, Cabbage] ) :-
```

```
    oneEq (Man, Goat, Wolf) ,
```

```
    oneEq (Man, Goat, Cabbage) .
```

# Solutions

- A solution is a starting configuration and a list of moves that takes you to **[e,e,e,e]**, where all the intermediate configurations are safe

```
solution([e,e,e,e],[ ]).  
solution(Config,[Move|Rest]) :-  
    move(Config,Move,NextConfig),  
    safe(NextConfig),  
    solution(NextConfig,Rest).
```

# Prolog Finds A Solution

```
?- length(X,7), solution([w,w,w,w],X).  
X = [goat, nothing, wolf, goat, cabbage, nothing, goat] .
```

- Note: without the **length(X,7)** restriction, Prolog would not find a solution
- It gets lost looking at possible solutions like **[goat,goat,goat,goat,goat...]**
- More about this in Chapter 20

# What Prolog Is Good For

- The program specified a problem logically
- It did not say how to search for a solution to the problem – Prolog took it from there
- That's one kind of problem Prolog is especially good for