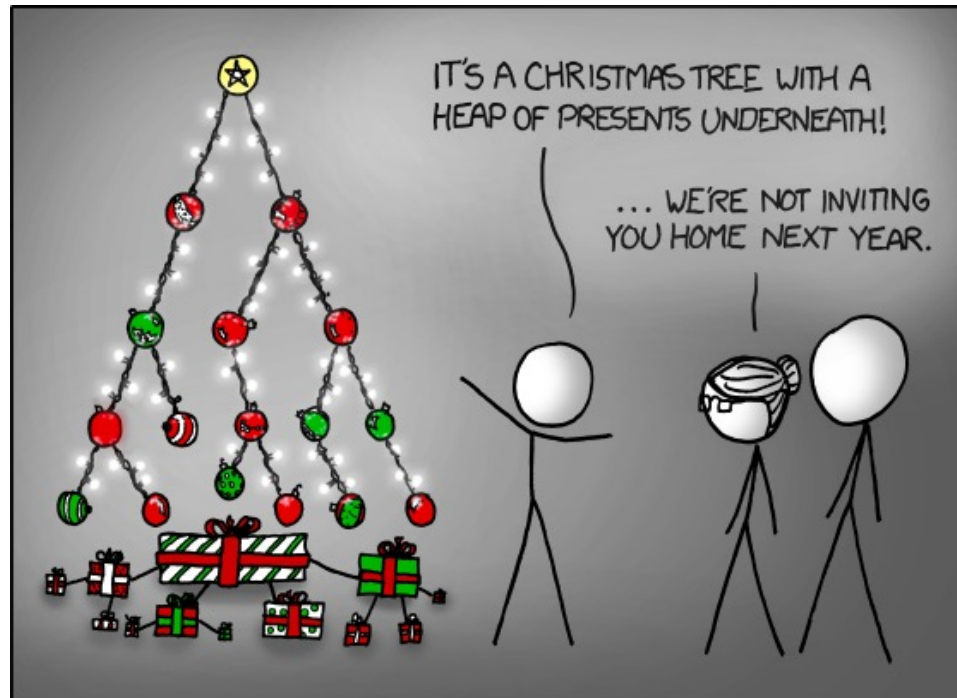


VE280 Programming and Elementary Data Structures

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Binary Search Tree and Set



Learning Objectives

- Understand what is a binary search tree
- Learn how to implement a binary search tree
- Learn how to implement a set using a binary search tree

Outline

- Binary Search Tree
- Implementing a set using a binary search tree

Binary Tree

- A binary tree
 - is empty, or
 - contains a root node, a left child, and a right child; both children are binary trees.
- Depth of a tree = height - 1

Implementation of Binary Tree

```
struct Node{  
    int val;  
    Node *left;  
    Node *right;  
}  
class BinaryTree{  
    Node *root;  
public:  
    ...  
}
```

Pre-order Tree Traversal

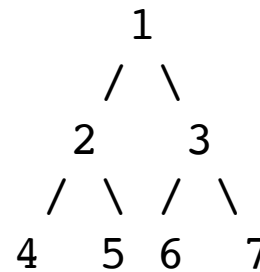
```
void BinaryTree::_preorder(Node *n) {  
    if (!n) {  
        // perform some operations on root  
        _preorder(root->left);  
        _preorder(root->right);  
    }  
}  
  
void BinaryTree::preorder() {  
    _preorder(root);  
}
```



Pre-order Tree Traversal

Assuming that the operation on the root node is printing its value, what is the output of pre-order tree traversal on this tree?

- **A.** 1234567.
- **B.** 1245367.
- **C.** 4251637.
- **D.** 4526731.



In-order Tree Traversal

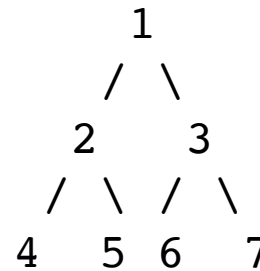
```
void BinaryTree::_inorder(Node *n) {  
    if (!n) {  
        _inorder(root->left);  
        // perform some operations on root  
        _inorder(root->right);  
    }  
}  
  
void BinaryTree::inorder() {  
    _inorder(root);  
}
```




In-order Tree Traversal

Assuming that the operation on the root node is printing its value, what is the output of in-order tree traversal on this tree?

- **A.** 1234567.
- **B.** 1245367.
- **C.** 4251637.
- **D.** 4526731.



Post-order Tree Traversal

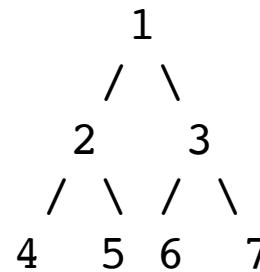
```
void BinaryTree::_postorder(Node *n) {  
    if (!n) {  
        _postorder(root->left);  
        _postorder(root->right);  
        // perform some operations on root  
    }  
}  
  
void BinaryTree::postorder() {  
    _postorder(root);  
}
```



Post-order Tree Traversal

Assuming that the operation on the root node is printing its value, what is the output of post-order tree traversal on this tree?

- A. 1234567.
- B. 1245367.
- C. 4251637.
- D. 4526731.



Binary Search Tree

- A binary search tree:
 - is a search tree,
 - if not empty, the key contained in its root node is larger than those contained in its left child and smaller than those contained in its right child,
 - Both left and right children are binary search trees.
- A binary search tree is balanced if the difference of the depths of its left and right subtrees is bounded by one.

Implementation of Binary Search Tree

```
struct Node{
    int key;
    Node *left;
    Node *right;
}
class BinarySearchTree{
    Node *root;
public:
    ...
    void add(int k);
    bool contains(int k);
    void delete(int k);
}
```

Implementation of add

```
Node *BinarySearchTree::_add(Node *n, int k);
    if (!n){
        Node *nn = new Node(k, NULL, NULL);
        return nn;
    }
    if (n->key > k){
        n->left = _add(n->left, k);
        return n;
    }
    if (n->key < k){
        n->right = _add(n->right, k);
        return n;
    }
}
void BinarySearchTree::add(int k);
    root = _add(root, k);
}
```

Implementation of contains

```
Node *BinarySearchTree::_contains(Node *n, int k);  
    if (!n){  
        return False;  
    }  
    if (n->key > k)  
        return _contains(n->left, k);  
    if (n->key < k)  
        return _contains(n->right, k);  
    if (n->key == k)  
        return True;  
}  
bool BinarySearchTree::contains(int k);  
    return _contains(root, k);  
}
```

Implementation of delete

```
Node *BinarySearchTree::_delete(Node *n,
int k){
    if (!n) return NULL;
    if (n->key > k){
        n->left = _delete(n->left, k);
        return n;
    }
    if (n->key < k){
        n->right = _delete(n->right, k);
        return n;
    }
    if (n->key == k){
        if (!n->left && !n->right)
            return NULL;
        if (!n->left){
            Node *right = n->right;
            delete n;
            return right;
        }
        if (!n->right){
            Node *left = n->left;
            delete n;
            return left;
        }
        Node *left = n->left;
        Node *right = n->right;
        delete n;
        return minKey(left, right);
    }
}
```

```
    if (!n->right){
        Node *left = n->left;
        delete n; return left;
    }
    n->key = minKey(n->right);
    n->right = _delete(n->right,
        n->key);
}

Node *BinarySearchTree::delete(int k){
    if (!root) return;
    root = _delete(root, k);
}

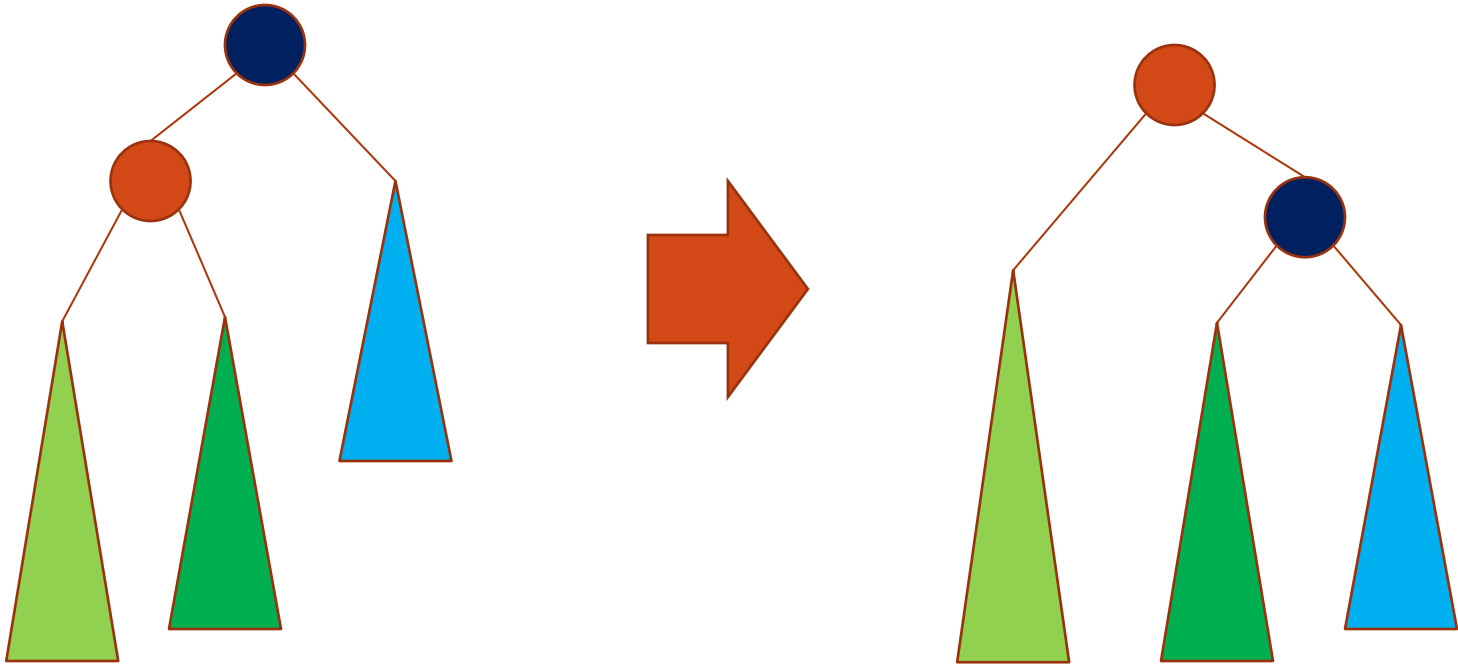
Node *BinarySearchTree::minKey(Node
*n){
    if (!n) throw exception;
    if (!n->left) return n->key;
    return minKey(n->left);
}
```


Complexity of add, contains, delete

- The number of nodes visited is in $O(\text{depth of tree})$
- If tree is balanced, depth is in $O(\log(n))$ where $n = \#$ values
- Therefore, for balanced binary search tree, the computational complexity is in $O(\log(n))$

Balance a binary search tree

- Possibly needed after a call of add or delete
- Example:



Implementation of Set

```
class Set{  
    BinarySearchTree bst;  
public:  
    ...  
    void add(int k);  
    bool contains(int k);  
    void delete(int k);  
}
```