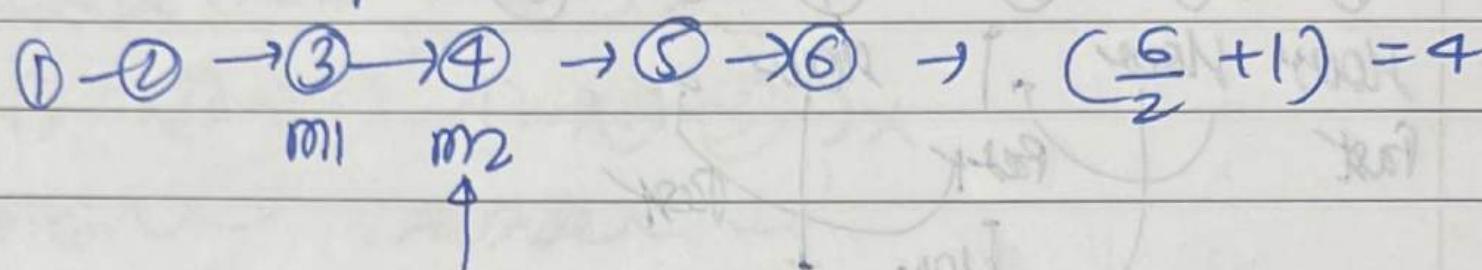
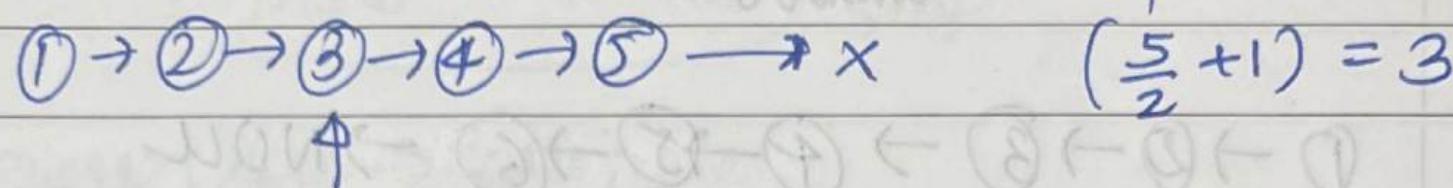


∅ } Middle element of the linkedlist $\left\{ \frac{N}{2} + 1 \right\}$



② Brute force :- Node* calculateHead() {

int cnt = 0;

$T \rightarrow O(N^2)$ Node* temp = head;

for (temp = head; temp != NULL; temp = temp->next);

while (temp) {

temp = temp->next;

cnt++;

}

int mid = cnt / 2 + 1;

temp = head;

while (temp) {

mid--;

if (!mid) return temp;

temp = temp->next;

}

return head;

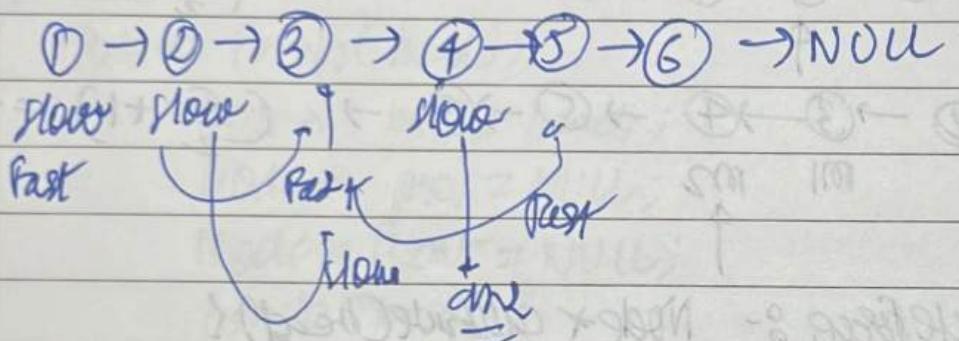
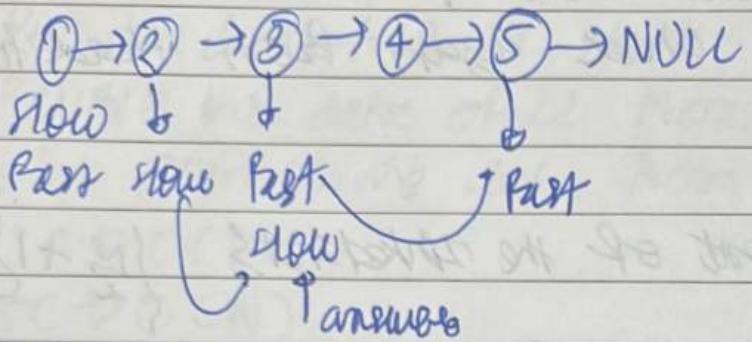
}

$O(N^2)$

COSTLY

{ Tortoise & Hare }

Optimal:- { Tortoise and Hare }



Node* callOptNode (head)?

Nodes slow = head, fast = head;

while (Fast != NULL || Fast->next != NULL) {

 slow = slow->next;

 fast = fast->next->next;



return slow;

TC → O(N/2)

SC → O(1)

⑥ { Reverse LL3

Φ Brute

Put the data of LL in stack

{ Rewrite the data of LL from Head to tail }

{ By popping the data from stack }

T C → O(2N)

S C → O(N)

Φ Optimal-O

Node* Revert(Head){

Node* temp = head;

Node* prev = NULL;

Node* front = NULL;

while (temp){

front = temp->next;

temp->next = prev;

prev = temp;

temp = front;

}

return prev;

T C → O(N)

S C → O(1)

T C → O(N)

S C → O(N) Recursive Node* newHead =

using Recursion

Node* Recursion(Node* Head){

if (Head == NULL || Head->next == NULL) return Head;

Node* newHead = Recursion(Head->next);

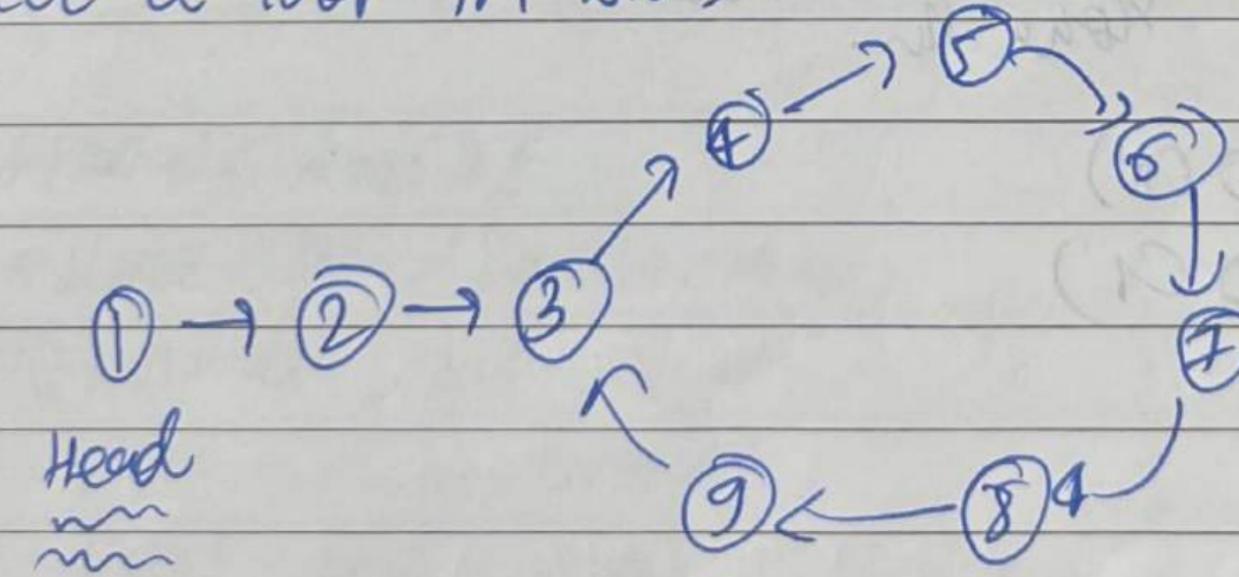
Node* front = Head->next;

front->next = Head;

Head->next = NULL;

return newHead;

⑩ { detect a loop in LL3

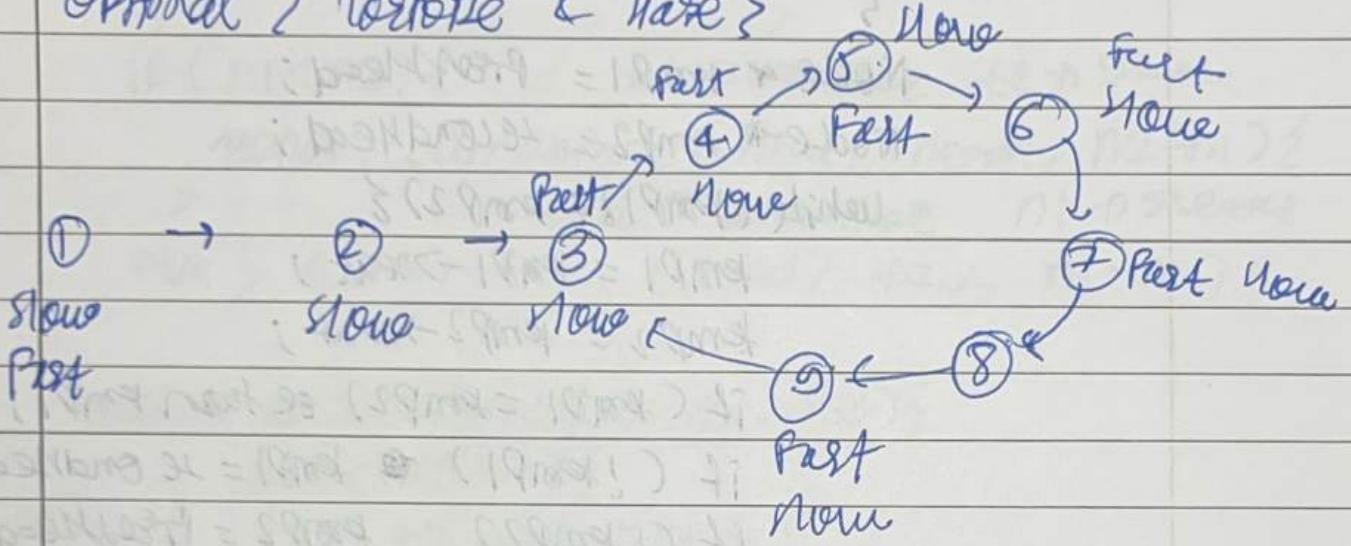


Brute 3- store in map if it already in map else?

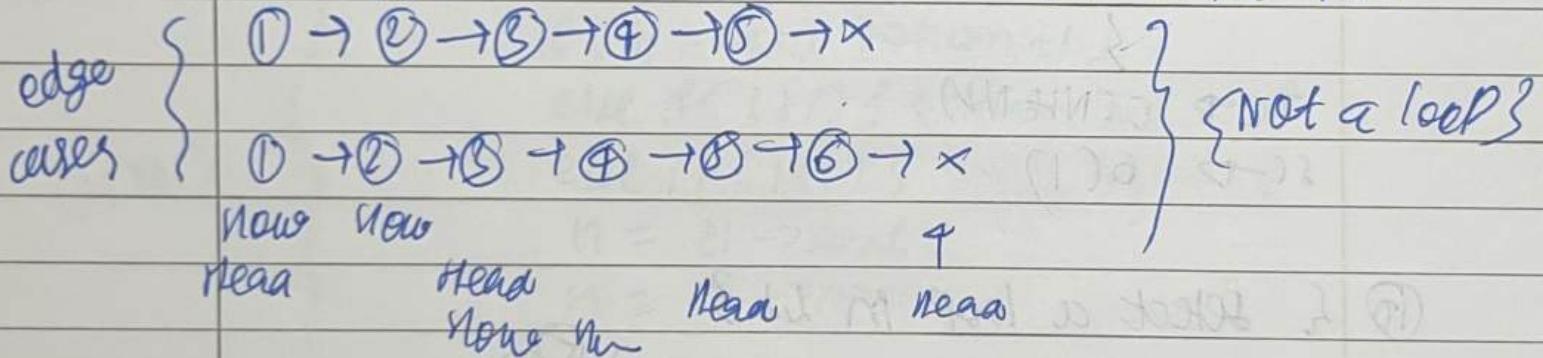
$T_C \rightarrow O(N^2)$ for unordered map?

$S_C \rightarrow O(N)$ } two operations { checking & storing in map

OPTIONAL { Tortoise & Hare }



head None If (they are standing
Slow None Head at same pts)



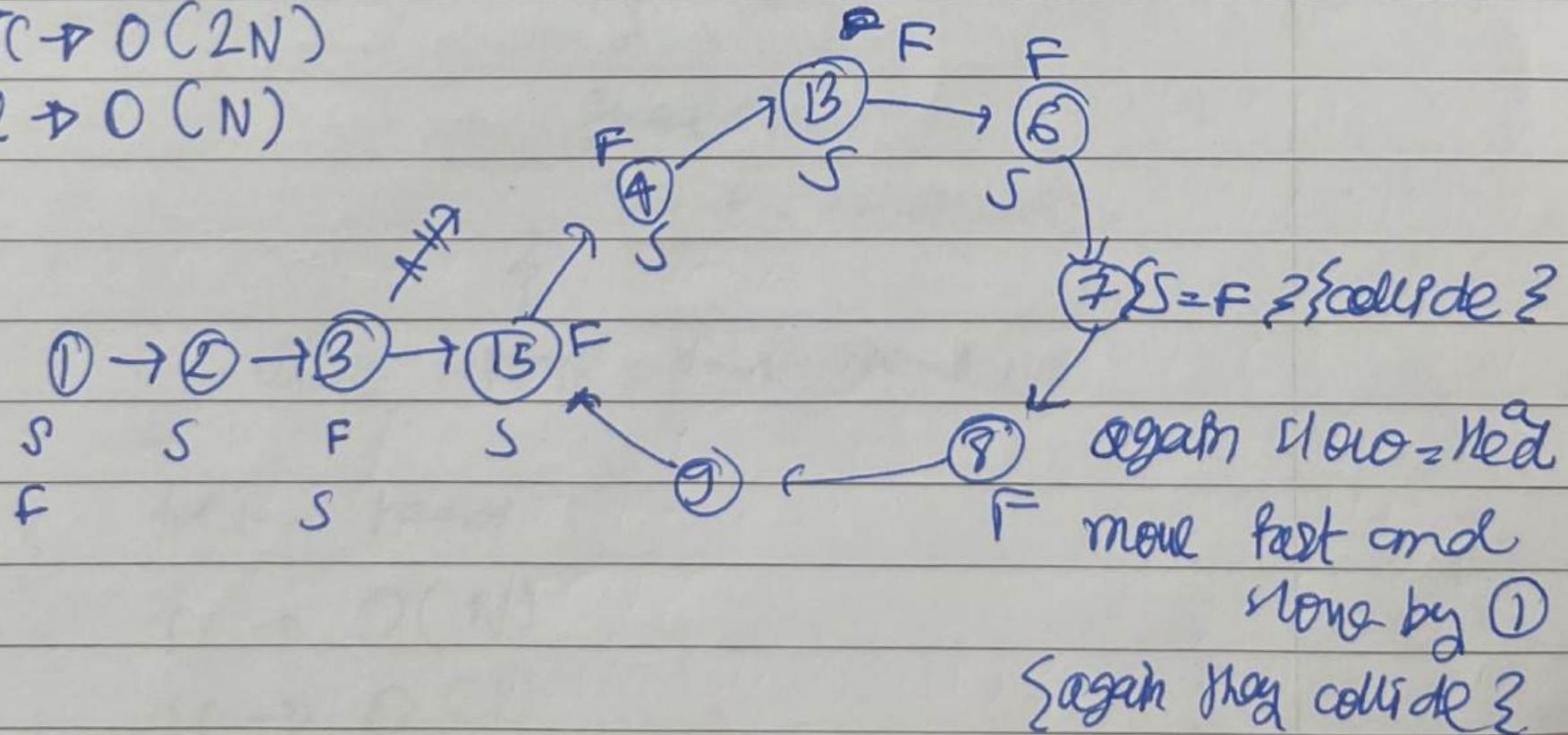
$T_C \rightarrow O(N)$

$S_C \rightarrow O(1)$

⑬ { starting node in a loop of LL3
back & start in the map }

$T C \rightarrow O(2N)$

$S Q \rightarrow O(N)$



slow = head; fast = head

while (!fast and !fast->next) {

slow = slow->next

fast = fast->next->next

if (slow == fast) {

slow = head

while (slow != fast) {

slow = slow->next

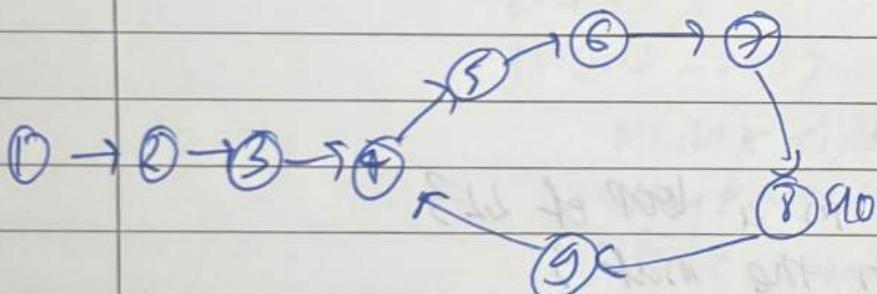
fast = fast->next

} return slow;

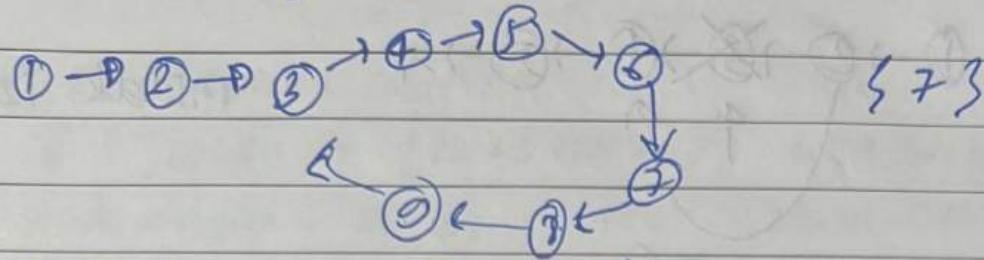
}

} return NULL;

}



⑪ { Find the length of a loop in the LL }



int breakLongLoop(Node* head) {

Node* kmp = head;

int len = 0;

unordered_map<Node*, int> mapp;

while (temp) {

if (mapp.find(temp) == mapp.end()) {

len = len + 1;

mapp[temp] = (len++);

temp = temp->next;

}
return 0;

TC \rightarrow O(2N) anode

SC \rightarrow O(N)

{ TC \rightarrow O(N * length) }
{ SC \rightarrow O(1) }

Optimal

int optimal(Node* head) {

\Rightarrow

Node* slow = head, * fast = head,

while (fast & fast->next) {

slow = slow->next;

fast = fast->next->next;

if (fast == slow) {

return length(slow, fast);

}
return 0; }

} int length(Node, fast) {

int cnt = 1;

fast = fast->next;

while (fast == slow) {

cnt++;

fast = fast->next;

} return cnt;

⑦ { Check if L2 P3 Palindrome }
∅ brute { Put all data in stack
and check again from Head to tail }

T C → O(2N)

S C → O(N)

∅ { Middle element of the linked lists } $\left\lfloor \frac{N}{2} + 1 \right\rfloor$

$$1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow x \quad \left(\frac{5}{2} + 1\right) = 3$$

$$1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow x \quad \left(\frac{6}{2} + 1\right) = 4$$

m1 m2
↑

⑧ Brute force :- Node * callBrute(head) {
int cnt = 0;

T C → O(2N * N/2) Node * temp = head;

S C → O(C1) while (temp) {
temp = temp → next;
cnt++; } } O(N)

int mid = cnt / 2 + 1;

temp = head;

while (temp) {

mid--;

if (!mid) return temp;

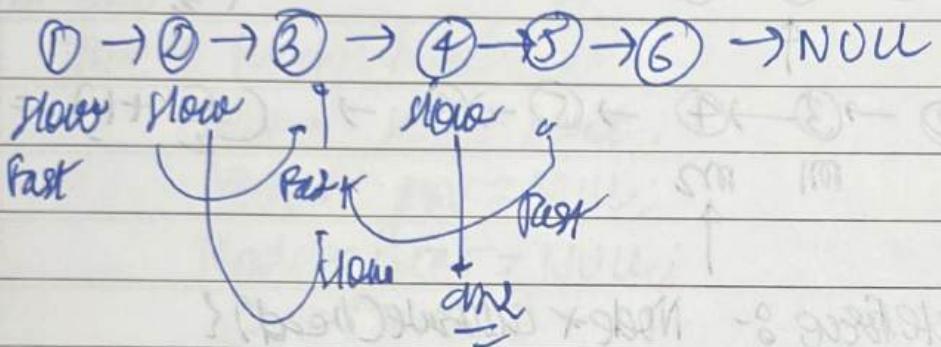
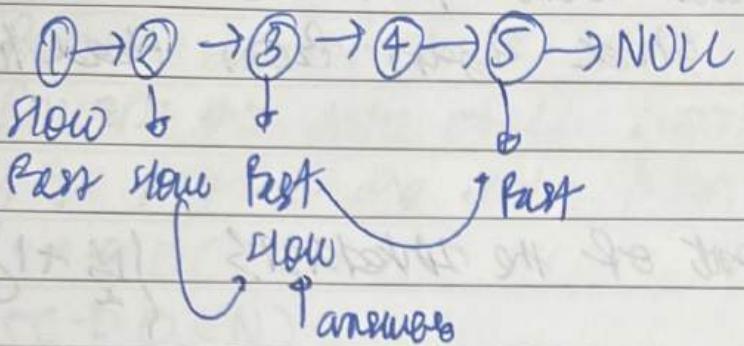
temp = temp → next; } } } O(N) ~~GEORGE~~

return head;

}

{ Tortoise & Hare }

Optimal :- { Tortoise and Hare }



Node* callOfNode (head)?

Nodes slow = head, fast = head;

while (Fast != NULL && Fast->next != NULL) {

slow = slow->next;

fast = fast->next->next;

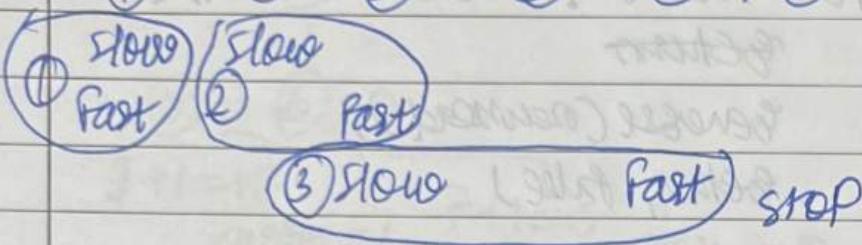
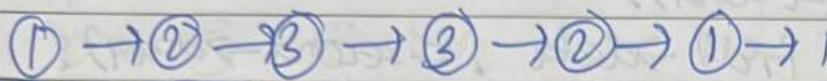
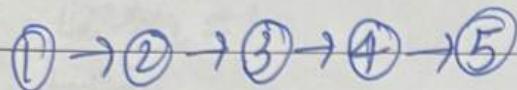


return slow;

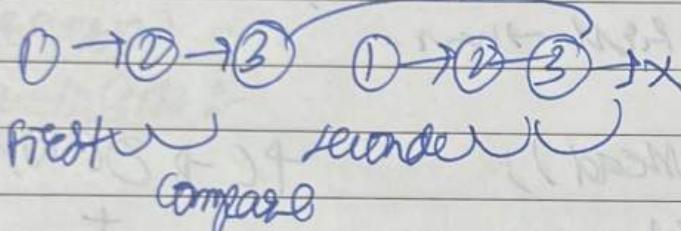
TC → O(N/2)

SC → O(1)

Palindrome?



reverse (slow \rightarrow next);



Now again reverse newHead

1 \rightarrow 2 \rightarrow 3 \rightarrow 2 \rightarrow 1 \rightarrow x } even length

1 \rightarrow 2 \rightarrow 3 \rightarrow 2 \rightarrow 1 \rightarrow x } odd length

1 \rightarrow 2 \rightarrow 3 \rightarrow 1 \rightarrow 2 \rightarrow x
newHead

Psuedo code:-

slow = head, fast = head

while (fast \rightarrow next != NULL and fast.next.next != NULL) {

slow = slow \rightarrow next;

fast = fast \rightarrow next.next;

\Rightarrow

Node x nowHead = reverse (slow \rightarrow next);

first = head

second = newHead

while(second)

if(first->data != second->data){

return

else

return false;

}

second = second->next;

first = first->next;

{

else

true;

{

+ $O(N^2)$

+

$O(N^2)$

+ $O(2N)$

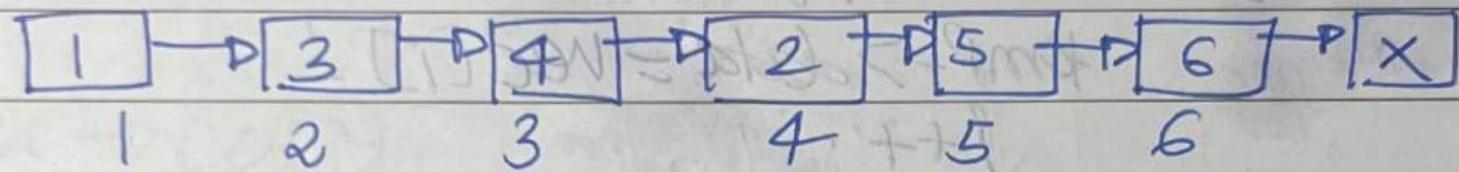
+ $O(Nh)$

$O(1)$

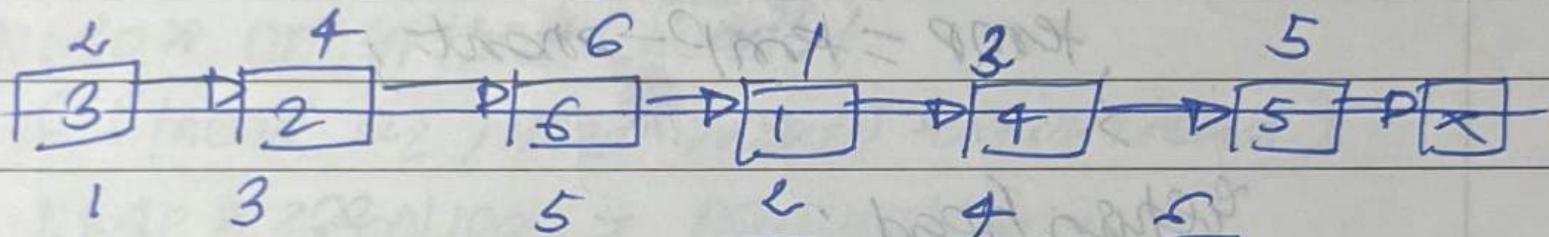
+ $O(Nh)$

⑧ { Legecgate Even and Odd }

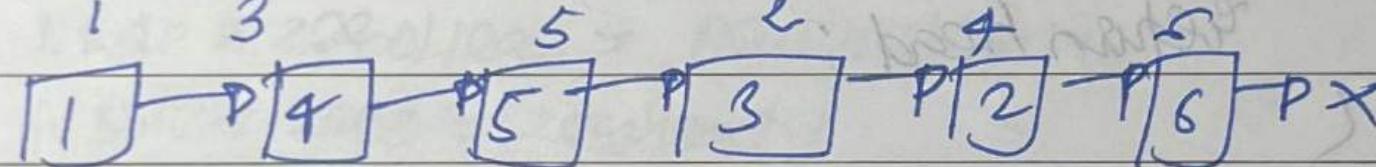
Input {



Output {



Output {



∅ create \rightarrow store in vector and can change the value in each mode

{Pseudo
code}

Node* callbrute(Node* head){

vector<int> vec;

Node* temp = head;

while(temp != NULL && temp->next != NULL)

{ vec.push_back(temp->data);

temp = temp->next->next;

}

if(temp){ vec.push_back(temp->data) }

TC \rightarrow O(2N)

SC \rightarrow O(N)

$\text{temp} = \text{head} \rightarrow \text{next}$;

while ($\text{temp} \neq \text{NULL}$ and $\text{temp} \rightarrow \text{next} \neq \text{NULL}$) {

 vec.push_back($\text{temp} \rightarrow \text{data}$);

 temp = $\text{temp} \rightarrow \text{next} \rightarrow \text{next}$;

}

if (temp) {

 vec.push_back($\text{temp} \rightarrow \text{data}$);

}

$\text{temp} = \text{head}$;

int i = 0;

while (temp) {

 temp->data = vec[i];

 i++;

 temp = temp->next;

return head;

}

optimal

$O(N)$

Node* callOPLinval (head) {

 Node* odd = head;

 Node* even = head->next;

 Node* memoiz = head->next;

 while (even != NULL and even->next != NULL) {

}

 odd->next = odd->next->next;

 even->next = even->next->next;

 odd = odd->next;

 even = even->next;

}

 odd->next = memoiz;

 return head;

⑤

{ Remove kth Node from end in 223

} Break;

Node* removekth(head, k) {

 int cnt = 0;

 Node* temp = head;

 while (temp) { cnt++; temp = temp->next; }

} edge } →

 if (cnt == k) { return head->next; }

 int one = cnt - k;

 temp = head;

 while (one)

while(temp){

ans--;

if (ans == 0) {

Node* del = temp->next;

temp->next = temp->next->next;

free(del);

break;

}

temp = temp->next;

3

return head;

5

TC $\rightarrow O(\log n) + O(\log n - N)$

SC $\rightarrow O(1)$

\Rightarrow OPTIMAL

Node* remove(head, k){

Node* fast = head;

Node* slow = head;

for (int i = 0 to k) { fast = fast->next; }

while(fast) { } if (fast == NULL) {

fast = fast->next,

return head->next; }

slow = slow->next,

{ }

if Node* del = slow->next)

slow->next = slow->next->next;

free(del);

return head;

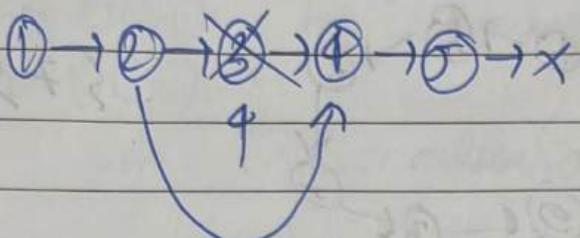
{ }

TC $\rightarrow O(kn)$

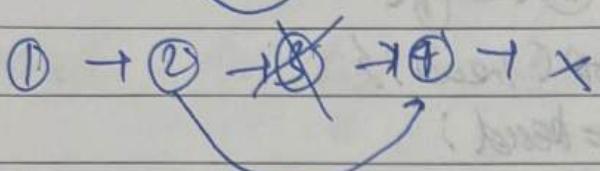
SC $\rightarrow O(1)$

① → ② → x

② } Delete the middle node of list 3



$$\text{middle node} \leq \left\lfloor \frac{N}{2} + 1 \right\rfloor$$



Node * calculate (Node * head) {

 if (!head || !head->next) {

 return NULL;

}

int count = 0;

O(N) {
 temp = head;
 while (temp) {
 cnt++;
 temp = temp->next;
 }

 cnt = cnt / 2;

 temp = head;
 while (temp) {

 cnt--;

 if (cnt == 0) {

 Node * middle = temp->next;

 temp->next = temp->next->next;

 free (middle);

 break;

}

 temp = last P. node;

{
 break; last = last->next;
}

 return head;

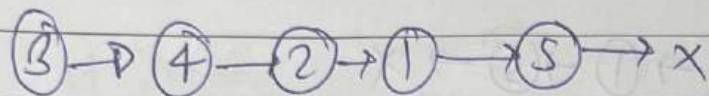
}

T -> O(N) + O(N/2)

S -> O(1)

②) Last LS

→



Node P call last(head){

arr = [];

temp = head;

while(temp){

arr.add(temp.data);

temp = temp.next;

}

last(arr))

temp = head, p = 0

while(temp){

temp = arr[p];

p++

temp = temp.next;

}

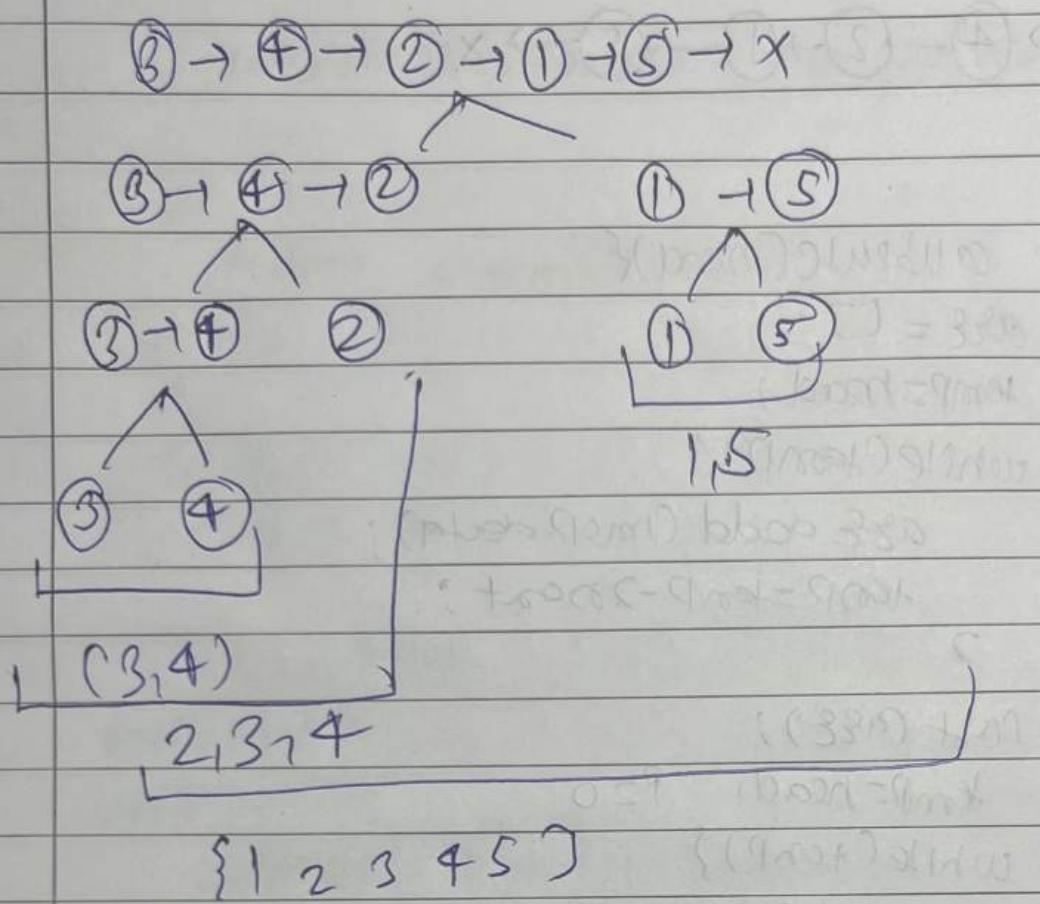
return head;

{

TC → O(N log N + N + N)

SC → O(N).

{ applying merge sort }



mergeSort(arr , low , $high$)

if ($low \leq high$) {

 return;

}

 mid = $(low + high) / 2$

 mergeSort(arr , low , mid);

 mergeSort(arr , $mid + 1$, $high$);

 merge(arr , low , mid , $high$);

}

if (head) {

 if (head == NULL || head->next == NULL) {
 return head;
 }

 middle = findMiddle(head);

 leftHead = head;

 rightHead = middle->next;

 middle->next = NULL;

 leftHead = f(leftHead);

 rightHead = f(rightHead);

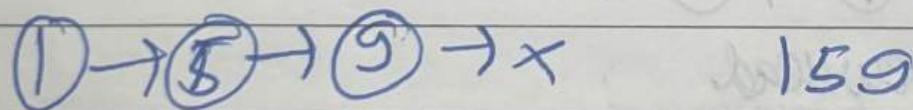
 return merge2(leftHead, rightHead);

}

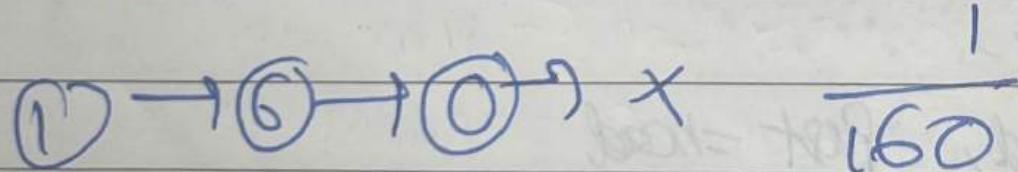
 └─ See me Platon LL Code

⑧ ↲ Add 1 to linkedList3

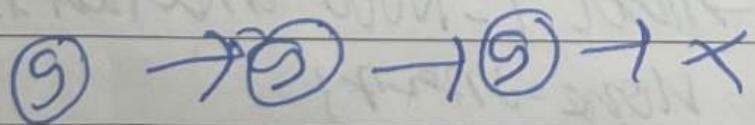
input



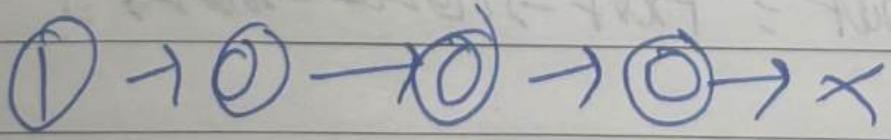
output



refugee



output



$0 \rightarrow 5 \rightarrow 5 \rightarrow *$

carry = 1

++.

$0 \rightarrow 5 \rightarrow 1 \rightarrow *$

$0 \rightarrow 6 \rightarrow$

$9+1=10$

carry = 1

$5+1=0$

carry = 0

break;

~~Step 2~~

Pseudo Code :-

Node* calculate (head)?

newhead = recursion(head)

int carry = 1

dummy node = new Node(-1),

temp = dummyNode,

while(newhead)

sum = carry;

if (newhead){

sum += newhead->data;

newhead = newhead->next;

}

dummyNode->next = new Node(sum % 10);

dummyNode = dummyNode->next;

sum / = 10;

carry = sum;

}

```
if (carry) {  
    dummyNode->next = new Node(carry);  
}
```

return recursion(temp->next);

}

```
int recursionOptimize (Node){
```

```
if (!temp) return 0;
```

```
int carry = recursionOptimize (temp->next);
```

```
temp->data = temp->data + carry;
```

```
if (temp->data < 0) return 0;  
return 1;
```

}

```
Node* callOptimal (head){
```

```
carry = recursionOptimize (head);
```

```
if (carry){
```

```
Node* newNode = new Node (carry);
```

```
newNode->next = head;
```

```
return newNode;
```

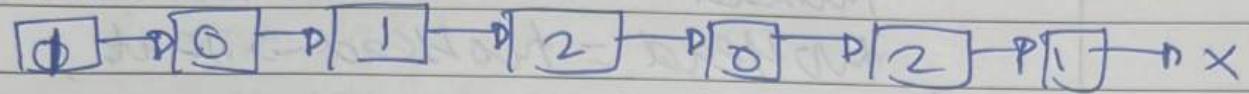
}

```
return head;
```

}

④ {Sort 0, 1, 2}

Brute:



$\text{cnt0} = 2$ } Replace data of LL as per }

$\text{cnt1} = 3$ } the count

$\text{cnt2} = 2$

$T C \rightarrow O(2N)$

$S E \rightarrow O(1)$

Optimal :-

$T C \rightarrow O(N)$

$S E \rightarrow O(1))$

Node * callOptimal (head) {

if (!head or !head->next) return head;

Node * zeroHead = new Node(-1);

Node * zero = zeroHead;

oneHead = new Node(-1);

one = oneHead;

twoHead = new Node(-1);

two = twoHead;

Node * temp = head;

while (temp) {

if (temp->data == 0) {

zeroHead->next = new Node(0);

zeroHead->next = zeroHead->next;

}

else if (temp->data == 1) {

oneHead->next = new Node(1);

oneHead = oneHead->next;

else {

twoHead->next = new Node(2),

twoHead = twoHead->next;

}

temp = temp->next;

zeroHead->next = (one->next) ? one->next :

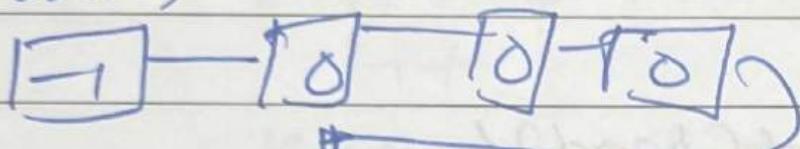
two->next,

oneHead->next = two->next;

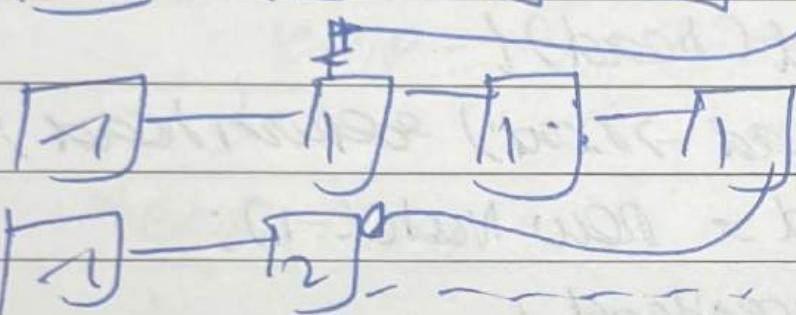
return zero->next;

#

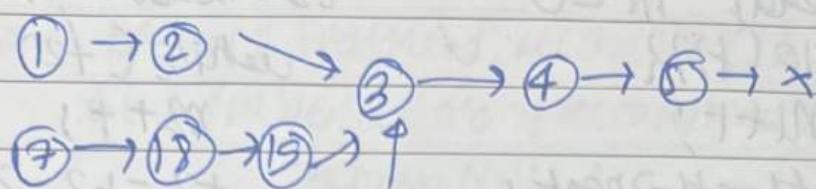
Approach



connect all 0



⑨ To find the intersection point of 4 linked-list



Brute Force :-

```
Node* calculateFirstHead, secondHead){
```

```
Node* ans = NULL;
```

```
unordered map<Node*, int> mapp;
```

```
while(firstHead){
```

```
    mapp[firstHead] = 1;
```

```
    firstHead = firstHead->next;
```

```
}
```

```
while(secondHead){
```

```
if(mapp.find(secondHead) != mapp.end()) {
```

```
    ans = secondHead;
```

```
    break;
```

```
}
```

```
secondHead = secondHead->next;
```

```
}
```

```
return ans;
```

```
}
```

Time Complexity :- $O(N_1 \times N_2) + O(N_2 \times 1)$

Space Complexity :- $O(N_1)$

To solve first heads

better
 { Optimal approach }

```

    t1 = head1, n1 = 0           t2 = head2, n2 = 0
    while (t1 != null)          while (t2 != null) {
        n1++;
        t1 = t1->next;
    }                            n2++;
    if (n1 < n2) {              t2 = t2->next;
        t1 = t1->next;          }
        return collisionPoint(head1, head2, n2 - n1);
    } else {                     n2 = n2 - n1;
        t2 = t2->next;          t1 = t1->next;
        return collisionPoint(head1, head2, n1 - n2);
    }
  
```

$O(N_2 - N_1)$

```

    Node* collisionPoint(t1, t2, diff) {
        while (diff--) {
            t2 = t2->next;
        }
        while (t1 != t2) {
            if (t1 == t2) return t1;
            else if (t1 == NULL) return NULL;
            else if (t2 == NULL) return NULL;
            t1 = t1->next;
            t2 = t2->next;
        }
        return NULL;
    }
  
```

$T.C \rightarrow O(N_1 + O(N_2) + O(N_2 - N_1) + O(N_1))$
 $\rightarrow O(N_1 + O(N_2))$
 $\rightarrow O(N_1 + 2N_2)$
 $SC \rightarrow O(1)$

{ Optimal Approach }

```
Node* Optimal( Prethead, SecondHead ) {  
    if ( C is Prethead or !SecondHead ) {  
        return NULL;  
    }
```

```
    Node* temp1 = Prethead;
```

```
    Node* temp2 = SecondHead;
```

```
    while ( temp1 != temp2 ) {
```

```
        temp1 = temp1->next;
```

```
        temp2 = temp2->next;
```

```
        if ( temp1 == temp2 ) return temp1;
```

```
        if ( !temp1 ) temp1 = SecondHead;
```

```
        if ( !temp2 ) temp2 = Prethead;
```

```
}
```

```
return temp1;
```

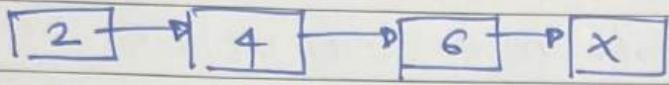
```
}
```

TC → O(N + N)

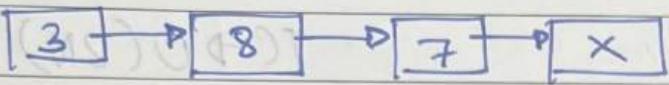
SC → O(1)

③ { Add two number in LL }

Input {

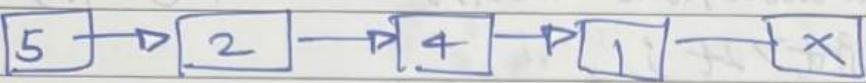


$$\begin{array}{r} 6 \ 4 \ 2 \\ + 7 \ 8 \ 3 \\ \hline \end{array}$$



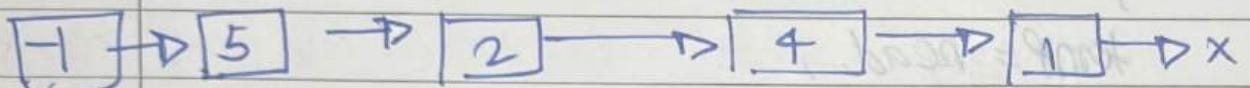
$$\begin{array}{r} 14 \ 2 \ 5 \\ \hline \end{array}$$

Output {



{ No need to reverse linked list }

$$\begin{array}{r} 2 \ 4 \ 6 \ x \\ + 3 \ 8 \ 7 \ x \\ \hline \end{array}$$



function {
dummy ↗ } carry=0 carry=1 carry=1 carry=0
{
 ↑
 function
 dummy ↗ }
 12+0 13+1
 =12 =14

Pseudo code:- Node* func(Node* n1, Node* n2) {

 int carry = 0 , sum = 0 ;

 Node* dummyNode = new Node(-1) ;

 Node* current = dummyNode ;

 while(num1 != NULL || num2 != NULL) {

 sum = carry ;

 if (num1) { sum += num1->data ; num1 = num1->next ; }

 if (num2) { sum += num2->data ; num2 = num2->next ; }

 Node* temp = new Node(sum % 10) ; { num2->next ; }

 current->next = temp ;

 current = current->next ;

 carry = sum / 10 ;

```
if (carry) {
```

```
    current->next = new Node(carry % 10);
```

```
    return dummyNode->next;
```

```
}
```

```
current->next = NULL;
```

```
return dummyNode->next;
```

```
}
```