

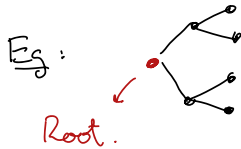
Lecture 10

Trees and searching:

Recall: Tree \Leftrightarrow cctd &

- ① All bridges
- ② $\exists!$ paths
- ③ No circuits
- ④ $V - E = 1$

def a root is a distinguished vertex:



\rightsquigarrow Outward flow assigns direction

$$d_{in} = 1, d_{out} = n - 1$$

$$d_{in} = 0 \Leftrightarrow v = r$$

leaf: $d_{out} = 0$, internal: $d_{out} \neq 0$.

binary for case $m = 2$.

def If $d_{out} \equiv m$ then m-ary tree

$$n = l + i$$

Prop for m-ary Tree, $n = mi + 1$, $i = |\text{internal vertices}|$.

Pf Each vertex is the child of unique parent so $n = mi + \underline{\text{root}}$.



Eg 56 people sign up for a bracket: How many matches?

Ans: $n = ?$, $m = 2$, $l = 56$, $i = ? \Rightarrow l + i = 2i + 1$

$$l = i + 1 \Rightarrow l - 1 = \boxed{i = 55.}$$

def height of tree: length of longest path from root.

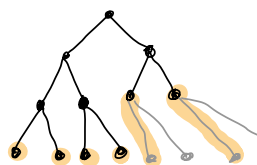
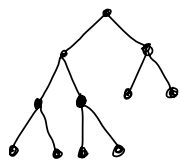
def balanced iff all leaves at height h or $h-1$. } accommodate m-ary.

Thm T an m-ary tree w/ height h and l leaves.

(a) $l \leq m^h$ (if all leaves at height $h \Rightarrow l = m^h$)

(b) $h \geq \log_m(l)$ If balanced then $h = \lceil \log_m(l) \rceil$

Pf: At each height there are $\leq m^h$ verts. and if the tree is complete then $l = m^h$. Otoh IF not then complete it and total leaves $\leq m^h$ by flowing down.



$$l \leq m^h. (a) \checkmark$$

Take log both sides, then balanced $\Rightarrow m^{h-1} \leq l \leq m^h \Rightarrow$
 $h-1 \leq \log_m l \leq h$
 $\Rightarrow h = \lceil \log_m l \rceil$.

Eg. How many binary tests to find a number? $\in \{1, \dots, N\}$.

Soln Searching gives binary tree w/ N verts (could stop at internal vertex!)

$$\Rightarrow \begin{matrix} 2i+1=n \\ i+l=n \end{matrix} \Rightarrow \frac{n-1}{2} = n-l \Rightarrow n-1 = 2n-2l$$

$$\Rightarrow 2l = n+1 \Rightarrow l = \frac{n+1}{2} \Rightarrow h = \lceil \log_2(n+1) - 1 \rceil = \lfloor \log_2(n+1) \rfloor$$

Qn How many trees are there on n labeled vertices?

(Meaning $i \rightarrow i$ does not give an isom)

Thm Cayley: n^{n-2} trees on n -labels.

Pf n -lab trees \longleftrightarrow sequences of length $n-1$.

Idea: Append to the sequence the parent of lowest labeled leaf then delete the leaf. Stop at 2 leaves.

Note length = #labels - 2.

Reconstruction: Note Leaves do not appear.

Smallest missing # is adj to 1st term.

Do (6, 2, 2, 3, 3, 3) $\leadsto n=8$
 $1, 2, 3, 4, 5, 6, 7, 8$

