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hw1_search_q13_a*_cscs

Question 13: A*-CSCS

2/2 points (ungraded)

Recall that a dictionary, also known as a hashmap, works as follows:

Inserting a key-value pair into a dictionary when the key is not already in the dictionary adds the pair to the dictionary:

```
\begin{array}{l} \textit{dict} \leftarrow \text{ an empty dictionary} \\ \textit{dict}[\text{``key''}] \leftarrow \text{``value''} \\ \\ \textit{print } \textit{dict}[\text{``key''}] \\ \rightarrow \text{``value''} \end{array}
```

Updating the value associated with a dictionary entry is done as follows:

```
dict["key"] \leftarrow "new value"
print dict["key"]
\rightarrow "new value"
```

We saw that for A^* graph search to be guaranteed to be optimal the heuristic needs to be consistent. In this question we explore a new search procedure using a dictionary for the closed set, A^* -graph-search-with-Cost-Sensitive-Closed-Set (A^* - CSCS).

```
function A*-CSCS-GRAPH-SEARCH(problem, fringe, strategy) return a solution, or failure

closed ← an empty dictionary
fringe ← INSERT(MAKE-NODE(INITIAL-STATE[problem]), fringe)

loop do

if fringe is empty then return failure

node ← REMOVE-FRONT(fringe, strategy)
if GOAL-TEST(problem, STATE[node]) then return node

if STATE[node] is not in closed or COST[node] < closed[STATE[node]] then

closed[STATE[node]] ← COST[node]

for child-node in EXPAND(node, problem) do
fringe ← INSERT(child-node, fringe)
end
end
```

Rather than just inserting the last state of a node into the closed set, we now store the last state paired with the cost of the node. Whenever A^* - CSCS considers expanding a node, it checks the closed set. Only if the last state is not a key in the closed set, or the cost of the node is less than the cost associated with the state in the closed set, the node is expanded.

For $oldsymbol{regular} oldsymbol{A^*}$ $oldsymbol{graph}$ search which of the following statements are true?
$\ lue{}$ If $m{h}$ is admissible, then $m{A^*}$ graph search finds an optimal solution.
$lacksquare$ If $m{h}$ is consistent, then $m{A^*}$ graph search finds an optimal solution.
Submit
✓ Correct (2/2 points)
problem 2/2 points (ungraded) In each of the following parts, select all true statements about A^* -cscs
$lacksquare$ If $m{h}$ is admissible, then $m{A^*}$ - CSCS finds an optimal solution.
$ extcolor{black}{ lack} extcolor{black}{ If h is consistent, then A^*- CSCS finds an optimal solution.}$
Submit
✓ Correct (2/2 points)
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2/2 points (ungraded)

If $m{h}$ is admissible, then $m{A^*}$ - CSCS will expand at most as many nodes as $m{A^*}$ tree search.
If h is consistent, then A^* - CSCS will expand at most as many nodes as A^* tree search.
Submit
✓ Correct (2/2 points)
problem
2/2 points (ungraded)
$lacksquare$ If $m{h}$ is admissible, then $m{A^*}$ - CSCS will expand at most as many nodes as $m{A^*}$ graph search.
If h is consistent, then A^* - CSCS will expand at most as many nodes as A^* graph search.
Submit
✓ Correct (2/2 points)
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