

Recall: Stopping times & Games

Interpretation: S_n = represent "amount of capital of the player after n rounds"

X_k = amount a player wins in round k .

* Expected "amount of capital" after n rounds
 $= E[S_n] = 0 \quad 0 \leq n \leq N$

Question: Is it possible to stop the game in a favorable moment?

Definition: An event $A \in \mathcal{F}$ is observable until time n when it can be written as a union of basic events of the form

$$\{\omega \in \Omega \mid \omega_1 = o_1, \dots, \omega_n = o_n\} \quad o_i \in \{-1, 1\} \quad 1 \leq i \leq n$$

i.e. A - can be determined from the outcome of the 1st n tosses.

$\mathcal{A}_n :=$ class of event A that can be observed by time n . [include \emptyset]

Definition: A map $T: \Omega \rightarrow \{0, 1, \dots, N\} \cup \{\infty\}$ is called a stopping time if

$$\{T = n\} \equiv \{\omega \in \Omega \mid T(\omega) = n\} \in \mathcal{A}_n \quad 0 \leq n \leq N$$

i.e. $\{T = n\}$ is an event observable until time n .

Example: For $a \in \mathbb{Z}$ let

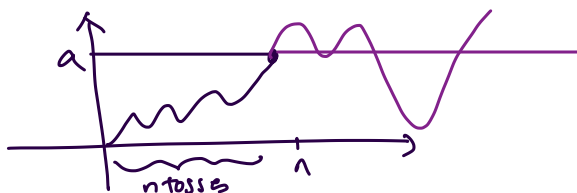
$$\sigma_a: \Omega \rightarrow \{0, 1, \dots, N\} \cup \{\infty\}$$

$$\sigma_a(\omega) = \min \{k \in [0, N] \mid S_k(\omega) = a\}$$

$$\min(\emptyset) = \infty$$

Ex:- $\sigma_a(\cdot)$ is a stopping time

$\{\omega \in \Omega : \sigma_a(\omega) = n\} = \dots$ "only depends on 1st n tosses"



Theorem 1 : For any stopping time

[Impossibility of a favorite stopping] $T: \Omega \rightarrow \{0, 1, \dots, N\}$

$$E[S_T] = 0$$

$S_T(\omega) = S_{T(\omega)}(\omega) \equiv$ outcome of the trajectory ω at stopping time $T(\omega)$

Proof :- T is a stopping time.

$$k=0, 1, 2, \dots, N \quad \{T \geq k\} \in \mathcal{A}_{k-1} \quad - (1)$$

$$[\because \{T \geq k\} = (\{T=1\} \cup \dots \cup \{T=k-1\})^c]$$

$\in \mathcal{A}_{k-1} \in \mathcal{A}_{k-1}$ is closed under complements

$$S_T = \sum_{k=1}^N X_k \mathbb{1}(T \geq k) \quad - (2)$$

$$[\because S_T = \sum_{k=1}^N S_k \mathbb{1}(T \geq k)]$$

$$= \sum_{k=1}^{N-1} S_k [\mathbb{1}(T \geq k) - \mathbb{1}(T \geq k+1)] + S_N \mathbb{1}(T=N)$$

$$= \sum_{k=1}^{N-1} S_k \mathbb{1}(T \geq k) - \sum_{k=1}^{N-1} S_k \mathbb{1}(T \geq k+1) + S_N \mathbb{1}(T=N)$$

$$= S_1 \mathbb{1}_{T \geq 1} + \sum_{k=2}^{N-1} S_k \mathbb{1}_{(T \geq k)} + S_N \mathbb{1}_{T=N} - \sum_{k=1}^{N-1} S_k \mathbb{1}_{(T \geq k+1)}$$

$$= S_1 \mathbb{1}_{T \geq 1} + \sum_{n=1}^{N-1} S_{n+1} \mathbb{1}_{(T \geq n+1)} + S_N \mathbb{1}_{T=N} - \sum_{k=1}^{N-1} S_k \mathbb{1}_{(T \geq k+1)}$$

$$= S_1 \mathbb{1}_{T \geq 1} + \sum_{\ell=1}^{N-1} (S_{\ell+1} - S_\ell) \mathbb{1}_{T \geq \ell+1} + (S_N - S_{N-1}) \mathbb{1}_{T=N}$$

$$\begin{aligned} \{T=N\} \\ \equiv \{T \geq N\} \end{aligned} = \left[\begin{array}{c} \downarrow \\ x_1 \mathbb{1}_{T \geq 1} \end{array} + \sum_{\ell=1}^{N-1} \left[\begin{array}{c} \downarrow \\ x_\ell \mathbb{1}_{T \geq \ell+1} \end{array} + \begin{array}{c} \downarrow \\ x_N \mathbb{1}_{T=N} \end{array} \right] \right]$$

$$E[S_T] \stackrel{(2)}{=} E\left[\sum_{k=1}^N x_k \mathbb{1}_{T \geq k}\right]$$

linearity of $E \leftarrow \sum_{k=1}^N E\left[x_k \mathbb{1}_{T \geq k}\right]$

\downarrow depends on the k^{th} toss $\downarrow \omega_k$
 $\underbrace{\mathbb{1}_{T \geq k}}_{\in A_{k-1}} \downarrow \text{"}\omega_1, \dots, \omega_{k-1}\text{"}$

independence $\xleftarrow{Ex} = \sum_{k=1}^N E[x_k] E[\mathbb{1}_{T \geq k}]$

\downarrow

$$= \sum_{k=1}^N 0 \cdot \mathbb{P}(T \geq k) = 0$$

□

Game System: $V_1, \dots, V_N: \Omega_N \rightarrow \mathbb{R}$

are random variables. Such that

$$\{V_k = c\} \in \mathcal{A}_{k-1} \quad c \in \mathbb{R}, \quad k=1, 2, \dots, N$$

$V_k \equiv$ amount of money you will place as a bet
in the k^{th} round
Result in the k^{th} round $= V_k X_k$

$$S_N^V = \sum_{k=1}^N V_k X_k$$

Theorem 2: For any V_1, V_2, \dots, V_N game system

then

$$E[S_N^V] = 0$$

Proof:

like before

$$E[S_N^V] = \sum_{k=1}^N E[V_k X_k] - (\#)$$

It is enough to show $E[V_k X_k] = 0$.

\mathcal{A}_{k-1}

depends on
 k^{th} toss

$$\text{Range}(v_k) = \{c_1, c_2, \dots, c_m\}$$

$$\Rightarrow X_k v_k = \sum_{i=1}^m c_i X_k \mathbb{1}(v_{1k} = c_i)$$

$$E[X_k v_k] = \sum_{i=1}^m c_i E[X_k \mathbb{1}(v_{1k} = c_i)]$$

\swarrow \nwarrow
 $"X_k"$ $"v_1, \dots, v_{k-1}"$

$$= \sum_{i=1}^m c_i E[X_k] P(v_{1k} = c_i)$$

$$= \dots 0 \dots = 0$$

\downarrow
 $\hookrightarrow \textcircled{x}$

Place \textcircled{x} into $\textcircled{\#}$ \Rightarrow to get result \square

Impact of Theorem 2:

$T: \Omega \rightarrow \{0, 1, \dots, N\}$ is a stopping time

$$v_k = \mathbb{1}_{\{T \geq k\}}$$

$$S_N^v = \sum_{k=1}^N X_k v_k = \sum_{k=1}^N X_k \mathbb{1}_{\{T \geq k\}}$$

$$\stackrel{\textcircled{2}}{=} S_T$$

[Theorem 2 \Rightarrow Theorem 1]

$$\bullet \quad v_k = \underbrace{S_{k-1}}_{(v_1, \dots, v_{k-1})} \underbrace{\perp_{T \geq k}}_{\textcircled{1}} \quad (\text{take this definition})$$

$\Rightarrow v_1, v_2, \dots, v_n$ is a game system.

$$v_k x_k = S_{k-1} x_k \perp_{T \geq k}$$

Observe $\therefore S_k^2 = (S_{k-1} + x_k)^2$

$$= S_{k-1}^2 - 2 \underline{S_{k-1} x_k} + x_k^2$$

$$\Rightarrow v_k x_k = \frac{(S_k^2 - S_{k-1}^2 - x_k^2)}{2} \perp_{T \geq k}$$

$$\sum_k^n v_k x_k = \sum_{k=1}^n v_k x_k$$

$$= \sum_{k=1}^n \frac{(S_k^2 - S_{k-1}^2 - x_k^2)}{2} \perp_{T \geq k}$$

$$\begin{aligned} (x_k^2 = 1) &= \sum_{k=1}^n \frac{S_k^2}{2} \perp_{T \geq k} - \sum_{k=1}^n \frac{S_{k-1}^2}{2} \perp_{T \geq k} \\ &\quad - \sum_{k=1}^n \frac{1}{2} \perp_{T \geq k} \end{aligned}$$

$$= \sum_{k=1}^n \frac{S_k^2}{2} \perp_{T=k} - \frac{T}{2}$$

$$S_n^V = \frac{(S_T^2 - T)}{2} - \textcircled{xxx}$$

Theorem 2 $\Rightarrow E[S_n^V] = 0$

\therefore By \textcircled{xxx} for any $T: \Omega \rightarrow \{0, 1, \dots, N\}$

we have

$$0 = E[S_n^V] = E\left[\frac{S_T^2 - T}{2}\right]$$

$$\Rightarrow E[S_T^2] = E[T] - \textcircled{\text{||||}}$$

Theorem 1: $E[S_T] = 0$

$\textcircled{\text{||||}}$
 \Rightarrow

$$\text{Var}[S_T] = E[T]$$

1.3 Ruin Problem:

Two players. Each with capital $\begin{cases} a \sim \text{I} \\ b \sim \text{II} \end{cases}$

Interpret

$S_n \equiv$ gain of player 1

Ruin (of player 1)

$$\{\sigma_a < \sigma_b; \sigma_a \leq N\}$$

where, $\sigma_u = \min \{ k : S_k = u \}$

Define: Ruin Probability

$$\delta'_N = \mathbb{P}(\sigma_{-a} \leq \sigma_b, \sigma_{-a} \leq N)$$

$$T_N = \min(\sigma_{-a}, \sigma_b, N)$$

Ex: T_N is a stopping time

$$\therefore \underline{\text{Theorem 1}} \Rightarrow E[S_{T_N}] = 0 \quad - (3)$$

$$T_N = \sigma_{-a} \Rightarrow S_{T_N} = -a$$

$$T_N = \sigma_b \Rightarrow S_{T_N} = b$$

$$T_N = N \Rightarrow S_{T_N} = S_N$$

- Use this in (3)

$$0 = -a \mathbb{P}(T_N = \sigma_{-a}) + b \mathbb{P}(T_N = \sigma_b) + E[S_N \mid T_N = N]$$

$$\text{As } T_N = \min(\sigma_{-a}, \sigma_b, N)$$

$$0 = -a \mathbb{P}(\sigma_{-a} < \sigma_b \text{ \& \> } \sigma_{-a} \leq N)$$

$$+ b P(\sigma_b < \sigma_{-a} \leq \sigma_b \leq n)$$

$$+ E S_n (1_{\min(\sigma_{-a}, \sigma_b) \geq n})$$

\Rightarrow

$$0 = -a P_n(\text{Ruin of Player 1})$$

$$+ b P_n(\text{Ruin of Player 2})$$

$$+ E [S_n 1_{\min(\sigma_{-a}, \sigma_b) \geq n}]$$

(4)

$$\bullet \quad r^1 = \lim_{n \rightarrow \infty} P_n(\text{Ruin of Player 1}) \quad \left. \vphantom{\lim} \right\} \text{Extinction}$$

$$\bullet \quad r^2 = \lim_{n \rightarrow \infty} P_n(\text{Ruin of Player 2})$$

$$(4) \quad \begin{matrix} E_x \\ \Rightarrow \\ \lim_{n \rightarrow \infty} \end{matrix} \quad 0 = -a r^1 + b r^2 + 0$$

$$\Rightarrow -a r^1 + b r^2 = 0 \quad (5)$$

$$r^1 + r^2 = 1 \quad (6)$$

Ex

$$(5) \text{ et } (6)$$

$$r^1 = \frac{b}{a+b} \quad r^2 = \frac{a}{a+b} !$$

□

