# Lecture 9 Graphs and trees

**COS10003 Computer Logic and Essentials (Hawthorn)** 



Semester 1 2021

# **Today**

- Foundation
- Paths and cycles
- 3 Algorithms
- 4 Search
- **5** Trees
- 6 Applications

Defining graphs and trees

Exploring graphs with special properties

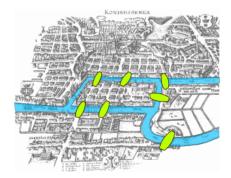
Algorithms for shortest paths

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#### Let's start with a problem

In 1736 a town in Prussia had seven bridges as shown. How is it possible to walk around the city crossing each bridge exactly once?



https://commons.wikimedia.org/wiki/File:Konigsberg\_bridges.png, CC SA 3.0

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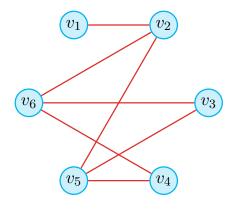
# What is a graph?

A graph G consists of two sets:

- ► A set V whose elements are called vertices, points or nodes.
- ► A set E of unordered pairs of distinct vertices, called edges or links.

We denote a graph by G(V, E) when we want to emphasise the two parts of G.

#### **Visualising**



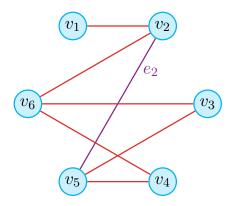
$$V = \{v_1, v_2, v_3, v_4, v_5, v_6\}$$

$$E = \{\{v_1, v_2\}, \{v_2, v_5\}, \{v_2, v_6\}, \{v_3, v_5\}, \{v_3, v_6\}, \{v_4, v_5\}, \{v_4, v_6\}\}$$

#### **More definitions**

- Vertices u and v are said to be adjacent if there is an edge e between them, that is (u, v).
- ► The edge e is said to be incident to u (and v).
- ► We can have multiple or parallel edges between the same vertices, however this then creates a multigraph.
- ► The degree of a vertex is the number of incident edges it has.

#### For you to do

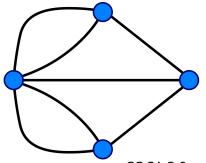


- ▶ The degree of  $v_4$ ? **2**
- ▶ The degree of  $v_5$ ?
- $ightharpoonup v_1$  is adjacent to?
- $ightharpoonup v_3$  is adjacent to?
- $ightharpoonup e_2$  is incident to?

# Revisiting our problem

In 1736 a town in Prussia had seven bridges. How is it possible to walk around the city crossing each bridge exactly once?

The problem becomes how to traverse every edge once. Can this be done?



https://commons.wikimedia.org/wiki/File:Konigsburg\_graph.svg, CC SA 3.0

#### Before we continue

Graphs are an example of how data can be represented for computation purposes.

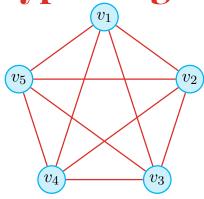
- Computers that are connected in some way in a network
- ▶ People who work together or are linked on a website
- Representation of possible paths to take for a robot or (autonomous) vehicle
- Allocation of resources such as timetabling or scheduling

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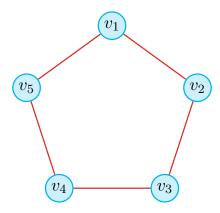
# More on degree

The degree sum formula says that the sum of degrees in a graph G is twice the number of edges, that is,  $\sum_{v \in V} \deg v = 2|E|$ .

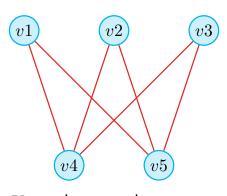
Types of graphs



 $K_n$ : the complete graph with n vertices ( $K_5$  shown)



 $C_n$ : the cycle on n vertices ( $C_5$  shown)



 $K_{m,n}$ : the complete bipartite graph with sets of m and n vertices ( $K_{3,2}$  shown)

# **Bipartite graphs**

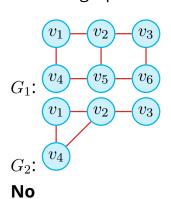
A bipartite graph is a graph for which it is possible to divide the vertices into two disjoint sets such that there are no edges between any two vertices in the same set.

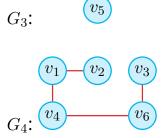
### **Subgraphs**

A subgraph G'(V', E') of a graph G(V, E) is a graph such that  $V' \subseteq V$  and  $E' \subseteq E$ , and the vertices incident to any edges in E' are in V'.

# For you to do

Which graphs are subgraphs of  $G_1$ ?





 $v_2$ 

 $v_3$ 

 $v_1$ 

### Isomorphic graphs

Two graphs are isomorphic if they are the same except for the naming of the vertices.

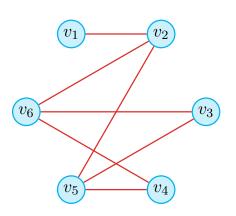
#### Look for:

- same number of vertices and edges
- same number of vertices with certain degrees
- same cycles

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#### **Paths**

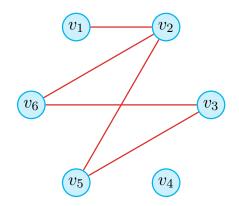
A path through a graph is an alternating sequence of n+1 vertices and n edges, so that edge  $e_i$  is incident to vertices  $v_{i-1}$  and  $v_i$ .



What is a path from  $v_1$  to  $v_3$ ? ( $v_1$ ,  $\{v_1, v_2\}$ ,  $v_2$ ,  $\{v_2, v_6\}$ ,  $v_6$ ,  $\{v_6, v_3\}$ ,  $v_3$ )
This is a path of length 3; count the number of edges.

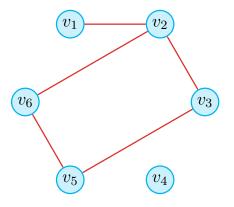
#### **Connected graphs**

A graph G is connected if for any vertices v and w in G, there is a path from v to w.



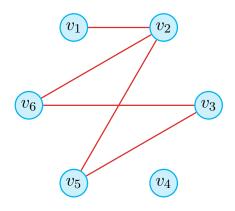
#### **Cycles**

A cycle is a path of non-zero length from vertex v to itself, with no repeated edges. A simple cycle has no repeated vertices apart from v at the start and end.



There is a cycle ( $v_2$ ,  $v_3$ ,  $v_5$ ,  $v_6$ ,  $v_2$ ).

#### For you to do



- Find a path between  $v_1$  and  $v_5$ . ( $v_1$ ,  $\{v_1, v_2\}, v_2, \{v_2, v_5\}, v_5$ )
- ▶ What is the length of the path?
- Find a path between  $v_2$  and  $v_4$ .
- Find the longest simple path between  $v_1$  and  $v_5$ ?
- ightharpoonup Find a cycle from  $v_3$ .

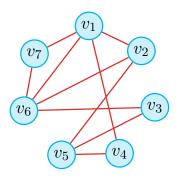
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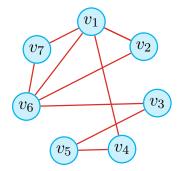
# **Eulerian cycle**

An Eulerian cycle is a cycle that uses all the edges in the graph exactly once and starts and ends at the same vertex.

For this to work, the graph must be connected and each vertex must have an even degree.

#### **Examples**





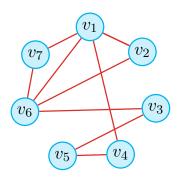
# Hamiltonian cycles

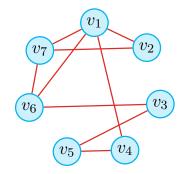
A Hamitonian cycle is a cycle where each vertex is visited exactly once, apart from the starting and ending vertex which is the same.

This is considered to be a very difficult problem to solve; there are no definite tips on properties to look for.

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#### **Examples**





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#### **Directed graphs**

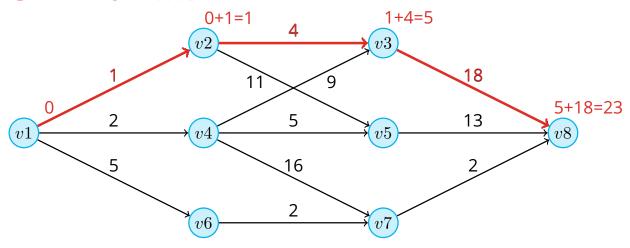
- In many case we want to work with edges that have a direction (a source vertex and a sink vertex).
- ► These are often weighted as well, signifying some sort of cost to traverse those links.
- ▶ This weight could be distance, cost, or time for example.



# **Shortest path**

Given any two nodes, find the path joining them of minimum weight. Assume all weights are positive.

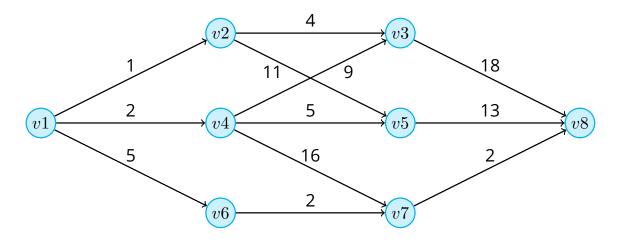
#### A greedy approach



# Dijkstra's algorithm

- 1. Mark all nodes unvisited. Create a set of all the unvisited nodes called the unvisited set.
- 2. Assign to every node a tentative distance value: set it to zero for our initial node and to infinity for all other nodes. Set the initial node as current.
- 3. For the current node, consider all of its unvisited neighbours and calculate their tentative distances through the current node. Compare the newly calculated tentative distance to the current assigned value and assign the smaller one. Otherwise, keep the current value.
- 4. When we are done considering all of the unvisited neighbours of the current node, mark the current node as visited and remove it from the unvisited set. A visited node will never be checked again.
- 5. If the destination node has been marked visited (when planning a route between two specific nodes) or if the smallest tentative distance among the nodes in the unvisited set is infinity (when planning a complete traversal; occurs when there is no connection between the initial node and remaining unvisited nodes), then stop. The algorithm has finished.
- 6. Otherwise, select the unvisited node that is marked with the smallest tentative distance, set it as the new "current node", and go back to step 3.

#### Dijkstra in action



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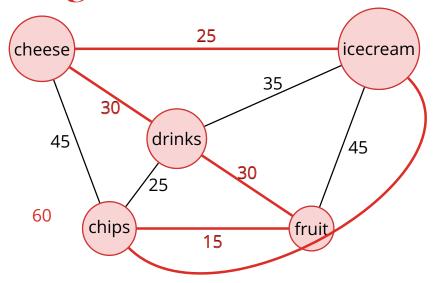
# **Another problem**

How can we visit each vertex exactly once and return to the starting point with minimum total weight?

This is known as the travelling salesman problem.

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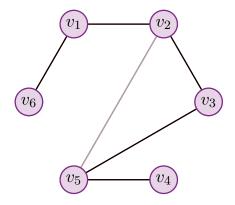
#### **Travelling Salesman Problem**



# Types of algorithms

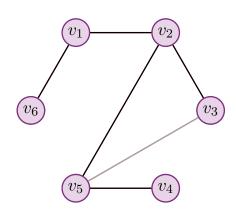
- Exact algorithms produce the correct solution, however tend to work only for small problems due to prohibitive runtime.
  - Enumerate every path (TSP)
  - Dijkstra's algorithm
- Heuristics have better runtimes and will produce a solution, but possibly not the correct solution.
  - Greedy shortest path
  - ▶ Nearest neighbour

# Depth-first search



- ightharpoonup start at  $v_1$ , see  $v_2$  and  $v_6$
- ightharpoonup move to  $v_2$ , see  $v_3$  and  $v_5$
- lacktriangle move to  $v_3$ , see  $v_5$
- ightharpoonup move to  $v_5$ , see  $v_4$
- ightharpoonup move to  $v_4$ , no more nodes to visit
- ightharpoonup move back to  $v_1$  and follow  $v_6$

#### **Breadth-first search**

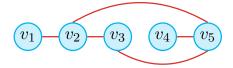


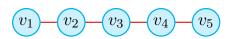
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- ightharpoonup see  $v_6$
- ightharpoonup move to  $v_3$
- ightharpoonup move to  $v_5$ , see  $v_4$
- ightharpoonup move to  $v_4$ , no more nodes to visit

#### What is a tree?

- A tree is a graph T where there is a unique simple path between all pairs of vertices u and v in T.
- ▶ Another way of saying this is that a tree with n vertices has n-1 edges and is connected.
- Trees are acyclic, that is, they have no cycles.
- ► Trees are commonly used in computer science for representing the organisation of data or for decision problems.

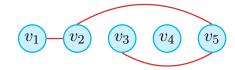
#### Which of these are trees?





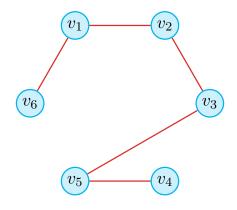
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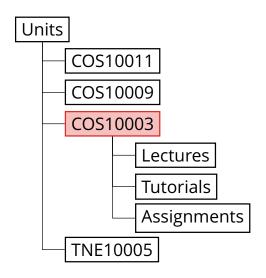
# **Spanning tree**

A spanning tree is a subset of a graph with n-1 edges.



#### **Rooted tree**

A rooted tree is a tree with one element named as the root element. The remaining elements are in a hierarchy.

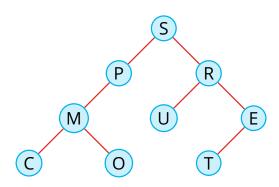


#### Tree traversal

- Preorder
- ▶ In-order
- Postorder

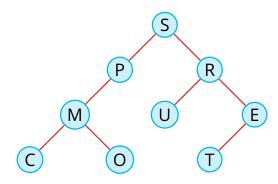
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#### **Pre-order**



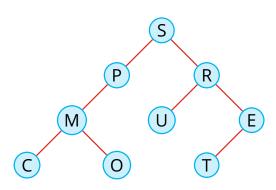
```
pre-order(node):
   if the current node is empty
     return
   else
     print current node
     pre-order(left node)
     pre-order(right node)
```

#### In order



```
in-order(node):
   if the current node is empty
     return
   else
     in-order(left node)
     print current node
   in-order(right node)
```

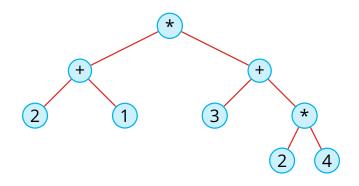
#### **Post-order**



```
post-order(node):
   if the current node is empty
     return
   else
     post-order(left node)
     post-order(right node)
     print current node
```

#### For you to do

What is the result of traversing this tree using post-order traversal?



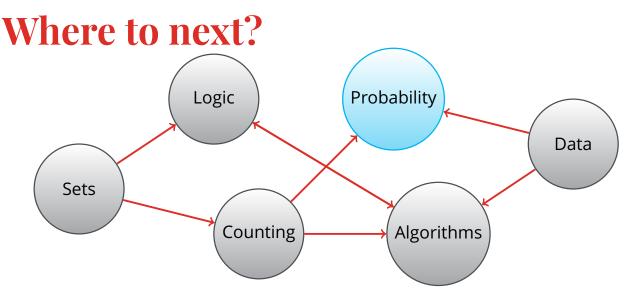
# Some applications

- ► Path planning and optimisation problems, e.g., for robotics, non-autonomous vehicles
- Puzzles and games
- Graph databases: more obhect-oriented than relational
- Medicine, biology, neuroscience: disease contagion, DNA, brain wiring

### Reflecting

- ▶ What are the main components of graphs?
- What are graphs useful for?
- ▶ How can you find the shortest path between two nodes?
- What are the specific characteristics of a tree?





In which we look at probability for computing.

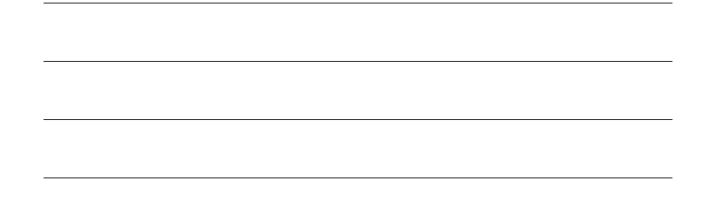
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# **Questions I still have**



# **Topics I need to review**

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