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## Q6: A\* Search: Batch Node Expansion

### Problem 6: A\* Search: Batch Node Expansion

#### Part 1

0.0/8.0 points (ungraded)

Recall that A\* graph search can be implemented in pseudo-code as follows:

```
1: function A*-GRAPH-SEARCH(problem, fringe)
2:   closed  $\leftarrow$  an empty set
3:   fringe  $\leftarrow$  INSERT(MAKE-NODE(INITIAL-STATE[problem]), fringe)
4:   loop do
5:     if fringe is empty then return failure
6:     node  $\leftarrow$  REMOVE-FRONT(fringe)
7:     if GOAL-TEST(problem, STATE[node]) then return node
8:     if STATE[node] is not in closed then
9:       add STATE[node] to closed
10:    child-nodes  $\leftarrow$  EXPAND(node, problem)
11:    fringe  $\leftarrow$  INSERT-ALL(child-nodes, fringe)
```

Now consider the following batch version of A\* graph search.

```

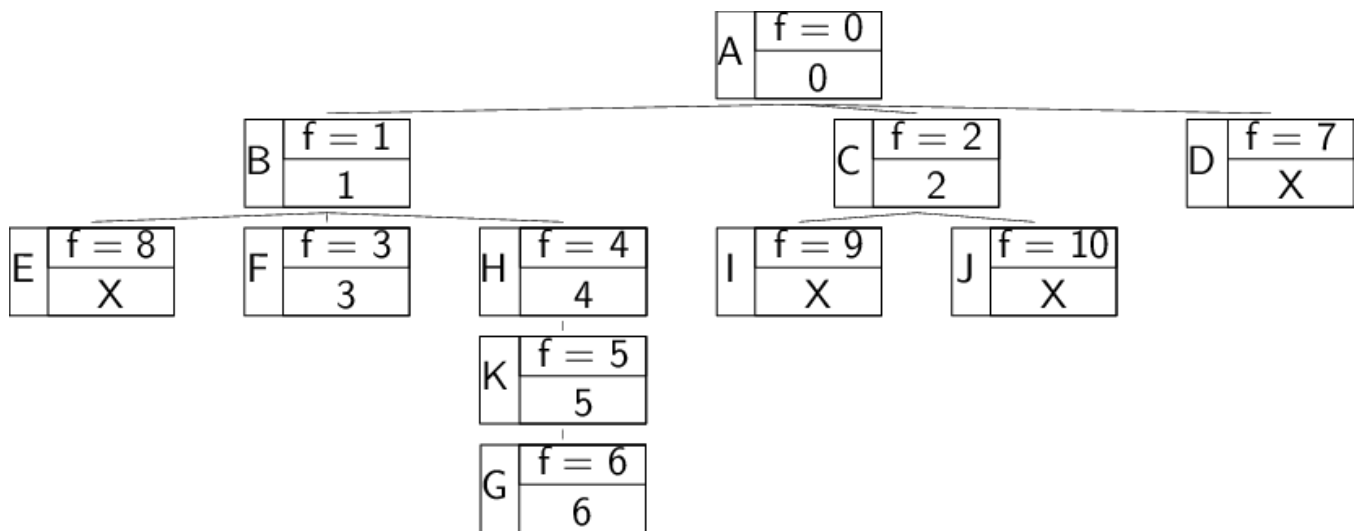
1: function A*-BATCH(problem, fringe, k)
2:   closed  $\leftarrow$  an empty set
3:   fringe  $\leftarrow$  INSERT(MAKE-NODE(INITIAL-STATE[problem]), fringe)
4:   loop do
5:     i  $\leftarrow$  0
6:     next-set  $\leftarrow$  an empty set
7:     while i < k and fringe is not empty do
8:       next-set  $\leftarrow$  INSERT(REMOVE-FRONT(fringe))
9:       i  $\leftarrow$  i + 1
10:    for node in next-set do
11:      if GOAL-TEST(problem, STATE[node]) then
12:        return node
13:      if STATE[node] not in closed then
14:        closed  $\leftarrow$  INSERT(STATE[node])
15:        child-nodes  $\leftarrow$  EXPAND(node, problem)
16:        fringe  $\leftarrow$  INSERT-ALL(child-nodes, fringe)

```

Rather than process the nodes from the fringe one at a time, this code instead pulls ***k*** nodes off the fringe, and processes them in a random order.

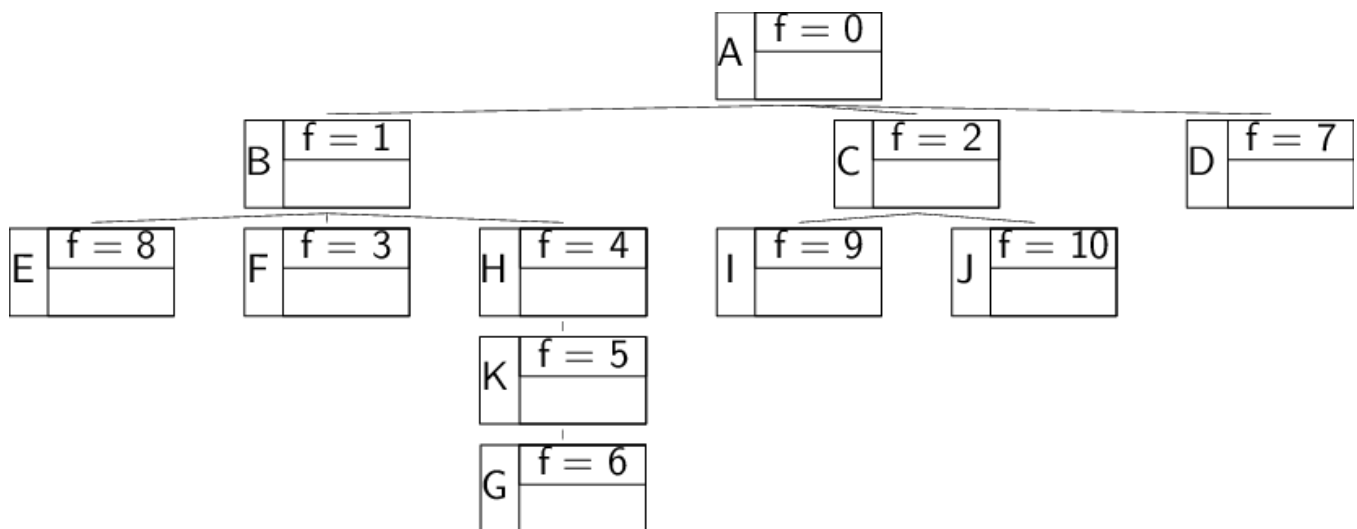
For this problem, we are interested in determining the iteration of the outermost loop at which each node is expanded.

Consider running this new A\*-**BATCH** algorithm on a search problem with the search tree shown in the diagram below with ***k* = 1**. Each node is drawn with the state at the left, the *f*-value at the top-right ( $f(n) = g(n) + h(n)$ ), and the iteration at which a node was expanded is at the bottom-right, with an 'X' if that node was not expanded. G is the unique goal node. In the diagram below, we can see that the start node A was expanded during iteration 0, then node B was expanded during iteration 1, node C was expanded during iteration 2, node F was expanded during iteration 3, node H was expanded during iteration 4, node K was expanded during iteration 5, and node G was expanded during iteration 6. Nodes D,E,I,J were never expanded.



In this question you'll complete similar diagrams by filling in the node expansion iterations for the cases of  $k = 2$  and  $k = 3$ . Note that now multiple nodes can (and typically will!) be expanded during any given iteration.

For each node below, fill in the node expansion iterations for the case of  $k = 2$  and fill in an 'x' for any node that is not expanded.



A:  Answer: 0

B:  Answer: 1

C:  Answer: 1

D:  Answer: 3

E:  Answer: x

F:  Answer: 2

H:  Answer: 2

I:  Answer: x

J:  Answer: x

K:  Answer: 3

G:  Answer: 4

Submit

You have used 0 of 2 attempts

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**i** Answers are displayed within the problem

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## Part 2

0.0/8.0 points (ungraded)

For the same tree, fill in the node expansion iterations for the case of  $k = 3$  and fill in an 'x' for any node that is not expanded.

A:  Answer: 0

B:  Answer: 1

C:  Answer: 1

D: 1

Answer: 1

E: 2

Answer: 2

F: 2

Answer: 2

H: 2

Answer: 2

I: 3

Answer: 3

J: 3

Answer: 3

K: 3

Answer: 3

G: 4

Answer: 4

Submit

You have used 0 of 2 attempts

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**i** Answers are displayed within the problem