

Today:

Ken

2.3 Valid & Invalid Arguments

Last time:

2.1 Logical Form & Equivalence

2.2 Conditional Statements

2.3 Valid & Invalid Arguments

⑥ C & D approach but only C speaks.

C says: Both of us are knaves.

What are C & D?

- subproof
- ① Suppose C is telling the truth. *assumption*
  - ② C & D are knaves. *what C said*
  - ③ C is a knave. *specialization*
  - ④ C is lying. *by def. of knave*
  - ⑤ C is telling the truth but lying. *conjunction*
  - ⑥ C is lying. *contradiction rule*
  - ⑦ C is a knave. *by def. of knave*
  - ⑧ Not both C & D are knaves. *negation of what C said*
  - ⑨ C is a knight or D is a knight. *DeMorgan's laws*
  - ⑩ Since C is a knave, D is a knight. *elimination*

$p \vee q$   
 $\neg p$

$\therefore q$

$\neg(p \vee \neg p)$

$\neg p$

$p \vee \neg p$

$\perp$

$\neg p$

$p \vee \neg p$

$\perp$

$\neg \neg(p \vee \neg p)$

$p \vee \neg p$

# Summary of Rules of Inference



Table 2.3.1 Valid Argument Forms

<p>Modus Ponens</p> $\begin{array}{l} p \rightarrow q \\ p \\ \therefore q \end{array}$	<p>Elimination</p> $\begin{array}{ll} \textcircled{a} \begin{array}{l} p \vee q \\ \sim q \\ \therefore p \end{array} & \textcircled{b} \begin{array}{l} p \vee q \\ \sim p \\ \therefore q \end{array} \end{array}$
<p>Modus Tollens</p> $\begin{array}{l} p \rightarrow q \\ \sim q \\ \therefore \sim p \end{array}$	<p>Transitivity</p> $\begin{array}{l} p \rightarrow q \\ q \rightarrow r \\ \therefore p \rightarrow r \end{array}$
<p>Generalization</p> $\begin{array}{ll} \textcircled{a} \begin{array}{l} p \\ \therefore p \vee q \end{array} & \textcircled{b} \begin{array}{l} q \\ \therefore p \vee q \end{array} \end{array}$	<p>Proof by Cases</p> $\begin{array}{l} p \vee q \\ p \rightarrow r \\ q \rightarrow r \\ \therefore r \end{array}$
<p>Specialization</p> $\begin{array}{ll} \textcircled{a} \begin{array}{l} p \wedge q \\ \therefore p \end{array} & \textcircled{b} \begin{array}{l} p \wedge q \\ \therefore q \end{array} \end{array}$	<p>Proof by Contradiction</p> $\begin{array}{l} \sim p \rightarrow \perp \\ \therefore p \end{array}$
<p>Conjunction</p> $\begin{array}{l} p \\ q \\ \therefore p \wedge q \end{array}$	

A set of premises and a conclusion are given. Use the valid argument forms listed in Table 2.3.1 to deduce the conclusion from the premises, giving a reason for each step as in Example 2.3.8. Assume all variables are statement variables.

- #41
- |     |                                      |     |
|-----|--------------------------------------|-----|
| (a) | $\neg p \vee q \rightarrow r$        | (3) |
| (b) | $s \vee \neg q$                      | (5) |
| (c) | $\neg t$                             | (2) |
| (d) | $p \rightarrow t$                    | (1) |
| (e) | $\neg p \wedge r \rightarrow \neg s$ | (4) |
| (f) | $\therefore \neg q$                  | (6) |

- |                            |      |                |
|----------------------------|------|----------------|
| $p \rightarrow t$          | (d)  |                |
| $\neg t$                   | (c)  |                |
| $\therefore \neg p$        | (i)  | modus tollens  |
| $\neg p$                   | (i)  |                |
| $\therefore \neg p \vee q$ | (ii) | generalization |

$$\neg p \vee q \rightarrow r$$

(a)

$$\neg p \vee q$$

(i)

$$\therefore r$$

(ii) modus ponens

$$\neg p$$

(i)

$$r$$

(ii)

$$\therefore \neg p \wedge r$$

(iv) conjunction

$$\neg p \wedge r \rightarrow \neg s$$

(e)

$$\neg p \wedge r$$

(iv)

$$\therefore \neg s$$

(v) modus ponens

$$s \vee \neg q$$

(b)

$$\neg s$$

(v)

$$\therefore \neg q$$

(vi) elimination

- #43
- |     |                                      |     |
|-----|--------------------------------------|-----|
| (a) | $\neg p \rightarrow r \wedge \neg s$ | (4) |
| (b) | $t \rightarrow s$                    | (5) |
| (c) | $u \rightarrow \neg p$               | (3) |
| (d) | $\neg w$                             | (2) |
| (e) | $u \vee w$                           | (1) |
| (f) | $\therefore \neg t$                  | (6) |

$u \vee w$	(e)	
$\neg w$	(d)	
$\therefore u$	(i)	elimination

$u \rightarrow \neg p$	(c)	
$u$	(i)	
$\therefore \neg p$	(ii)	modus ponens

$\neg p \rightarrow r \wedge \neg s$	(a)	
$\neg p$	(ii)	
$\therefore r \wedge \neg s$	(iii)	modus ponens

$r \wedge \neg s$	(iv)	
$\therefore \neg s$	(iv)	specialization

$t \rightarrow s$

$\neg s$

$\therefore \neg t$

⑥

④

⑤

modus tollens



## 3.1 Predicates & Quantifiers I

Some important sets:

$\emptyset$  or  $\{\}$  empty set

$\mathbb{N}$  natural numbers  $\{1, 2, 3, \dots\}$  or  $\{0, 1, 2, \dots\}$

$\mathbb{Z}$  integers  $\{0, 1, -1, 2, -2, 3, -3, \dots\}$

$\mathbb{Z}^-$  negative integers  $\{-1, -2, -3, \dots\}$

$\mathbb{Z}^+$  positive integers  $\{1, 2, 3, \dots\}$

$\mathbb{Z}^{\text{nonneg}}$  nonnegative integers  $\{0, 1, 2, \dots\}$

$\mathbb{Q}$  rational numbers  $\{m/n : m, n \in \mathbb{Z}; n \neq 0\}$

$\mathbb{Q}^-$  negative rational numbers  $\{r \in \mathbb{Q} : r < 0\}$

$\mathbb{Q}^+$  positive rational numbers  $\{r \in \mathbb{Q} : r > 0\}$

$\mathbb{Q}^{\text{nonneg}}$  nonnegative rational numbers  $\{r \in \mathbb{Q} : r \geq 0\}$

$\mathbb{R}$  real numbers

$\mathbb{R}^-$  negative real numbers

$\mathbb{R}^+$  positive real numbers

$\mathbb{R}^{\text{nonneg}}$  nonnegative real numbers

$\mathbb{C}$  complex numbers  $\{a+bi : a, b \in \mathbb{R}; i = \sqrt{-1}\}$

$\mathbb{H}$  quaternions

## some set theory and operations

Suppose  $A, B$  are sets.

$a \in A$

$a$  is an element of  $A$

$a \notin A$

$a$  is not an element of  $A$

$\{a_1, \dots, a_n\}$

set with elements  $a_1, \dots, a_n$

$A^c$

$A$  complement

$A \subset B$  or  $A \subseteq B$

$A$  is a subset of  $B$

$A \not\subset B$  or  $A \not\subseteq B$

$A$  is not a subset  $B$

$A = B$

$A$  equals  $B$

$A \cup B$

$A$  union  $B$

$A \cap B$

$A$  intersect  $B$

$A - B$

set difference  $A$  minus  $B$

$(x, y)$

ordered pair

$(x_1, x_2, \dots, x_n)$

ordered  $n$ -tuple

$A \times B$

Cartesian product of  $A$  &  $B$

$A_1 \times \dots \times A_n$

Cartesian product of sets  $A_1, \dots, A_n$

$\mathcal{P}(A)$

power set of  $A$

$$A = \{1, 2\}$$

$$2^2 = 4$$

$$\mathcal{P}(A) = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\}$$

$$A \times B = \{(a,b) : a \in A \text{ and } b \in B\}$$

$$p \rightarrow q \equiv \neg p \vee q$$

Theorem \*

(formerly "Awesome law" in Spring 2021)