

Introduction:-

A singly linked list can be effectively used to implement a queue data structure. This approach leverages the dynamic memory allocation capabilities of linked lists, which can be advantageous over array-based queues that have a fixed size.

Key Concepts:-

→ FIFO (First-In, First-Out):

Like a traditional queue, elements are added (enqueued) at the rear and removed (dequeued) from the front.

→ Nodes:

- Each element in the queue is represented by a node in the linked list.

- Each node contains:

- Data: The value of the element.

- Next: A pointer to the next node in the queue.

→ Front and Rear pointers:

- We maintain two pointers:

- Front: Points to the first node in the queue (where dequeuing occurs).

- Rear: Points to the last node in the queue (where enqueueing occurs).

Operations:

→ Enqueue (Insertion):

- A new node is created and added to the rear of the list.

- The rear pointer is updated to point to the new node.

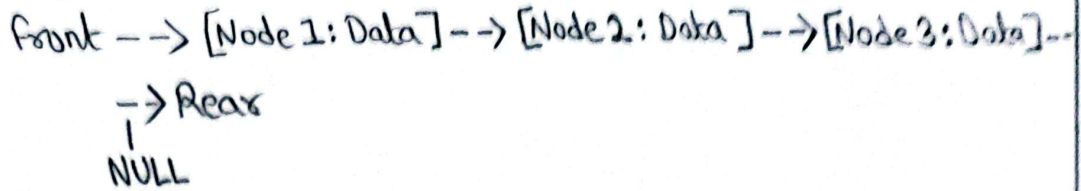
- If the queue is empty, both front and rear pointers are set to the new node.

→ Dequeue (Deletion):

- The node at the front of the list is removed.

- The front pointer is updated to point to the next node.

- If the queue becomes empty, both front and rear pointers are set to NULL.

Diagram:Explanation of the Diagram:

- "Front" indicates the beginning of the queue and "Rear" indicates the end.
- Each "[Node: Data]" represents a node in the linked list, storing a data element.
- The arrows represent the "next" pointers, connecting the nodes.
- The rear node next pointers points to NULL, indicating the end of the list.

Algorithm for Singly linked List Queue Implementation:

1) Structure Definition:

\rightarrow Define a structure 'Node' with two fields:

- 'info': An integer to store the data.
- 'next': A pointer to the next 'Node'.

\rightarrow Initialize two global pointers:

- 'front': Points to the front of the queue (initialized to 'NULL').
- 'rear': Points to the rear of the queue (initialized to 'NULL').

2) 'main()' Function:

\rightarrow Enter an infinite loop ('while(1)').

\rightarrow Display a menu with the following options:

- 1) Enqueue operation
- 2) Dequeue operation
- 3) Display
- 4) Exit

\rightarrow Read the user's choice ('ch')

\rightarrow Use a 'switch' statement to execute the corresponding operation:

• Case 1 (Enqueue):

- Prompt the user to enter data ('num').
- Call the 'enqueue(num)' function.

- break;
 - Case 2 (Dequeue):
 - Call the 'dequeue()' function.
 - break;
 - Case 3 (Display):
 - Call the 'display()' function.
 - break;
 - Case 4 (Exit):
 - Exit the program ('exit(0)').
 - break;
 - Default:
 - (Optional) Print an error message for invalid input.
- End of 'switch' statement.
- End of 'while' loop.
- (Optional) Wait for a key press ('getch()').
- Return 0.

- 3) 'enqueue(num)' Function:
- Allocate memory for a new 'Node' ('newnode').
 - Store the input data ('num') in 'newnode->info'.
 - If 'front' is 'NULL' (queue is empty):
 - Set both 'front' and 'rear' to 'newnode'.
 - Set 'front->next' and 'rear->next' to 'NULL'.
 - Else (queue is not empty):
 - Set 'rear->next' to 'newnode'.
 - Set 'rear' to 'newnode'.
 - Set 'rear->next' to 'NULL'.
 - Print a message indicating that 'num' was enqueued.

- 4) 'dequeue()' Function:
- If 'front' is 'NULL' (queue is empty):
 - Print "Queue is empty - - -".
 - Else (queue is not empty):
 - Store the current 'front' node in a temporary pointer 'temp'.
 - Update 'front' to point to the next node ('front = front->next').
 - Print the dequeued element ('temp->info').
 - Free the memory allocated for the dequeued node ('free(temp)').

5) 'display()' Function:

→ If 'front' is 'NULL' (queue is empty):

- print "Queue is empty".

→ Else (queue is not empty):

- Initialize a temporary pointer 'temp' to 'front'.

- Print "Contents are":

→ While 'temp' is not equal to 'rear':

- Print 'temp->info' followed by " - ->".

- Move 'temp' to the next node ('temp = temp->next').

→ Print the 'temp->info' (which is the last element of the queue).