

- * Exam 2 : Median 70 + 17
- * HW 4 graded
- * HW 5 out, due Nov 4 (One more HW with 4+ pbs after this)
- * Recap : Convex fns and Convex Optim

Gradient Descent (assume no constraints K)

$$\min_{x \in \mathbb{R}^n} f(x) \quad \leftarrow \text{we want } x^* \text{ s.t. } \|\nabla f(x^*)\| \approx 0$$

Since we want to find a local opt,

take a step in $-\nabla f(x) \leftarrow$ direction of steepest descent

(we are not given ∇f fn explicitly to solve directly)

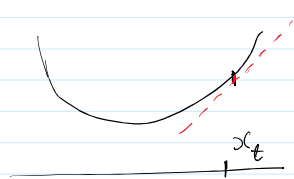
Grad Descent Alg:

- 1) Start with x_0
- 2) In t -th step:

$$x_t = x_{t-1} - \underbrace{\eta}_{\text{step size}} \nabla f(x_{t-1})$$

Exercise:

$$\lim_{\epsilon \rightarrow 0} \max_{y: \|y\|=\epsilon} \frac{f(x+y) - f(x)}{\epsilon} = \|\nabla f(x)\|_2$$



\rightarrow We will show that after $\text{poly}(1/\epsilon)$ steps, G.D. finds ϵ -optimal soln.

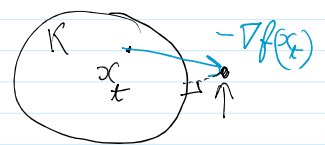
Properties: (1) G.D. is fast when ϵ is not very small.

(2) Only requires gradient oracle

(3) Well-defined (and works) even for non-convex fns to give local-opt.

(4) Similar results even hold with convex constraints K , but then we need to project on to K

Thm 1: Assume $\|\nabla f(x)\|_2 \leq G$ and $\|x_0 - x^*\| \leq D$.



After T steps of G.D with $\eta_t = \frac{D}{G\sqrt{T}}$,

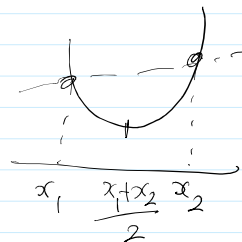
$$f\left(\frac{\sum_{t=0}^{T-1} x_t}{T}\right) - f(x^*) \leq \frac{GD}{\sqrt{T}} \Rightarrow \text{For } \epsilon \text{ error, } T \geq \frac{G^2 D^2}{\epsilon^2}.$$

$$f\left(\frac{\sum_{t=0}^T x_t}{T}\right) - f(x^*) \leq \frac{GD}{\sqrt{T}} \Rightarrow \text{For } \epsilon \text{ error, } T \geq \frac{G^2 D^2}{\epsilon^2}.$$

Obs: Suffices to prove $\sum_{t=0}^{T-1} \left[f(x_t) - f(x^*) \right] \leq \frac{GD}{\sqrt{T}} \times T$

Regret at step t

Since $f\left(\frac{\sum x_t}{T}\right) \stackrel{\text{Convexity}}{\leq} \sum \frac{f(x_t)}{T}$



Pf: Define a potential function $\phi_t = \|x_t - x^*\|^2 \cdot \frac{1}{2\eta}$

We will show that

$$\left[f(x_t) - f(x^*) \right] \leq \left(\phi_t - \phi_{t+1} \right) + \frac{\eta G^2}{2} \quad \text{--- (1)}$$

This suffices since $\sum_{t=0}^{T-1} \left[f(x_t) - f(x^*) \right] \leq \underbrace{\phi_0}_{\text{Regret}} - \underbrace{\phi_T}_{\geq 0} + \frac{\eta G^2}{2} T$

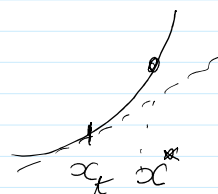
$$\leq \frac{D^2}{2\eta} + \frac{\eta G^2 T}{2} = DG\sqrt{T}, \quad \boxed{\eta = \frac{D}{G\sqrt{T}}}$$

Finally, to prove (1),

$$2\eta \cdot \phi_{t+1} = \|x_{t+1} - x^*\|^2$$

$$= \|x_t - \eta \nabla f(x_t) - x^*\|^2$$

$$= \underbrace{\|x_t - x^*\|^2}_{= \phi_t \cdot 2\eta} + \underbrace{\eta^2 \|\nabla f(x_t)\|^2}_{\leq G^2} + 2\eta \underbrace{\langle \nabla f(x_t), x^* - x_t \rangle}_{\leq f(x^*) - f(x_t)}$$



$$\Rightarrow f(x_t) - f(x^*) \leq \phi_t - \phi_{t+1} + \frac{\eta G^2}{2}$$



Takeaways and Remarks

1) Grad descent gets ϵ -approx in $\approx 1/\epsilon^2$ steps

- 1) Grad descent gets ϵ -approx in $\approx 1/\epsilon^2$ steps
- 2) This can be shown to be optimal for 1st-order methods (grad-based)
- 3) we can do better if more assumptions on gradient.

Smoothness and Strong-Convexity

Roughly, smoothness is an upper bound on 2nd derivative
 strong-convexity is a lower " " "

β -Smoothness:

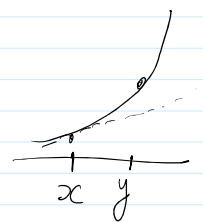
$$(1) \quad \nabla^2 f(x) \preceq \beta \quad \leftarrow \text{all e.v.s} \leq \beta$$

$$(2) \quad \|\nabla f(x) - \nabla f(y)\| \leq \beta \cdot \|x - y\|$$

$$(3) \quad f(y) \leq f(x) + \langle \nabla f(x), y - x \rangle + \frac{\beta}{2} \|y - x\|^2$$

Recall 1st-order defn of convexity:

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle$$



α -Strong Convexity:

$$(1) \quad \nabla^2 f(x) \succeq \alpha \quad \leftarrow \text{all e.v.s} \geq \alpha$$

$$(2) \quad \|\nabla f(x) - \nabla f(y)\| \geq \alpha \cdot \|x - y\|$$

$$(3) \quad f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle + \frac{\alpha}{2} \|y - x\|^2$$

E.g.: $f(x) = \|Ax - b\|^2 / 2$
 $\nabla f(x) = (A^T A x - A^T b)$
 $\beta = \text{Largest e.v. of } A^T A$
 $\alpha = \text{Min e.v. of } A^T A$

What can be shown:

	Function class	# iterations (ignoring non- ϵ dependency)
1)	Convex	$1/\epsilon^2$
2)	Strongly-convex	$1/\epsilon$
3)	Smooth	$1/\epsilon$
4)	Str-convex + Smooth	$\text{polylog}(1/\epsilon)$

$$4) \text{ Str-convex} + \text{Smooth} \left\{ \begin{array}{l} 1/\epsilon \\ \text{polylog}(1/\epsilon) \end{array} \right.$$