

Problem 1.

Please check "Problem 1. ipynb"

Problem 2:

$$\lambda = 50 \text{ ^{th}/month} \quad K_- = 50 \quad i = C_j = \frac{200}{12}$$

$$Q_j^* = \sqrt{\frac{2k\lambda}{h}} = \sqrt{\frac{2 \times 50 \times 50}{200/12}} = 17.3.$$

Only Q_1^* is reasonable, and the cost is:

$$g_1(17) = 50 \times 50 + \frac{50 \times 50}{17} + \frac{\frac{200}{12} \times 17}{2} = 25788.7$$

$$g_2(65) = 495 \times 50 + \frac{50 \times 50}{65} + \frac{\frac{200}{12} \times 65}{2} = 25330.1$$

$$g_3(129) = 485 \times 50 + \frac{50 \times 50}{129} + \frac{\frac{200}{12} \times 129}{2} = 25344.4$$

Therefore, the optimal order quantity is $Q = 65$.

the cost per month is: \$ 25330.1

$$(b) \quad g(65) = 520 \times 50 + \frac{0.50}{65} + \frac{0.65}{2} = 26000 > 25330.1$$

\Rightarrow Zens should not accept the offer.

Problem 3

(a) $K = 1000, h = 1.2$

$$Q_8 = 0$$

$$Q_7 = K + h(0 \cdot d_7) + \theta_8 = 1000, \quad [S(7) = 8]$$

$$Q_6 = \min \{ K + h(0 \cdot d_6) + \theta_7, K + h(0 \cdot d_6 + 1 \cdot d_7) + \theta_8 \}$$

$$= \min \{ 2000, 1348 \}$$

$$= 1348 \quad [S(6) = 8]$$

$$Q_5 = \min \{ K + h(0 \cdot d_5) + \theta_6, K + h(0 \cdot d_5 + 1 \cdot d_6) + \theta_7, K + h(0 \cdot d_5 + 1 \cdot d_6 + 2 \cdot d_7) + \theta_8 \}$$

$$= \min \{ 2348, 2252, 1948 \}$$

$$= 1948 \quad [S(5) = 8]$$

$$Q_4 = \min \{ K + h(0 \cdot d_4) + \theta_5, K + h(0 \cdot d_4 + 1 \cdot d_5) + \theta_6, K + h(0 \cdot d_4 + 1 \cdot d_5 + 2 \cdot d_6) + \theta_7, \\ K + h(0 \cdot d_4 + 1 \cdot d_5 + 2 \cdot d_6 + 3 \cdot d_7) + \theta_8 \}$$

$$= \min \{ 2948, 2552, 2708, 2752 \}$$

$$= 2552 \quad [S(4) = 6]$$

$$Q_3 = \min \{ K + h(0 \cdot d_3) + \theta_4, K + h(0 \cdot d_3 + 1 \cdot d_4) + \theta_5, K + h(0 \cdot d_3 + 1 \cdot d_4 + 2 \cdot d_5) + \theta_6, \\ K + h(0 \cdot d_3 + 1 \cdot d_4 + 2 \cdot d_5 + 3 \cdot d_6) + \theta_7, K + h(0 \cdot d_3 + 1 \cdot d_4 + 2 \cdot d_5 + 3 \cdot d_6 + 4 \cdot d_7) + \theta_8 \}$$

$$= \min \{ 3552, 3056, 2864, 3272, 3664 \}$$

$$= 2864 \quad [S(3) = 6]$$

$$Q_2 = \min \{ K + h(0 \cdot d_2) + \theta_3, K + h(0 \cdot d_2 + 1 \cdot d_3) + \theta_4, K + h(0 \cdot d_2 + 1 \cdot d_3 + 2 \cdot d_4) + \theta_5, \\ K + h(0 \cdot d_2 + 1 \cdot d_3 + 2 \cdot d_4 + 3 \cdot d_5) + \theta_6, K + h(0 \cdot d_2 + 1 \cdot d_3 + 2 \cdot d_4 + 3 \cdot d_5 + 4 \cdot d_6) + \theta_7, \\ K + h(0 \cdot d_2 + 1 \cdot d_3 + 2 \cdot d_4 + 3 \cdot d_5 + 4 \cdot d_6 + 5 \cdot d_7) + \theta_8 \}$$

$$= \min \{ 3864, 3678, 3290, 3302, 3962, 4702 \}$$

$$= 3290 \quad [S(2) = 5]$$

$$Q_1 = \min \left\{ K + h(0 \cdot d_1) + \theta_2, K + h(0 \cdot d_1 + 1 \cdot d_2) + \theta_3, K + h(0 \cdot d_1 + 1 \cdot d_2 + 2 \cdot d_3) + \theta_4, \right. \\ K + h(0 \cdot d_1 + 1 \cdot d_2 + 2 \cdot d_3 + 3 \cdot d_4) + \theta_5, K + h(0 \cdot d_1 + 1 \cdot d_2 + 2 \cdot d_3 + 3 \cdot d_4 + 4 \cdot d_5) + \theta_6, \\ K + h(0 \cdot d_1 + 1 \cdot d_2 + 2 \cdot d_3 + 3 \cdot d_4 + 4 \cdot d_5 + 5 \cdot d_6) + \theta_7, \\ \left. K + h(0 \cdot d_1 + 1 \cdot d_2 + 2 \cdot d_3 + 3 \cdot d_4 + 4 \cdot d_5 + 5 \cdot d_6 + 6 \cdot d_7) + \theta_8 \right\}$$

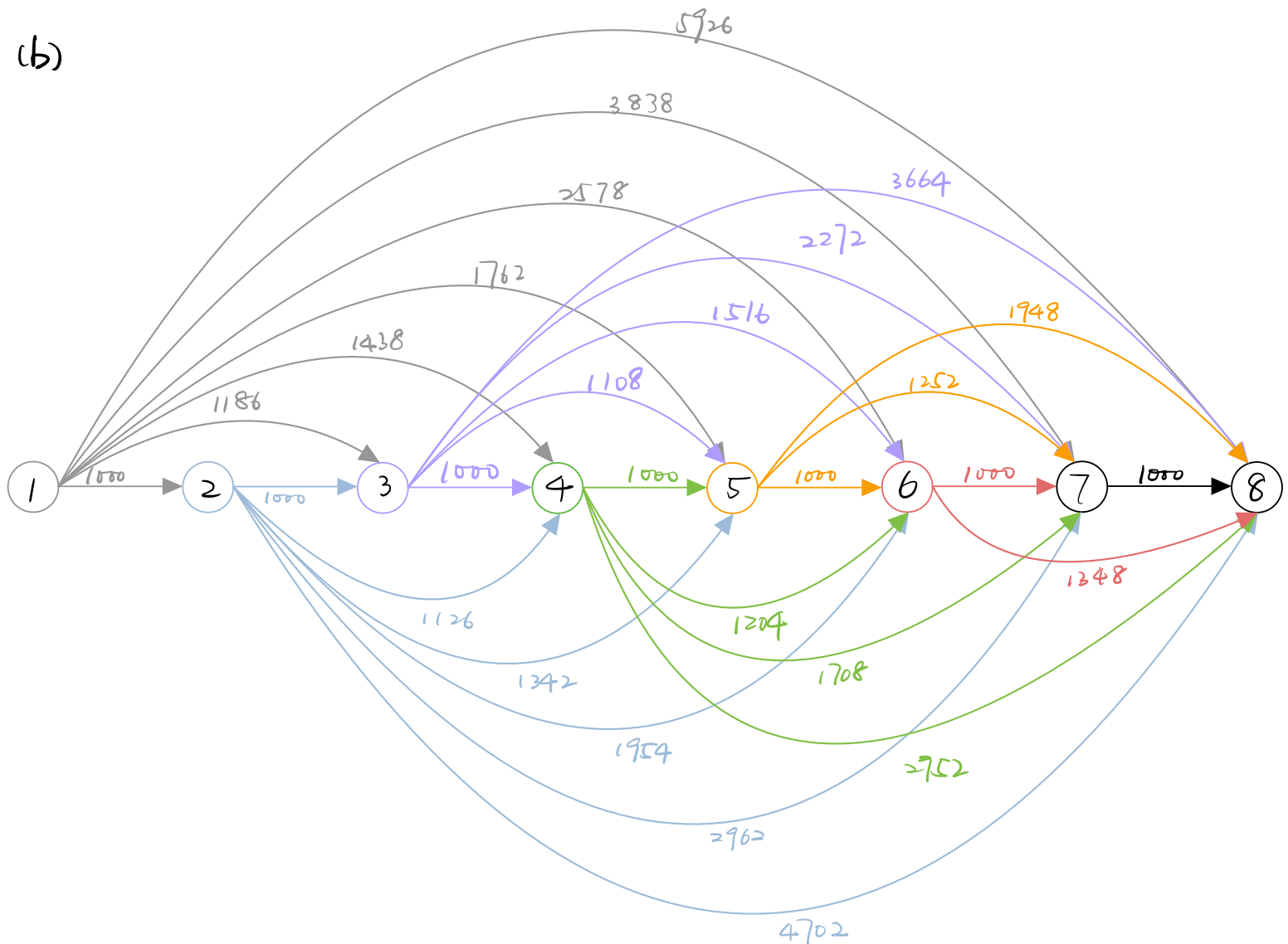
$$= \min \{ 4290, 4030, 3990, 3710, 3726, 4838, 5726 \}$$

$$= 3710 \quad [S(1) = 5]$$

Therefore, we should order 570 on Day[#] 1,
and order 670 on Day[#] 5.

The total cost is \$3710.

(b)



path $\emptyset \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow 5 \rightarrow 7 \rightarrow 5 \rightarrow 8$

cost 0 1000 1186 1438 1762 2578 3014 3710

Therefore, we should order 570 on Day[#] 1,
and order 670 on Day[#] 5.

The total cost is \$3710.

(c) Denote q_t : number of units ordered in period t .
 y_t : 1 for order in period t , 0 otherwise
 x_t : the inventory level at the end of period t .

$$\min \sum_{t=1}^T (1000 y_t + 1.2 x_t)$$

$$\text{s.t. } x_t = x_{t-1} + q_t - d_t$$

$$q_t \leq 10000 y_t$$

$$x_t \geq 0$$

$$\text{for } \forall t \in \{1, \dots, T\}$$

$$q_t \geq 10$$

$$y_t \in \{0, 1\}$$

The code is in 'Problem 3c.ipynb'

Problem 4

$$(a) \quad A = \begin{pmatrix} 3 & 2 & 1 \\ 1 & 2 & 4 \end{pmatrix} \quad x = \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} \quad c = \begin{pmatrix} 6 \\ 14 \\ 13 \end{pmatrix}$$

$$A^T = \begin{pmatrix} 3 & 1 \\ 2 & 2 \\ 1 & 4 \end{pmatrix} \quad b^T = (24 \quad 60)$$

$$\min (24y_1 + 60y_2)$$

$$\begin{aligned} \text{s.t.} \quad & 3y_1 + y_2 \geq 6 \\ & 2y_1 + 2y_2 \leq 14 \\ & y_1 + 4y_2 = 13 \\ & y_1 \geq 0 \\ & y_2 \leq 0 \end{aligned}$$

$$(b) \quad c = \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix} \quad c.size() = |I| \times |J| \quad b = \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix} \quad b.size() = |I| + |J|$$

$$A = \begin{pmatrix} 1 & 1 & 1 & \dots & 0 & 0 & \dots & 0 \\ 0 & 0 & 0 & \dots & 1 & 1 & \dots & 1 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 0 & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 0 & 0 & \dots & 1 \end{pmatrix}$$

$$A^T = \begin{pmatrix} 1 & 0 & \dots & 0 & 1 & 0 & \dots & 0 \\ 1 & 0 & \dots & 0 & 0 & 1 & \dots & 0 \\ 1 & 0 & \dots & 0 & 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 1 & \dots & 0 & 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & \dots & 0 & \dots & 1 & 0 & \dots & 0 & 1 \end{pmatrix}$$

For A^T , we can find that there're two 1 in each row. one in $|I|$, one in $|J|$.

The dual is:

$$\begin{aligned} \min \quad & \sum_{i \in I} u_i + \sum_{j \in J} v_j \\ \text{s.t.} \quad & u_i + v_j \geq 1 \quad \text{for } \forall i \in I, \forall j \in J. \\ & u_i \geq 0 \quad \text{for } \forall i \in I \\ & v_j \geq 0 \quad \text{for } \forall j \in J. \end{aligned}$$