

7/10 LECTURE 5: MORE EXAMPLES ON "EQUALLY LIKELY OUTCOMES".

For us prob. space, finite set Ω & pmf p or PD P .

"In school" — finite set Ω & $P(A) = \frac{|A|}{|\Omega|} = \frac{\# \text{ favourable outcomes}}{\# \text{ outcomes}}$
 Prob. of A

Such a P as above is what we call as uniform PD or "Equally likely outcomes".

For eg. by saying Ω is a prob. space with equally likely outcomes means that the PD P on Ω is uniform PD.

Why go from "Eq. lik. outcomes" to more gen. prob. spaces?

— In reality, not all outcomes are equally likely.

Eg: $\Omega = \{W, L\}$ of "eq. lik. out." then $p(W) = p(L) = 1/2$.

Eg 5.2 [sampling with replacement]

A box contains ' n ' ^{labelled} objects. We pick an object, note down the label & put it back. Then we repeat this process $(k+)$ times more.

— choosing k objects out of n (with repetition/replacement).

sample space $\Omega =$ set of k -tuples from $[n] := \{1, \dots, n\}$

$$= \{(r_1, \dots, r_k) : r_i \in [n] \forall i\}$$

[k labelled particles into ' n ' cells]

$$= [n] \times [n] \times \dots \times [n] \quad (A \times B \text{ — Cartesian product})$$

$$[A \times B = \{(a, b) : a \in A, b \in B\}, \text{ } A \text{ \& B are sets}]$$

$$=: [n]^k$$

$$(A^k = A \times \dots \times A \text{ } k\text{-times})$$

$$A^k = \{(a_1, \dots, a_k) : a_i \in A \forall i\}$$

Equally likely outcomes $\rightarrow p(w) = \frac{1}{|\Omega|} = \frac{1}{n^k}$, $w \in \Omega$
 $w = (w_1, \dots, w_k)$
 (Ω, p) or (Ω, P) is said as "sampling with replacement from $[n]$ k -many times".

Ex 3: [hashing] There are k individuals & each individual is assigned a code. There are n distinct codes/labels.

For proper identification, codes of k ind. "should" be distinct.

Keeping track of prev. codes isn't easy.

Say a 'lazy' person assigns (hash) codes randomly.

Sample space = set of all ^{possible assignment of} (hash) codes for k individuals.

$$= \{ (h_1, \dots, h_k) : h_i \in [n] \forall i \}$$

$$= [n]^k \quad h_i - \text{hash code of } i^{\text{th}} \text{ person.}$$

All assignments are equally likely.

$$\Rightarrow p(w) = \frac{1}{n^k} \quad w = (w_1, \dots, w_k) \in [n]^k$$

(Error) $E = \{h_1, \dots, h_k \text{ are not distinct}\}$

(Good) $G = E^c = \{h_i \text{'s are distinct}\}$

$$= \{ (h_1, \dots, h_k) \in \Omega : h_i \neq h_j \forall i \neq j \}$$

Prob. (lazy hash code is good) $= P(G)$ (defn of G)

$$= \frac{|G|}{|\Omega|} \quad (\text{Eq. likely outcomes})$$

$$= \frac{|G|}{n^k}$$

$$|G| = n(n-1) \cdot \dots \cdot (n-k+1) = \prod_{i=0}^{k-1} (n-i)$$

$$P(G) = \prod_{i=0}^{k-1} \left(1 - \frac{i}{n}\right)$$

Eg 5.4 (Birthday problem) — if there are k people in a room
 what is the chance a pair shares same birthday? [Von Mises & others?]
 label days as $1, \dots, 365$; Set $n = 365$.

b_i = birthday of i^{th} person.

Ω = all possible birthdays of k people

$$= \{ (b_1, \dots, b_k) : b_i \in [n] \forall i \}$$

(Same as sampling with repl. or hashing)

Assume any ordering of b_i 's is equally likely.

$$P(w) = \frac{1}{n^k} \quad w = (w_1, \dots, w_k) \in \Omega$$

$$IP(\exists_{\text{at least one pair}} \text{ sharing same birthday}) = IP(\{ (b_1, \dots, b_k) : b_i = b_j \text{ for some } i \neq j \})$$

(from computation of $P(B)$ in Eg 5.3)

$$= 1 - \prod_{i=0}^{k-1} \left(1 - \frac{i}{n}\right) = c(k).$$

$n = 365$

$c(2) = 0.002$, $c(5) = 0.027$, $c(10) = 0.116$, $c(23) = 0.507$
 if there are 23 people, $IP(\text{at least 2 people share same birthday}) > 1/2$.

$$c(40) > 0.8 \quad c(50) \approx 1.$$

$IP(\text{person 1 shares a birthday with one another})$

$$= 1 - IP(\text{person 1 doesn't share a birthday with anybody})$$

$$= 1 - IP(\{ (b_1, \dots, b_k) : b_j \neq b_1 \forall j \neq 1 \})$$

$$= 1 - \frac{|\{ (b_1, \dots, b_k) : b_j \neq b_1 \forall j \neq 1 \}|}{n^k}$$

$$= 1 - \frac{\underbrace{(n-1) \cdot (n-1)^{k-1}}_{b_1}}{n^k} = 1 - \left(1 - \frac{1}{n}\right)^{k-1} = b(k)$$

$$b(40) \leq 0.11, \quad c(40) \geq 0.89.$$

$$\approx \frac{40}{365}$$

THM 4.1 Let (Ω, \mathcal{P}) be a fin. prob. space. Then

(i) $P(A) \leq P(B) \quad \forall A \subseteq B \subseteq \Omega$. [Monotonicity]

(ii) $P(A) \leq 1 \quad \forall A \subseteq \Omega$

(iii) $P(\emptyset) = 0$

(iv) $P(A_1 \cup \dots \cup A_n) \leq \sum_{i=1}^n P(A_i) \quad \forall A_1, \dots, A_n \subseteq \Omega$. [finite subadditivity]

(v) $P(A \cup B) = P(A) + P(B) - P(A \cap B)$ [Inclusion-Exclusion]

Proof: (i) $P(A) = \sum_{\omega \in A} p(\omega) \leq \sum_{\omega \in B} p(\omega) = P(B) \quad (A \subseteq B \text{ \& } p(\omega) \geq 0 \quad \forall \omega)$
= $P(B)$ (defn)

(ii) Set $B = \Omega$. (i) $\Rightarrow P(A) \leq P(\Omega) = 1$
 \hookrightarrow (prop. of P)

(iii) $P(A \cup \emptyset) = P(A) + P(\emptyset) \quad \text{as } A \cap \emptyset = \emptyset$.
 $\Rightarrow P(A) = P(A) + P(\emptyset) \quad \text{as } A \cup \emptyset = A$.
 so $P(\emptyset) = 0$.

(iv) we'll prove by induction.

$n=1$ is trivially true. $P(A_1) \leq P(A_1) \quad \checkmark$

Let $n=2$. Then $A_1 \cup A_2 = A_1 \sqcup (A_2 \setminus A_1)$

$P(A_1 \cup A_2) = P(A_1 \sqcup (A_2 \setminus A_1))$
 $= P(A_1) + P(A_2 \setminus A_1) \quad \left(\begin{array}{l} \text{finite} \\ \text{additivity of } P \end{array} \right) \quad \left(\begin{array}{l} \text{disjoint union} \\ (A_1 \cap (A_2 \setminus A_1) = \emptyset) \end{array} \right)$
 $\leq P(A_1) + P(A_2)$
(monotonicity as $A_2 \setminus A_1 \subseteq A_2$)

Assume finite subadd. holds for $n-1 \geq 2$.

Consider A_1, \dots, A_n . We've to prove fin. subadd for n .

$$A_1 \cup \dots \cup A_n = A_1 \cup (A_2 \cup \dots \cup A_n / A_1)$$

$$\begin{aligned} P(A_1 \cup \dots \cup A_n) &= P(A_1) + P(A_2 \cup \dots \cup A_n / A_1) \quad (\text{fin. add.}) \\ &\leq P(A_1) + P(A_2 \cup \dots \cup A_n) \quad (\text{mon.}) \\ &\leq P(A_1) + P(A_2) + \dots + P(A_n) \quad (\text{fin. subadd. holds for } n-1) \end{aligned}$$

$$(v) \quad A \cup B = A \cup (B / A \cap B)$$

$$\begin{aligned} P(A \cup B) &= P(A) + P(B / A \cap B) \quad (\text{fin. add.}) \\ &= P(A) + P(B) - P(A \cap B) \quad (A \cap B \subseteq B \text{ \& use Lemma 4.4}) \end{aligned}$$

□