

Assignment3

Question1

Consider the directed acyclic graph G in Figure 3.10. How many topological orderings does it have?

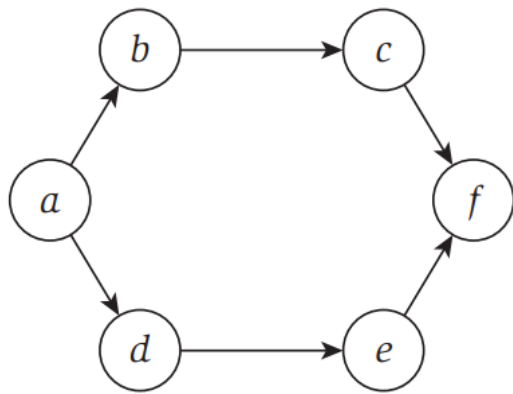


Figure 3.10 How many topological orderings does this graph have?

There 6 topological orderings this graph have

1. a b d c e f
2. a b d e c f
3. a b c d e f
4. a d b c e f
5. a d b e c f
6. a d e b c f

Question3

The algorithm described in Section 3.6 for computing a topological ordering of a DAG repeatedly finds a node with no incoming edges and deletes it. This will eventually produce a topological ordering, provided that the input graph really is a DAG.

But suppose that we're given an arbitrary graph that may or may not be a DAG. Extend the topological ordering algorithm so that, given an input directed graph G , it outputs one of two things: (a) a topological ordering, thus establishing that G is a DAG; or (b) a cycle in G , thus establishing that G is not a DAG. The running time of your algorithm should be $O(m + n)$ for a directed graph with n nodes and m edges.

Here is the pseudocode of the judgement:

```
1 // G: A directed graph
2 DAG_Judge(G){
3     Create queue Q; // for topological order
4     Create ArrayList A; // record the topological answer
5     // O(n)
6     for (all Nodes in G){
7         if (there is not nodes which in-degree is 0){
8             return("Not a DAG");
9         }else{
10             if(node in-degree == 0){
11                 Q.add(node);
12             }
13         }
14     }
15     // O(n + m)
16     while(Q is not null){
17         node = Q.pop(); // will probably pop n nodes
18         A.add(node);
19         count += 1;
20         // will probably pass m edges
21         for(outnode in all the outnodes of node){
22             outnode in-degree -= 1;
23             if(outnode in-degree == 0){
24                 Q.add(outnodes);
25             }
26         }
27     }
28     // Judgement
29     if(A.size() == G node number){
30         return(A);
31     }else{
32         return("Not a DAG");
33     }
34 }
```

