Complex Networks: Quiz #8

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## Problem 1

Make a program of breadth-first search (BFS).

#### Answer 1

The BFS made by myself:

```
import networkx as nx
# note the iterated nodes as 1
# push the current node into list2(FIF0)
# -> push its adjanct nodes into list2
# ->dequeue current node
def my_BFS(G,s):
    count = 0
    statement = []
   queue = []
   n = len(G)
   result = []
   # use one list to save the statement
   # initial all nodes as zero (haven't iterate)
   for node in range(0,n):
        statement.append(0)
   #visit the source node
   queue.append(s)
   count = count + 1
    statement[s] = 1
    while (queue!=[]):
        current = queue[0]
        result.append(current)
        #if haven't reach the node, push it into the queue
        if (count-1<n):</pre>
            for i in G[count-1]:
                if(statement[i] == 0):
                    queue.append(i)
        # dequeue current node
        queue.remove(current)
        count = count + 1
        #change the statement of nodes in queue
        for v in queue:
            statement[v] = 1
    return result
G = nx.karate_club_graph()
result = my_BFS(G,source)
print(result)
```

and the result is:

```
[0, 1, 2, 3, 4, 5, 6, 7, 8, 10, 11, 12, 13, 17, 19, 21, 31, 30, 9, 27, 28, 32, 16, 33, 25, 29, 23, 24]
```

# Problem 2

Show the final status of distance array and queue after BFS is done from vertex 0 of Karate club network.

### Answer 2

After BFS has been done, the queue is empty.

The distance array is:

# Problem 3

Explain why BFS is not good for networks with varying edge lengths.

#### Answer 3

In BFS, it will search all of the nearest nodes from the source firstly, and after that expand the nodes connected to the current node. So the depth of searching is small and need a lot of memory to storage the array.

If the edge length is varying, which means the graph is deeper, BFS may not have a good performance because of the huge cost of space. In terms of mathematics, the space complexity of BFS is  $O(B^M)(from\ wiki)$ , in which B is the max number of branches and M is the max length of the tree. Therefore, with more branches and when the graph is deeper, the space cost will increase.