

1. Recall the extension of Euclid's algorithm that we discussed in class.

```

1  def extended_Euclid(a,b):
2      """
3      a, b are non-negative integers
4      The function returns (u, v, d) such that d = gcd(a,b)
5      and d = ua + vb
6      """
7
8      if b == 0: return (1, 0, a)
9      (u, v, d) = extended_Euclid(b, a % b)
10     return (v, u - v * (a//b), d)

```

We argued in class that the algorithm correctly returns  $(u, v, d)$  as stated in the comment in the beginning of the code. Suppose for a certain input  $(a, b)$ , where  $a > b \geq 1$ , the call to `extended_Euclid(a, b)` executes line 9 a total of  $t$  times (where  $t \geq 1$ ). Let the value of  $(a, b)$  in the  $i$ -th call to `extended_Euclid(a, b)` be  $(a_i, b_i)$ ; let  $(a_0, b_0) = (a, b)$ . Let the value  $(u, v)$  returned by the  $i$ -th call be  $(u_i, v_i)$ , so that  $u_i a_i + v_i b_i = d$ ; thus  $(u_t, v_t) = (1, 0)$ . Then, for  $i = 1, 2, \dots, t$ , we have

$$\begin{bmatrix} a_{i-1} \\ b_{i-1} \end{bmatrix} = \begin{bmatrix} q_i & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} a_i \\ b_i \end{bmatrix};$$

$$\begin{bmatrix} u_{i-1} & v_{i-1} \end{bmatrix} = \begin{bmatrix} u_i & v_i \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & -q_i \end{bmatrix},$$

where  $q_i$  is the quotient obtained on dividing  $a_{i-1}$  by  $b_{i-1}$ .

- Show that  $|u_i| \leq a_i/d$  and  $|v_i| \leq b_i/d$ , where  $d = \gcd(a, b)$ . You may use induction to show that the claim holds for  $i - 1$  assuming it holds for  $i$ ; what is the base case?
- Suppose  $a$  and  $b$  are  $n$ -bit integers. Show that the total number of bit operations needed for `extended_Euclid(a, b)` is  $O(n^3)$ , assuming that integer division of  $\ell$ -bit integers can be done in using  $O(\ell^2)$  bit operations.

2. Consider the following modification to Euclid's algorithm.

```

1  def modified_Euclid(a,b):
2      """
3      a, b are non-negative integers
4      The function returns (u, v, d) such that d = gcd(a,b)
5      and d = ua + vb
6      """
7
8      if b == 0: return a
9      r = a % b
10     if r < b/2:
11         return modified_Euclid(b, r)
12     else:
13         return modified_Euclid(b, b-r)

```

- Argue that for integers  $a > b > 0$ , `modified_Euclid(a, b)` returns the gcd of  $a$  and  $b$ .

- (b) How many times is `modified_Euclid` is called recursively after `modified_Euclid(a, b)` is called with Fibonacci numbers  $a = F_{t+1}$  and  $b = F_t$ ?
3. Describe an algorithm to determine if a given positive number  $N \geq 2$  can be written in the form  $N = Q^E$ , where  $Q$  and  $E$  are both integers at least 2. For  $n$ -bit numbers  $N$ , your algorithm should run in time  $O(n^k)$  for some small constant  $k$  (fixed independent of  $n$ ).
4. Suppose  $x$ ,  $y$  and  $\ell$  are  $n$ -bit numbers, such that  $x > y$ . Suppose the binary expansion of the fraction  $x/y$  is

$$0.b_0b_1b_2b_3 \dots = \sum_{i \geq 1} b_i 2^{-i},$$

which in general may not terminate. Describe an algorithm to determine  $b_\ell$ , given  $x$ ,  $y$  and  $\ell$ . Your algorithm should run in time  $O(n^k)$  for some small constant  $k$  (fixed independent of  $n$ ).

**(Due 30 Aug 2023)**