CSA0389 - DATA STRUCTURE FOR STACK PMPLEMENTATION

ASSIGNMENT-2

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BHUVANESHWARI. S 192311375 BF-CSE \$ 40 delete 201

Describe the Abstract data type (ADT) and how they differ from concrete data structures. Design on ADT for a stack and implement it using arrays and Linked 18st in c. Include operations Like Push. Pop, Peek, 1s empty, is full and Peek.

Abstract Data Type (ADT):

An Abstract Data Type (ADT) is a theoretical model that defines a set of operation and the behavior of there operations on a data structure, without specifying how the data structure is should be implemented. It provides a high level description of what operations can be performed on the data and what constraints apply to those operations.

characteristics of ADT's:

* operations: Defines a set of operations that can be performed on the data structure.

* semantics: especifies the behaviour of each operation.

* Enapsulation: Hides the implementation details, forwing on the interface provided to the wer.

ADT for stack:

A stack is a functionental data structure that follows the Last 2h, first out (L8FO) principle. It supports the following operations.

Push: Adds on element to the top of the stack

Pop: Removes and meturn the element from the top of the stack.

Peek: Return the element from the top of the stack without removing it.

is empty: checks if the stack is empty

is tull: checks if the astack is full.

Concrete data structure:

The implementation using arrays and Linked 19sts are specific ways of implementing the stack ADT inc.

How ADT differ from concrete data structures:

ADT focuses on the operations and their behavior, while concrete data structures focus on how these operations are realized using speafic programming constructs larrays are Minked 18sts).

Advantages of ADT:

By reparating the ADT from its implementation, you achieve modularity, encapsulation and blexibility in designing and using data structures in programs. This reparation allows for easier maintenance, code neme and abstraction of the complex operations.

Implementation in c wing Array.

include / studio. hz

int main () of Stack-Array stack; Stack . top = -19

stack. [tems [++ stack . top] = 10;

Stack . I tems [++ Stack, top] = 20; stack. Ptems [++ stack. top]=300,

```
(stack . top 1 = -1) {
     print + ("Top element: % d \n", stack. Ffems [stack.top]);3
 else d
      print f (" stack Ps empty: \n"); }
 [ ( stack , top 1 = -1) of
      print+ ("popped element: %d \n", stack. "tem [stack.top=]);
else of
     printy ("stack underflow: \n");
1 (stack . top! = -1) {
       printf ["popped element: "/. d(h", stack. Ptem [stack. top-]);
 3 ches
   print ("stack undeflow: \n");
1 (stack . top 1 = -1) of
      printf (" top element after pops: "din", stack. Thems[stack =
                                                 f ([got
elie 9
     point { [ stack is empty: (h");
```

```
Implementation in c using
 # include < studio. h>
  int main () &
      Nocle * top = NULL;
      Node * newnode = (Nocle+) malloc (size of (Node));
  9) (newNode = = NULL){
      prints ("Memory allocation failed: \n"):
      return 1;
  new Node -) data = 10°,
  new Node -> next = top;
 top = new Node ;
  newnode = (Node+) malloc (size of (Node));
 96 (nownade == NULL) 5
    prints ("Memory allocation foiled: \n");
    refun 1;
 top = top > next ;
 Free (temp);
4 else of
    printf ("stack underflow: \n");
if 140p! = NULL) &
     printf [" Top element after pops: "/.d \n", top-) data); 3
else q
    printf["stack is empty: \n"); g
while (top ! = NULL) &
     Node * temp = tor;
     top = top -) next 1
     tree (temp)?
   return og
```

for placement training. The student xxx, reg. no. 20142010 whites to check whether Ps name Ps Proted or not. The 18st is not sorted in any order.

Linear Search:

Linear wearch works by checking each element in the list one by one until the desired element is tound or the end of the list is oneached. It is a simple wearching technique that doesn't require any prior sorting to the data.

Steps for Linear rearch.

- 1) start the first element
- 2) check of the current element is equal to the target element.
- 3) If the current element is not the target, move to the next element in the list.
- to und or you neach the end of the 18st.
- 6) Its the tanget Ps found, one furn Pts position. If the end of the 1Pst Ps neached and the element has not been tound, Indicate that element is not present.

Procedure

Given the 19st:

- D start at the first element of the fist.
- a) compare '20142010' with '20142015' (first element), 20142033 (second element) these are not equal.

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```
3) compane '20142010' with '20142010' [fifth element).
 4) The element '20142010' Ps found at the bilth position
They are equal.
(Pindex 4) Pn the 18st.
c code for linear search:
# Producte 2stdPo. hs
 Put main () 5
     int regno [] = { 3;
     int tanget = 20142010;
     Put n = size of (regno) /size of (regno[o]);
  int found = 0;
 for (i=0; " cn , i++) q
    if [regno[t]= = tanget) &
       printf ("kogistration number %. d found at index 4.d./
                                    taget ,i);
    break ;
if (!found) &
     printf ["Rogistration number %. d not tound in Pst. In",
                                            target) 9
  Jeturn o;
```

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Pseudocode for stack operations: Puh (elements): Pb stack is full: print "stack overflow" add element to the top of the wtack Increment top pointer 2) POP (): Pl stack is empty: print (" stack underflow") return null for appropriate onor value) olse ? remove and return element from the top of stack decrement and pointer. 3) Peck (): If stack is empty: Print " stack is empty" return null (or appropriate error value) ohe: ale return element at the top of the stack (without removing it) 4) 8 Empty (): return true of top 95 -1 (stack is empty) otherwise, netwon take F) & Full: return true, 9 top 95 equal to maxsize -1 (stack is full) otherwise, neturn take.

Explanation ob Pseudocade: # Adds an element to the top of the stack checks if the stack is full before pushing. * Removes and returning the element from the top of the stack checks if the stack is empty before popping. * Returns the element at the top of the stack without removing of checks of the stack or empty before * checks 1/ the stack is empty by inspecting the top peeking. pointer or equivalent variable * checks of the stack is full by comparing the top pointer or equivalent variable to the max size of stack