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Lecture I. Basics of Graph Theory.

Q1: Does a town always contain two people with the same number of friends?

Idea: Find comb objects which model the town & friendships, then translate and some the problem.

Obs: No if town contains one person. Q: if population >, 2.

del: A graph is a collection of vertices V and edges E s.t.

- Each citge connects two distinct vertices
- Each pair of vertices is cold by I or O edges.

Variations:

"Multi-graph" — Verties may be cefd by more than one edge.
"General graph" — Edges may cet a vertex to itself.

Examples:
$$K_1$$
 K_2 K_3 K_4 $V = \{1,2,3,4\}$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad$$

Multi-graph: General Graph: •-

Notation: Given $e \in E$, $\partial e := \{V_1, V_2\}$ the set of vertices cctd by e.

Information about graphs:

- * Number of Verts, Number of Edges. ∈ Z7.0
- * For $V \in V$, how many edges leave V = 0 deg (V)Make a function deg $: V \to \mathbb{Z}_{> 0}$

def
$$v, v' \in V$$
 are adjacent \iff there is an edge $e \in E$ s.t. $\partial e = \{v, v'\}$

Already enough to ans Q1!

Idea: turn town/friendships into a graph.

Eg. 3 people, 3 friendships:



Person 1 has 2 friends 2 edges leane P1.

friends = deg

Q1: (=> Are there 2 vertices w/ same degree?

(Assume n Vertices) Minimum = 0 edges, Maximum = n-1

"Isolated put" "Cold to all verts except itself." Thus $\forall v \in V$, deg $v \in \{0, ..., n-1\}$.

Obs: If deg vo = O then vo is adj to nothing.

If deg $V_1 = n-1$ then V_1 is adj to everything.

=> There cannot be a deg O and deg n-1 vertex together.

Thus Either:

deg
$$v \in \{1, ..., n-1\}$$
 OR deg $v \in \{0, ..., n-2\}$
Size $n-1$.

 $deg: \bigvee \longrightarrow Size n-1 \Rightarrow \exists v_i \neq v_j \mid deg(v_i) = deg(v_j)$ Sizen => Two people have some # of friends. Note: is a multi-graph w/ degrees (1,3,2) all different is a general graph w/ degs (1,3,2).

More examples:

Kn := Graph with n vertices and an edge cetting any point of nexts.

Kn,m == (V = Vn U Vm, Edges cct each vertex of Vn to Vm).

def: G is Bipartite (V splits as VI UVz with VI n Vz = Ø and VI, Vz \$\frac{1}{2} \partial e (\text{YeeE}).

def: A matching of a graph is a set MCE s.t.

 $\bigcirc \forall m_1, m_2 \in M \mid \partial m_1 \cap \partial m_2 = \emptyset.$

2 EVEDM | MEM } = V.

dy: An independent set is I⊆V s.t. ∀e∈E/ ∂e ≠ I.

"pw. non-adj vertices"

dy An Edge Couer 'N 6 € V s.t. Ye € E, de n 6 ≠ Ø.

" has at least one nertex of each edge!"

Prop: The smallest E.C. corresps to the largest IS.

End Lecture 1.