MySql Concrete Architecture (Storage Management)

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Team

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Introduction

- This presentation illustrates the divergences between the conceptual and concrete architecture of a system.
- Our topic of interest is the storage management component of the MySQL Server.
- The learning outcome of this presentation is a better understanding of why certain differences exist between architectures.

Layout

• Subsystems & Interactions

Derivation Process

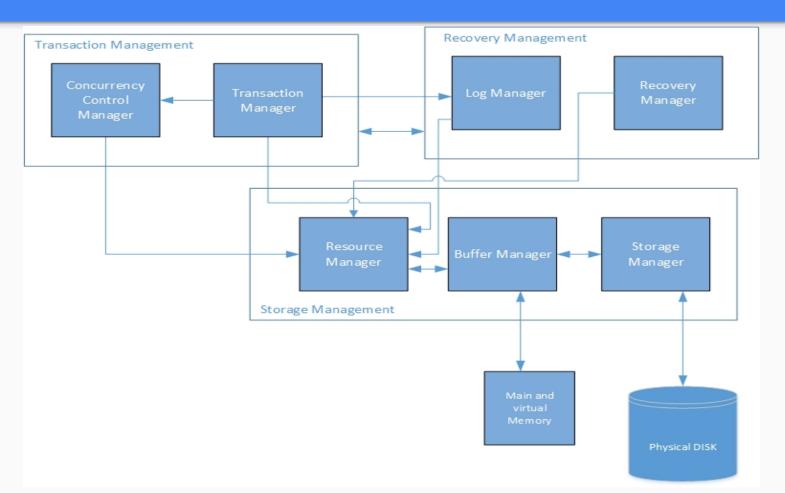
Architecture Styles & Design
 Patterns

Storage Management

- Architecture Differences
- LsEdit Screenshots
- Use Case

Lessons Learned

Subsystems and Interactions



Derivation Process

- Visualized interactions using Isedit.
- Simplified the view of Isedit by removing unneeded file dependences.
- Analyzed the code structure and file names and tried to map file names to components.
- If filename wasn't clear we went inside the code and analyzed the comments.
- When we found the right files, we contained and visualized them.

Some Important files in MySQL Server Folder

client - command line client utility code.

sql - main folder of MySQL it contains the parser and optimizer, execution engine and handler to choose a storage engine and many other things.

Sql-common - files that used by client and server.

Storage - The storage engines like Inndob and MyISSAM implemented here.

Vio - virtual input output contains implementation of low level network I/O code.

Architecture Styles and Design Patterns

- Layered architecture (top-level and Storage management internals)
- Pipe-and-filter architecture (query processing)
- Client-server architecture (application layer)
- Facade design pattern. The file mysqld.cc act as an interface between the client and the other components.

Storage Management - Storage Manager

Storage Manager

- Responsible for providing fast read/write access to tablespaces and logs of the database.
- fil0fil.cc located in the fil directory contains the implementation of the lowlevel file system. Contains the methods and structures needed to manipulate them.
- fsp0fsp.cc located in the fsp directory takes care of the file space management, which handles the allocation of pages and keeps track of which files are used, open, and closed.

Storage Management - Buffer Manager

Buffer Manager

- The Buffer Manager is responsible for efficiently storing the data in memory for manipulation
- it accepts requests from the Resource Manager and decides how much memory to allocate.
- The buffer manager is located in the buf directory.
- buf0buf.cc file is the buffer pool, buf0flu.cc flush the buffer pool and buf0lru.cc contains the replacement algorithm that decides which blocks should be shifted back to disk.
- buf0rea.cc calls the storage manager to initiate a file read.

Storage Management - Resource Manager

Resource Manager

- The responsibility of the Resource Manager is to accept requests from the higher level layers and translate them into an appropriate format that can be understood by the Buffer Manager.
- The resource manager is found inside the dict directory.
- It can be seen from the diagram that there is a direct interaction between storage manager and resource manager where the resource manager skips the buffer and accesses the storage manager directly.

Architecture Differences: Sticky Notes

Which: dict0dict.cc (ResourceManager) -> fil0fil.cc (StorageManager)

fil_make_filepath

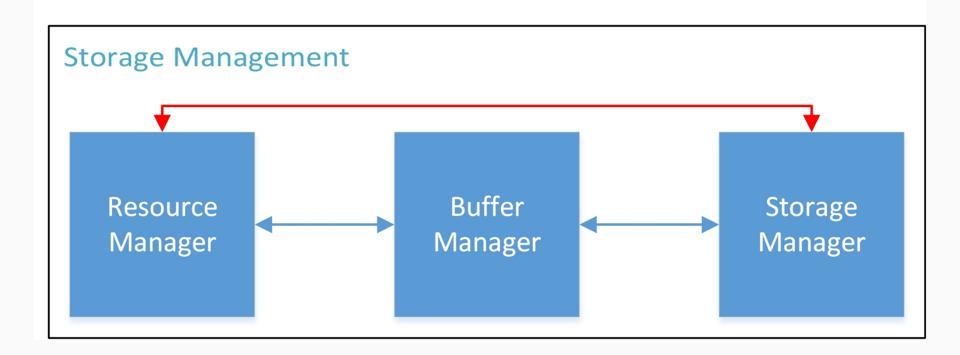
Who: Kevin Lewis

When: Jan 17, 2014

Why: Refactoring. Consolidating several functions that make a filename.

Legend

- → Convergence
- → Divergence



3751		<pre>- filepath = os_file_make_remote_pathname(</pre>
3752		<pre>- table->data_dir_path, table->name, "ibd");</pre>
	3752	+ filepath = fil_make_filepath(
	3753	+ table->data_dir_path,
	3754	+ table->name, IBD, true);
3753	3755	<pre>} else if (table->dir_path_of_temp_table) {</pre>
3754		- filepath = fil_make_ibd_name(
3755		<pre>- table->dir_path_of_temp_table, true);</pre>
	3756	+ filepath = fil_make_filepath(
	3757	+ table->dir_path_of_temp_table,
	3758	+ NULL, IBD, false);
3756	3759	} else {
3757		<pre>- filepath = fil_make_ibd_name(tablename, false);</pre>
	3760	+ filepath = fil_make_filepath(
	3761	+ NULL, tablename, IBD, false);
2750	2761	1
		14

3/50 3/51

Architecture Differences: Sticky Notes

Which: fil0fil.cc (StorageManager) -> dict0mem.cc (ResourceManager) dict_mem_create_temporary_tablename

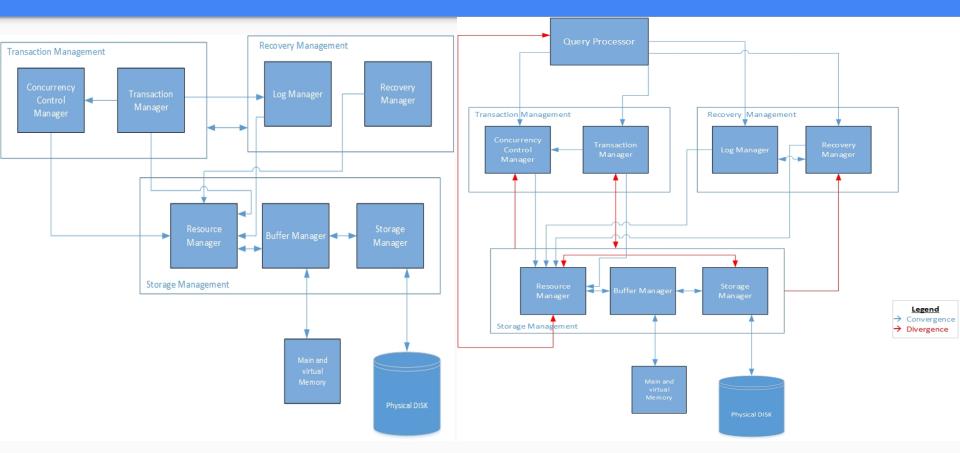
Who: Jimmy Yang

When: Sept 5, 2012

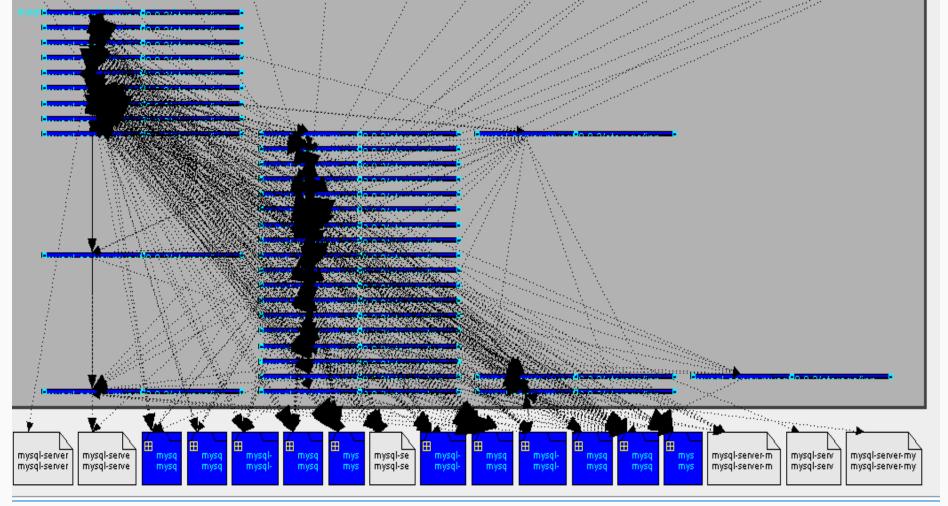
Why: Bug fixing

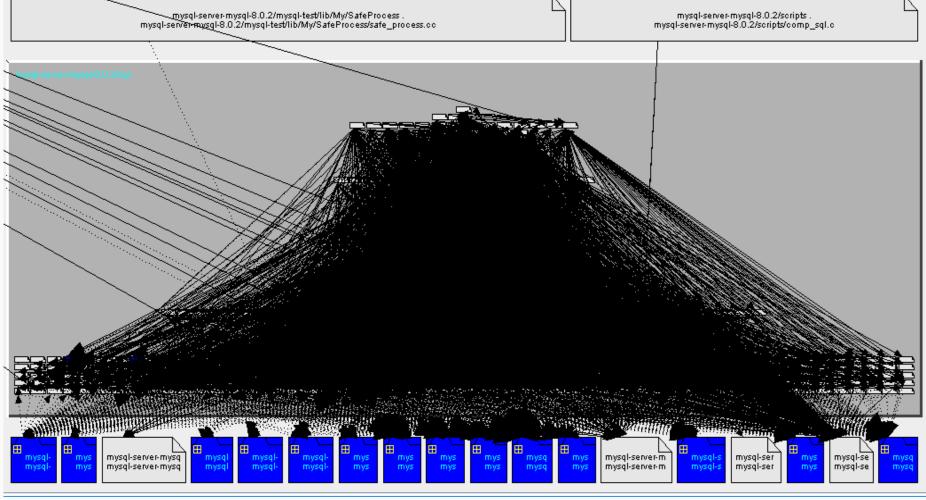
```
604
     +dict_mem_create_temporary_tablename(
     +/*======*/
     +
            mem heap t* heap, /*!< in: memory heap */
607
            const char* dbtab, /*!< in: database/table name */</pre>
    +
608
           table_id_t id) /*!< in: InnoDB table id */
    +
    +{
610 +
            const char* dbend = strchr(dbtab, '/');
611 +
            ut_ad(dbend);
612 +
            size t dblen = dbend - dbtab + 1;
613 +
            size_t
                         size = tmp_file_prefix_length + 4 + 9 + 9 + dblen;
614 +
615
            char* name = static_cast<char*>(mem_heap_alloc(heap, size));
    +
616 +
            memcpy(name, dbtab, dblen);
617
            ut_snprintf(name + dblen, size - dblen,
618
                      tmp file prefix "-ib" UINT64PF, id);
   +
619 +
            return(name);
620
     +}
```

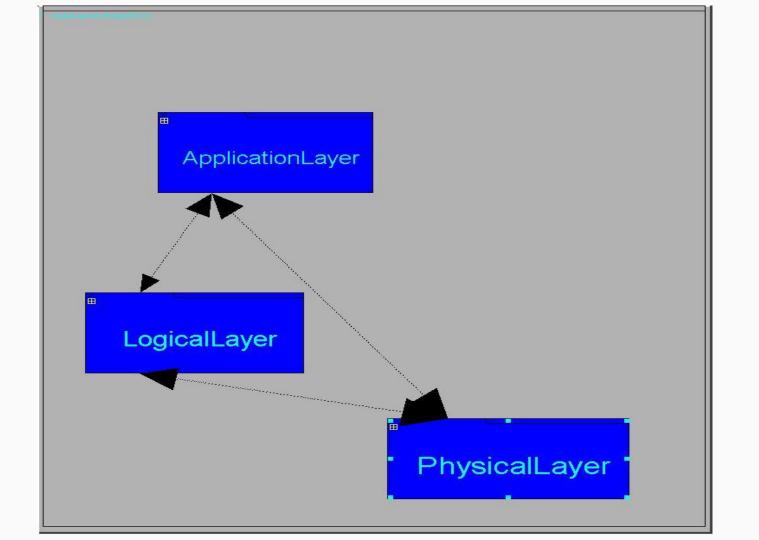
Side by side comparison: Conceptual and Concrete architectures

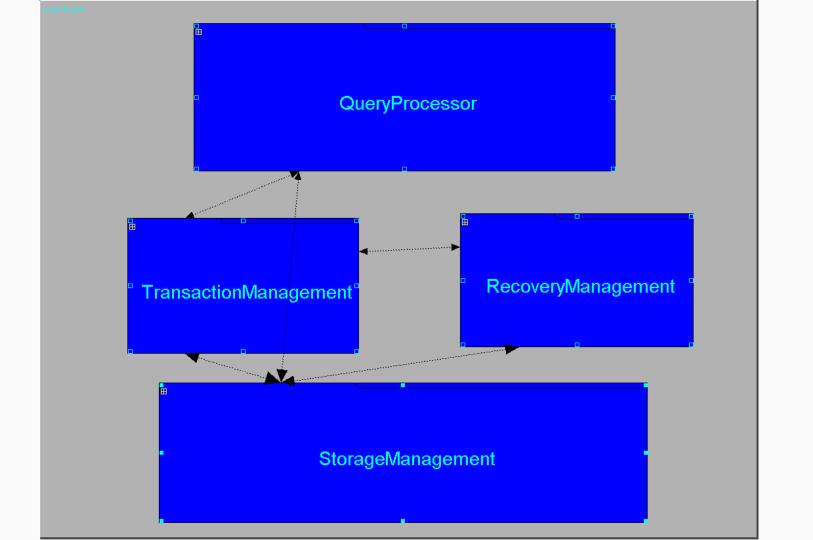


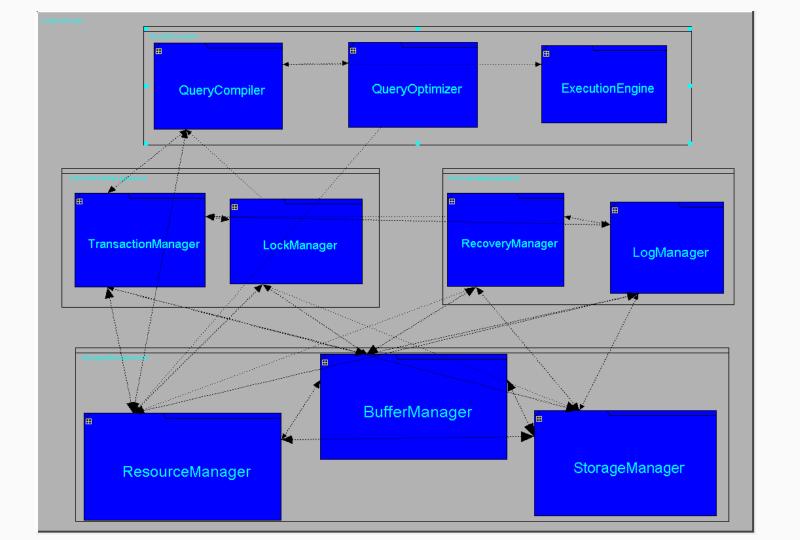
LsEdit Screenshots



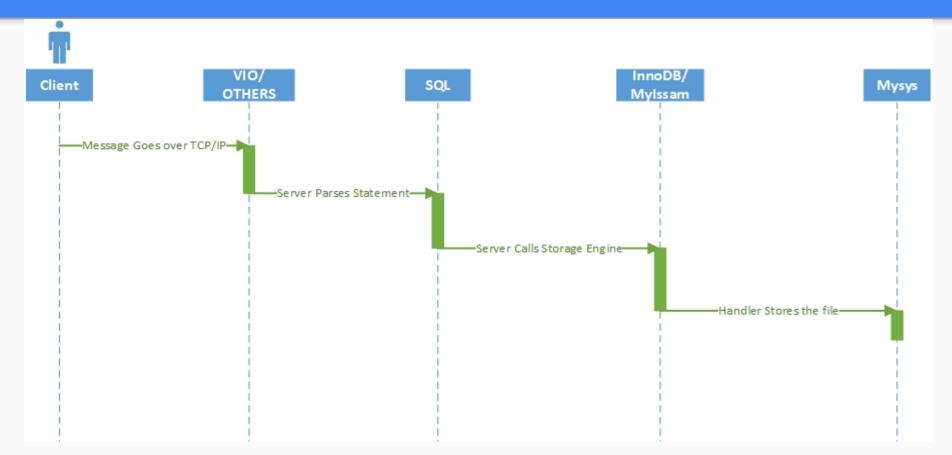








Use Case - INSERT Statement



Lessons Learned

- Retrieving Concrete architecture of a good documented software is possible to achieve with less problems.
- Concrete architecture may contain components that don't exist in the conceptual architecture.
- There are more interactions between components in the concrete architecture than the conceptual.
- Lsedit is a powerful tool but not as user friendly.