





gasgas4@gmail.com

咨訊主百 答!



Malware Must Die!

Semper legerent "Salve Regina" ante venatione malware

Thursday, September 17, 2015

MMD-0043-2015 - Polymorphic in ELF malware: Linux/Xor.DDOS

Background

A share of knowledge I have, hopefully to make internet safer - @unixfreaxjp

The threat of **Linux/XOR.DDoS**, a China-made ELF backdoor & ddoser malware, a rather specific threat compares to other Chinese ELF ddosers, and it's still on going. I just received a good question (from I assumed from a victim of infection or a researcher) about why the found malware binary is not the same as what was firstly executed one. Well, this writing is short and covering the answer for the asked question only. But, the information maybe important for the mitigation and detection, and also various methodology I use for the sharing to other NIX mates, so I write this post with three processes I conduct to every ELF malware investigation: in *reversing*, *debugging* and *forensics* ways. Please bear with the poor english since I had few time to check, or to the lack of the explanation.

Polymorphic is a behavior of malware during self-reproduction constantly changes ("morphs") the file characteristic (size, hash, etc), and it may not be the same with the previous copy or as previous pre-infection state. The goal of this changes is **to makes it difficult** for signature-based antivirus software programs to recognize and detect the polymorphed malware.

Polymorphic method in malware is an usual practise in windows malware. In UNIX malware maybe it is not as commonly heard as in Windows; but since the nature of NIX malware are coming from networking, either to be "extracted" from encoder/infector files, downloaded or dropped by other malware from the beginning, so... I guess we have many hashes by default. But in this post, we are actually dealing with a polymorphic behavior malware just like ones infecting Windows during the self-copy method.. so I guess it is worth to write a bit.

The reported case was a real infection, a case of known gang/crooks, I am allowed to post the the attack log as per following:

```
1 2015-09-17 12:04:39+0900 New connection: 43.229.53.90; Session: 133]
2 2015-09-17 12:04:39+0900 [session=133,ip=43.229.53.90] Remote SSN version: SSN-2.0-PUTTY
3 2015-09-17 12:04:39+0900 [session=133,ip=43.229.53.90] outgoing: aes128-ctr hmac-shal none
4 2015-09-17 12:04:39+0900 [session=133,ip=43.229.53.90] incoming: aes128-ctr hmac-shal none
6 2015-09-17 12:04:39+0900 [session=133,ip=43.229.53.90] incoming: aes128-ctr hmac-shal none
6 2015-09-17 12:04:39+0900 [session=133,ip=43.229.53.90] Satring service ssh-userauth
8 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] Toot trying auth none
9 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] root trying auth password
10 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] root trying auth password
11 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] root trying auth password
12 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] root authenticated with password
12 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] starting service ssh-connection
13 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] cannel open
14 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] channel open
15 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] channel open
16 2015-09-17 12:04:40+0900 [session=133,ip=43.229.53.90] channel open
17 2015-09-17 12:09:42+0900 [session=133,ip=43.229.53.90] channel open
18 2015-09-17 12:09:42+0900 [session=133,ip=43.229.53.90] channel open
19 2015-09-17 12:09:42+0900 [session=133,ip=43.229.53.90] channel open
20 2015-09-17 12:09:42+0900 [session=133,ip=43.229.53.90] channel open
21 2015-09-17 12:09:42+0900 [session=133,ip=43.229.53.90] channel open
22 2015-09-17 12:09:42+0900 [session=133,ip=43.229.53.90] channel open
23 2015-09-17 12:09:42+0900 [session=133,ip=43.229.53.90] channel open
24 2015-09-17 12:09:42+0900 [session=133,ip=43.229.53.90] channel open
25 2015-09-17 12:09:43+0900 [session=133,ip=43.229.53.90] channel open
26 2015-09-17 12:09:43+0900 [session=133,ip=43.229.53.90] channel open
27 2015-09-17 1
```

Yes, it is a recent attack, please block the IP addresses.

The above log is typical Linux/Xor.DDOS hostasa.org ssh brute attack pattern. I announced the case not so long ago here (different cases, same attacker)-->[link] and the recent incident was reported too in here-->[link]. I uploaded this ELF malware sample into Virus Total w/the link is here-->[link].

Polymorphic PoC

When Linux/XOR.DDoS malware was executed, it will come to the stage that it seeks the place to self-copy it self, in my case the linux system call can show us the effort to write file like:

open("/usr/bin/lgjgjmkkgd", O_WRONLY|O_CREAT, 0777); depends, in mine is ? open("/bin/lgjgjmkkgd", O_WRONLY|O_CREAT, 0777); depends, in mine is -1

In a well-hardened linux system and if the malware is not executed as root you should see the same

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Since malware and its evolution is becoming the serious threat in our internet and computer industry. We are now coming to the stage to admit the fact that malware is actually "winning" this longest 15+ years historical "battle" by keep on its existence, infecting and lurking us....[Read More]

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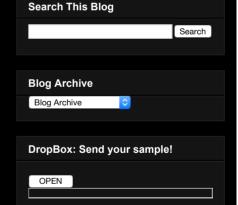
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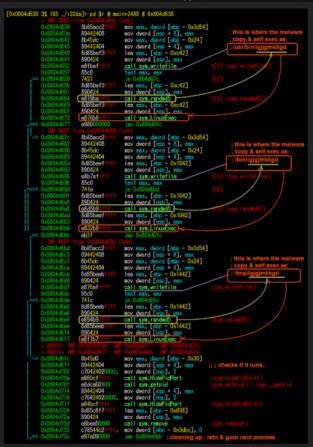
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result as per pasted above. And that time the malware will aim to the only their favorite heavenly place to copy: /tmp:

By reverse engineering the ELF malware, after seeking for a while, the assembly procedure below is responsible for the above operation: (the bigger picture click-->>THIS)



You can see the cascade of jumps during each error that might occur until it ends up to the accessed one for the self-copy purpose, starting from /usr/bin to /bin, and in my case it is ended with /tmp/[randomname]. The filename is random and the full path with the directory aimed is to be "fired" via an original API to execute the execve(), but we will go to this topic later on.

In Linux memory forensics the blob data copied can be seen clearly with some beautify effort, a good old *hexdump* is still a favorite in dealing with raw hex data:

```
## Copy process illustration (read and write of copy process) in the end on [...]

3 00098bd0 6d 65 00 5f 64 6c 5f 6d 61 70 5f 6f 62 6a 65 63 | me._dl_map_obje |

4 00098be0 74 5f 64 65 70 73 00 5f 6e 6c 5f 43 5f 4c 43 5f | t_deps._nl_C_LC |

5 00098bf0 49 44 45 4e 54 49 46 49 43 41 54 49 4f 4e 00 5f | IDENTIFICATION. |

6 00098c00 64 6c 5f 6e 73 00 5f 6e 6c 5f 6c 6f 61 64 5f 6c | dl_ns._nl_load_ |

7 00098c10 6f 63 61 6c 65 5f 66 72 6f 6d 5f 61 72 63 68 69 | ocale_from_arch |

8 00098c20 76 65 00 77 63 74 72 61 6e 73 00 | ve.wctrans.| |

9 00098c2b
```

And the copy process was ended gracefully, as per debug check shows in the system call below:

```
1 read(3, "", 4096): ; EO/termination w/no space
2 close(3); ; end of copy (reading)
3 close(4); ; end of copy (writing)
```

Nothing so special about operation above, but it is related to the next steps, let's go forward.. Now, we can see up to here that the malware was self copied! But why the file gets different?

The next system's call showing the effort to open the written file afterward with flag to write.. What's going on?

```
open("/tmp/lgjgjmkkgd", O_WRONLY); ; opening the copied file
lseek(3, 0, SEEK_END) = 625707 <==size ; set LSET to the EOF for writing
; SEEK_END = *) ; note the size of original malware</pre>
```

It looks like the pointer of *LSET* used to write is pointing to the end of the file itself, noted the *SEEK_END* flag. For the illustration see the paste "*)" position below:

```
1 | ## Illustration of the LSET set in the end of file.. | ?
```

```
MMD-0037-2015 - A bad Shellshock & Linux/XOR.DDoS CNC "under the hood"
```

MMD-0036-2015 - KINS (or ZeusVM) v2.0.0.0 tookit (builder & panel source code) leaked.

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DDoS'er as Service - a camouflage of legit stresser/booter/etc

Most read analysis

Oute Modified Mar 27, 2014 1:49 PM Mar 27, 2014 1:49 PM Mar 27, 2014 1:49 PM Mar 27, 2014 2:11 PM Mar 27, 2014 2:11 PM Mar 27, 2014 1:49 PM Mar 27, 2014 2:11 PM

MMD-0020-2014 -Analysis of infection ELF malware: libworker.so - A shared (DYN) malicious llibrary by LD_PRELOAD



MMD-0028-2014 - Fuzzy reversing a new China ELF "Linux/XOR.DDoS"



The Evil Came Back: Darkleech's Apache Malware Module: Recent Infection, Reversing, Prevention & Source

And then we have these two operation called **timeoftheday()** and writing the specific strings in the end of the file:

```
gettimeofday({1442479267, 397488}, NULL); for randomid() seed..

write(3, "wlpvpovdvi\0", 11); 'size is set to 11'

; write string "wlpvpovdvi\0"-
; in the LSET position (EOF)
```

So this is what happened for BEFORE and AFTER the writing:

```
74 51 64 65 70 73 00 5f 6e 6c 5f 43 5f 4c 43 5f
49 44 45 4e 54 49 46 49 43 41 54 49 4f 4e 00 5f
64 6c 5f 6e 73 00 5f 6e 6c 5f 6c 6f 61 64 5f 6c
6f 63 61 6c 65 5f 66 72 6f 6d 5f 61 72 63 68 69
76 65 00 77 63 74 72 61 6e 73 00
00098be0 74 5f 64 65 70 73 00 5f 00098bf0 49 44 45 4e 54 49 46 49
                                                                                                                                  t_deps._nl_C_LC_
                                                                                                                                   IDENTIFICATION.
00098c00
                                                                                                                                  dl_ns._nl_load_ll
locale_from_archil
00098c10
00098c20
                                                                                                                                  ve.wctrans.
00098c2b
;; AFTER
00098be0
                    74 5f 64 65 70 73 00 5f
49 44 45 4e 54 49 46 49
64 6c 5f 6e 73 00 5f 6e
6f 63 61 6c 65 5f 66 72
76 65 00 77 63 74 72 61
                                                                          6e 6c 5f 43 5f 4c 43 5f
43 41 54 49 4f 4e 00 5f
6c 5f 6c 6f 61 64 5f 6c
6f 6d 5f 61 72 63 68 69
                                                                                                                                    t_deps._nl_C_LC_
IDENTIFICATION._
00098bf0
00098c00
                                                                                                                                   dl_ns._nl_load_l
ocale_from_archi
00098c10
00098c20
                                                                           6e 73 00 77 6c 70 76 70
                                                                                                                                   ve.wctrans.wlpvp
00098c30
                   6f 76 64 76 69 00
                                                                                                                                   ovdvi.
00098c36
```

So we see the file was added to 11 characters, which means we should have 11 bytes bigger for the size of file after this self-copy process, we'll get there..hang on!

Following the calls of the malware process, we can see the new file was saved:

```
1 | close(3) ; end of writing process..
```

And executed! Noted: execve() function is used to spawn the shell command.

```
1 | execve("/tmp/lgjgjmkkgd", ..); ; main running process of XOR.DDOS in new Pp ; with new size (& hash)
```

You can see how it was executed in the saved process data in the **/proc** :-), so believe me, it doesn't really any fancy tools for UNIX forensics, since UNIX gods already provided us openly with everything:

```
lgjgjmkkg 14881 MMD
lgjgjmkkg 14881 MMD
                                  DIR
                                                  4096
                                                                  7209106 /TESTDIR
                                        8,6
                                        8,1
8,1
                          rtd
                                                  4096
lgjgjmkkg 14881 MMD
                                                 "625718 <== NEW SIZE"
                                                                           829 /tmp/lgjgjmk
lgjgjmkkg
            14881 MMD
                                        1,3
                                                   0t0
lgjgjmkkg 14881 MMD
lgjgjmkkg 14881 MMD
                             1u
                                                    ata
                                                                      1028 /dev/null
                                  CHR
                                                                      1028 /dev/null
                             2u
                                                    0t0
```

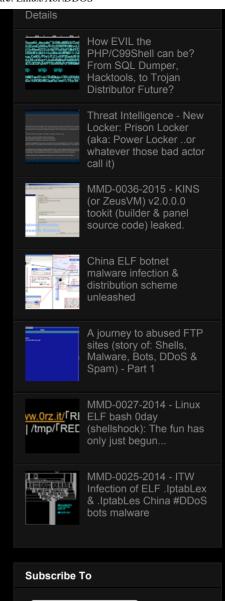
..as per seen here it runs in new PID, not clone nor *forking/threading* since execution used the shell spawning. See the new size, it gets bigger by 11 bytes.

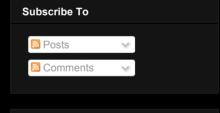
Below is the illustration of malware samples original and after copy-injected.

We have different hash and size.

Okay, we're done with the debugging and forensics. Let's see how the reverse engineering goes for this ELF malware binary for the above processes.

This is the part where the **malware self-copy process was executed** in my sample case. Noted: there are so many cases to trail with the similar codes in copying, write files and randomizing them, I counted about more than 4 scenarios prepared for this operation and the author really calculate every possibilities in his code to make sure the malware will run.





the jump to 0x804dfc2 will take you to the next process.

The assembly snip below is explaining the writing process to the done-copied file by the malware, it is not using the randomizing **11 characters** but the malware was picking a hard coded *xor crypt strings* that is saved in **0x080cf120** (symbol: *str.__Ff3VE.__*7).

```
e99c000000 jmp 0x804e06a
38 0855beebffff lea eax, [ebp - 0x1442] 890424 mov dword [esp], eax call sym.randmd5 lea eax, [ebp - 0x842]
                                   mov dword [eax], 0
mov dword [eax + 4], 0
 c700000000000
                                  mov dword [eax + 4], 0
mov word [eax + 8], 0
call sym.getpid ;sym.getpid()
mov dword [esp + 0xc], eax
mov dword [esp + 8], 0xa
lea eax, [ebp - 0x842] ;preparing print
mov dword [esp], eax
call sym.snprintf ;sym.snprintf
mov dword [esp], 0
call sym.randomid ;where the timeofday() cal
movzx eax, ax
 8944240c
c74424081132.
c74424040a00.
 e8a49c0000
c74424041700.
 c70424000000.
 e8e81e0000
0fb7c0
8945d4
8b55d4
                                    movzx eax, ax
mov dword [ebp - 0x2c], eax
mov edx, dword [ebp - 0x2c]
 89d0
c1e002
01d0
                                    shl eax 2
                                     add eax, edx
 c1e002
0520f10c08
                                     add eax, str.__Ff3VE._7; this lead to 11 chars to be added
                                   add eax, str._risve.,_
mov edx, eax
lea eax, [ebp - 0x842]
mov dword [esp + 8], eax
mov dword [esp + 4], edx
lea eax, [ebp - 0x1442]
mov dword [esp], eax
call sum linuxfxec Argv2
                                    call sym_LinuxExec_Argv2
  8345cc01
                                     add dword [ebp - 0x34], 1
837dcc04 cmp dword [ebp - 0x34], 4
```

The **snprintf()** is an *API function* that will lead (in the VERY end) to **SYS_write** at sys/syscall, since we deal with the statically compiled ELF many libc trails will appear in reversing the function, we may see more of these, sorry to say, unnecessary codes.

The **timeoftheday()** result which was shown during debugging is caused by the function which was called, named function **randomid()**.

```
98
; arg int arg_2; arg int arg_3; var int local_0_1; var int local_2; var int local_3; var int local_6; var int local_6.
                                                @ ebp+0x8
                                                 @ ebp+0xc
                                                @ ebp-0x1
@ ebp-0x8
                                                 @ ebp-0xc
                                                @ ebp-0x18
@ ebp-0x1c
                                                                      ebp
                                                            push ebx
sub esp, 0x24
                                                            sub esp, 0x24
mov eax, dword [ebp+arg_2]
mov edx, dword [ebp+arg_3]
mov word [ebp-local_6], ax
mov word [ebp-local_7], dx
                                                                                           local_3]
                                                             mov dword [esp], eax
                                                           call sym._gettimeofday_internal ;
mov eax, dword_lebp-local_2]
          04 f f 3
                            e8a0710100
0x0804ff43
                             890424
                                                             mov dword [esp], eax
```

↑Obviously, is a self-explanatory that the **timeoftheday()** is fetching the system time as the seed needed in **randomid()** function.

There is an additional information too actually: I think maybe it is good for our community to know too: Linux/XOR.DDoS ELF malware is using a uncommon seen function to execute the shell command, it was called: LinuxExec_Argv() and LinuxExec_Argv2(), which was called to act as an API to execute non direct syscall basis commands by the malware (well, this is a static compiled binary), these functions are typical in characteristic, it is a very simple in use, easy to spot (smile) and these are responsible to call execve(), a linux system call commands (with the environment parameter

parsed) to be executed during an infection, and also to call **execvp()** for the file execution purpose (with parsing the file path), i.e. shown in the code below:

```
8b4508
8945e8
8b450c
                                           mov eax, dword [ebp+arg_2]
0x08049050
0x08049053
                                           mov dword [ebp-local_6], eax
                                           mov eax, dword [ebp+arg_3]
                                           mov dword [ebp - 0x14], eax
mov eax, dword [ebp+arg_4]
mov dword [ebp - 0x10], eax
0x08049056
0x08049059
0x0804905c
                     8945ec
                     8b4510
                     8945f0
0x0804905f
0x08049062
0x08049066
                     8d45e8
89442404
                                           lea eax, [ebp-local_6]
mov dword [esp + 4], eax
mov eax, dword [ebp+arg_2]
                     8b4508
0x08049069
0x0804906c
                     890424
                                           mov dword [esp], eax
                   e85fe90100
                                           call sym.exe
0x08049071
0x08049078
                     c70424000
                                           mov dword [esp], 0
                   e8e3cd0000
                                           call sym.exit
0x0804907d
                     8b45f8
                                            mov eax, dword [ebp-local_2]
                    c9
0x08049081
                     сЗ
```

You may want to see the reference of exec method with UNIX C library (libc) on execve, execvp at man(2) pages, and yes, UNIX gods are also providing us with good reference too.

Conclusion & reference

Yes, Linux/XOR.DDoS malware after copied and executed (read: successfully infecting us) will have a different size (11 bytes bigger..depends.. I only check one binary for this), and have a different hash. So this means that the malware spotted in the panel may not be detected by the scanner used inside of the Linux box if only detecting by the hash.

Many of us still think, Yeah..ELF malware..won't harm us or end users much.. But remember, IoT are mostly linux basis, take a look of the most of router's OS now. Also, the infection method and volume of ELF malware is getting better and bigger by days. As proof: We have about 6 of new ELF malware for 2 and half years span only! As MMD (read: MalwareMustDie, NPO), we suggest to be prepared to update the ELF malware detection quality as earliest as possible, once an ELF malicious binary hit a server the impact can be way much bigger than a PE hit a PC.

Below are links to the previous Linux/XOR.DDoS analysis:.

http://blog.malwaremustdie.org/2015/07/mmd-0037-2015-bad-shellshock.html http://blog.malwaremustdie.org/2014/09/mmd-0028-2014-fuzzy-reversing-new-china.html

The "new" CNC of the threat:



Oh btw,the CNC is very alive even now...and sending the download/payload too. here's the pcap snips for a hard proof:

