

Supporting Protocols

ARP and ICMP

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Recap

- Forwarding needs IP to MAC address mapping
 - Service provided by ARP protocol
- Network layer needs to provide means for debugging (error signaling) and for router-host communication (determine MTU size, indicate better routes, provide netmask info etc)
 - Service provided by ICMP protocol

Problem Statement

- IP layer forwarding is based on IP addresses
 - Next-hop delivery based on Link addresses (MAC)
 - Need to perform IP to MAC address translation
 - Answer: Address Resolution Protocol (ARP)
- what layer ?*
- How do you ensure ARP process gets the relevant packets? → demux
- what address should the frame carry ?
- what messages would you send & how do you act on a message received message ?

Address Resolution Protocol (ARP)

- Operates at Link layer (Frame type = 0x0806)

but it doesn't do any framing, MAC, it is handled by ethernet.
so ARP is like a companion protocol

- Based on broadcast: What is the MAC address corresponding to given IP address?"



- Host with matching IP address replies
- Each host maintains a cache with IP to MAC translations
 - Entries in cache timed out periodically (15 min)

to account for changes in ip-MAC mappings

ARP Packet Format

hardware : ethernet
Protocol : ip

0	8	16	31
Hardware Type (=1)		Protocol Type (=0x0800)	
HLEN (=48) <small>MAC</small>	PLEN (=32) <small>IP</small>	Operation <small>request, reply</small>	
Source Hardware Address (Bytes 0-3)			
Source Hardware Address (Bytes 4-5)		Source Protocol Address (Bytes 0-1)	
Source Protocol Address (Bytes 2-3)		Target Hardware Address (Bytes 0-1)	
Target Hardware Address (Bytes 2-5) <small>- ?</small>			
Target Protocol Address (Bytes 0-3) <small>✓</small>			

Numbers in brackets capture mapping
IP addresses to Ethernet addresses

Address Resolution Protocol (ARP)

- Originator: Add entry to cache corresponding to target
↳ ARP request
- Target: Add entry to cache corresponding to the originator (sender)
ARP reply
- Intermediate hosts: Refresh existing entries
- When forwarding a datagram, check cache, if no mapping, invoke ARP

Example

(A): MAC corresponding to IP of B

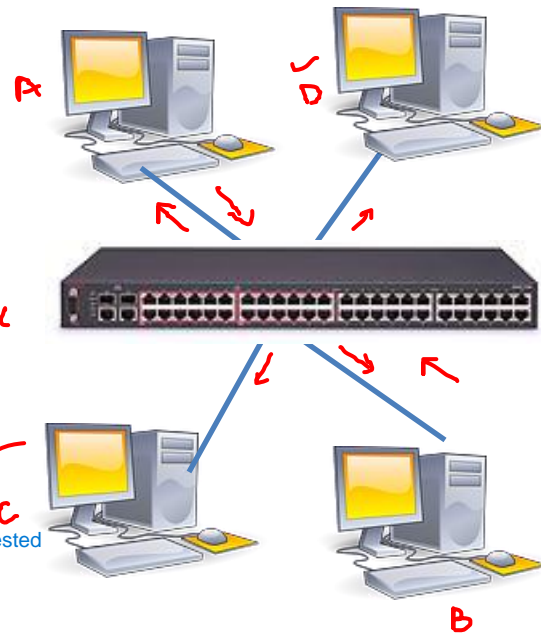
ARP request
↓
broadcast

MAC header
↓
Payload

src add	dst add		
src MAC	src IP	Target MAC	Target IP
A	A	?	B

ff:ff:...

Cache A
BIP → B MAC



ARP reply
from B
↓
"unicast"

MAC →
↓
Payload

src add	B	dst add	A
src MAC	src IP	Target MAC	Target IP
B ← B	B	A	A

no reason that other hosts would be interested in B's address

Cache B
AIP → A MAC

C, D ARP request
↓
A → refresh (next time)
no entry A, ignore ARP request

C, D won't have anything to do with ARP reply

C, D

Address Resolution Protocol (ARP)

- Originator: Add entry to cache corresponding to target *IP address, MAC*
- Target: Add entry to cache corresponding to the originator (sender)
- Intermediate hosts: Refresh existing entries
- When forwarding a datagram, check ^{ARP} cache, if no mapping, invoke ARP

Gratuitous ARPs

- Generated by a host to inform others of its IP to MAC mapping

→ broadcast
MAC header
src MAC
↑
host
dest
ff:ff:ff:ff:ff:ff
own

- Could be a request or reply

- If request, no reply will occur
- If reply, there was no preceding request
- Source IP = destination IP = IP of machine generating gratuitous ARP

Ethernet payload i.e, ARP data

ARP data
src, Target
IP=?, IP=?
src → MAC
MAC

- Target MAC = ?

ff:ff:ff → request
MAC → reply } window

Uses of Gratuitous ARPs

- Issued whenever IP or MAC address of an interface changes or brought up from down state
 - Help rectify cached ARP entries
 - Report IP address conflicts (duplicate IP)
if u manually configure a dup ip, then a grat ARP, will reach the host with that ip, and it replies that u have used a dup ip... so u better change it.
 - Inform bridges of the location of new host

ICMP: Internet Control Message Protocol

- Used by hosts & routers to communicate network-level information
 - Error reporting: unreachable host, network, port, protocol
 - Diagnostic purposes: Echo request/reply (used by ping)
 - Routing: Source quench

↗ data
↖

if source sends too much data rate than it can handle.

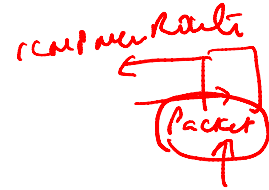
ICMP Packet Format

ICMP protocol operates at Network Layer,,,

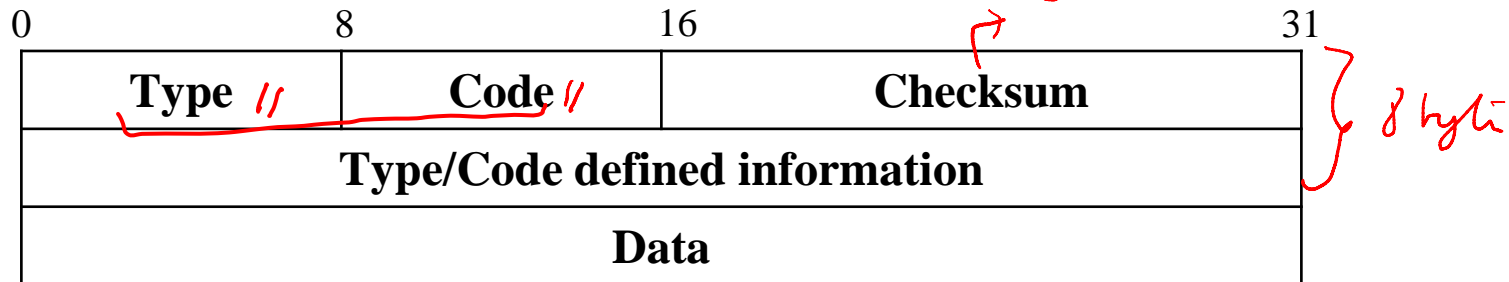
- ICMP messages carried in IP datagrams
- 8 bytes of header followed by data.
- Data field in error messages carry
 - entire IP header and first 8 bytes of data of IP packet that caused the error

not all ICMP msgs contain data, it is contained only in the error messages


demux TCP = 6
IP → ICMP
demux key: 1
if demux key is 1, then pass packet to the ICMP protocol



entire ICMP message



Select ICMP Messages



Type	Code	Description
0	0	Echo Reply (Ping)
3	0	Destination network unreachable
3	1	Destination host unreachable
3	3	Destination port unreachable
3 /	4 /	Fragmentation required, DF flag set
3	6	Destination network unknown
3	7	Destination host unknown

if we send a large packet, and still say DontFragment ... then it drops the packet and sends us error message

Select ICMP Messages

Type	Code	Description
4	0	Source Quench
5	0	Redirect datagram for the network
8	0	Echo request (<u>Ping</u>)
11	0	<u>TTL</u> expired
12	0	Bad IP header
13	0	Timestamp
14	0	Timestamp reply
17	0	Address mask request
18	0	Address mask reply ←

these not error messages

Example: Fragmentation Required

0	8	16	31
Type=3	Code=4	Checksum	
Unused		Next hop MTU ✓	
IP header and first 8 bytes of original datagram's payload			

Traceroute

- Source sends series of UDP segments to destination one after another

it keeps on sending till it gets a reply.

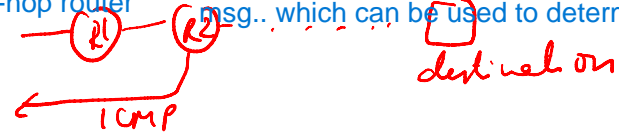
- First has TTL=1

The packet is dropped by the very next-hop router



and the next-hop-router sends back an ICMP msg.. which can be used to determine the route

- Second has TTL=2, etc.

- Destination port is set to an unlikely number



Traceroute

- When n^{th} datagram arrives to n^{th} router:
 - Router discards datagram
 - Sends to source an ICMP message (type 11, code 0)
 - Message includes name of router & IP address
- For each ICMP message, sending host notes router id and RTT time 
- Sending host stops when it gets ICMP message (type 3, code 3) 

Summary

- Studied two useful protocols: ARP and ICMP
- ARP is needed for forwarding
 - Performs IP to MAC address translation
- ICMP helps with error reporting and host signaling