

Design and Analysis of Algorithms

BFS, DFS, and topological sort

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Graph Traversal

Two techniques of graph traversing for both directed or undirected graphs:

- ❖ Breadth First Search (BFS) – Uses Queue
- ❖ Depth First Search (DFS) – Uses Stack

What are BFS and DFS?

- Two ambiguous terms: search, traversal
- Visit each of vertices once
 - E.g.: tree walks of a binary search tree
 - Traversal
 - Search
- Start from a vertex, visit all the **reachable** vertices
 - Search

What are BFS and DFS?

- To eliminate the ambiguity, in my class
- Search indicates
 - Start from a vertex, visit all the **reachable** vertices
- Traversal indicates
 - Visit each of vertices once
- However, in other materials, you may see some time “search” is considered as “traversal”

What are BFS and DFS?

- BFS
 - Start from a vertex, visit all the reachable vertices in a breadth first manner
- DFS
 - Start from a vertex, visit all the reachable vertices in a depth first manner
- BFS or DFS based traversal
 - Repeat BFS or DFS for unreachable vertices

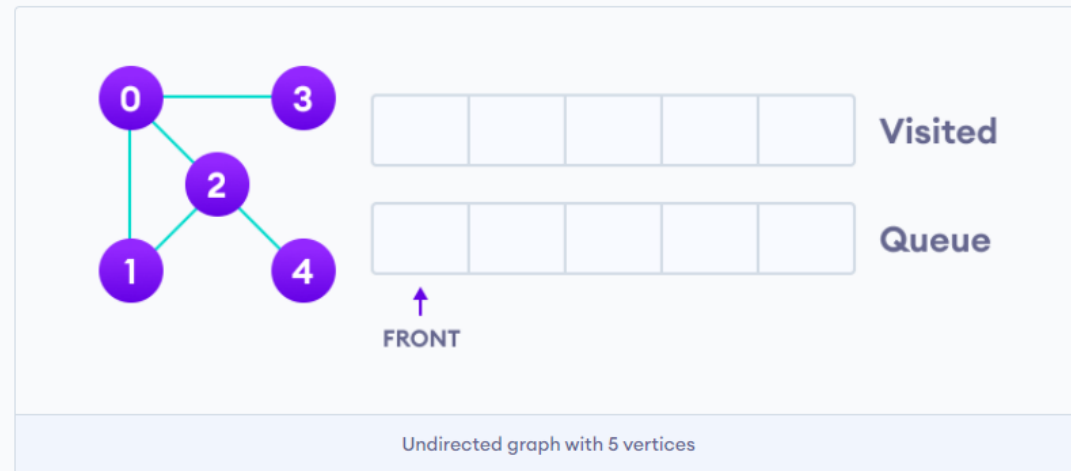
BFS, BF tree and shortest path

- Breadth-first search
 - From a source vertex s
 - Breadth-firstly explores the edges to discover every vertex that is reachable from s
- **BFS**(s)

```
visit(s);
queue.insert(s);
while( queue is not empty ){
    u = queue.extractHead();
    for each edge <u, d>{
        if(d has not been visited)
            visit(d);
            queue.insert(d);
    }
}
```

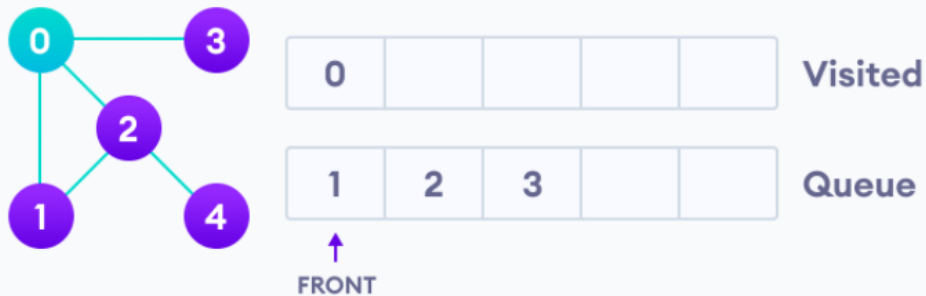
Graph Traversal : BFS

Let's see how the Breadth First Search algorithm works with an example. We use an undirected graph with 5 vertices.



Graph Traversal : BFS

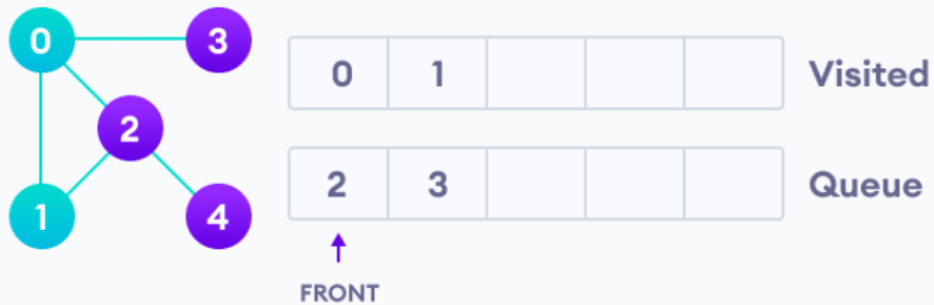
We start from vertex 0, the BFS algorithm starts by putting it in the Visited list and putting all its adjacent vertices in the stack.



Visit start vertex and add its adjacent vertices to queue

Graph Traversal : BFS

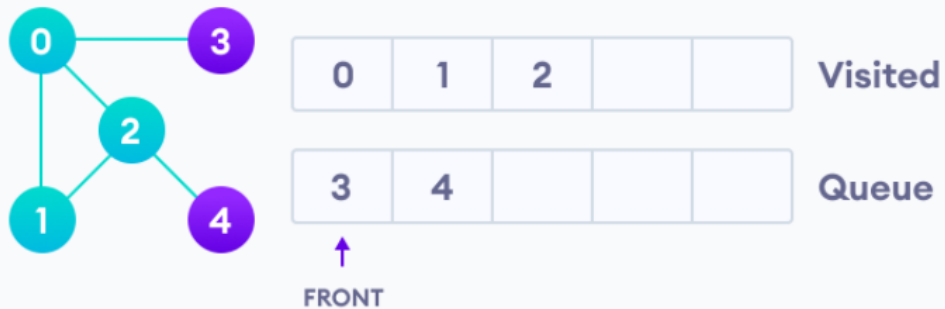
Next, we visit the element at the front of queue i.e. 1 and go to its adjacent nodes. Since 0 has already been visited, we visit 2 instead.



Visit the first neighbour of start node 0, which is 1

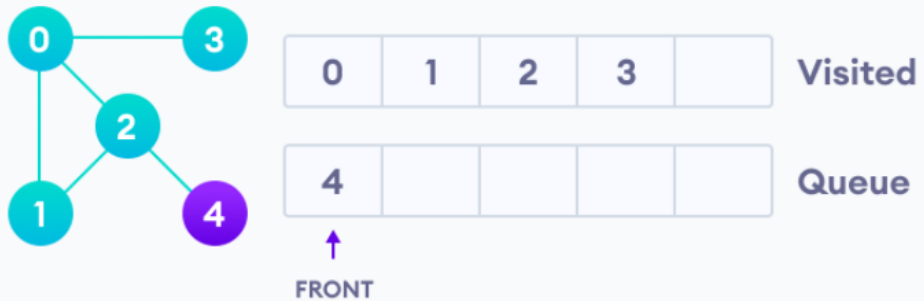
Graph Traversal : BFS

Vertex 2 has an unvisited adjacent vertex in 4, so we add that to the back of the queue and visit 3, which is at the front of the queue.



Visit 2 which was added to queue earlier to add its neighbours

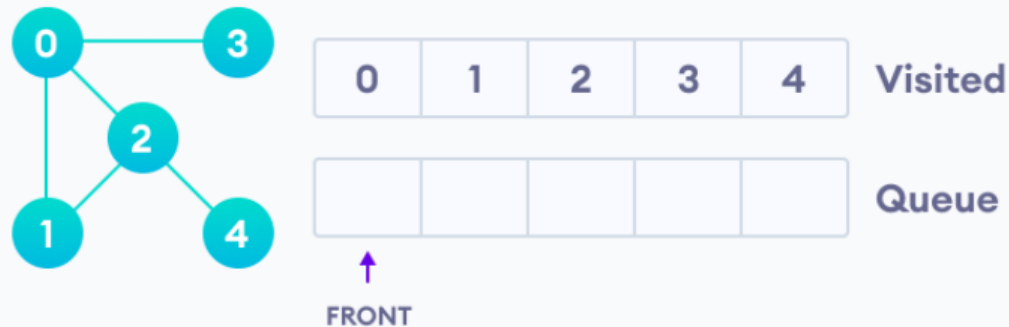
Graph Traversal : BFS



4 remains in the queue

Only 4 remains in the queue since the only adjacent node of 3 i.e. 0 is already visited. We visit it.

Graph Traversal : BFS

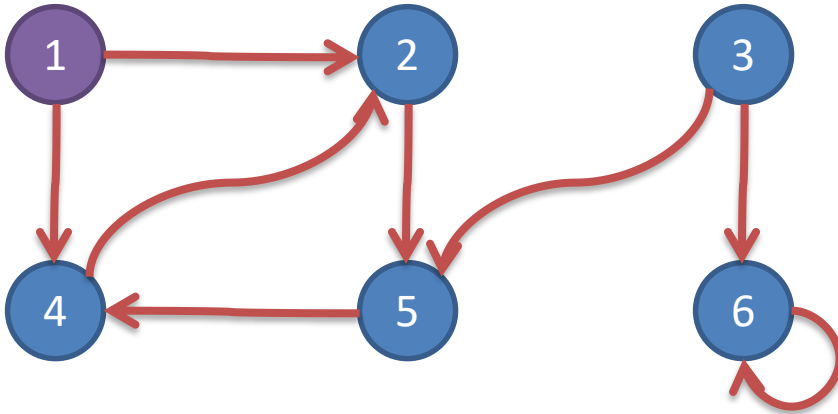


Visit last remaining item in the stack to check if it has unvisited neighbors

Since the queue is empty, we have completed the Breadth First Traversal of the graph.

BFS, BF tree and shortest path

BFS(1)



Queue: 1 4 2 5

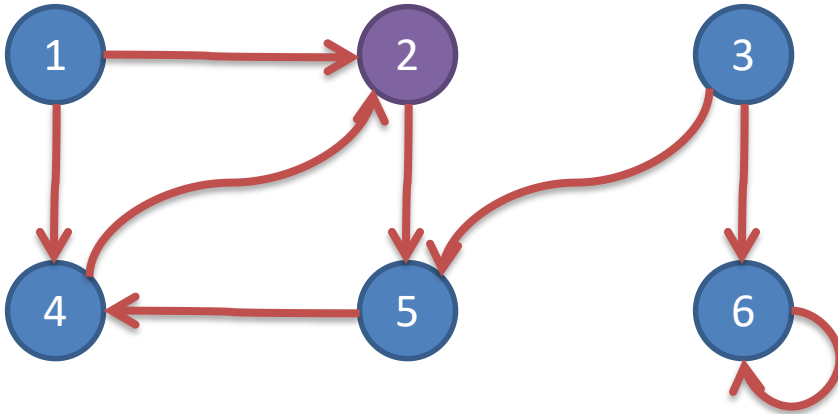
Visit order: 1 4 2 5

BFS(s)

```
visit(s);
queue.insert(s);
while( queue is not empty ){
    u = queue.extractHead();
    for each edge <u, d>{
        if(d has not been
           visited)
            visit(d);
            queue.insert(d);
    }
}
```

BFS, BF tree and shortest path

BFS(2)



Queue: 2 5 4

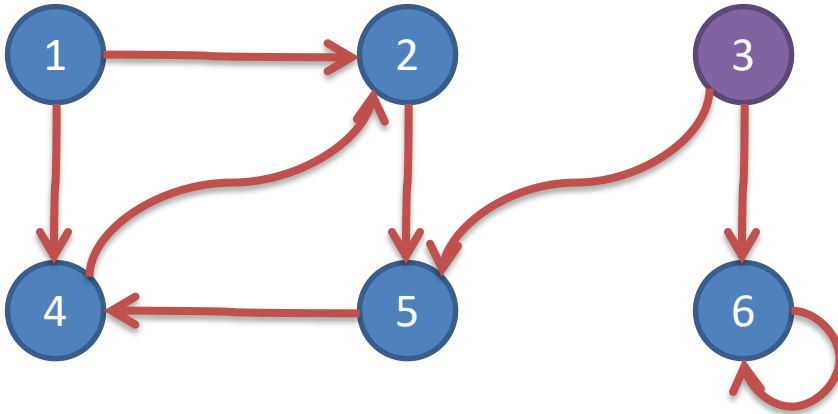
Visit order: 2 5 4

BFS(s)

```
visit(s);
queue.insert(s);
while( queue is not empty ){
    u = queue.extractHead();
    for each edge <u, d>{
        if(d has not been
           visited)
            visit(d);
            queue.insert(d);
    }
}
```

BFS, BF tree and shortest path

BFS(3)



Queue: 3 6 5 4 2

Visit order: 3 6 5 4 2

BFS(s)

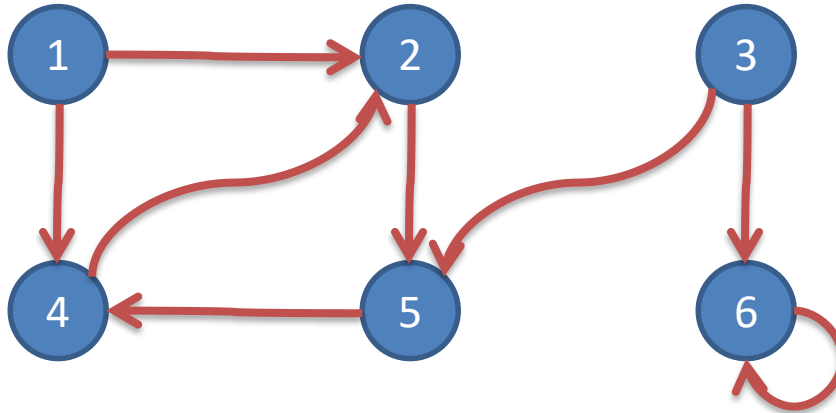
```
visit(s);
queue.insert(s);
while( queue is not empty ){
    u = queue.extractHead();
    for each edge <u, d>{
        if(d has not been
           visited)
            visit(d);
            queue.insert(d);
    }
}
```

Note that: no matter visit 5 first or visit 6 first, they are BFS

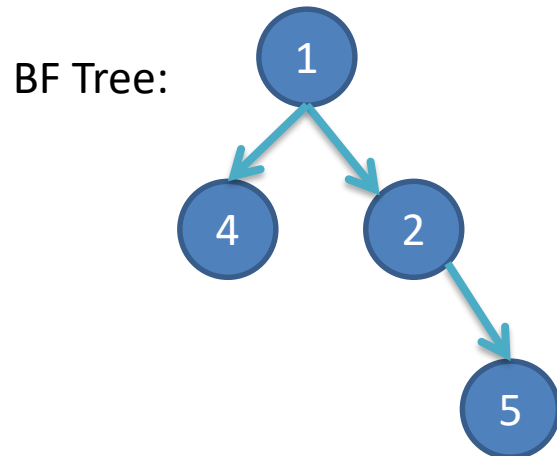
BFS, BF tree and shortest path

- Byproducts of BFS(s)
 - Breadth first tree
 - The tree constructed when a BFS is done
 - Shortest path
 - A path with minimum number of edges from one vertex to another
 - BFS(s) find out all the shortest paths from s to all its reachable vertices

BFS, BF tree and shortest path

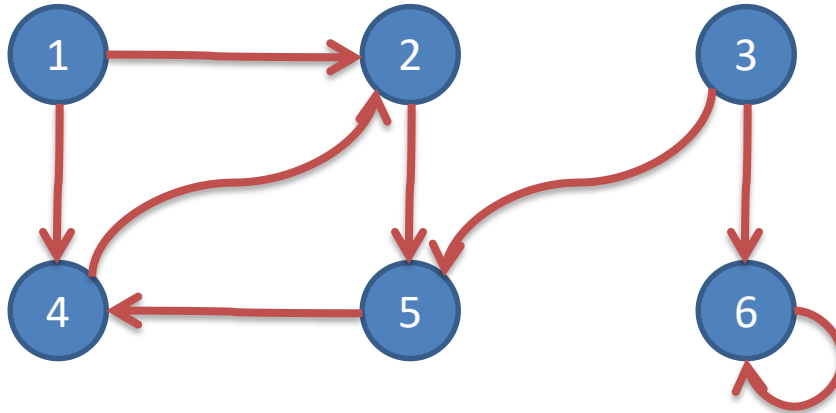


BFS(1): 1 4 2 5



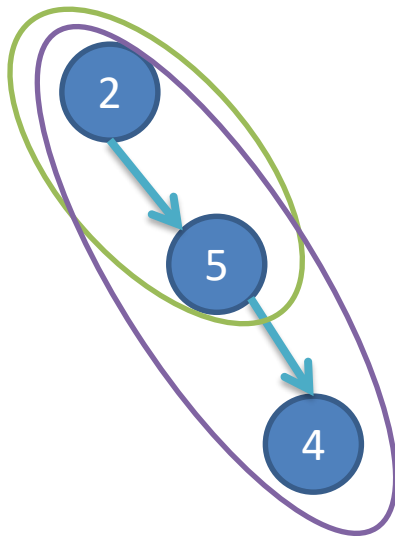
All shortest paths started
from vertex 1 are found
e.g. 1 to 5

BFS, BF tree and shortest path



BFS(2): 2 5 4

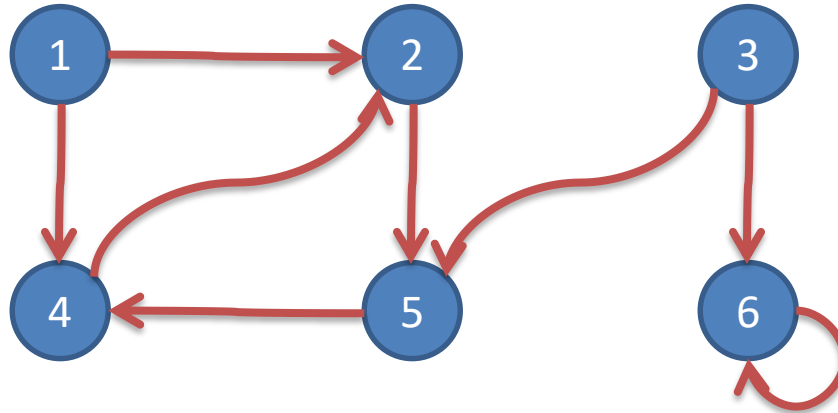
BF Tree:



Shortest 2 to 5

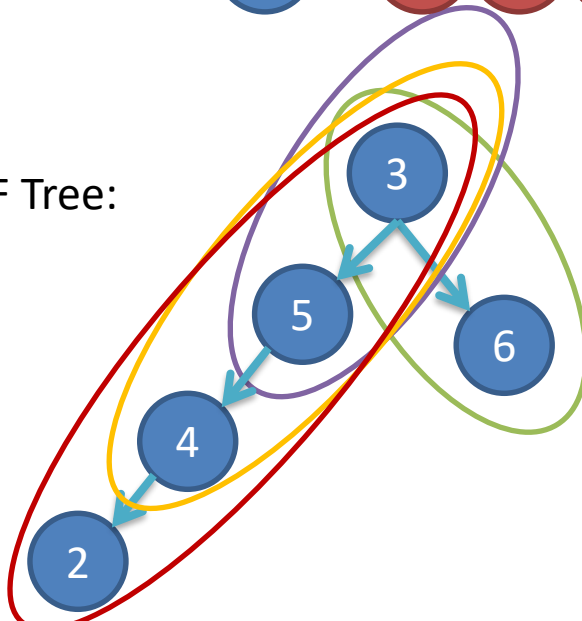
Shortest 2 to 4

BFS, BF tree and shortest path



BFS(3): 3 6 5 4 2

BF Tree:



Shortest 3 to 6

Shortest 3 to 5

Shortest 3 to 4

Shortest 3 to 2

BFS, BF tree and shortest path

- BFS Traversal
- **BFS_Traversal(G)**

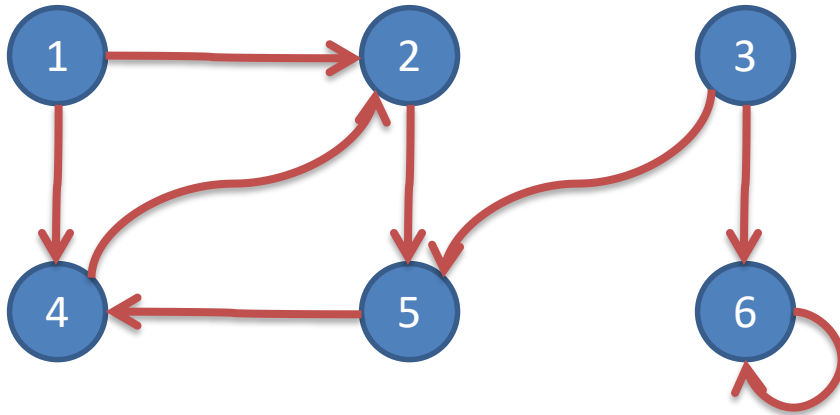
for each v in G {

 if (v has not been visited)

BFS(v);

}

BFS, BF tree and shortest path



Queue: 1 4 2 5 3 6

Visit order: 1 4 2 5 3 6

```
BFS_Traversal(G)
for each v in G{
    if (v has not been visited)
        BFS(v);
}
```

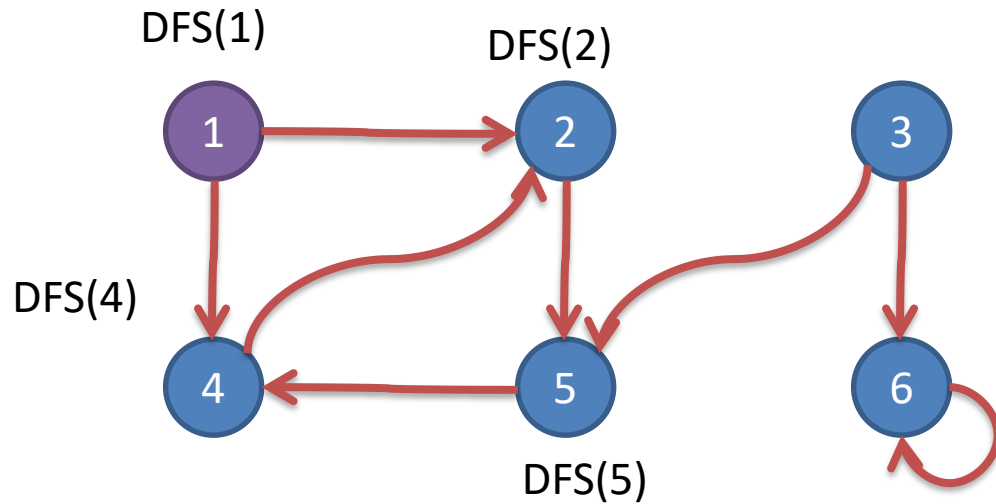
DFS, DF tree

- Depth-first search
 - From a source vertex s
 - Depth-firstly search explores the edges to discover every vertex that is reachable from s
- **DFS**(s):
 - $s.\text{underDFS} = \text{true};$ // grey
 - for each edge $\langle s, d \rangle \{$
 - if(! $d.\text{underDFS}$ and d has not been visited)
 - DFS**(d);
 - }
 - Visit(s); // black

From the deepest one to
the current one

DFS, DF tree

DFS(1)



DFS(s):

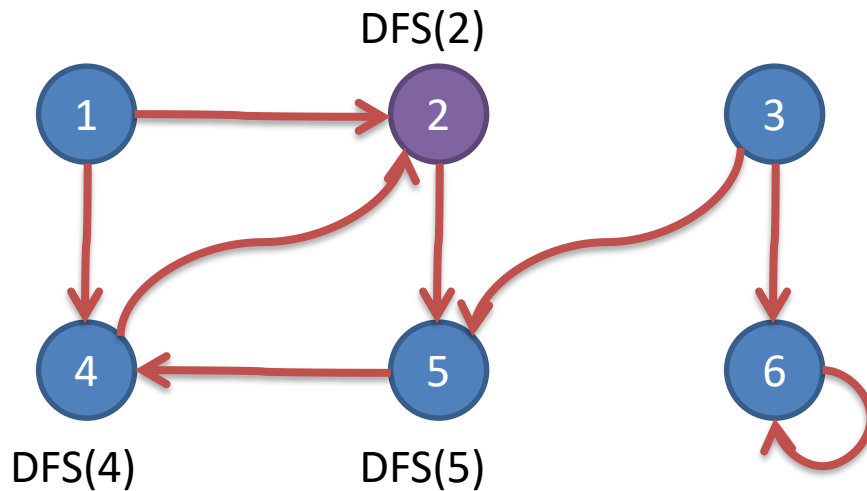
```
s.underDFS = true;  
for each edge <s, d>{  
    if((! d.underDFS and d has  
        not been visited)  
        DFS(d);  
}  
Visit(s);
```

Visit order:



DFS, DF tree

DFS(2)



DFS(s):

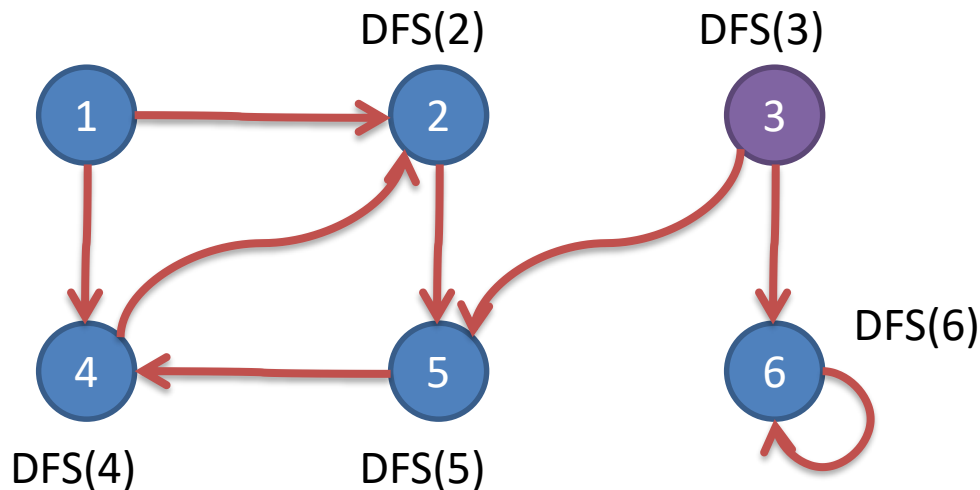
```
s.underDFS = true;  
for each edge <s, d>{  
    if((! d.underDFS and d has  
        not been visited)  
        DFS(d);  
}  
Visit(s);
```

Visit order:



DFS, DF tree

DFS(3)



DFS(s):

```
s.underDFS = true;  
for each edge <s, d>{  
    if((!d.underDFS and d has  
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        DFS(d);  
}  
Visit(s);
```

Visit order:



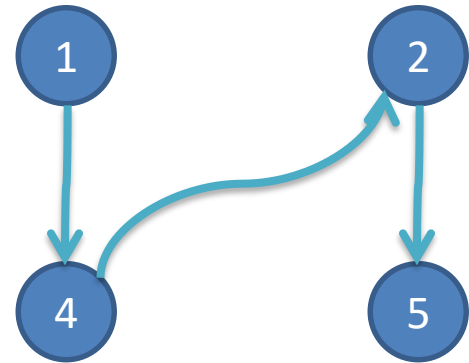
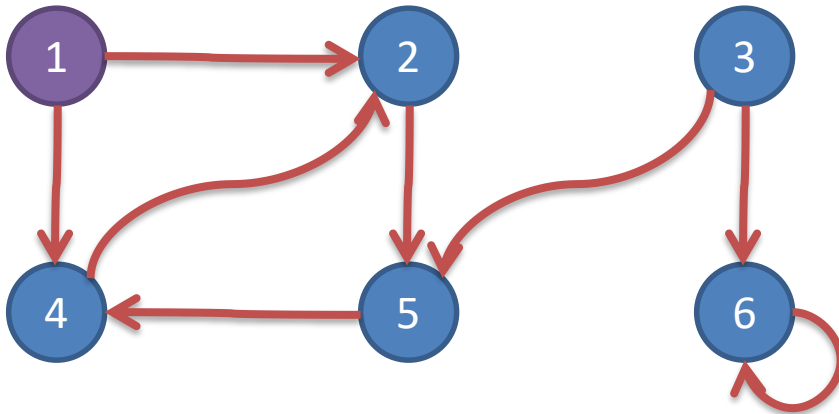
The reachable vertices are exactly the same with
BFS, but with a different order

DFS, DF tree

- Depth first tree
 - The tree constructed when a DFS is done

DFS, DF tree

DF Tree of **DFS(1)**



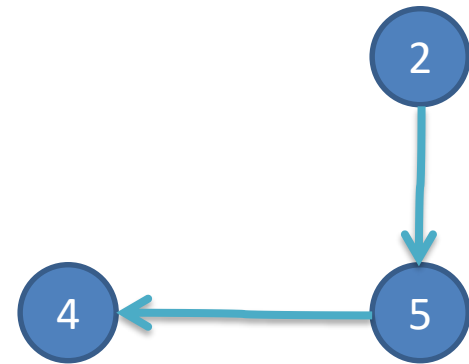
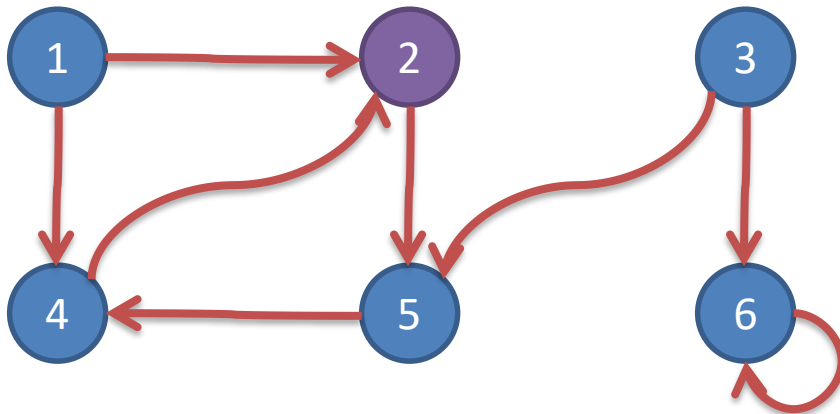
DF Tree

Visit order:



DFS, DF tree

DF Tree of **DFS(2)**



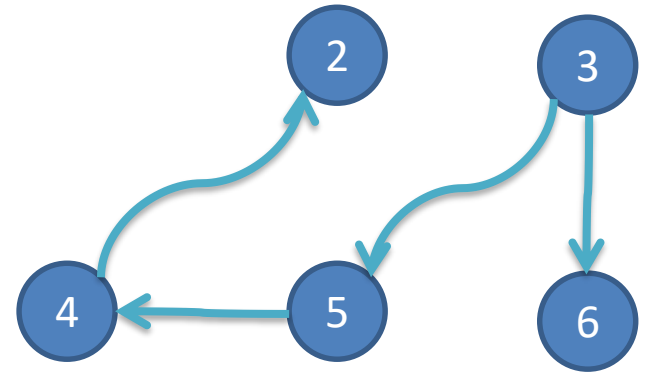
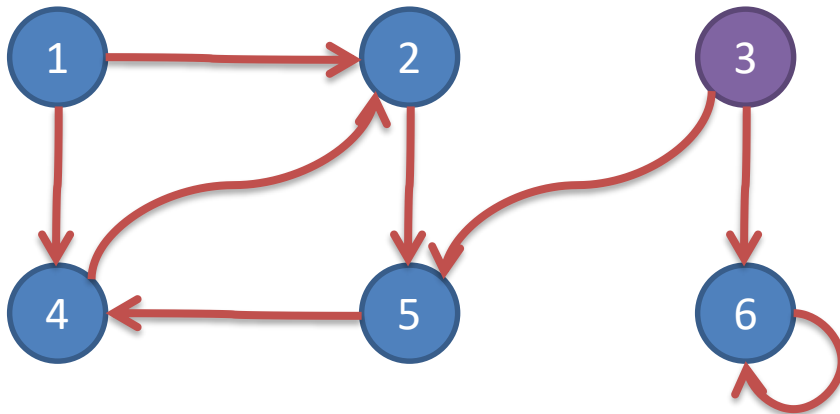
DF Tree

Visit order:



DFS, DF tree

DF Tree of **DFS(3)**



DF Tree

Visit order:



DFS, DF tree

- DFS Traversal // (The DFS in the textbook)
- **DFS_Traversal(G)**

for each v in G {

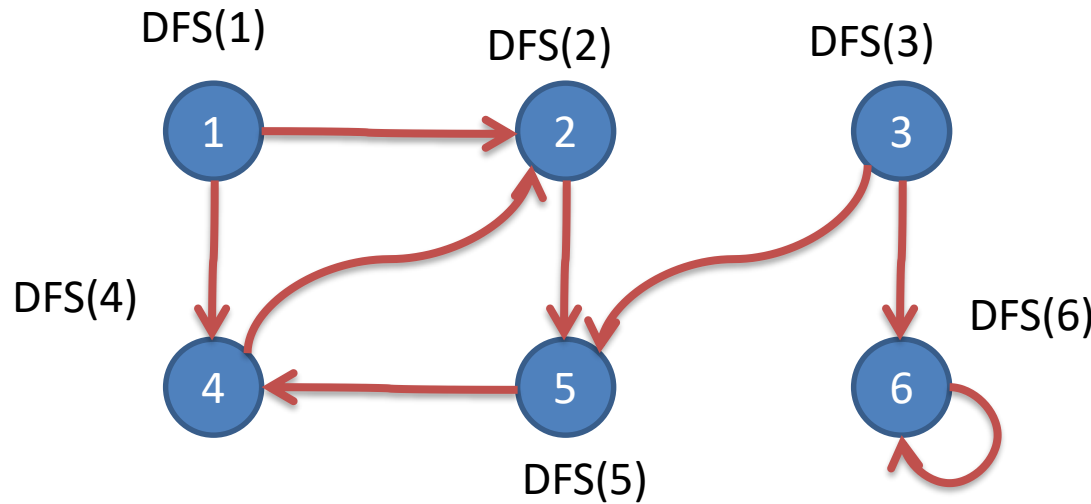
 if (v has not been visited)

DFS(v);

}

DFS, DF tree

DFS_Traversal(G)



DFS_Traversal(G)

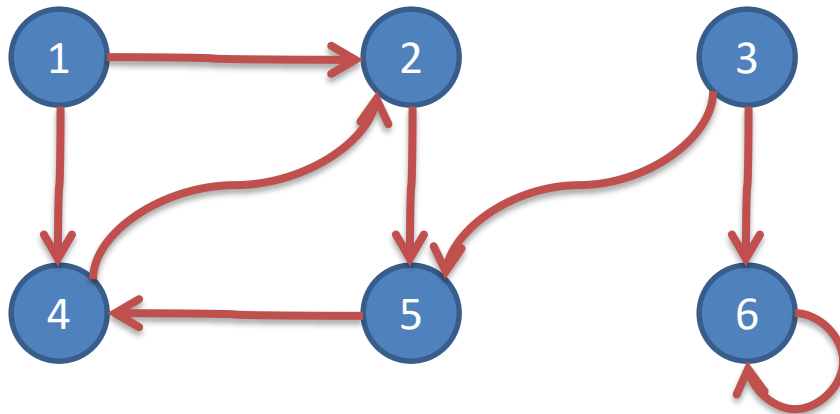
```
for each v in G{  
    if (v has not been visited)  
        DFS(v);  
}
```

Visit order:



Topological sort

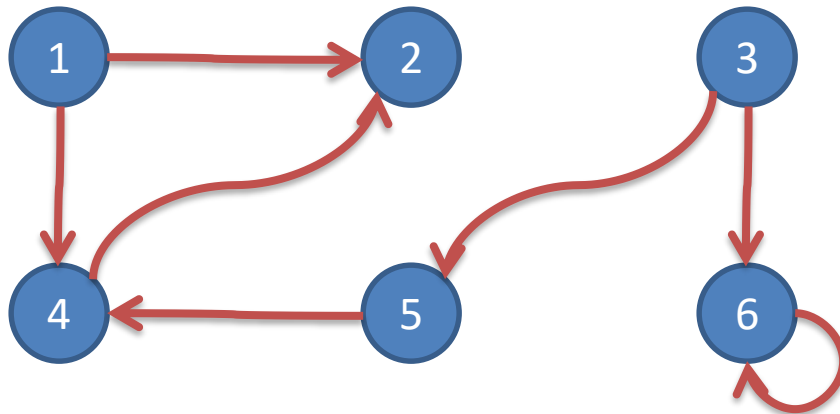
- DAG: directed acyclic graph
 - A graph without cycles



Dag? No

Topological sort

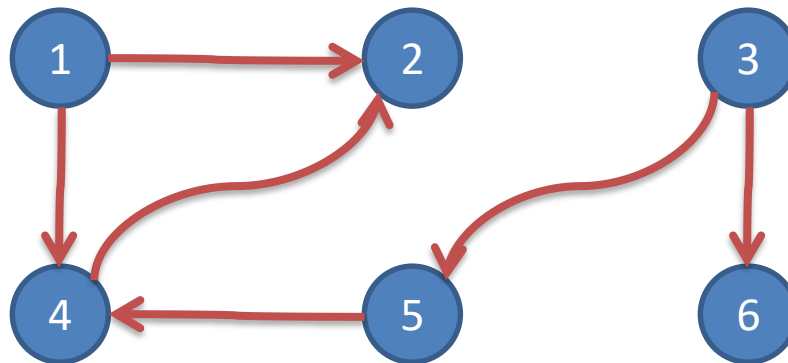
- DAG: directed acyclic graph
 - A graph without cycles



Dag? No

Topological sort

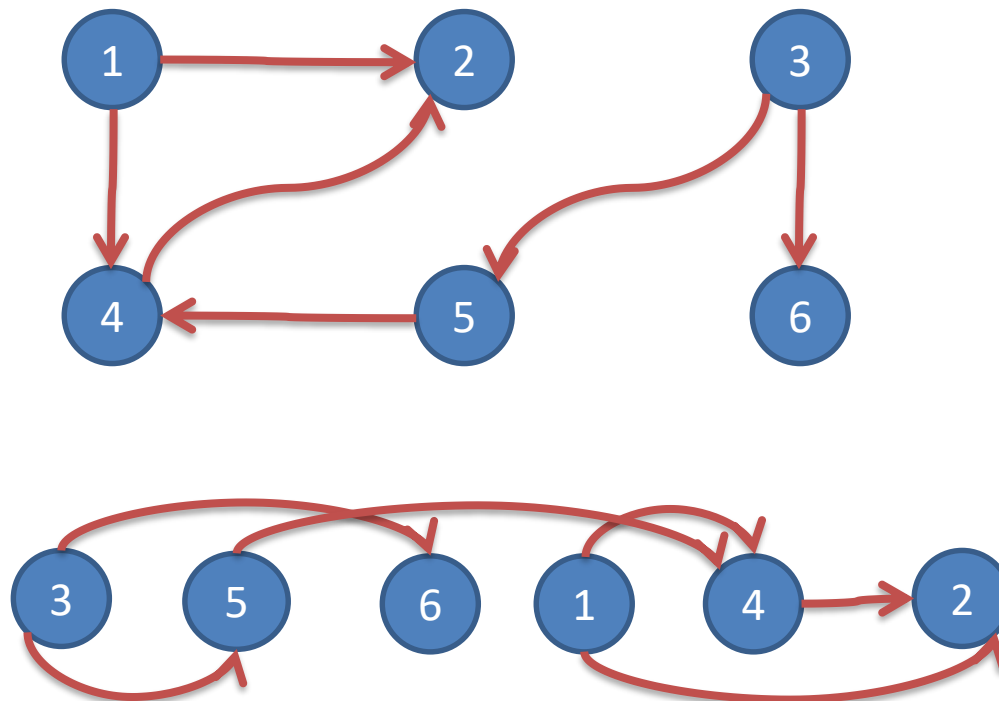
- Ordering in DAGs
 - If there is an edge $\langle u, v \rangle$, then u appears before v in the ordering



Dag? Yes

Topological sort

- Example

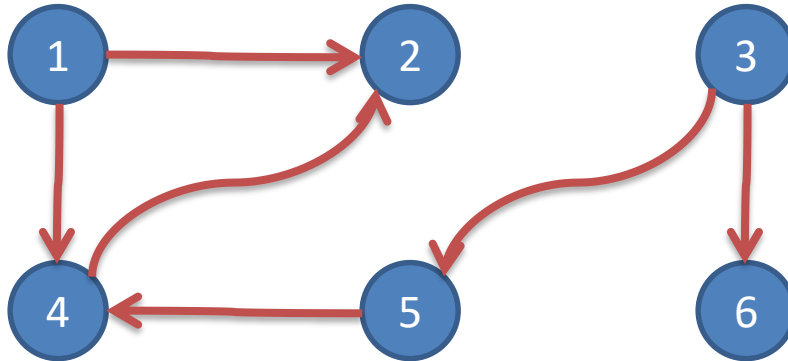


Put all the topological sorted vertices in a line, all edges go from left to right

Topological sort

- How to topological sort a dag?
- Just use **DFS_Traversal**
- The reverse order of DFS_Traversal is a topological sorted order

Topological sort



DFS_Traversal(G): 2 4 1 6 5 3

