习题: 假定全局关系和分段模式如下:

全局关系

Doctor (D#, Name, SLDEPT)

Patient (P#, Name, SLDEPT, Treat, D#)

Care (P#, Drug, QTY)

分段模式

Doctor1 = SLDEPT= 'Surgery' (Doctor)

Doctor2 = SLDEPT= 'Pediatrics' (Doctor)

Doctor3 = SLDEPT= 'Surgery' and SLDEPT= 'Pediatrics' (Doctor)

Patient1 = SLDEPT= 'Surgery' and treat= 'intensing' (Patient)

Patient2 = SLDEPT= 'Surgery' and treat = ' intensing' (Patient)

Patient3 = SLDEPT= 'Surgery' (Patient)

Care1 = Care SJ p#=p# Patient1

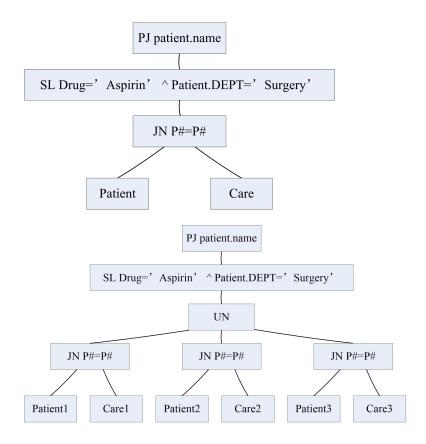
Care2 = Care SJ p#=p# Patient2

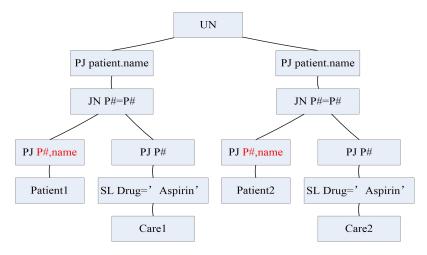
Care3 = Care SJ p#=p# Patient3

使用变换准则,将下面的全局查询变化为分段查询,并对它们加以简化。当需要时用限 定关系代数消除查询中不用的段。

(a) 列出在 Care 中使用'Aspirin'的 Patient 名字;

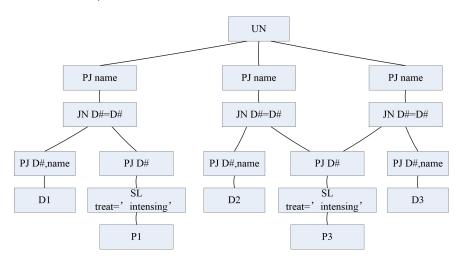
$$\pi_{\textit{Patient.Name}} \left(\sigma_{\textit{Care.Drug}='\textit{Aspirin}' \land \textit{Patient.DEPT}='\textit{Surgery}'} \left(\textit{Patient} \quad \textit{Care} \right) \right)$$





(b) 列出给病人开出接受精细护理的医生名字。

$$\pi_{{\scriptscriptstyle Doctor.Name}}\left(\sigma_{{\scriptscriptstyle Patient.Treat}\, = {\scriptscriptstyle Intensing}} \left({\scriptscriptstyle Doctor} \propto_{{\scriptscriptstyle Doctor.D}\, \# = {\scriptscriptstyle Patient}\, {\scriptscriptstyle D\#}} {\scriptscriptstyle Patient} \right) \right)$$



1.假设两个事务 T 和 U 的 log 记录如下所示:

<T,Start>

<T,A,10,25>

<U,Start>

<U,B,20,15>

<T,C,30,20>

<U,D,40,30>

<U,Commit>

<T,E,50,70>

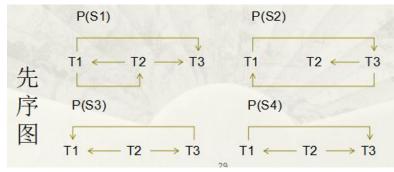
<T,Commit>

如果系统故障时,磁盘上记录的 Log 记录如下,请描述数据库恢复管理器的动作. a) < U, Start>

- b) < U, Commit>
- c) <T,E,50>
- d) <T,Commit>
- (a)undo U, undo T (不能是 undo T,undo U)
- A 还原成初始值
- (b)undo T, redo U
- AC 还原成初始值, BD 分别为 20、40
- (c)undo T, redo U
- ACE 还原成初始值, BD 分别为 20、40
- (d)redo T, redo U
- ACE 分别为 10、30、50, BD 分别为 20、40
- 1. Undo 是逆向完成
- 2. Redo 是顺序完成
- 3. 必须先 Undo 后 Redo
 - 1. 指出下面的冲突等价调度(忽略提交C和夭折A动作)
 - S1=
 - $W_2(x),W_1(x),R_3(x),R_1(x),C_1,W_2(y),R_3(z),C_3,R_2(x),C_2$ \$2=
 - $R_3(z), R_3(y), W_2(y), R_2(z), W_1(x), R_3(x), W_1(x), R_1(x), C_1, C_2, C_3$ S3=
 - $R_3(z), W_2(x), W_2(y), R_1(x), R_3(x), R_2(z), R_3(y), C_3, W_1(x), C_2, C_1$ S4=
 - $\mathsf{R}_3(z), \mathsf{W}_2(x), \mathsf{W}_2(y), \mathsf{C}_2, \mathsf{W}_1(x), \mathsf{R}_1(x), \mathsf{A}_1, \mathsf{R}_3(x), \mathsf{R}_3(z), \mathsf{R}_3(y), \mathsf{C}_3$

一个冲突等价调度如下:

- S1' = W2(x),W1(x),R1(x),R3(x),C1,W2(y),R3(z),C3,R2(x),C2
- S2' = R3(z), R2(z), R3(y), W2(y), W1(x), R3(x), W1(x), R1(x), C1, C2, C3
- S3' =W2(x),W2(y),R2(z),C2,R3(z),R3(x),R3(y),C3,R1(x),W1(x),C1
- S4' = W2(x), W2(y), C2, W1(x), R1(x), A1, R3(z), R3(x), R3(z), R3(y), C3



S3 和 S4 为可串调度