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Q4: Value Functions

Problem 4: Value Functions

Consider a general search problem defined by:

- A set of states, S.
- A start state s_0 .
- A set of goal states G, with $G \subset S$.
- A successor function Succ(s) that gives the set of states s' that you can go to from the current state s.
- For each successor s' of s, the cost (weight) W(s,s') of that action.

As usual, the search problem is to find a lowest-cost path from the start state s_0 to a goal $g \in G$. You may assume that each non-goal state has at least one successor, that the weights are all positive, and that all states can reach a goal state in a finite number of steps.

Define C(s) to be the *optimal cost* of the state s; that is, the cost of the lowest-cost path from s to any goal state. For $g \in G$, clearly C(g) = 0.

Part 1

4/4 points (ungraded)

Select the Bellman-style (one-step lookahead) equation that expresses $C\left(s\right)$ for a non-goal state \boldsymbol{s} in terms of the optimal costs of other states.

$$igcup C\left(s
ight) = \sum\limits_{s' \in \mathrm{Succ}\left(s
ight)} W\left(s, s'
ight)$$

$$igcup C\left(s
ight) = \min_{s' \in \mathrm{Succ}\left(\mathrm{s}
ight)} W\left(s, s'
ight)$$

$$egin{aligned} & C\left(s
ight) = \sum\limits_{s' \in \mathrm{Succ}\left(s
ight)} \left[W\left(s, s'
ight) + C\left(s'
ight)
ight] \end{aligned}$$

$$ullet C\left(s
ight) = \min_{s' \in \mathrm{Succ}\left(\mathrm{s}
ight)} \left[W\left(s, s'
ight) + C\left(s'
ight)
ight] imes$$

$$igcup C\left(s
ight) = \max_{s' \in \operatorname{Succ}\left(s
ight)} \left[W\left(s, s'
ight) + C\left(s'
ight)
ight]$$

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✓ Correct (4/4 points)

Part 2

2/2 points (ungraded)

Consider a heuristic function $h\left(s\right)$ with $h\left(s\right)\geq 0$. What relation must hold between $h\left(s\right)$ and $C\left(s\right)$ for $h\left(s\right)$ to be an admissible heuristic?

$$lacksquare h\left(s
ight)=C\left(s
ight), orall s\in S$$

$$ullet h\left(s
ight) \leq C\left(s
ight), orall s \in S$$

$$igcup h\left(s
ight) \geq C\left(s
ight) ,orall s\in S$$

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✓ Correct (2/2 points)

Part 3

By analogy to value iteration, define $C_k\left(s\right)$ to be the minimum cost of any plan starting from s that is either of length k or reaches a goal state in at most k actions. Define $C_0\left(s\right)=0$ for all s. $C\left(s\right)$ is defined as before to be the cost of the lowest-cost path from s to any goal state, without any restriction on the number of actions starting from s.

Suppose we use C and C_k as heuristic functions.

Part 3.1

2/2 points (ungraded)

Which of the following statements regarding the admissibility of $m{C}$ and $m{C_k}$ are true? Check all that apply.

- lacksquare C_k might be inadmissible for any given value of k.
- $ule{C}_{k}$ is admissible for all k.
- lacksquare Let L be the length of the longest path from a state to a goal. C_k is only guaranteed to be admissible if $k \geq L$.



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Correct (2/2 points)

Part 3.2

2/2 points (ungraded)

Which of the following statements regarding the consistency of ${\it C}$ and ${\it C}_k$ are true? Check all that apply.

- lacksquare C_k might be inconsistent for any given value of k.
- lacksquare Let L be the length of the longest path from a state to a goal. C_k is only guaranteed to be consistent if $k \geq L$.
- $ule{\hspace{-0.1cm} \hspace{-0.1cm} \hspace{-0.1cm} \hspace{-0.1cm} \hspace{-0.1cm} \boldsymbol{C}}$ is consistent.



Submit

✓ Correct (2/2 points)

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