

IPv6

Kameswari Chebrolu

All the figures used as part of the slides are either self created or from the public domain with either 'creative commons' or 'public domain dedication' licensing. The public sites from which some of the figures have been picked include: <http://commons.wikimedia.org> (Wikipedia, Wikimedia and workbooks); <http://www.sxc.hu> and <http://www.pixabay.com>

IPv4 - 1981

Background

- Early 1990's CIDR, NAT proposed
 - Temporary fixes; Not possible to achieve 100% efficiency
- Mid 1990's: Next Generation IP (IPng) – IPv6
 - Apart from addressing, fix other aspects of the protocol based on experience

IPv5

Desirable Features

- Support billions of hosts in a scalable fashion
- Allow fast processing at routers
- Support real-time applications
- Provide security
- Multicast support
- Mobility support
- Need to be backward compatible

VC



Extension Headers

$$\begin{array}{r} \text{HL} \\ 2^4 \times 4 \\ = 64 \text{ by } 6 \\ 20 \\ = 44 \end{array}$$

- Next Header field replaces both options and ‘upper-layer protocol field’ of IPv4

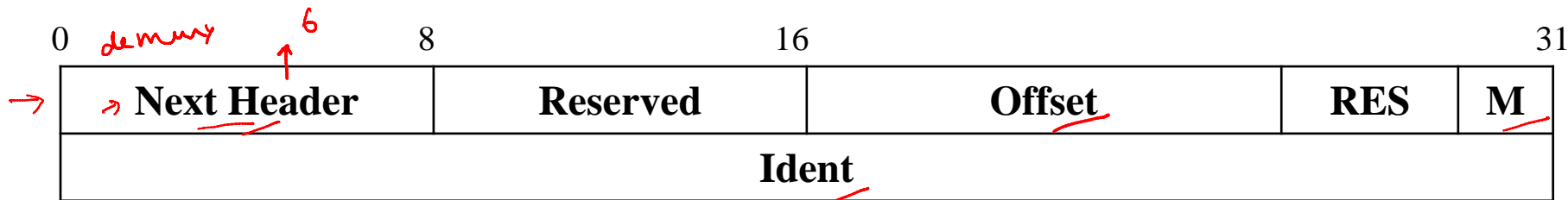
In IPv4 what is the maximum length the options field can occupy? Specify in Bytes.

- Structure improves router performance
- Can support arbitrary length options (IPv4 restricted to be under 44 bytes)

HL counts in multiples of 4 ... 4bits ==>
16 combinations x 4 = 64 bytes -
20byte header = 44

- Each option has an ‘extension header’
 - Next Header field within indentifies the header following it

IPv6 Fragmentation Extension Header



40 bytes | Next header = 44
 → Next header/data

- Assume only one option that of fragmentation
- Next header field in Ipv6 header will take value 44 to indicate fragmentation header
- Next header in fragmentation header will take the value 6 to indicate pass to TCP
- * Only source host does fragmentation, not routers

Points to Note

- 128 bit addresses can support $3 * 10^{38}$ hosts
- Fast router processing
 - Streamlined header of 40 bytes
 - No checksum, no fragmentation
- Support for real-time applications via traffic class and flow label

Points to Note

- Other features handled via options field
- ICMP extended for IPv6
 - Packet too big
 - Multicast, mobility support

Intermission



Addressing

- 128 bits $\rightarrow 3 * 10^{38}$ nodes $\leftarrow 2^{128}$
 - Consider entire surface of earth; $7 * 10^{23}$ IP addresses per square foot
 - $4.354 \pm 0.012 \times 10^{23}$ micro seconds since Big Bang
- Notation: x:x:x:x:x:x:x:x $202.13.5.6$
 - X is hexadecimal representation of 16 bit piece of address
 - E.g: 2001:0DB8:0000:0000:95CD:BBE0:000B:0001
 - Short form: 2001:DB8::95CD:BBE0:B:1

short form leading zeros can be eliminated

- Classless addressing /41 132 . . .
- Number of addresses with special meaning

Prefix	Usage
<u>00...0</u> (128 bits)	Unspecified
00...1 (128 bits)	loopback
<u>1111 1111</u>	Multicast
1111 1110 10	<u>Link local unicast</u>
<u>::ffff:0:0/96</u>	IPv4 mapped IPv6 addresses

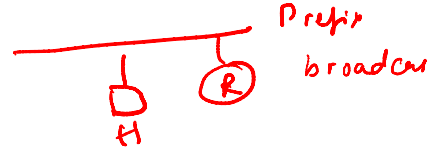
0's ← 1's

Sample Set

- Routing very similar to IPv4 except for some new extension routing header
 - Can specify which provider network to use for which packets

flow → PN1 cheap
 flow → PN2 E

Autoconfiguration



- In IPv4 done via DHCP servers
- IPv6: Stateless auto configuration without servers
 - Need unique IP address, need correct address prefix *netmask*
- Solution: Routers announce prefix; Host autoconfigures address as: prefix 00..00 Ethernet-MAC-addr
globally routable addr *few bytes* *6 bytes* *128 bits*
- Globally not routable: 1111 1110 10 0...0 Ethernet-MAC-Addr
not routable ... it is a link local addr, can be used to communicate within ur physical network *128 bits*

Transition from IPv4 to IPv6

- Impossible for a flag-day
- Incremental deployment of IPv6
 - IPv4 nodes don't know anything about IPv6
 - IPv4 nodes should be able to talk with other IPv4 nodes and IPv6 nodes ↗ IPv4
.....here v6 nodes can also support v4 .. bcz v6 already know v4
 - IPv6 nodes should be able to talk with other IPv6 nodes over intermediate IPv4 nodes
- Solution: Dual stack operation and Tunneling

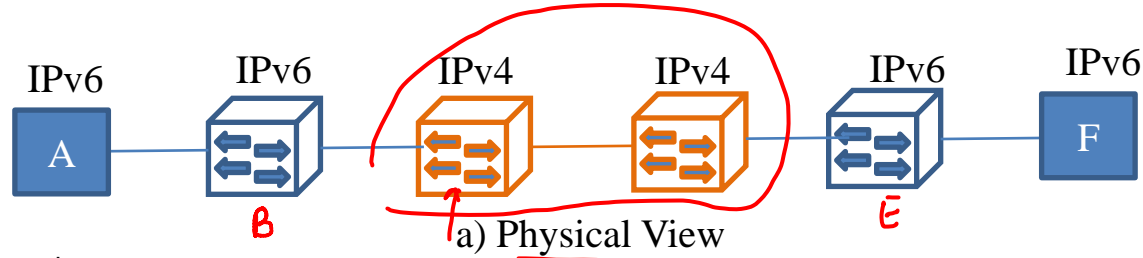
Dual Stack Operation

- IPv6 nodes run both IPv4 and IPv6 and use version field to call the right process

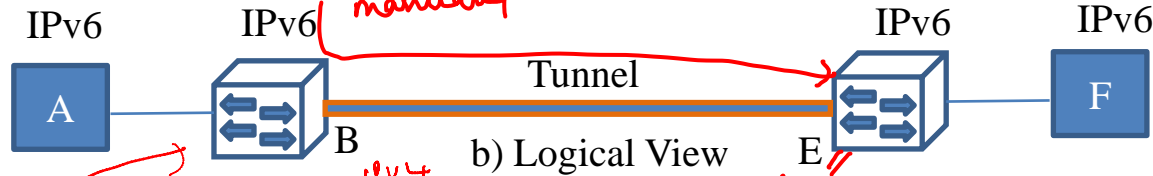


Tunneling

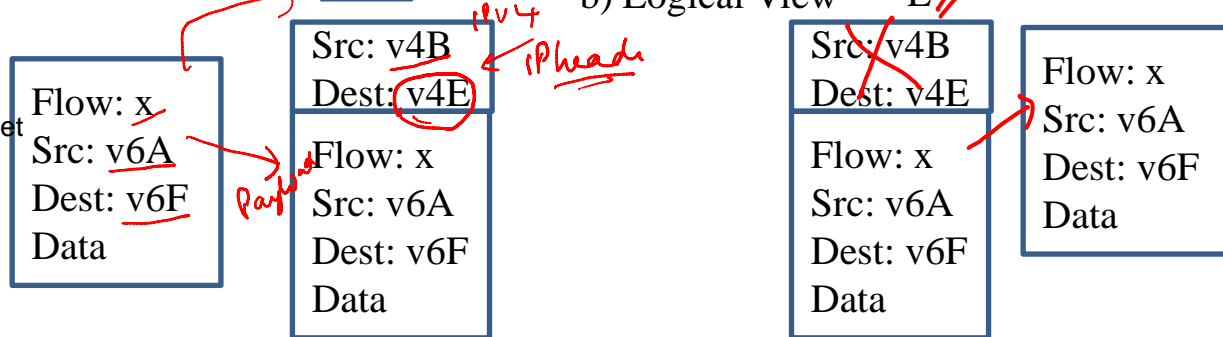
- Send IP packet as a payload of another IP packet



All this tunnel info is manually configured.
like how does B know to give Dest as E ??



IPv4 cannot interpret these fields like Flow : x etc etc



Summary

- IPv6 long term solution to IPv4 address exhaustion
- Addresses other shortcomings of IPv4
- Many interesting features
- Migration via Dual-stack operation/Tunneling
- As of 2011, few RIRs have exhausted their IPv4 address space
- As of Nov 2012, IPv6 share of Internet traffic is 1%