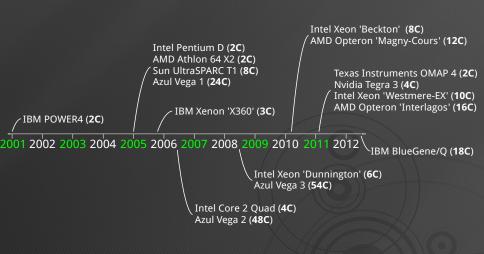
A Software-based Method-Level Speculation Framework for the Java Platform

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2001 2002 2003 2004 2005 2006 2007 2008 2009 2010 2011 2012 Many applications still single or lightly-threaded Not feasible to rewrite many existing applications to work in parallel

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→ Automatic Parallelization

Thread-Level Speculation

Employs speculation

• No need to prove that parallelization is valid

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- Correctness dynamically ensured at runtime
- → Uses some kind of in-memory transactions to buffer and validate tentative changes

Our proposal: JaSPEx-MLS

• Software-based speculative parallelization framework for Java



- Software-based speculative parallelization framework for Java
- Java Bytecode rewriting



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...on top of the OpenJDK Java VM

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Target applications

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- Existing single-threaded JVM bytecode
- Object-oriented / method-heavy
 - Limited support for loops
- Static analysis is hard or impossible

• Modern open-source VM by Oracle and community contributors



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- → JaSPEx-MLS works on commonly available hardware
- → Same codebase used to run regular Java applications

Method-Level Speculation

```
int example1() {
   int x = computeValue();
   int y = 0;
   for (...) y += ...;
   return x+y;
}
```

• Run method call in parallel with code following its return

Method-Level Speculation

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int example1() {
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for loop

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Classloader

• Prepares code for speculative parallelization

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Runtime Orchestration Library

- Coordinates all runtime events
 - o Creation, validation, commit, abort, ...

JaSPEx-MLS Classloader

Transactification

• Intercept access to slots and arrays



JaSPEx-MLS Classloader

Transactification

Intercept access to slots and arrays

```
object.field = 100l;
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Transaction.storeLong(object, field, 1001);

Transactification

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object.field = 100l;
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- Static and type-specific
 - → Designed to be inlined by the VM



Handling Non-Transactional Operations

→ Operations outside the control of the transactional system



- → Operations outside the control of the transactional system
 - native methods

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```
System.out.println("Hello World!");
...
```

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```
SpeculationControl.nonTransactionalActionAttempted();
System.out.println("Hello World!");
...
```

Modifications for MLS

• Replace method calls with call to spawnSpeculation()



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- Replace method return value with Future



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int mlsExample1() {
    return compute(0) + compute(1);
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int mlsExample1() {
    Future temp1 = spawnSpeculation(...); // compute(0)
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```

Usually needs control flow analysis and code duplication

STM support for Futures

- Allow Futures to be "written" to memory locations
- Delay obtaining value until it is needed (read or commit)

```
STM support for Futures
// Original Method
void doCompute(Object[] results) {
  for (int i = 0; i < results.length; i++) {
    results[i] = compute(i);
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// *Attempted* parallelization
void doCompute(Object[] results) {
  for (int i = 0; i < results.length; i++) {
    Future f0 = spawnSpeculation(...); // compute(i)
    TM.storeObjectArray(results, i, f0.get());
  }}
```

```
... let's try this again
// Original Method
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// *Successful* parallelization
void doCompute(Object[] results) {
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    TM.storeFutureObjectArray(results, i, f0); // f0 handed
  }}
                                                 // to the STM
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What happens inside spawnSpeculation?

Check threadpool state

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- Capture first-class continuation

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 - Clean current stask
 - Execute compute(0)
 - Write result to child task
 - Return to poll
- Child task:
 - Start transaction
 - Resume continuation
 - Return from spawnSpeculation

Committing a speculation:

• Task completed work



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- Need to obtain value from a Future, and oldest-running task



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- Tried to execute non-transactional operations
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- If exception → abort
- Else try to commit

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• Program-order mode: Reads and writes directly to memory



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Transaction commit is simpler, because only one transaction may commit at a time

→ No synchronization needed

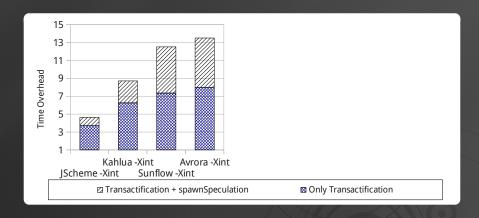


Benchmarks

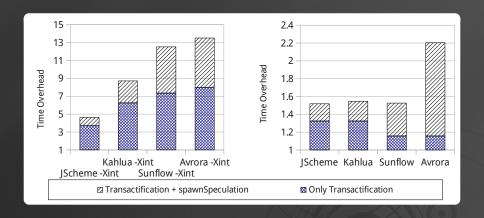
Benchmarks

- Test overheads from bytecode changes
- Early results from full system
- Test system: Intel Core i5 750, 8GB RAM, Ubuntu Linux 12.04 64-bit

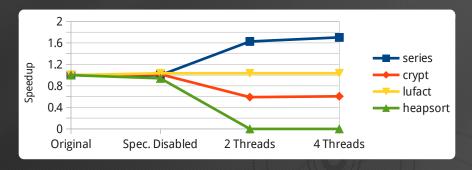
Benchmarks — Overhead with Interpreter



Benchmarks — Overhead with JIT



Benchmarks — Parallelization



- Benchmarks from Java Grande Framework
- Source was **not** modified

JaSPEx-MLS: Automatic parallelization framework for Java



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- Optimizing VM makes *all* the difference
 - o Able to hide much of the added overhead
- Already able to extract parallelism in some benchmarks

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- Re-use waiting threads by returning them to the thread-pool

Thank you!