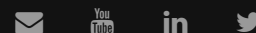


breaking



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Moker, Part 1: dissecting a new APT under the microscope

Recently, we came across Moker, an advanced malware residing in a sensitive network of a customer. Since the malware did not try to access an external server, but rather tamper with the system inner workings, we decided to give this malware a second look.

Indeed, we found that there is more to this malware than meets the eye.

First, it was obvious that the malware authors placed many anti-research measures, beyond those of the so-called usual anti-debugging techniques.

Second, once we started analyzing the malware itself, we were also able to analyze its capabilities and the advanced techniques it uses to remain stealthy, hook itself into the operating system, and its sophisticated inter-process communication.

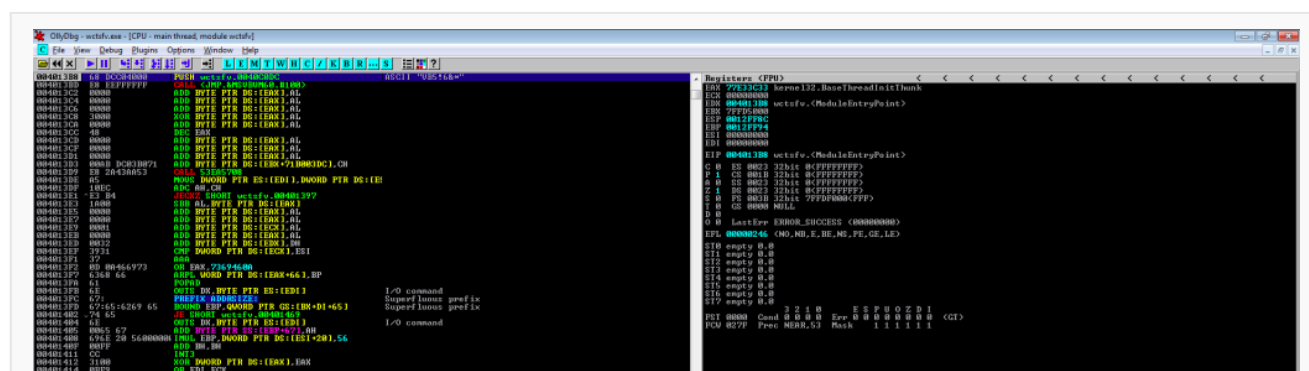
We decided to call this malware Moker as it's the file description that the malware author gave to the malware's executable file.

In this first of a 2-series blog we'll present the challenges that Moker placed in front of researchers to avoid detection and anti-dissection; and show how we overcame these measures one-by-one until we arrived at the stripped-down malware sample.

In an upcoming second blog entry, we'll present the comprehensive list and different advanced techniques that Moker utilized. Since we're still analyzing the malware, we'll present the technical details within the next few days.

Challenge #1: initial anti-debugging and unpacking

We run the malware inside our virtual environment to first understand what type of executable we are looking at. In fact, we are able to recognize that the malware (or at least its packer), is a Visual 5 or 6 binary:



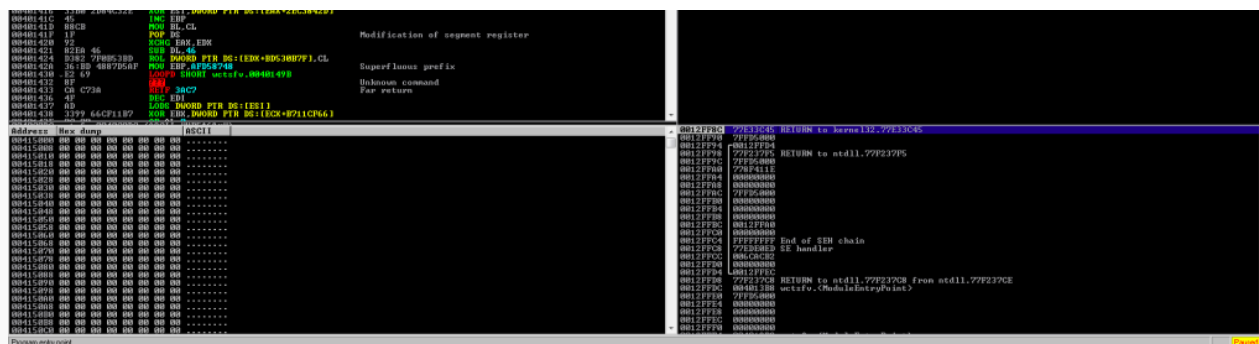


Figure 1: Recognizing that it's a Visual 5 or 6 Binary (first line, highlighted in blue)

Understanding the type of binary, we run the sample inside the appropriate debugger but very quickly the program crashes on Access Violation.

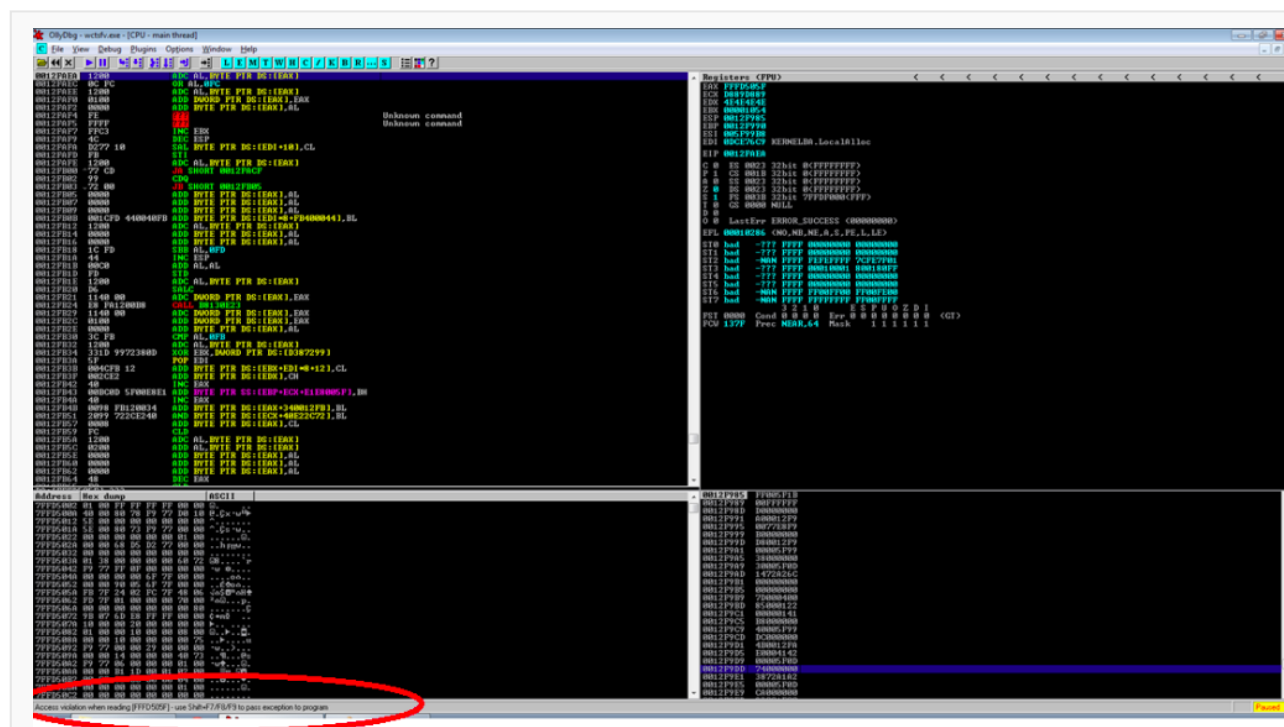


Figure 2: Access Violation in the stack memory space

The malware somehow knows that it is running inside the context of a debugger and jumps to some junk code.

The next step is to look for the appropriate debugger state flag – the PEBBeingDebugged flag.

The corresponding byte is in the Process Environment Block (PEB). The byte is read by the API IsDebuggerPresent, and can be read manually by accessing:

FS:[0x30]+2

Since this byte has no other meaning apart from the debug status, we are able to safely reset it to

zero.

Run again. Crash again. On a different address this time.

Our packer has another anti-debugging trick up its sleeve. This time it turns out to be a slightly less common method (though still in use): the use of PEB!NtGlobalFlags. Its corresponding DWORD can be read using an API, or by directly accessing this address:

FS[0x30]+0x68

Reset the DWORD to zero. Run. No crash this time. However, the process creation terminated.

This fact alone means little. Anti-debugging can still be present and the process could have called ExitProcess or TerminateProcess before it does what it is designed to do.

In this specific case, however, anti-debugging is defeated.

Challenge #2: Process Hollowing

Beginning to understand what the process is doing, let's set breakpoints at some interesting APIs, one of these being process creation.

Sure enough, a breakpoint at CreateProcessW is hit.

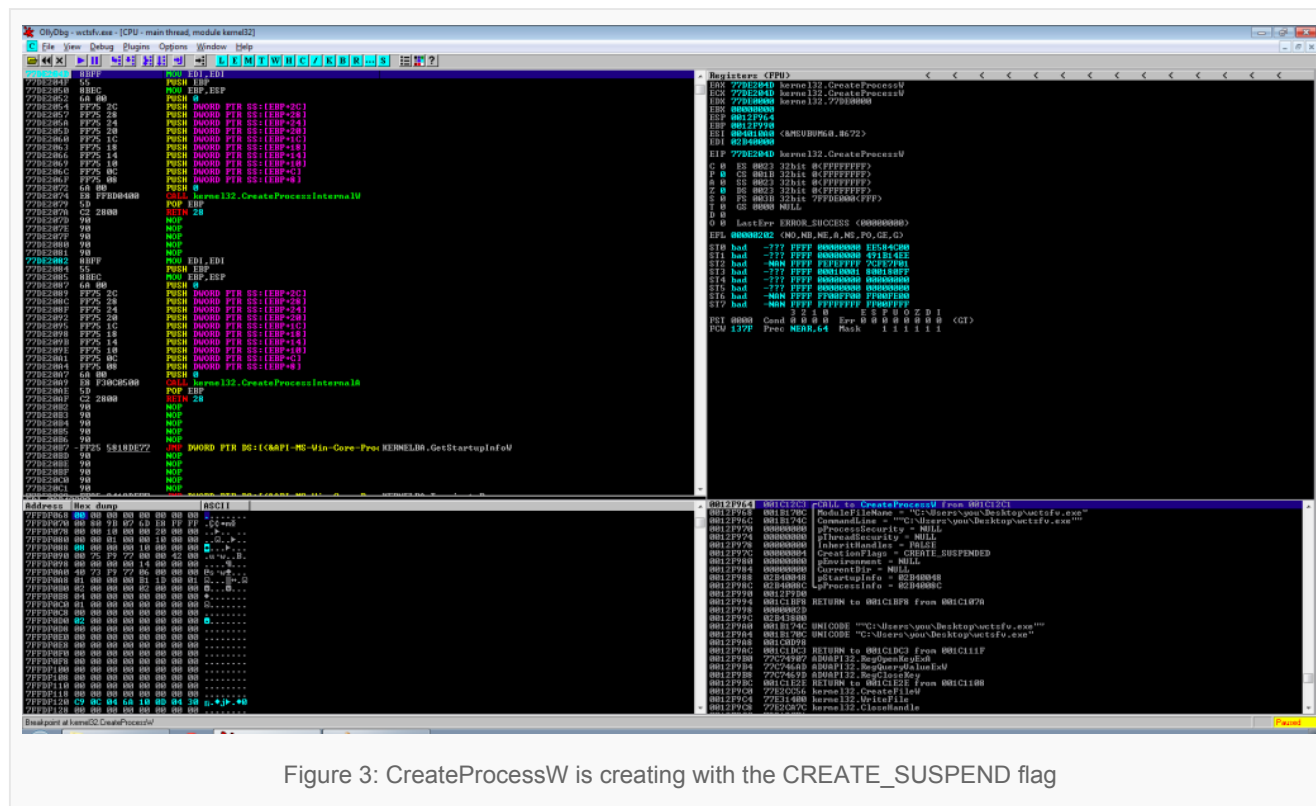


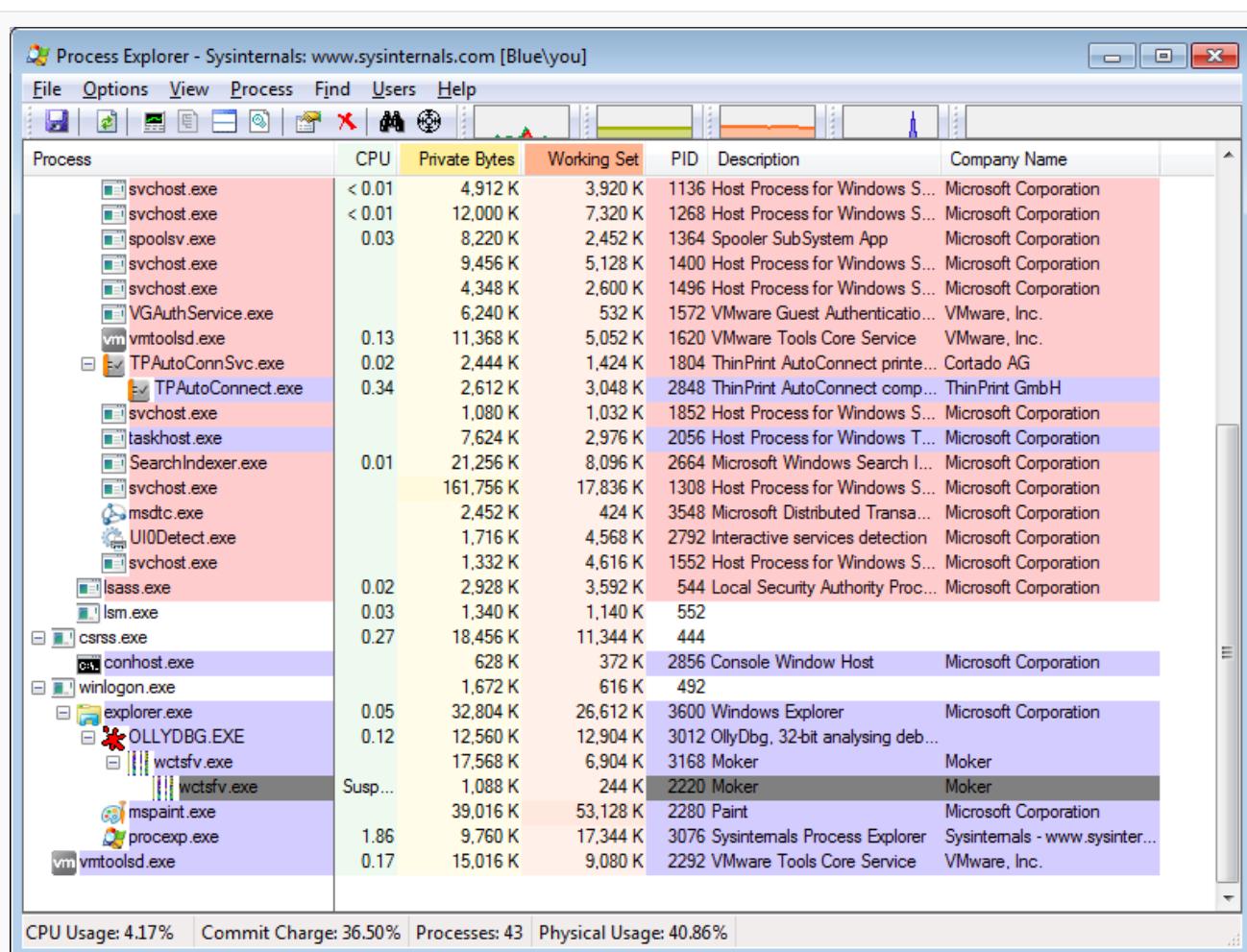
Figure 3: CreateProcessW is creating with the CREATE_SUSPEND flag

As indicated by the code above, the Moker sample tries to run its own executable without arguments,

and with the `CREATE_SUSPENDED` flag. This means that when the call returns, a new process will be created in a suspended state. The process will remain suspended until killed or until one of its process' threads is resumed.

Continuing to run Moker, we see that on the created process at its base address (0x400000) the sample uses `UnmapViewOfSection` and `MapViewOfSection`. It then calls `NtResumeThread`, but let's wait with that one for now.

As we can see, at this stage we have a process running, but its code is not the original code in the file. In fact, the sample completely overwrites the new process' PE header, giving it a new entry point and a whole section of new code.



The screenshot shows the Process Explorer window with the following table of processes:

Process	CPU	Private Bytes	Working Set	PID	Description	Company Name
svchost.exe	< 0.01	4,912 K	3,920 K	1136	Host Process for Windows S...	Microsoft Corporation
svchost.exe	< 0.01	12,000 K	7,320 K	1268	Host Process for Windows S...	Microsoft Corporation
spoolsv.exe	0.03	8,220 K	2,452 K	1364	Spooler Sub-System App	Microsoft Corporation
svchost.exe		9,456 K	5,128 K	1400	Host Process for Windows S...	Microsoft Corporation
svchost.exe		4,348 K	2,600 K	1496	Host Process for Windows S...	Microsoft Corporation
VGAuthService.exe		6,240 K	532 K	1572	VMware Guest Authenticatio...	VMware, Inc.
vmtoolsd.exe	0.13	11,368 K	5,052 K	1620	VMware Tools Core Service	VMware, Inc.
TPAutoConnSvc.exe	0.02	2,444 K	1,424 K	1804	ThinPrint AutoConnect printe...	Cortado AG
TPAutoConnect.exe	0.34	2,612 K	3,048 K	2848	ThinPrint AutoConnect comp...	ThinPrint GmbH
svchost.exe		1,080 K	1,032 K	1852	Host Process for Windows S...	Microsoft Corporation
taskhost.exe		7,624 K	2,976 K	2056	Host Process for Windows T...	Microsoft Corporation
SearchIndexer.exe	0.01	21,256 K	8,096 K	2664	Microsoft Windows Search I...	Microsoft Corporation
svchost.exe		161,756 K	17,836 K	1308	Host Process for Windows S...	Microsoft Corporation
msdtc.exe		2,452 K	424 K	3548	Microsoft Distributed Transa...	Microsoft Corporation
UIODetect.exe		1,716 K	4,568 K	2792	Interactive services detection	Microsoft Corporation
svchost.exe		1,332 K	4,616 K	1552	Host Process for Windows S...	Microsoft Corporation
lsass.exe	0.02	2,928 K	3,592 K	544	Local Security Authority Proc...	Microsoft Corporation
lsim.exe	0.03	1,340 K	1,140 K	552		
csrss.exe	0.27	18,456 K	11,344 K	444		
conhost.exe		628 K	372 K	2856	Console Window Host	Microsoft Corporation
winlogon.exe		1,672 K	616 K	492		
explorer.exe	0.05	32,804 K	26,612 K	3600	Windows Explorer	Microsoft Corporation
OLLYDBG.EXE	0.12	12,560 K	12,904 K	3012	OllyDbg, 32-bit analysing deb...	
wctsfv.exe		17,568 K	6,904 K	3168	Moker	Moker
wctsfv.exe	Susp...	1,088 K	244 K	2220	Moker	Moker
mspaint.exe		39,016 K	53,128 K	2280	Paint	Microsoft Corporation
procepx.exe	1.86	9,760 K	17,344 K	3076	Sysinternals Process Explorer	Sysinternals - www.sysinter...
vmtoolsd.exe	0.17	15,016 K	9,080 K	2292	VMware Tools Core Service	VMware, Inc.

At the bottom of the window, the status bar shows: CPU Usage: 4.17%, Commit Charge: 36.50%, Processes: 43, Physical Usage: 40.86%.

Figure 4: Moker as a suspended process

Challenge #3: Planned exceptions, or anti-debugging revisited

Before resuming the process and letting it run, let's create a dump in order to analyze it.

Dump. Fix sections.

As a first, we need to look at the import table to get a general idea about our unpacking status

As a first, we need to look at the import table to get a general idea about our unpacking status.

[ImportTable]					
DllName	OriginalFirstThunk	TimeDateStamp	ForwarderChain	Name	FirstThunk
KERNEL32.DLL	000D60DC	00000000	00000000	000D608C	000D60FC
USER32.DLL	000D61A0	00000000	00000000	000D609A	000D61AC
ADVAPI32.DLL	000D61D0	00000000	00000000	000D60A6	000D61D8
COMCTL32.DLL	000D61F0	00000000	00000000	000D60B4	000D6200
COMDLG32.DLL	000D6250	00000000	00000000	000D60C2	000D6260
ThunkRVA	ThunkOffset	ThunkValue	Hint	ApiName	
000D60DC	000D60DC	000D611C	0000	ExitProcess	
000D60E0	000D60E0	000D612A	0000	GetComputerNameA	
000D60E4	000D60E4	000D613E	0000	GetModuleHandleA	
000D60E8	000D60E8	000D6152	0000	LoadLibraryA	
000D60EC	000D60EC	000D6162	0000	VirtualAlloc	
000D60F0	000D60F0	000D6172	0000	VirtualFree	
000D60F4	000D60F4	000D6180	0000	AddVectoredExceptionHandler	
Number Of Thunks: 7h / 7d (OriginalFirstThunk chain)					
					<input type="checkbox"/> View always FirstThunk

Figure 5: Import table shows that no fixing needs to be made

Since the import table seems normal, we can open it with IDA.

When we open the dumped code in IDA, we see something strange. IDA finds the start address and parses the code, but the code cannot run very far as it attempts to run “*in al,dx*” at 4CEA08 – a privileged instruction. Since this is a user-mode program, it cannot natively run this instruction without generating an exception.

```

.text:004CEA00 ;
.text:004CEA00 public start
.text:004CEA00 start: ; DATA XREF: .text:start o
.text:004CEA00 lea     ebx, start
.text:004CEA00 call   sub_4CEC99
.text:004CEA00 in      al, dx
.text:004CEA00 les     eax, [edx+7468016Ah]
.text:004CEA12 pop     ss
.text:004CEA12 inc     eax
.text:004CEA14 add     ah, ah
.text:004CEA16 pusha
.text:004CEA17 push    0
.text:004CEA19 push    0
.text:004CEA1B push    1
.text:004CEA1D push    offset unk_401774
.text:004CEA22 in      al, 61h ; PC/XT PPI port 8 bits:
; 0: Tmr 2 gate OR 03H=spkr ON
; 1: Tmr 2 data e AND 0fCH=spkr OFF
; 3: I=read high switches
; 4: 0=enable RAM parity checking
; 5: 0=enable I/O channel check
; 6: 0=hold keyboard clock low
; 7: 0=enable kbdr
.text:004CEA22
.text:004CEA22
.text:004CEA22
.text:004CEA22
.text:004CEA24 lea     edx, unk_401774
.text:004CEA2A mov     dword_401788, 0Ch
.text:004CEA34 mov     dword_40178C, edx
.text:004CEA3A push    4
.text:004CEA41 push    1C1h
.text:004CEA41 loc_4CEA41: ; CODE XREF: .text:004CEA53 j
.text:004CEA41 push    offset unk_40136C
.text:004CEA46 call   sub_4D4AD8
.text:004CEA48 in      al, dx
.text:004CEA4C pop     esp
.text:004CEA4D xor     esp, ebp
.text:004CEA4F add     al, 4Fh
.text:004CEA51 in      al, dx
.text:004CEA52 int     3 ; Trap to Debugger
.text:004CEA53 ja     short loc_4CEA41
.text:004CEA55 int     3 ; Trap to Debugger
.text:004CEA58 add     bh, al
.text:004CEA5D add     eax, 402148h
.text:004CEA5D pushf
.text:004CEA5D
.text:004CEA5E dw     0

```

Figure 6: IDA Dump shows that the code attempts to run a privileged op-code

The only thing we can run is the call prior to “*in al,dx*” – sub_4CEC99. What’s in that call? A small function that registers a Vectored Exception Handler. IDA has already identified that and named a “Handler” function on the list.

At this stage, when trying to step through the program with a debugger, the program attempts to execute the privileged instruction and an exception is raised. When we return from the exception Handler, the program calls NtContinue and the debugger loses track of the program.

This is an interesting anti-debugging technique that our new in-memory module uses.

It was time to investigate that Handler.

We found out that the Handler reads the byte at the address where the exception occurred and then enters a state machine.

This set of rules reads one or more bytes further, depending on the first byte's value, then alters the CONTEXT fields of the thread in order to simulate a call instruction.

A return address will also be inserted to the stack – to the precise location where our special opcode ends.

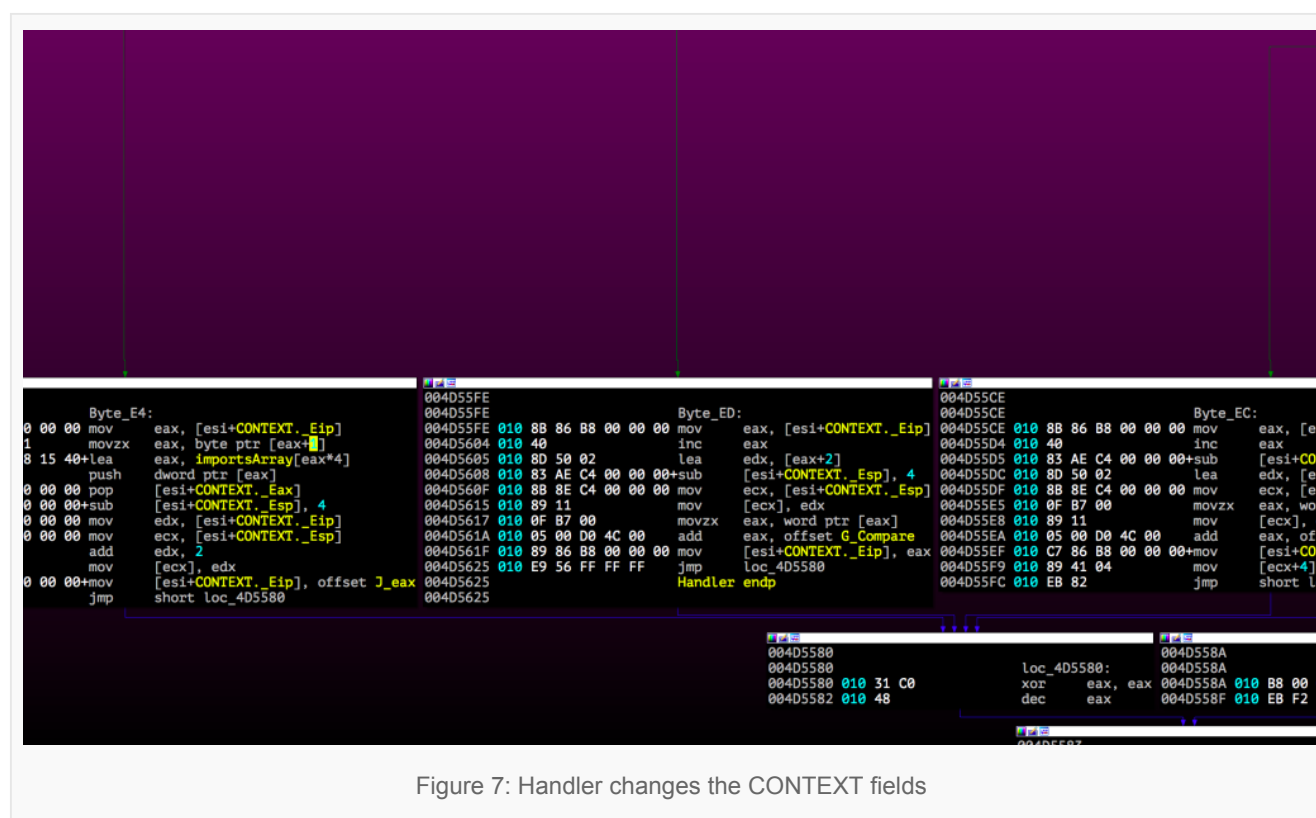


Figure 7: Handler changes the CONTEXT fields

The rules are:

- E4 <byte> – Call a Windows API from the array of API addresses by index <byte>.
- ED <word> – Call a local function by offset Base+<word>
- EC <word> – Call a specific function with an offset Base+<word> as a parameter.

This allows us to find system calls and local functions. This is enough to figure out the intentions of this

binary but the analysis cannot be complete without knowing what the EC “call” does.

Challenge #4: Decrypted code

The EC opcode throws an exception to the Handler, as expected. The Handler then changes `_eip` in the `CONTEXT` struct of the thread to make the program jump to a function with an offset as a parameter.

The data at this offset does not look useful at its raw form.

Now that we know where the function is, there is no problem running it with a debugger and letting it decrypt the data blob into code and jump to it. However, the algorithm was simple enough to write a decoder script for it.

Finally, we were able to decode the whole program.

We use IDA's structs and enums to make the exception-handler-calls more readable:

```

text:004CEA00 ; -----
text:004CEA00
text:004CEA00
text:004CEA00
text:004CEA00
text:004CEA00 8D 1D 00 EA 4C 00
text:004CEA06 E8 8E 02 00 00
text:004CEA06
text:004CEA0B EC C4 82
text:004CEA0E
text:004CEA0E 6A 01
text:004CEA10 68 74 17 40 00
text:004CEA10
text:004CEA15 E4 60
text:004CEA17
text:004CEA17 6A 00
text:004CEA19 6A 00
text:004CEA1B 6A 01
text:004CEA1D 68 74 17 40 00
text:004CEA1D
text:004CEA22 E4 61
text:004CEA24
text:004CEA24 8D 15 74 17 40 00
text:004CEA2A C7 05 88 17 40 00+
text:004CEA34 89 15 8C 17 40 00
text:004CEA3A 6A 04
text:004CEA3C 68 C1 01 00 00
text:004CEA41 68 6C 13 40 00
text:004CEA46 E8 8D 60 00 00
text:004CEA46
text:004CEA4B EC 5C 31
text:004CEA4E EC 04 4F
text:004CEA51 EC CC 77
text:004CEA54 EC CC 00
text:004CEA57
text:004CEA57 C7 05 48 21 40 00+
text:004CEA61 68 48 21 40 00
text:004CEA61
text:004CEA66 E4 17
text:004CEA68
text:004CEA68 83 3D 2C 16 40 00+
text:004CEA6F 74 09
text:004CEA71 68 30 21 40 00
text:004CEA76 6A FF
text:004CEA76
text:004CEA78 E4 25

; -----
public start
start:
; DATA XREF: sub_150000+2'o
; seg006:001B0004'o ...
    lea     ebx, start
    call    Call_AddVectoredExceptionHandler

; -----
DecryptCall <0ECh, LoadNtdll>

; -----
    push    1
    push    offset unk_401774
; -----
    CallAPI <0E4h, InitializeSecurityDescriptor>

; -----
    push    0
    push    0
    push    1
    push    offset unk_401774
; -----
    CallAPI <0E4h, SetSecurityDescriptorDacl>

; -----
    lea     edx, unk_401774
    mov     dword_401788, 0Ch
    mov     off_40178C, edx
    push    4
    push    1C1h
    push    offset aNtdll_dll ; "ntdll.dll"
    call    sub_4D4AD8

; -----
DecryptCall <0ECh, VolumeInfoOrExit>
DecryptCall <0ECh, CreateFileAndMap>
DecryptCall <0ECh, CheckGlobalFlag>
DecryptCall <0ECh, EnumKeys>

; -----
    mov     dword_402148, 9Ch
    push    offset dword_402148
; -----
    CallAPI <0E4h, GetVersionExA>

; -----
    cmp     dword_4015B4+78h, 0
    jz      short loc_4CEA7A
    push    offset dword_402130
    push    0FFFFFFh
; -----
    CallAPI <0E4h, IsWow64Process>
  
```

Figure 8: pseudo-code for exception handler calls

Challenge #5: Process injection

As usual, the malware searches through processes, with known APIs and injects itself into

Explorer.exe as well as to the following three typical Windows' background processes:

- Svchost.exe
- lsass.exe
- csrss.exe

The injected thread tries to load a library using LoadLibraryW with a file inside the user's temp folder, the name is similar to this:

C:\Users\user\AppData\Local\Temp\4FE53.

Notice the "." at the end. Naming a file as anything that ends with a dot causes it to be hidden in Windows Explorer. This is done to prevent users from finding this file by listing its folder's contents.

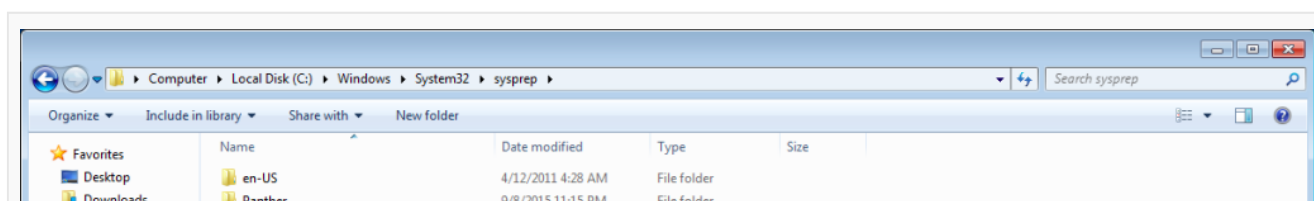
Challenge #6: Elevation of privileges

To recall, in Windows, some system operations are only available to programs that request elevation. Requesting elevation requires Administrator rights and also pops up the User Account Control (UAC) mechanism that asks the user for permissions to elevate (a.k.a "Allow this program to make changes to this computer?").

How was Moker able to ensure that it gains system-privileges, without requiring the user's consent?

1. **Leeches on to an elevated program.** In order to enable proper functioning of the Windows system, there are exceptional programs that are always granted elevation permission – without a user prompt. One of these exceptional processes that Moker found is the Microsoft System Preparation Tool that resides in C:\Windows\system32\sysprep.
2. **Elevates the privilege of a DLL.** There is a Windows design vulnerability which enables loading unauthorized DLLs by authorized applications. This vulnerability is found in the way Windows loads DLLs upon request. When an application loads a DLL by its name, Windows first looks through the application's current folder, and then proceeds to search at system directories (and other paths). In order to avoid a predicted chaos using this technique, some DLLs are always loaded from the system directory regardless of any other same named DLLs found in its own path. Monitoring sysprep shows it loading one DLL which does not follow this restriction: ActionQueue.dll.

Putting it all together, Moker writes a file named "ActionQueue.dll" into the "sysprep" directory and then runs sysprep so that the DLL is able to run with elevated privileges.



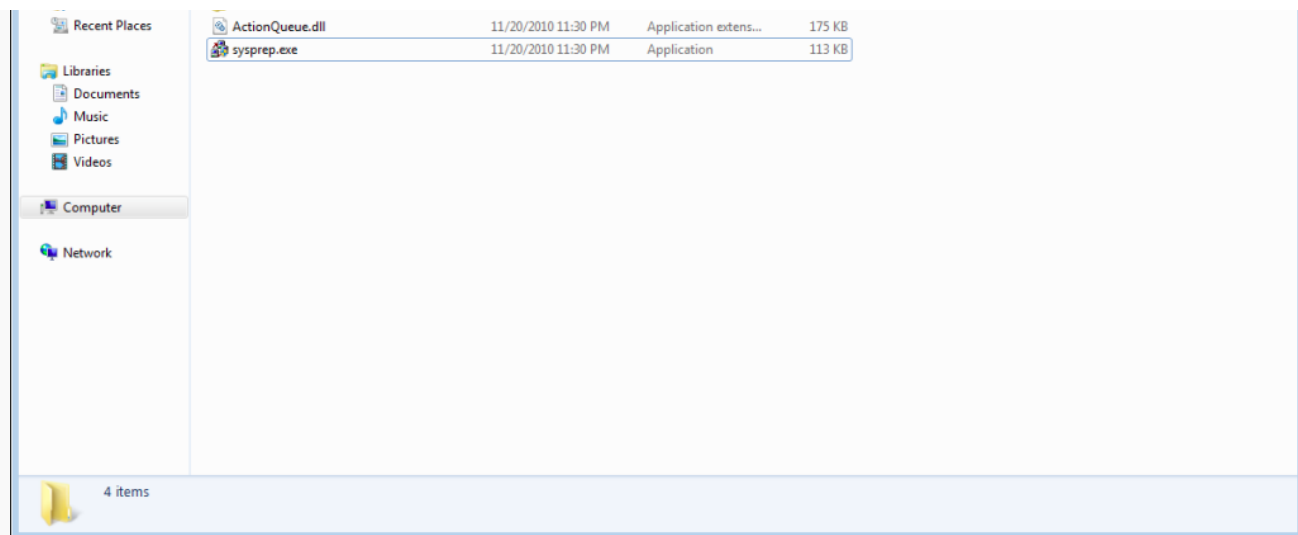


Figure 9: In the sysprep directory, Moker runs sysprep whereas its executable ActionQueue.dll receives elevated privileges

Challenge #7: Downloading the payload

We have skipped another part of the story where Moker loaded itself into the sysprep process and injects itself into other processes. That actually happened earlier (challenge #5).

By analyzing the running code, we realize that we actually have a downloader in our hands.

The downloader itself uses one of its injected threads to communicate with some unknown server. Interestingly, at the time of analysis, the address for the server was empty. This could suggest that:

1. The executable was changed since it was first installed
2. The payload is loaded in different ways, not only via the network
3. There is another layer of protection protecting specifically that part of information

At this point let's try to emulate a server response, just to see its program flow.

It uses `InternetReadFile`, reads the response to a buffer, and saves it to the registry as-is.

This leads us to understand that the payload is in the registry.

Setting a breakpoint at `RegQueryValueEx` takes us the next step.

At this stage the program reads the registry value, manipulates it (decrypts it), validates that it contains a valid PE header ("MZ", "PE"), and then loads it as a library.

Extracting the registry value from the infected machine and feeding it to the decryption code results in a valid DLL with relatively readable code.

Moker is finally installed, without any external connections.

Coming Up: analyzing Moker's capabilities

As this stage, we are ready to look under the hood of Moker.

A first look reveals that Moker features RAT capabilities, a few different techniques to obfuscate strings, a rather sophisticated inter-process communication and synchronization between its components (including the downloader from previous parts).

A distinct feature is its ability to start the Remote Desktop service and allow external access to the machine using Remote Desktop clients (like Windows client or Linux' 'rdesktop').

Since we're still analyzing the malware, we'll present the technical details within the next few days. In the meanwhile, you can find a high-level overview of the capabilities here:

<http://blog.ensilo.com/moker-a-new-apt-discovered-within-a-sensitive-network>



 Posted October 6, 2015 by [Yotam Gottesman](#)

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