

## Lecture 2 : Basics of Graphs Continued.

HW1 delayed Access issues  
O.H. Poll in CW  
1-2 BkL vs 3-4 Manh. ?

Review: def of IS / def of EC.

Prop  $I$  is indep  $\iff V \setminus I$  is an edge cover.

Proof:

Assume  $I$  indep but  $V \setminus I$  is not an EC.

$$\Rightarrow \exists e \in E \mid \partial e \cap V \setminus I = \emptyset$$

$$\Rightarrow \exists e \in E \mid \partial e \subseteq I \Rightarrow I \text{ NOT independent.}$$

Assume  $V \setminus I$  an edge cover &  $I$  NOT indep.

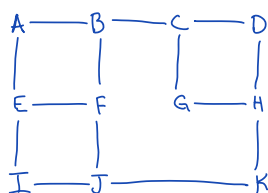
$$\Rightarrow \exists e \in E \mid \partial e \subseteq I \Rightarrow \exists e \in E \mid \partial e \not\subseteq V \setminus I$$

$$\Rightarrow V \setminus I \text{ not an edge cover.}$$

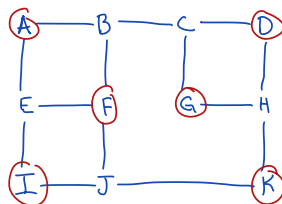


Cor The largest independent set corresponds to the smallest edge cover.

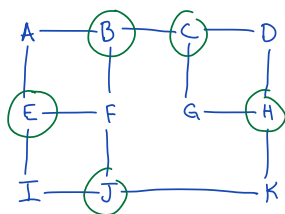
Eg. Find the smallest edge cover of:



Find largest IS.



Notice IS can be built from balls of size 2.



$\rightsquigarrow$  EC of size 5

def: A directed graph every edge has an arrow,  
two fixed vertices are connected by  
at most one arrow each way.

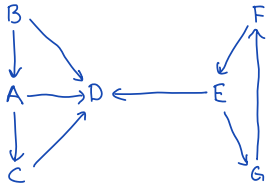
def: A (directed) path is a sequence of distinct vertices w/ each  
consecutive pair joined by an edge.  
↳ All edges flow in the direction of the sequence.

def: a circuit is a path whose final vertex is adj to its  
initial vertex.

def  $G$  is ctd if there's a path between any two vertices.

def vertex basis is  $B \subseteq V$  s.t. there's a directed path to  
any vertex from one in  $B$ .

Eg:



$\{B, F\}$  form a v.b. (Notation for a path?)  
↳ b/c source  $p = v_1 - v_2 - v_3 - \dots - v_k$   
↳ Then remove its image.  $p: v_1 \mapsto v_k$

Note loops act on the vertex bases.

End lecture 2